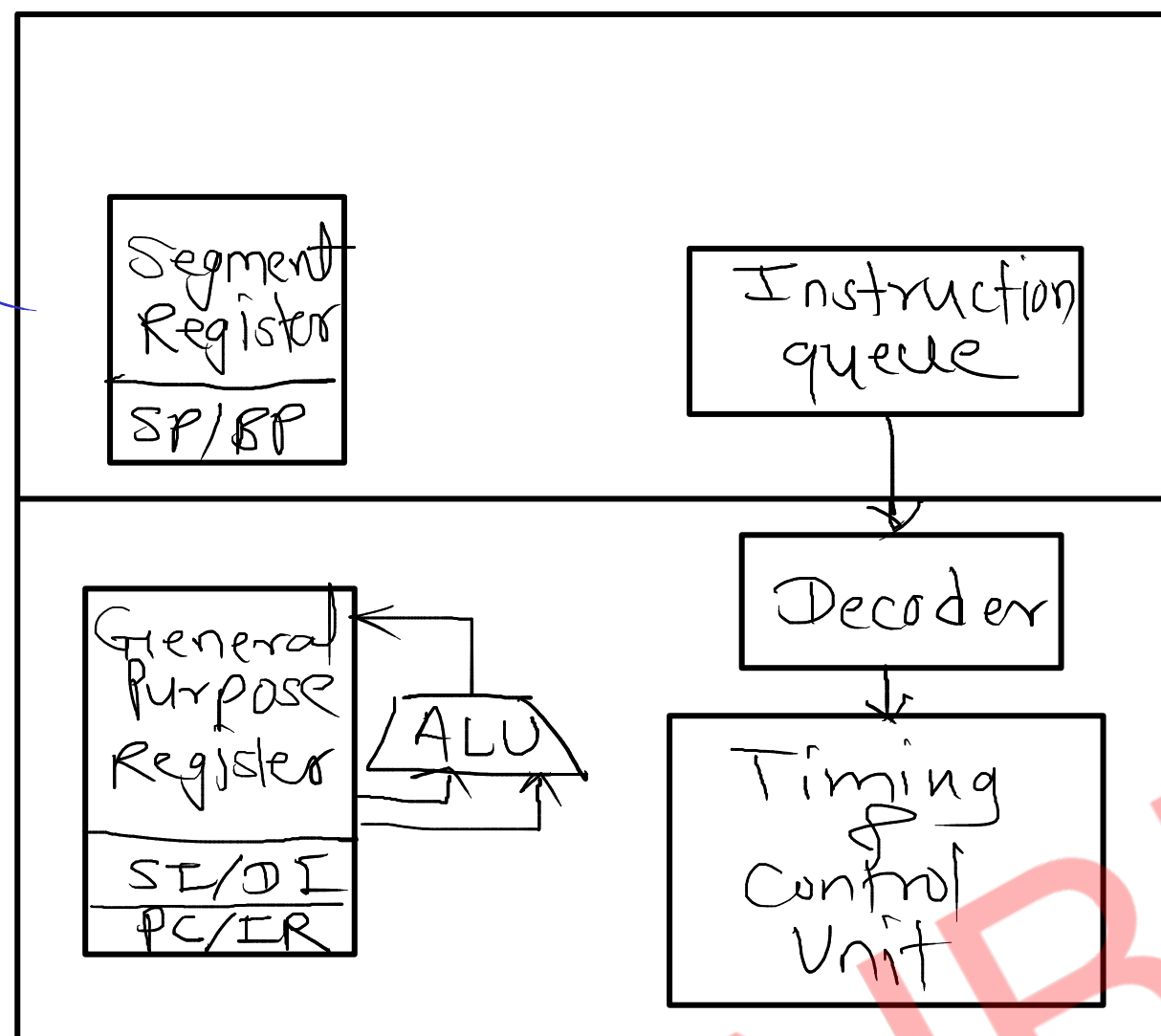


Bus
Interface
Unit



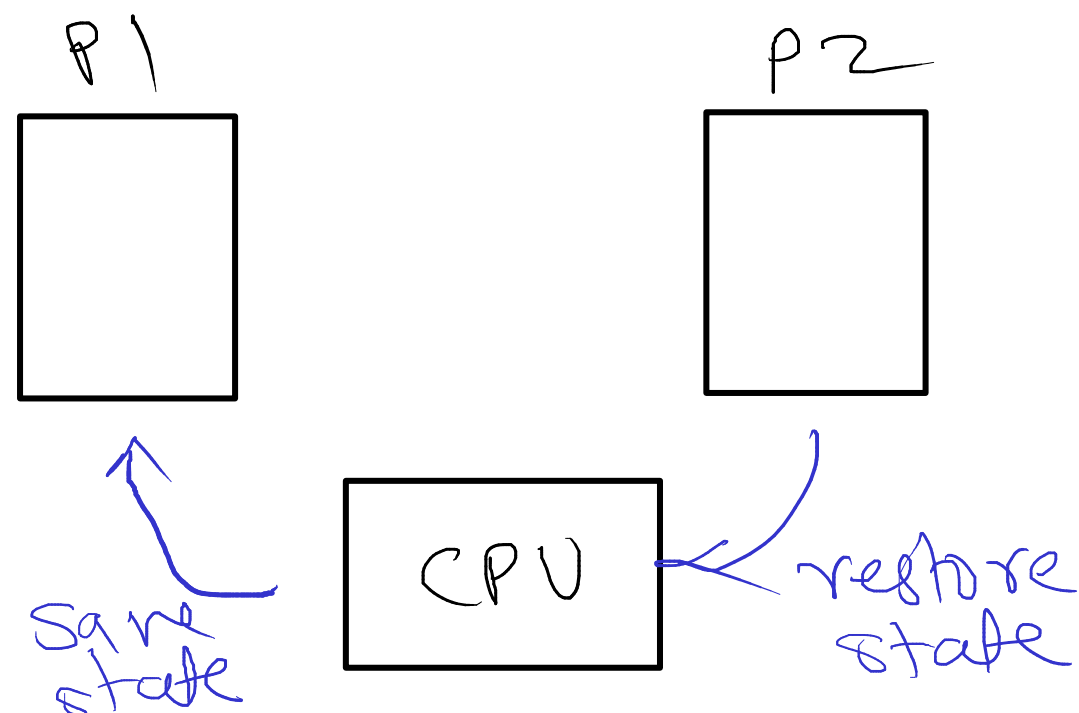
Execution
Unit

Execution Context:

— values of CPU registers

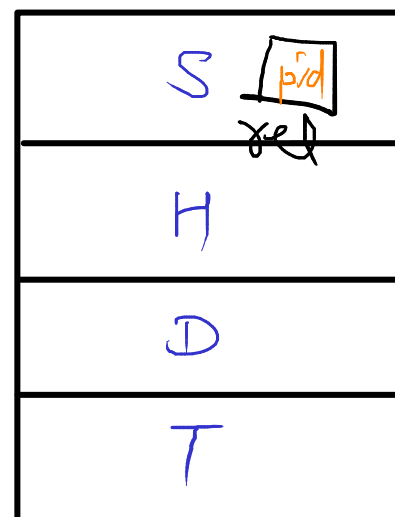
Context switching:

— save execution context of one process & restore execution context of another process.



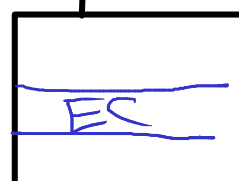
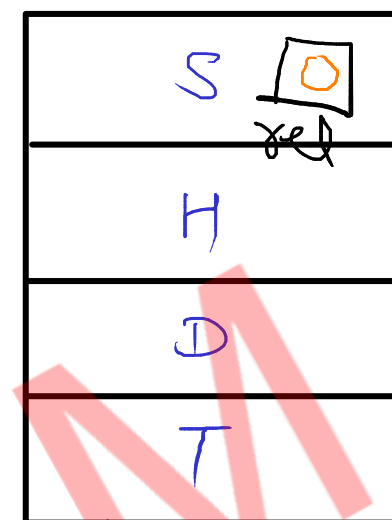
Parent
process

/a.out

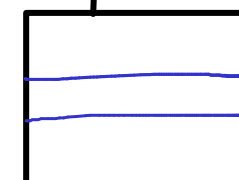


fork() →

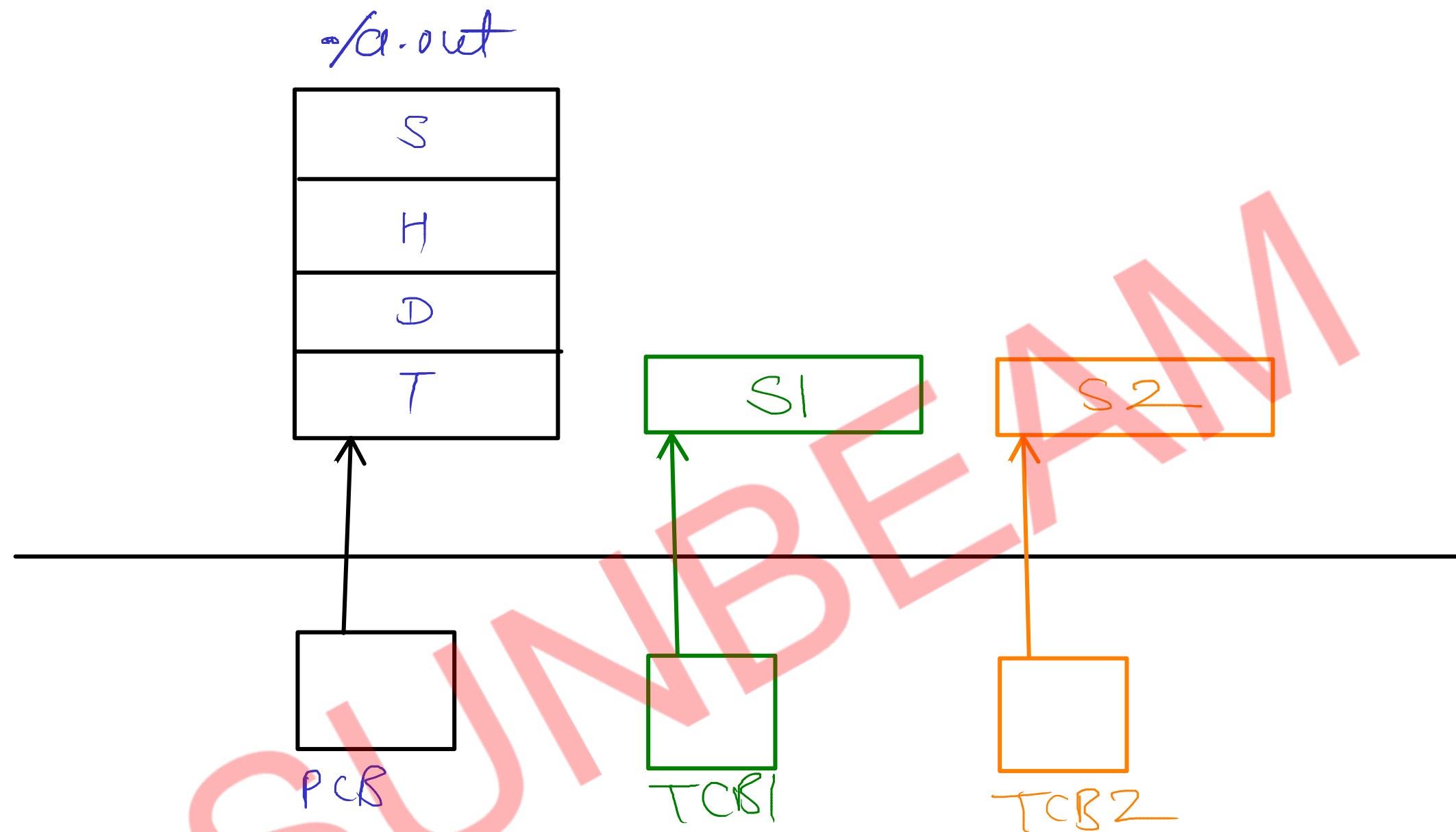
child
process



PCB



PCB



Semaphore

- semaphore is internally a counter

Dec:

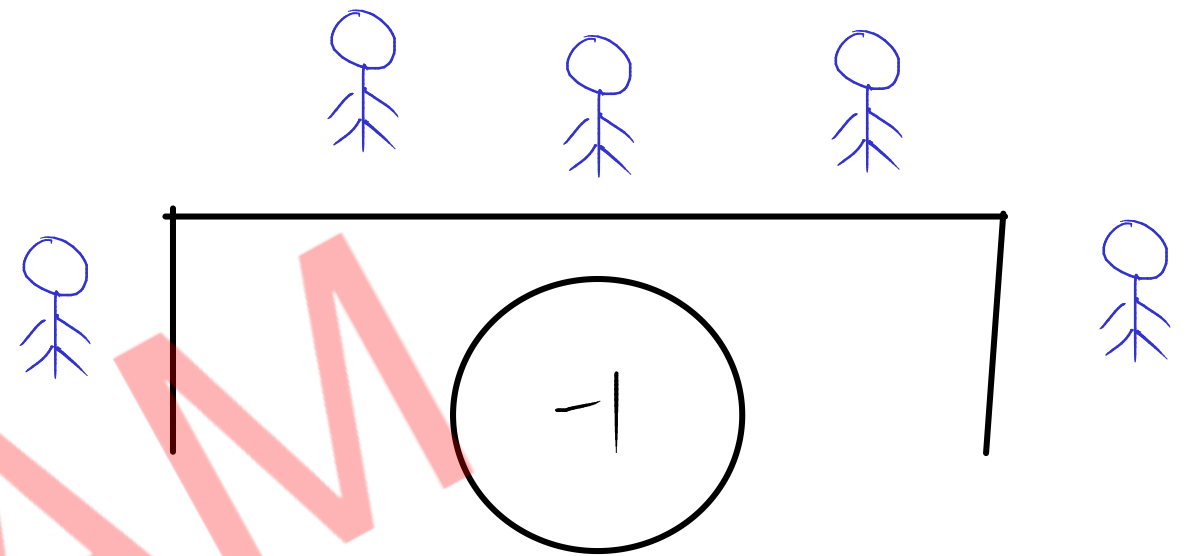
- 1) decrement counter
- 2) if counter < 0 , then block current thread/process.

Inc:

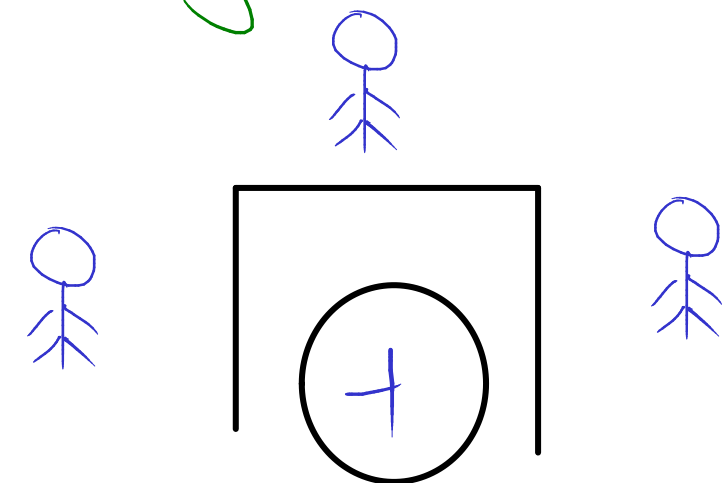
- 1) increment counter
- 2) if any process/thread is blocked on this semaphore, wake up one.

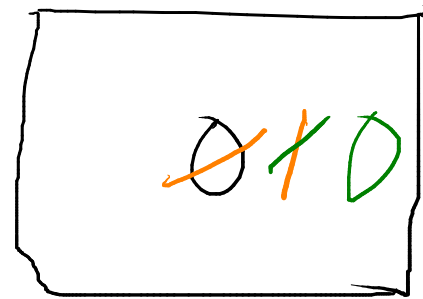
create	-	sem_init()
dec	-	sem_wait()
inc	-	sem_post()
destroy	-	sem_destroy()

Counting semaphore



Binary semaphore





count

sem_init(&s, 0, 1)

```
mid inc(void)
{
```

```
for( )
{
```

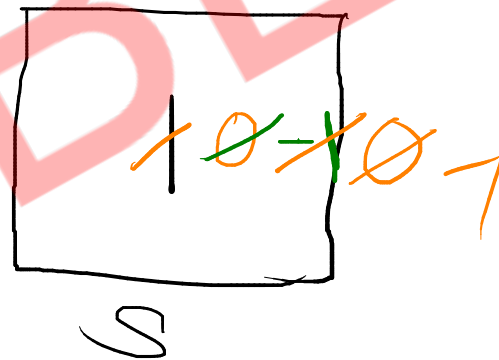
```
sem_wait(&s)
```

```
count++;
```

```
sem_post(&s)
```

```
}
```

```
}
```



```
mid dec(void)
{
```

```
for( )
{
```

```
sem_wait(&s)
```

```
count--;
```

```
sem_post(&s)
```

```
}
```

```
}
```

sem_destroy(&s)

Mutex

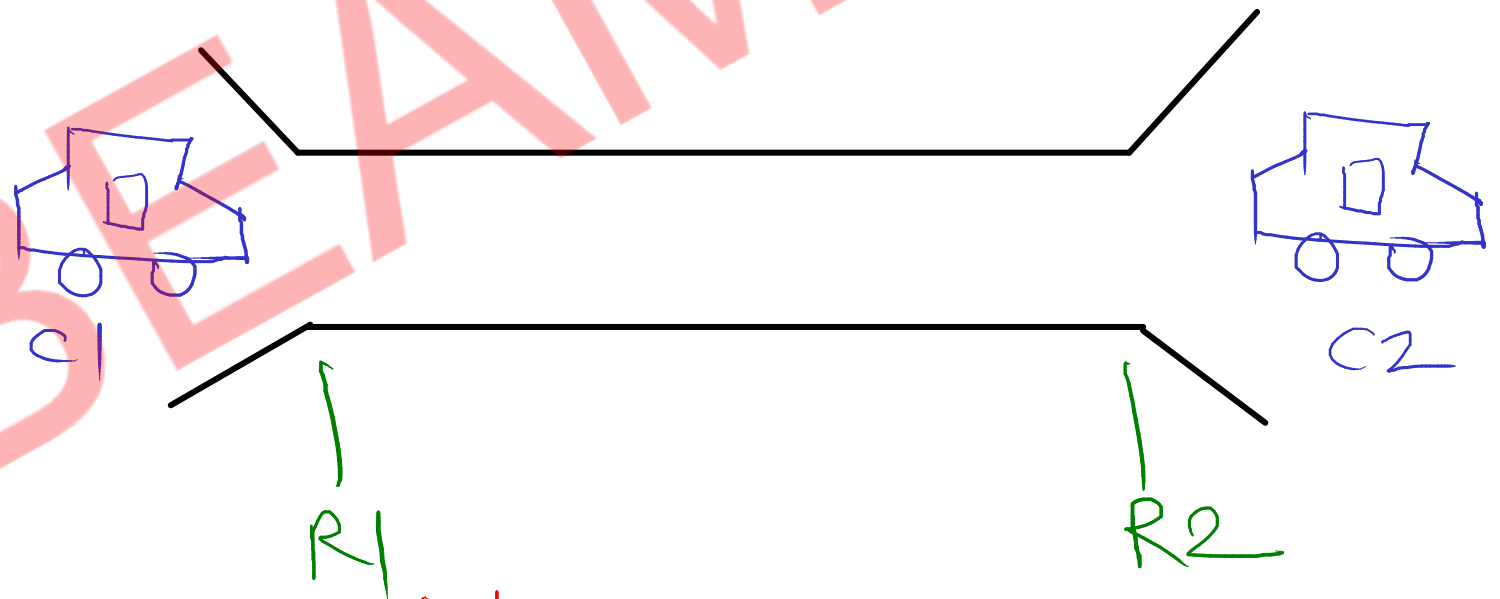
- Mutual Exclusion — one at a time
- Mutex is similar to binary semaphore.
- Mutex operations — lock(), unlock()
- The process who will lock the mutex, will become owner of that mutex.
- Only owner can unlock the mutex.

pthread_mutex_create()	— create
pthread_mutex_lock()	— lock
pthread_mutex_unlock()	— unlock
pthread_mutex_destroy()	— destroy.

DeadLock

- indefinit waiting of process for some common resource.
- there are four conditions of dead lock. & when all four conditions will be true ^{at same time}, deadlock will occur.

- 1) Mutual exclusion
- 2) No preemption
- 3) Hold & wait
- 4) Circular wait.



Prevention:

- while implementing code of OS, it is ensured that, any one condition out of 4 is always false

Avoidance:

- 1) resource allocation graph
- 2) Banker's algorithm
- 3) Safe state algorithm

Recovery:

- 1) resource preemption
- 2) Process termination