

Tutorial-4

① $T(n) = 3T(n/2) + n^2$

$a = 3, b = 2$ $f(n) = n^2$

$$n \log_b^a = n \log_2^3$$

comparing $n \log_2^3$ and n^2

$$n \log_2^3 < n^2 \text{ (case 3)}$$

\therefore acc. to master's Theorem

$$T(n) = O(n^2)$$

② $T(n) = 4T(n/2) + n^2$

$a = 4, b = 2$

$$n \log_b^a = n \log_2^4 = n^2 = f(n) \text{ (case 2)}$$

\therefore acc. to master's theorem $T(n) = O(n^2 \log n)$

③ $T(n) = T(n/2) + 2^n$

$a = 1, b = 2$

$$n \log_2^1 = n^0 = 1$$

$$1 < 2^n \text{ (case 3)}$$

\therefore Acc. to master's theorem $T(n) = O(2^n)$

④ $T(n) = 2^n T(n/2) + n^n$

\therefore Master's theorem is not applicable as a w.f.ⁿ of n

$$\textcircled{5} \quad T(n) = 16T(n/4) + n$$

$$a=16, b=4, f(n)=n$$

$$n \log_b a = n \log_4 16 = n^2$$

$$n^2 > f(n) \text{ (case 1)}$$

$$T(n) = O(n^2)$$

$$\textcircled{6} \quad T(n) = 2T(n/2) + n \log n$$

$$a=2, b=2, f(n)=n \log n$$

$$n \log_b a = n \log_2 2 = n$$

$$\text{Now } f(n) > n$$

$$\therefore \text{Acc. to master's } T(n) = O(n \log n)$$

$$\textcircled{7} \quad T(n) = 2T(n/2) + \frac{n}{\log n}$$

$$a=2, b=2, f(n) = \frac{n}{\log n}$$

$$n \log_b a = n \log_2 2 = n$$

$$n > f(n)$$

$$\therefore \text{Acc. to master's theorem } T(n) = O(n)$$

$$\textcircled{8} \quad T(n) = 2T(n/4) + n^{0.51}$$

$$a=2, b=4, f(n)=n^{0.51}$$

$$n \log_b a = n \log_4 2 = n^{0.5}$$

$$n^{0.5} < f(n)$$

$$\therefore \text{Acc. to master's theorem } T(n) = O(n^{0.51})$$

$$(9) \quad T(n) = 0.5T(n/2) + \frac{1}{n}$$

\therefore Master's Not applicable as $a < 1$

$$(10) \quad T(n) = 16T(n/4) + n!$$

$$a = 16, b = 4, f(n) = n!$$

$$n^{\log_b a} = n^{\log_4 16} = n^2 \quad n^2 < n!$$

\therefore Acc. to master's, $T(n) = O(n!)$

$$(11) \quad T(n) = 4T(n/2) + \log n$$

$$a = 4, b = 2, f(n) = \log n$$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$

$$n^2 > f(n)$$

\therefore Acc. to master's, $T(n) = O(n^2)$

$$(12) \quad T(n) = \sqrt{n} T(n/2) + \log n$$

\therefore Master's not applicable as a is not const.

$$(13) \quad T(n) = 3T(n/2) + n$$

$$a = 3, b = 2, f(n) = \sqrt{n}$$

$$n^{\log_b a} = n^{\log_2 3} = n$$

$$n > \sqrt{n}$$

\therefore Acc. to master theorem, $T(n) = O(n)$

(14)

$$T(n) = 3T(n/3) + \sqrt{n}$$

$$a = 3, b = 3, f(n) = \sqrt{n}$$

$$n^{\log_b a} = n^{\log_3 3} = n$$

$$n > \sqrt{n}$$

\therefore Acc. to master theorem, $T(n) = \Theta(n)$

(15)

$$T(n) = 4T(n/2) + Cn$$

$$a = 4, b = 2, f(n) = C * n$$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$

$$n^2 > C * n$$

\therefore Acc. to master's theorem, $T(n) = \Theta(n^2)$

(16)

$$T(n) = 3T(n/4) + n \log n$$

$$a = 3, b = 4, f(n) = n \log n$$

$$n^{\log_b a} = n^{\log_4 3} = n^{0.79}$$

$$n^{0.79} < n \log n$$

$$\Theta(n) = \Theta\left(\frac{n}{2}\right)$$

\therefore Acc. to master's theorem

$$T(n) = \Theta(n \log n)$$

$$(17) \quad T(n) = 3T(n/3) + n/2$$

$$a=3, b=3, f(n) = n/2$$

$$n^{\log_b a} = n^{\log_3 3} = n$$

$$\Theta(n) = \Theta(n/2)$$

\therefore Acc. to master's theorem $T(n) = \Theta(n^2 \log n)$

$$(18) \quad T(n) = 6T(n/3) + n^2 \log n$$

$$a=6, b=3; f(n) = n^2 \log n$$

$$n^{\log_b a} = n^{\log_3 6} = n^{1.63}$$

$$n^{1.63} < n^2 \log n$$

\therefore Acc. to master's theorem $T(n) = \Theta(n^2 \log n)$

$$(19) \quad T(n) = 4T(n/2) + n/\log n$$

$$a=4, b=2, f(n) = n/\log n$$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$

$$n^2 > n/\log n$$

\therefore ~~Acc.~~ Acc. to master's theorem $T(n) = \Theta(n^3)$

$$(20) \quad T(n) = 64T(n/8) - n^2 \log n$$

Master's theorem is not applicable as $f(n)$ is not increasing function.

Q1

$$T(n) = \cancel{64} 7T(n/3) + n^2$$

$$a=7, b=3, f(n)=n^2$$

$$n^{\log_b a} = n^{\log_3 7} = n^{1.7}$$

$$n^{1.7} < n^2$$

\therefore Acc. to master's, $T(n) = O(n^2)$

Q2

$$T(n) = T(n/2) + n(2 - \cos n)$$

Master's theorem isn't applicable since regularity condition is violated in case 3.