BHCA ⇒ no. of calls in 1 hour

n calls = 1 hour

1 call = $\frac{1 \text{ hour}}{n}$ ⇒ Call Processing Time

tp ⇒ time taken to process 1 call

n time = 1 call

1 hour = $\frac{1 \text{ hour}}{n}$ ⇒ no. of calls

Tute-1- Hax

① CCR = Call Completion Rate = $\frac{\text{no. of succesfull calls}}{\text{total calls}} = \frac{\text{succesful traffic}}{\text{total traffic}}$

② BHCA = call attempts in busy hours (1hr by default)

③ BHCR = Busy Hour Calling Rate = no. of succesfull calls made by each sub $= \frac{\text{no. of succesfull calls}}{\text{subscribers}}$

④ Traffic = $\frac{\text{occupied/busy time of server}}{\text{total time}}$

⑤ Hax : (i) $BHCA = \frac{1\,hr}{t_p} \Rightarrow t_p = \frac{1\,hour}{BHCA}$

(ii) $BHCA = \frac{BHCR \times BCOS}{CCR}$

(iii) offered/total traffic = $\frac{\text{no. of call} \times \text{call duration}}{\text{observation time}}$ unit $\Rightarrow E$

$= \text{call rate} \times \text{call duration}$

(iv) $GOS = \frac{\text{loss traffic}}{\text{offered traffic}} = \frac{\text{lors}}{\text{offer}} \Rightarrow$ probability of lors

1. Sequential Write/Random Read:
   X x Y —> Consider X for Adress Lines and Data Memory. For Control Memory, consider Y
2. Sequential Read/Random Write:
   X x Y —> Consider Y for Address Lines and Data Memory. Control Memory, consider X
3. Diagram: given x....y —>
   inlets.....outlets

If Read —> first wale bahar

If Write —> second wale bahar

Edited 8:45 AM

TELECOMMUNICATION NETWORKS

MODERATE ASSIGNMENT

NAME : ANUNAY DWIVEDI

ROLL : 9922102048

BATCH : E2

① Total no. of subscriber = 1000

Peak hour = 40 % of 1000

= 400

External Traffic = $\frac{400}{5} = 80$

② $SAR = \frac{\sqrt{N}}{3}$

(a) $N = 228$, $SAR = \frac{\sqrt{228}}{3} = 3.72$

(b) $N = 32768$, $SAR = \frac{\sqrt{32768}}{3} = 60.34$

③ $M = N = 512$, $P = q = 16$, $\alpha = 0.7$

~~$P = \frac{\alpha}{S(1-}$~~ $P_B = \frac{(M-S)(N-S)\alpha}{MN}$

(a) $S = 16 \Rightarrow P_B = \frac{172262}{262144} = 0.657$

(b) $S = 24 \Rightarrow P_B = \frac{160366}{262144} = 0.635$

(c) $S = 31$, $P_B = \frac{161211}{262144} = 0.615$

Error = $\frac{0.615 - 0.61}{0.61} \times 100 = 0.820$

(4) Inlet utilization = 0.1

$$P_B = 0.002 = \frac{(M-S)(N-S)\alpha}{MN}$$

$$S = N - \sqrt{N \ln(1/P_B)}$$

(a) N = 128

$$S = 128 - \sqrt{128 \ln(500)}$$

$$= 100$$

(b) N = 2048

$$S = 2048 - \sqrt{2048 \ln(500)}$$

$$= 1935$$

(c) N = 8192

$$S = 8192 - \sqrt{8192 \ln(500)}$$

$$= 7967$$

Tute - 3 - 4 - Mark

① Binary encoding : (i) Dataword (iv) error ⇒ make any "1" to "0"
(ii) Divisor
(iii) Aug. dataword = Dataword + (no. of divisor bits - 1) 0's
(iv) Divide

② Polynomial encoding : (i) dataword (v) error ⇒ add any term a/c 'ly
(ii) divisor
(iii) Aug dataword = dataword × (divisor's max power term)
(iv) Divide

③ Flag : 0 1111 1110 ⇒ zero triple one triple one zero

④ Bit Stuffing : add '0' after consecutive '1' that apears 5 times
↳ add flag befor & after the final code

⑤ $\text{Transmission Time} = \frac{\text{Frame Size}}{\text{Frame Rate}} = \frac{L}{R}$ → bits, → kbps = msec

⑥ $\text{Total Time} = \text{Transmission Time} + 2(\text{Propagation Time})$ unit msec

⑦ $\text{Data Link Utilization} = \frac{T_{frame}}{T_{frame} + 2T_p} = \frac{\frac{L}{R}}{\frac{L}{R} + 2T_p}$
∪ ⇒ :

⑧ $U\ (\text{w/o errors}) = \frac{1}{1+2A}$ where $A = \frac{t_p}{t_f}$

⑨ If "w" is given ⇒ $U = \frac{W}{1+2A}$ unit o/o

⑩ If error "p" is given ⇒ $U = \frac{1-p}{1+2A}$ unit o/o

⑪ if "w" is given ⇒ $U = \frac{W(1-p)}{1+2A}$ unit o/o

ChatGPT
Explore GPTs

Today
Telecomms
Internship Form Help

Yesterday
DHD

Previous 7 Days
Literature
Photovoltaic
VLSI
GitHub

Upgrade plan
More access to the best models

🌟 Key Difference:

Feature	Stop-and-Wait	Sliding Window
Frames sent at a time	1 frame at a time	Multiple frames at a time
Efficiency	Low	High
Waiting time	Idle time while waiting for ACK	Continuous frame transmission
Ideal for	Small-distance networks	Long-distance, high-speed networks

POSSESION OF MOBILES IN EXAM IS UFM PRACTICE

Name Anurag Dwivedi Enrollment No. 9922102048

Jaypee Institute of Information Technology, Noida
Test-2 Examination, 2025
B.Tech. VI Semester

Course Title: Telecommunication Networks **Maximum Time: 1 Hr**
Course Code: 15B11EC611 **Maximum Marks: 20**

CO1	**To understand the basic concepts of Telecommunication network model, Traffic engineering and switching technology. Also, to understand various mechanisms involved in OSI model, TCP/IP and LAN access protocols, ATM and ISDN.**	**Understanding Level (C2)**
CO2	**To apply the concepts of traffic engineering, switching technologies and various network protocols for serving network related problems.**	**Applying Level (C3)**
CO3	**To analyze the link utilization and data packet generated after incorporation of data link error control and flow control mechanisms.**	**Analyzing Level (C4)**
CO4	**To apply the concept of subnetting for creating and assigning address blocks in a network. Applying various routing algorithms to create routing table for communication between two nodes.**	**Evaluating Level (C5)**

NOTE – Attempt all the questions.

Q.1. (a) Derive the utilization factor for Stop-and-Wait ARQ. Demonstrate it with diagram. **[3]**

[CO4 - Analyzing Level]

(b) A channel has a data rate of 3 Kbps and propagation delay of 10 msec. For what range of frame sizes does Stop-and-Wait give an efficiency of at least 70%. **[3]** [CO4 - Analyzing Level]

Q.2. Explain the methods of flow control with appropriate diagrams. **[3]** [CO4 - Analyzing Level]

Q.3. The following character encoding is used in a data link protocol:

A : 01100111

B : 11100010

FLAG : 01111110

ESC : 11100001

Show the bit sequence transmitted (in binary after bit stuffing for the four characters frame:A B ESC FLAG. **[3]**

[CO4 - Analyzing Level]

Q.4. Given the data word 1010111011 and the divisor 1001:

(a) Show the generation of the code word at the sender's side.

(b) Show how to determine an error at the receiver's side (if error is introduced in the 2nd bit from LSB of the transmitted code). **[2.5 + 2.5]** [CO4 - Analyzing Level]

Q.5. State differences between Virtual circuits and Datagram networks. **[3]** [CO3 - Applying Level]

Date .....

Tute-5-Hax

$$s = \frac{d}{t} \Rightarrow t = \frac{d}{s}$$

① General Stuff:

(i) N = no. of stations

(ii) Frame Size = ~~no. of frames~~ no. of bits in one frame

(iii) Frame Rate = no. of frames / sec being sent

(iv) Bandwidth = kbps of Mbps = speed of data transfer

(v) Throughput = Successfully delivered kbps

(vi) Frame time = Time taken for frames to be sent in sec or μsec

(vii) G = Channel Load

② Pure ALOHA:

$$S_{max} = 18.4\% = 0.184 = Ge^{-2G}$$ = Throughput

③ Slotted ALOHA:

$$S_{max} = 36.8\% = 0.368 = Ge^{-G}$$ =

④ Data speed = $\frac{\text{frame size}}{\text{frame time}}$

⑤ Load = G = frame rate × time for one frame → (frame size / bandwidth)

⑥ Amount of frames traffied = S × frame rate

S → throughput

Date .....

Tute-5

① Given: N stations, bandwidth = 5 kbps, pure = $18.4 \times 10 = 0.184$
frame size = 1000 bit every 100 sec

To find: N

Sol^n: 1 frame every 100 sec from 1 station
amount of data = $\frac{1000}{100}$ = 10 bps ——— (1)
from all frames
pure ALOHA = 0.184
Speed = 5×0.184

$$N = \frac{5 \times 0.184}{10} = 1030$$

② Given: N = 10000, Slotted = 36.8% To find: G = Normal load
one station = 18 req / hour
time = 125 usec / slot

Sol^n: Amount of req = $10000 \times \frac{18}{60 \times 60} = 50$ req/sec ——— (1)

$$\text{Load} = 125 \times 10^{-6} \times 50 = 0.625\%$$

③ Given: Slotted ALOHA = S_{max} = 36.8%, frame size = 200 bit
bandwidth = 200 kbps

Sol^n:

$$S = G e^{-G}$$ where: G = load

(a) G = frame rate × time $= 1000 \times 10^{-3} = 1$

$S = e^{-1} = 0.367 = 36.7\%$

amount of frames = $S \times 1000 = 368$

(b) $G = 500 \times 10^{-3}$

$S = 0.5 e^{-0.5} = 0.303 = 30.3\%$

amount = $0.303 \times 500 = 152$

(c) $G = 250 \times 10^{-3} = 0.25$

$S = 0.25 e^{-0.25} = 0.19$

amount = $0.19 \times 250 = 49$

④ N = 100, ban = 1 Mbps, fs = 1000

fs = 10 f-r/sec

$S = G e^{-G}$

$G = 10 \times 10^{-3} \times 100 = 1$

$S = e^{-2}$
$= 13.53\%$

Spiral

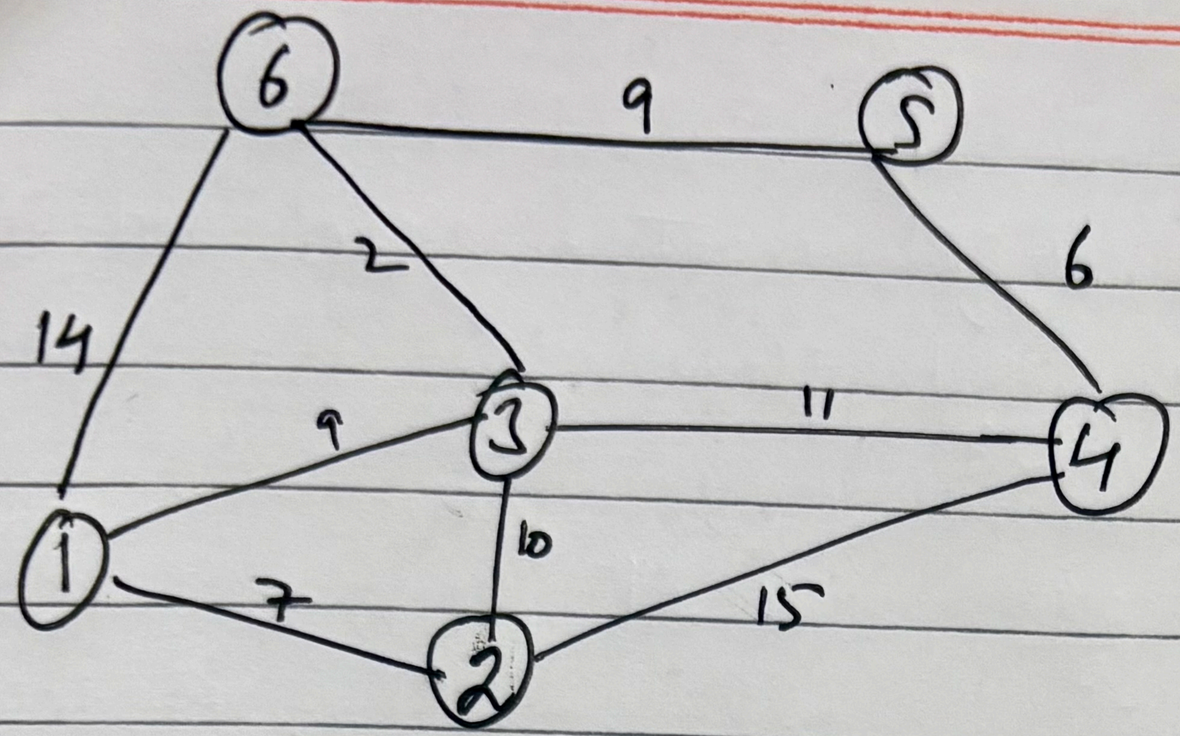
Telecom T3

(1) T1 ⇒ 5 marks (switching, ARQ) only theory
(2) T2 ⇒ 5 marks (
(3) T3 ⇒ IPB4, IPB6, routing algo, ALOHA and all
↳ CSMA, congestion

TCP, Quality of Service (3 marks) → not told but in syllabus

Logical addressing

(4) Tute: 4, 5, 6 ⇒ numericals



(i) Select root node

(ii) Calculate all direct distances

(iii) Select min and don't change it

(iv) Look through all possible paths

Source (1)	2	3	4	5	6
	∞	∞	∞	∞	∞
	7	9	∞	∞	14
1,2	7	9	22	∞	14
1,2,3	7	9	20	∞	11
1,2,3,6	7	9	20	23	11
1,2,3,6,4	7	9	20	20	11