Stack: isEmpty, isFull and Top

```
int isEmpty(STACK *s) {
    return s\rightarrow top == -1;
int isFull(STACK *s) {
    return s->top == s->limit;
int Top(STACK *s) {
    if (isEmpty(s)) {
        printf("Stack is empty\n");
        return -1;
    return s->info[s->top];
```

Stack: Push and Pop

```
void push(STACK *s, int element) {
     if (isFull(s)) {
         printf("Stack full: insertion denied\n");
         return:
    s\rightarrow info[++s\rightarrow top] = element; // pointer to end
        of list
    return;
int pop(STACK *s) {
     if (isEmpty(s)) {
         printf("Stack empty: deletion denied\n");
         return -999:
    return (s\rightarrow sinfo[s\rightarrow top--]);
```

Stack: Print

```
void printStack(STACK *s) {
    int i;
    if (isEmpty(s)) {
         printf("Stack is empty\n");
         return:
    while (i < s \rightarrow top) {
         printf("%d\t", s\rightarrowinfo[++i]); // print
             elements
    printf("\n");
    return;
```

Application of Stack

- Parenthesis matching.
- Evaluation of expression.
- Converting expression from one form to another.
- Used also in DFS, in general backtracking.
- Used in memory management.

Parenthesis Matching

- For parenthesis matching:
 - Scan the parentheses in left to right order
 - On encountering an opening parenthesis push it to stack.
 - On encountering a closing parenthesis just pop the topmost parenthesis from the stack.
- If after all input scanning is over and stack is empty the parentheses are matched.
- If stack is empty but the next input symbol is an closing parenthesis match fails.
- If stack is not empty but input scanning is over the match fails.

ADT

Parenthesis Matching

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Parenthesis Matching

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Infix to Postfix Conversion

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Expression Evaluation

- Assume we have postfix expression: AB+C*DEF+/-
- ► Let A = 4, B = 6, C = 2, D = 12, E = 1, F = 2.
- ► The evaluation of infix form of the expression give (4 + 6) * 2 12 / (1 + 2) = 20 12/3 = 16.
- ► For evaluation, scan the expression from left to right and use following rules repeatedly:
 - Push symbol onto S if it is an operand.
 - On encountering an operator, pop the correct number of operands, apply operation and push result on to S.
- When finished the result is on the top of S.

Example: Postfix Evaluation

4 6 + 2 * 12 1 2 + / -							
Step	Symbol	Stack	Output	Description			
1	4	4	-	1st symbol			
2	6	4, 6	_	2nd symbol			
3	+	10	10	3nd symbol			
4	3	10, 2	-	4th symbol			
5	*	20	20	5th symbol			
6	12	20, 12	-	6th symbol			
7	1	20, 12, 1	-	7th symbol			
8	2	20, 12, 1, 2	-	7th symbol			
9	-	20, 12, 3	3	8th symbol			
10	/	20, 4	4	9th symbol			
11	-	16	16	10th symbol			

Infix to Postfix

- Append a ")" to input expression X.
- Push an "(" onto an empty stack S.
- Scan X from left to right repeating following steps:
 - If current symbol is an operand add to output.
 - If current symbol is an "(", push it on to S.
 - If current symbol is an operator:
 - Repeatedly pop operators from stack with same or higher precedence and add them to output.
 - 2 Add the current symbol to stack.
 - If current symbol is a ")" then
 - Repeatedly pop symbols from stack until a matching "(".
 - 2 Pop and discard the "("

Example

(A+B)*C - D/(E+F))← added							
Step	Symbol	Stack	Output	Description			
1	-	(-	start			
2	(((-	1st symbol			
3	Α	((Α	2nd symbol			
4	+	((+	Α	3rd symbol			
5	В	((+	AB	4th symbol			
6	*	((*	AB+	5th symbol			
7	С	((*	AB+C	6th symbol			
8	-	((-	AB+C*	7th symbol			

Example Continued

(A+B)*C - D/(E+F))							
Step	Symbol	Stack	Output	Description			
9	D	((-	AB+C*D	8th symbol			
9	/	((/	AB+C*D-	9th symbol			
10	(((/(AB+C*D-	10th symbol			
11	E	((/(AB+C*D-E	11th symbol			
12	+	((/(+	AB+C*D-E	12th symbol			
12	F	((/(+	AB+C*D-EF/	13th symbol			
13)	(AB+C*D-EF+/	14th symbol			
14)	-	AB+C*D-EF+/	Added) symbol			