

# CS201A:Mathematics for Computer Science - I

## Assignment 2

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**Question 1: [9 points] Write a pseudocode to find the next natural number in base b.**

We assume that the given number as well as its next natural number can be stored in  $k$  digits in  $b$  base notation.

Since  $b$  can be greater than 10, we might not be able to write it as a decimal number, therefore, we use an array 'd' to input the  $k$  digits of  $n$ , where  $d[0]$  will contain the least significant digit of  $n$  in base  $b$  and  $d[k-1]$  store the most significant digit of  $n$ .

Following is the pseudocode of my algorithm to find the next natural number in base  $b$ .

```
find_next_number( $d, b$ )  
next[ $k$ ]  
incremented  $\leftarrow 0$   
 $i \leftarrow 0$   
while  $i < k$  do  
    if not (incremented) then  
        if  $d[i] = b - 1$  then  
            next[ $i$ ]  $\leftarrow 0$   
        else  
            next[ $i$ ]  $\leftarrow d[i] + 1$   
            incremented  $\leftarrow 1$   
        end if
```

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    else
        next[i] ← d[i]
    end if
    i ← i + 1
end while
return next

```

**Question 2: [6 points]** Learn what a graph is. Then, show that the sum of degree of each vertex is twice the number of edges.

Consider a Graph  $G$  with vertices  $V$ ,  
Consider that initially

$$E(G) = \phi,$$

Therefore,

$$\deg(v) = 0 \quad \forall v \in V,$$

Also,

$$|E| = 0$$

Hence the give proposition holds for an arbitrary edge set  $E \subseteq V \times V$ .

Then if we add a new edge  $(u, v)$  where  $u, v \in V$  the number of edges increase by 1.

Also,  $\deg(u)$  and  $\deg(v)$  also increase by 1 each.

So the following holds

$$\sum_{v \in V} \deg(v) = 2|E|$$

Thus by the principle of mathematical induction, the proposition is proved.

**Question 3: [6+9 points]**

1. Find the total number of  $r$ -ary (i.e. in  $r$  arguments) functions from  $[n]$  to  $[k]$ .

**2. Show that the number of surjective (or, onto) maps from  $[n]$  to  $[k]$  is,**

$$\sum_{i=0}^k (-1)^i \binom{k}{i} (k-i)^n$$

**1. Total number of r-ary functions from  $[n]$  to  $[k]$**

First we find the number of single argument functions from  $[m]$  to  $[k]$   
As there is no condition on the type of function, for each element in  $[m]$ , we just need an element belonging to  $[k]$  as the image for the element  
Therefore, each element of  $[m]$  has exactly  $k$  choices.

Hence, we can apply product rule in this situation which will give us the total number of distinct functions to be:  $k \times k \times \dots \times k (m \text{ times}) = k^m$

Now each  $r$ -argument function can be seen as a single argument function where the argument is a  $r$ -length tuple.

Again using product rule, the number of  $r$ -tuples that can be made using elements from  $[n]$  is equal to  $n^r$ . Therefore, the set of all the possible  $r$ -tuples from  $[n]$  is in bijection with the set  $[n^r]$

Implies the number of  $r$ -ary functions from  $[n]$  to  $[k]$  will be equal to the number of single argument functions from  $[n^r]$  to  $[k]$  which we can calculate using the formula derived above to be  $k^{n^r}$

**2. Number of surjections from  $[n]$  to  $[k]$**

Consider  $A_1, A_2, A_3 \dots A_k$  denote the set of functions in which range of  $A_i$  does not contain  $i$

Using the principle of inclusion-exclusion on set  $A_i$ 's gives

$$\left| \bigcup_{i=1}^k A_i \right| = \sum_{S \subseteq [k], S \neq \emptyset} (-1)^{|S|+1} \left| \bigcap_{j \in S} A_j \right| \quad (1)$$

Now if we simplify the summation in RHS by taking  $|S| = 1, 2, 3 \dots k$ .

For  $|S| = p$ , the number of terms will be  $\binom{k}{p}$

Also, for any given  $S$  with  $|S| = p$ , the term  $\left| \bigcap_{j \in S} A_j \right|$  can be calculated

as follows:

We know that range of any function in  $\bigcap_{j \in S} A_j$  will not contain atleast  $p$  elements (the elements which are there in  $S$ )

So each element in domain  $([n])$  will have will be left with only  $(k - p)$  choices.

Now, using product rule,

$$\left| \bigcap_{j \in S} A_j \right| = (k - p)^n$$

Therefore, the summation in RHS of equation 1 can be simplified to

$$\left| \bigcup_{i=1}^k A_i \right| = \sum_{p=1}^k (-1)^{p+1} \binom{k}{p} (k - p)^n$$

Now, we see that the LHS of above equation denotes the number of functions which have do not have atleast one element in  $[k]$  in their range.

So, the number of surjections will be just total number of functions from  $[n]$  to  $[k]$  - LHS of above eqn

$$\#Surjections = k^n - \left| \bigcup_{i=1}^k A_i \right| = k^n - \sum_{p=1}^k (-1)^{p+1} \binom{k}{p} (k - p)^n$$

Which on some rearrangement gives ...

$$\#Surjections = (-1)^0 \binom{k}{0} (k - 0)^n + \sum_{i=1}^k (-1)^i \binom{k}{i} (k - i)^n$$

$$\#Surjections = \sum_{i=0}^k (-1)^i \binom{k}{i} (k - i)^n$$

Hence proved,

$$\#Surjections \text{ from } [n] \text{ to } [k] = \sum_{i=0}^k (-1)^i \binom{k}{i} (k - i)^n$$

**Question 4: [7 points]** Find all possible solutions for a sequence  $S_n$  which satisfies,

$$S_n = S_{n-1} + 6S_{n-2}$$

Let  $S_0$  and  $S_1$  be respectively the first and second term of the sequence. (As we are starting indexing from zero, here the  $n^{th}$  term of sequence will be denoted by  $S_{n-1}$ )

Now assume the generating function for this sequence as  $G(t)$  given by

$$\begin{aligned} G(t) &= \sum_{i=0}^{\infty} S_i t^i \\ \implies tG(t) &= \sum_{i=0}^{\infty} S_i t^{i+1} \\ \implies tG(t) &= \sum_{i=1}^{\infty} S_{i-1} t^i \end{aligned} \tag{2}$$

Further,

$$\begin{aligned} t^2 G(t) &= \sum_{i=0}^{\infty} S_i t^{i+2} \\ \implies t^2 G(t) &= \sum_{i=2}^{\infty} S_{i-2} t^i \end{aligned} \tag{3}$$

From (1)+6×(2), we get ...

$$\begin{aligned} (t + 6t^2)G(t) &= \sum_{i=1}^{\infty} 6S_{i-1} t^i + \sum_{i=2}^{\infty} S_{i-2} t^i \\ \implies (t + 6t^2)G(t) &= S_0 t + \sum_{i=2}^{\infty} (S_{i-1} + 6S_{i-2}) t^i \\ \implies (t + 6t^2)G(t) &= S_0 t + \sum_{i=2}^{\infty} (S_i) t^i \end{aligned}$$

$$\implies (t + 6t^2)G(t) = S_0t - S_0 - S_1t + \sum_{i=0}^{\infty} (S_i)t^i = (S_0 - S_1)t - S_0 + G(t)$$

Which on simplifying gives,

$$G(t) = \frac{S_0 + (S_1 - S_0)t}{1 - t - 6t^2} \quad (4)$$

Now as we have got the generating function, we can simplify it using partial fractions as:

$$\begin{aligned} G(t) &= \frac{S_0 + (S_1 - S_0)t}{(1 - 3t)(1 + 2t)} \\ G(t) &= S_0 \left[ \frac{\frac{3}{5}(1 + 2t) + \frac{2}{5}(1 - 3t)}{(1 + 2t)(1 - 3t)} \right] + (S_1 - S_0) \left[ \frac{\frac{1}{5}(1 + 2t) - \frac{1}{5}(1 - 3t)}{(1 + 2t)(1 - 3t)} \right] \\ G(t) &= \frac{S_0}{5} \left[ \frac{3}{1 - 3t} + \frac{2}{1 + 2t} \right] + \frac{S_1 - S_0}{5} \left[ \frac{1}{1 - 3t} - \frac{1}{1 + 2t} \right] \\ G(t) &= \frac{2S_0 + S_1}{5} \left[ \frac{1}{1 - 3t} \right] + \frac{3S_0 - S_1}{5} \left[ \frac{1}{1 + 2t} \right] \end{aligned}$$

Now, using binomial theorem

$$\begin{aligned} G(t) &= \frac{2S_0 + S_1}{5} \sum_{i=0}^{\infty} (3t)^i + \frac{3S_0 - S_1}{5} \sum_{i=0}^{\infty} (-2t)^i \\ \implies G(t) &= \sum_{i=0}^{\infty} \left[ \frac{2S_0 + S_1}{5} (3)^i + \frac{3S_0 - S_1}{5} (-2)^i \right] t^i = \sum_{i=0}^{\infty} S_i t^i \end{aligned}$$

Now comparing the coefficients of  $t^n$  on both sides,

$$\mathbf{S_n} = \frac{2\mathbf{S_0} + \mathbf{S_1}}{5}(\mathbf{3})^n + \frac{3\mathbf{S_0} - \mathbf{S_1}}{5}(-\mathbf{2})^n$$

Hence, there can be infinitely many sequences be made by varying the values of  $S_0$  and  $S_1$

**Question 5: [8 points]** Given 9 vertices, join all pairs of vertices by either red or blue edge. Show that there is always either a red *triangle* or a blue *quadrilateral*.

For 9 vertices and all possible edges, each vertex has 8 edges connected to it. By pigeonhole principle, comparing edges to pigeons and colour of edge to holes, we can claim that for every vertex, either there exists at least 4 red edges or 4 blue edges connected to it.

Consider the following 2 possibilities :

**1. There exists a vertex with at least 4 red edges.**

In this case, let this vertex be  $v_1$  and let  $v_2, v_3, v_4, v_5$  are 4 vertices such that edges  $(v_1, v_2), (v_1, v_3), (v_1, v_4), (v_1, v_5)$  are red.

Now if,

$$\exists (u, v) : u, v \in \{v_2, v_3, v_4, v_5\} \text{ \& } (u, v) \text{ is red,}$$

then, there exists a red triangle with vertices  $(v_1, u, v)$ .

Otherwise, if there exists no such red edge, which implies all these edges are blue thus making a blue quadrilateral with vertices  $(v_2, v_3, v_4, v_5)$ .

Hence, in case 1, there will always exist a red triangle or blue quadrilateral.

**2. There is no such vertex with 4 or more red edges.**

$\implies$  Every vertex has at least 5 blue edges connected to it.

Now consider any vertex  $v_1$ , and say  $v_2, v_3, v_4, v_5, v_6$  are blue.

Now take another vertex  $v_7$ .

Then since there are only 9 vertices and at least 5 vertices from  $v_7$  must be blue, there has to be 2 distinct vertices,

$$u, v \in \{v_2, v_3, v_4, v_5, v_6\} : (u, v_7) \text{ \& } (v, v_7) \text{ are blue.}$$

But then  $(v_1, u, v_7, v)$  form a blue quadrilateral.

Hence in case 2, there always exists a blue quadrilateral.

So in all the possible situations, we have seen that there either exists a red triangle or a blue quadrilateral.

**Question 6: [5+8+2 points]**

1. Suppose  $\alpha$  is a rational. Show that there exists an  $n_0 \in \mathbb{N}$  such that for every rational number  $\frac{p}{q}$  with  $1 \leq q < n_0$ ,

$$\left| \alpha - \frac{p}{q} \right| \geq \frac{1}{n_0 q}$$

2. Suppose  $\alpha$  is an irrational real. Show that for any  $n \in \mathbb{N}$ , there is a rational number  $\frac{p}{q}$  with  $1 \leq q \leq n$ , s.t.,

$$\left| \alpha - \frac{p}{q} \right| < \frac{1}{nq}$$

**Finally, what do above two properties characterize ?**

1. Since  $\alpha$  is given to be rational, we can consider it of the form  $a/b$  where  $a$  and  $b$  are co-prime integers and  $b > 0$   
Multiplying by  $bq$  on both sides on given inequality reduces to

$$|qa - pb| \geq \frac{b}{n_0}$$

Now, if we take  $n_0 = b$ , then RHS of inequality becomes 1

And,  $q < n_0 \implies q < b$ , so  $\frac{a}{b} \neq \frac{p}{q} \implies |aq - bp| \neq 0$  (because  $a$  and  $b$  are co-prime, so for  $a/b = p/q$ , we need  $a=p$  and  $b=q$ )

Further,  $|qa - bp|$  has to be an integer, We can claim that

$$|qa - bp| \geq 1$$

Therefore, we can say for any rational  $\alpha$ , where  $\alpha = \frac{a}{b}$ ,  $a, b$  are co-prime,  
 $\exists n_0 = b$  such that for every rational  $p/q$  with  $1 \leq q < n_0$ ,  $\left| \alpha - \frac{p}{q} \right| \geq \frac{1}{n_0 q}$



2. Multiplying by  $q$  gives the simplified inequality

$$|q\alpha - p| < \frac{1}{n}$$

Now, consider the fractional part of  $q\alpha$ ,  $1 \leq q \leq n$

If there is some  $q$  for which fractional part of  $q\alpha < 1/n$ , then we have proven the inequality, we can assign  $p = [q\alpha]$  where  $[.]$  represents greatest integer function.

Now, we prove that there always exist such a  $q$ , we will use method of contradiction

Assuming  $\nexists q : q\alpha - [q\alpha] < 1/n$

Now consider the intervals of form  $(\frac{i}{n}, \frac{i+1}{n})$ ,  $1 \leq i < n$  as  $(n-1)$  holes and the fractional parts of  $q\alpha$  as  $n$  pigeons

Applying pigeonholes principle, we get atleast one interval in which there are two fractional parts lie for distinct  $q$ , say  $q_1$  and  $q_2$

But then  $|(q_1 - q_2)\alpha| < 1/n$  because the length of each interval is  $1/n$ , which is a contradiction.

Hence there always exist a  $q_0$ ,  $1 \leq q_0 \leq n$  for which  $q\alpha - [q\alpha] < 1/n$

Hence assigning,  $\frac{p}{q} = \frac{[q_0\alpha]}{q_0}$  will ensure that,

$$\left| \alpha - \frac{p}{q} \right| < \frac{1}{nq}$$

Property 1 characterizes that we can always quantify the difference between any two rationals using a natural number.

Property 2 characterizes that we can have a rational approximation for any irrational which can be made arbitrarily close to the irrational, as in the inequality we can take  $n$  to be arbitrarily large.

## 1 Acknowledgements

1. Professor Rajat Mittal's notes uploaded on Course Website, used in combining ideas of question 6.2.