

# Regular Languages

## Summary

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## 1 Quick Summary

### 1.1 DFA

#### Overview

##### What it contains

1. Set of states ( $Q$ )
2. An alphabet ( $\Sigma$ )
3. Transition function( $\delta$ )
4. Starting state ( $q_0$ )
5. Set of accepting state( $f$ )

The DFA starts at the initial state and follows the transition function consuming one letter at a time and accepts the string if and only if one of the accepting states is reached.

### 1.2 NFA

#### Overview

##### What it contains

1. Set of states ( $Q$ )
2. **An alphabet along with ' $\epsilon$ '** ( $\Sigma \cup \epsilon$ )
3. **Transition Relation** ( $\delta \subset Q \times \Sigma \times Q$ )
4. Starting state ( $q_0$ )
5. Set of accepting state( $f$ )

The NFA starts at the initial state and follows the transition relation consuming at most one letter at a time and accepts the string if and only if at least one of the runs end in an accepting state.

**Acceptance:** If at least one run terminates on an accepting state.

#### Languages recognized by NFA

The set of languages recognized by NFA is in fact, the set of regular languages.

### 1.3 Regular Languages

#### Closed under Complementation

##### How

Set all the non-accepting state as accepting states and accepting states as non-accepting states.

##### Why

Every string that terminated in a previously non-accepting state is accepted and vice versa.

##### Examples

1.  $L = \{a^n b^m \mid n, m \geq 0\}$  over  $\{a, b\}$
2.  $L = \{a^n b^n \mid n \geq 0\}$  over  $\{a, b\}$

#### Closed under Intersection and Union

##### Intersection

Do a "cross product" of the two machines (i.e Create a new machine that runs both the machines simultaneously) and accept iff both the machines accept.

##### Union

Do a cross product of the two machines and accept iff either of the two machines accept.

#### Closed under Concatenation and Star

##### Concatenation

Construct a new NFA where the accepting states of the first NFA have an epsilon transition to the start state of the second NFA.

##### Star

Construct a new NFA where the accepting states of the initial NFA have an epsilon transition to the start state of itself.

### 2 Tips and Tricks to solve problems

#### Quick tips and tricks

1. Any computation (decision) requiring finite memory can be performed in a DFA.
2. To show that a language is regular, prefer constructing an NFA over a DFA.
3. Instead, you can also use closure properties on known Regular language examples.
4. The idea of simultaneously running two Automatas can be simulated by taking the cross-product of states. This can be useful in several scenarios.