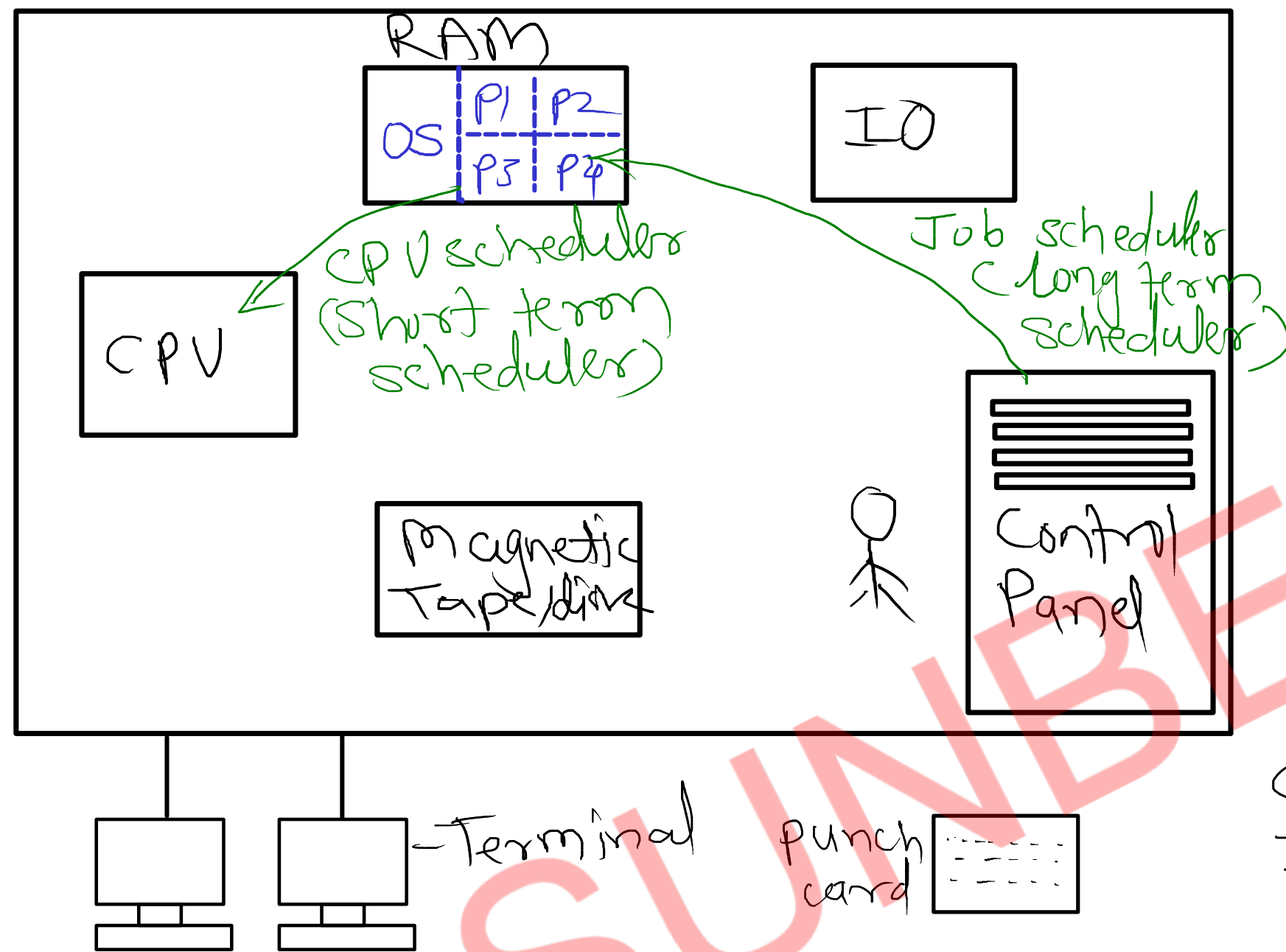


Types of Operating System



1) Resident Monitor

2) Batch systems

3) Multiprogramming system
- multiple programs are loaded into RAM.

Degree Multiprogramming -

- number of processes loaded into RAM.

CPU burst - time spent on CPU

IO burst - time spent on IO

$\text{CPU burst} > \text{IO burst}$ - **CPU bound**

$\text{IO burst} > \text{CPU burst}$ - **IO bound**

- mixture of CPU & IO bound processes is load into RAM

4) Time sharing system /

Multitasking system

- CPU time is shared into all the processes of RAM

Response time $< 1\text{sec}$

5) **Mult users system**

- multiple users are connected to the same system through different terminals.

- whoami, tty, who

6) Multiprocessing system

- multiple CPUs are fitted on single chip. such processor is known as "multiprocessor" / "multi core".

- OS can take advantage of this to schedule multiple processes at same time on different cores.

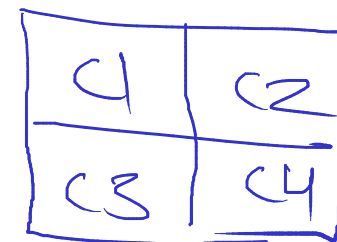
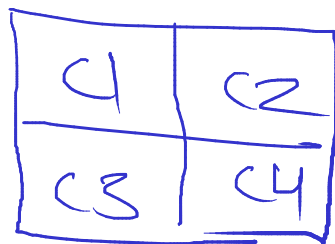
- processes can be executed parallelly, that's why it is also known as 'parallel systems'.

Linux - kernel 2.5+

Windows Vista

There are two types of multiprocessing

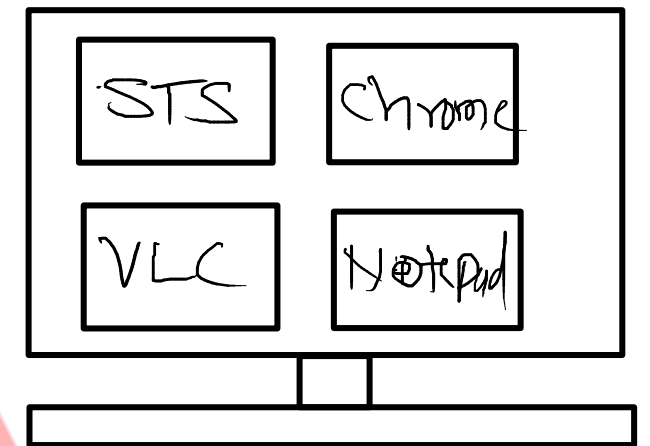
i) Asymmetric multiprocessing ii) Symmetric multiprocessing



Types of multitasking

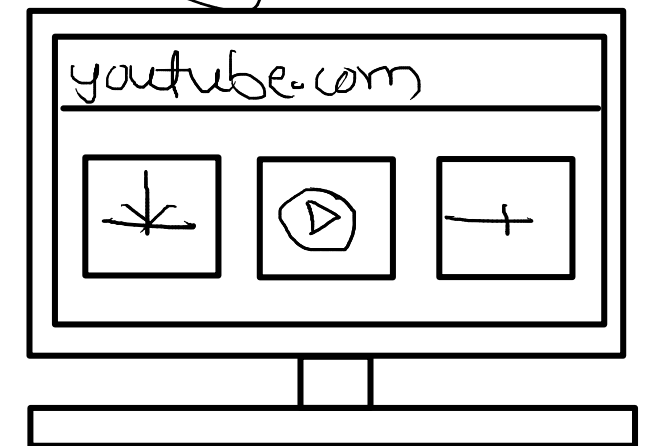
i) Process based multitasking

- system wide multitasking



ii) Thread based multitasking (multithreading)

- process wide (within process) multitasking



OS's Data Structures

1. Job queue / Process list

- all process of RAM are kept into this queue

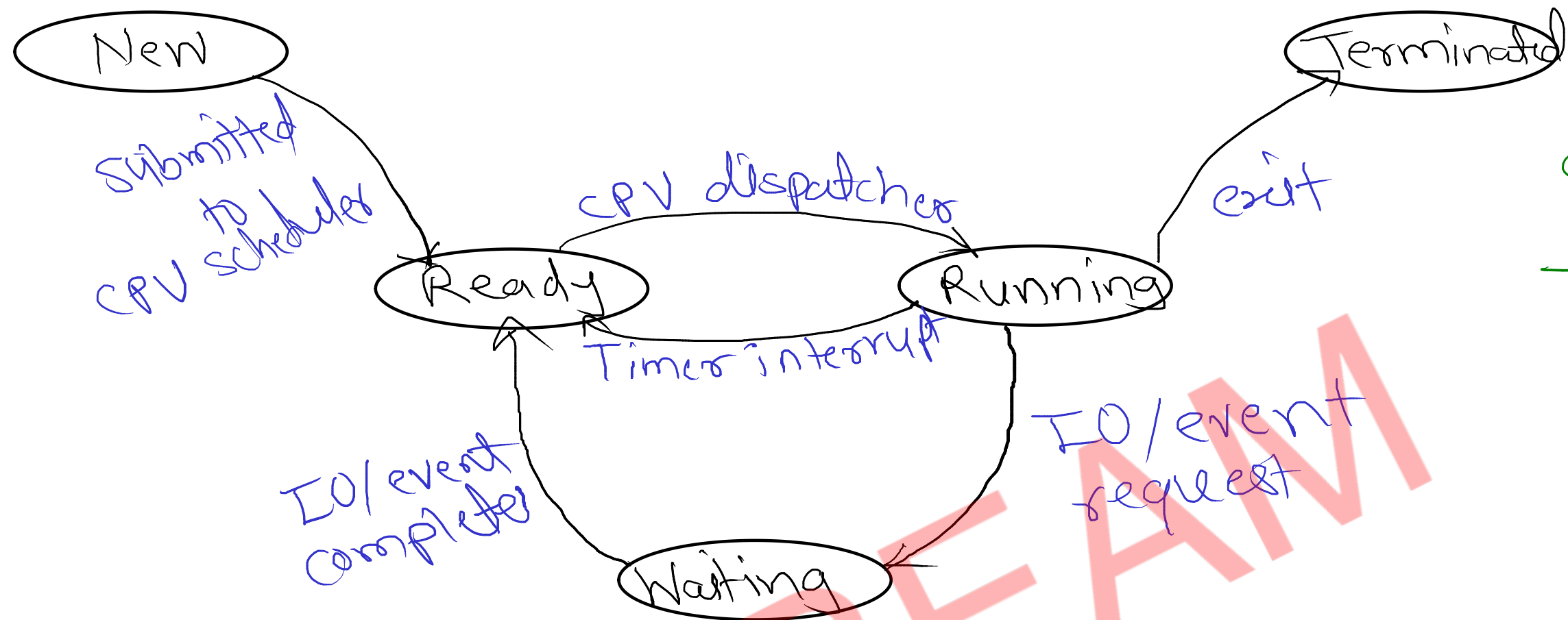
2. Ready queue

- processes which are ready for execution on CPU
- CPU scheduler is going to select one process from ready queue

3. Waiting queues

- processes which are waiting for I/O/event
- per device one waiting queue.

Process Life cycle



CPU sched Algorithm

- 1) FCFS
- 2) SJF
- 3) Priority
- 4) RR
- 5) Fair Share

- 1) Running → Terminated
- 2) Running → Waiting
- 3) Running → Ready
- 4) Waiting → Ready

} voluntarily
} forcefully

Types of scheduling

- 1) Non preemptive
CPU access is given to another process voluntarily
- case 1 & 2
- 2) preemptive
CPU access is given to another process forcefully.
- case 1 to 4

CPU Scheduling Criterias

1. CPU Utilization (Ideally : Max)

- Server OS \rightarrow Utilization - 90%
- Desktop OS \rightarrow Utilization - 70%

2. Through put (Ideally : Max)

- Amount of work done in unit time

3. Waiting time (Ideally : Min)

- time spent by process into ready queue to get access of CPU

4. Response time (Ideally : Min)

- time from arrival of process into ready queue upto first time getting scheduled/executed on CPU

5. Turn Around Time(TAT) (Ideally : Min)

- total time spent by process into RAM (memory)

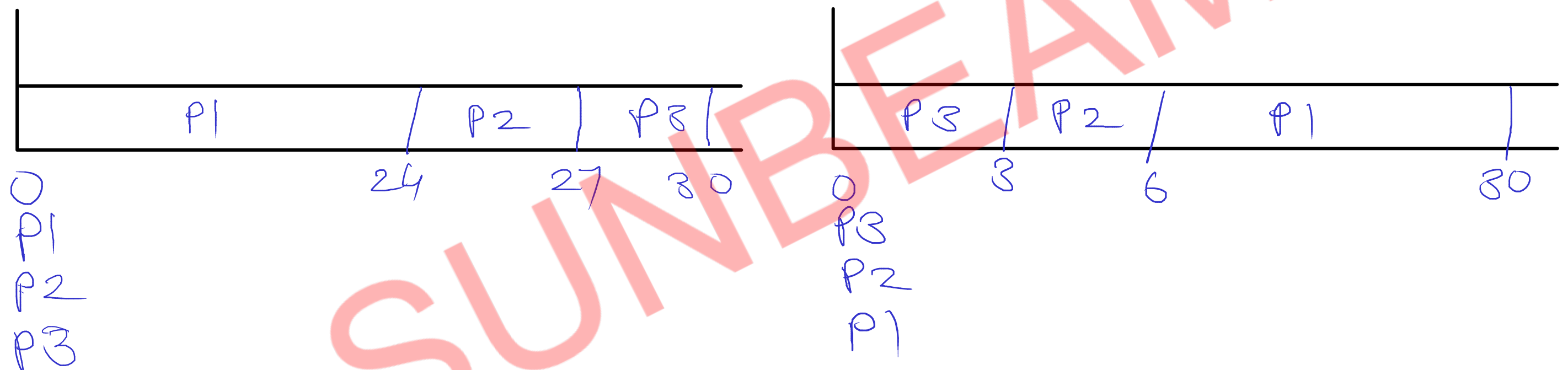
$$TAT = \text{CPU waiting time} + \text{CPU burst} + \text{IO waiting time} + \text{IO burst}$$

FCFS (First Come First Serve) (Non-preemptive)

Process	Arrival	CPU Burst	WT	RT	TAT
P1	0	24	0	0	24
P2	0	3	24	24	27
P3	0	3	27	27	30

Process	Arrival	CPU Burst	WT	RT	TAT
P3	0	3	0	0	3
P2	0	3	3	3	6
P1	0	24	6	6	30

Gantt's chart



Convo effect

- due to arrival of longer process early, all other processes has to wait for longer time.

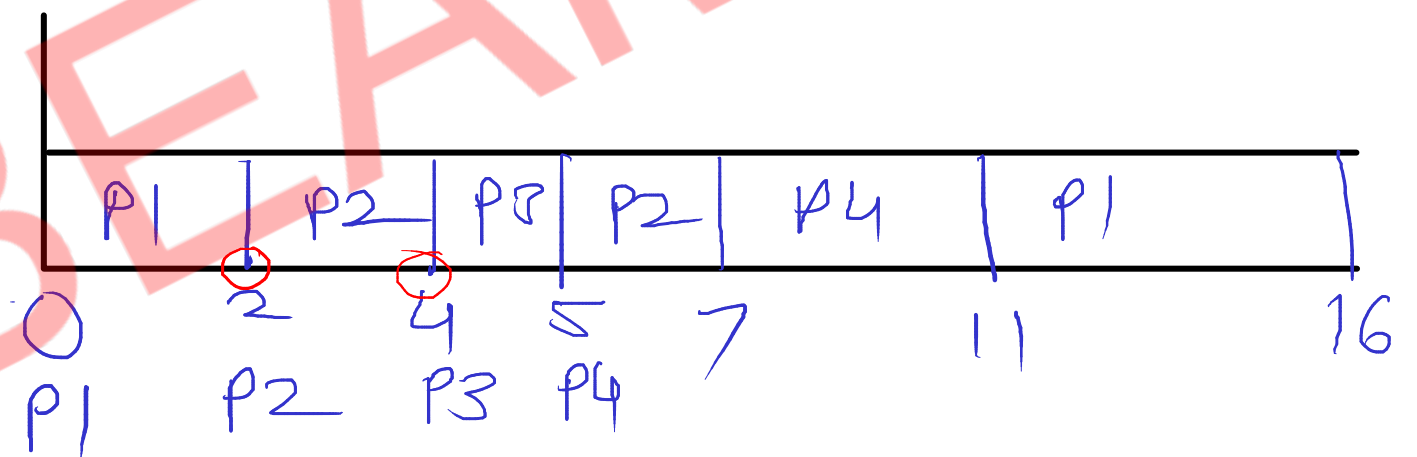
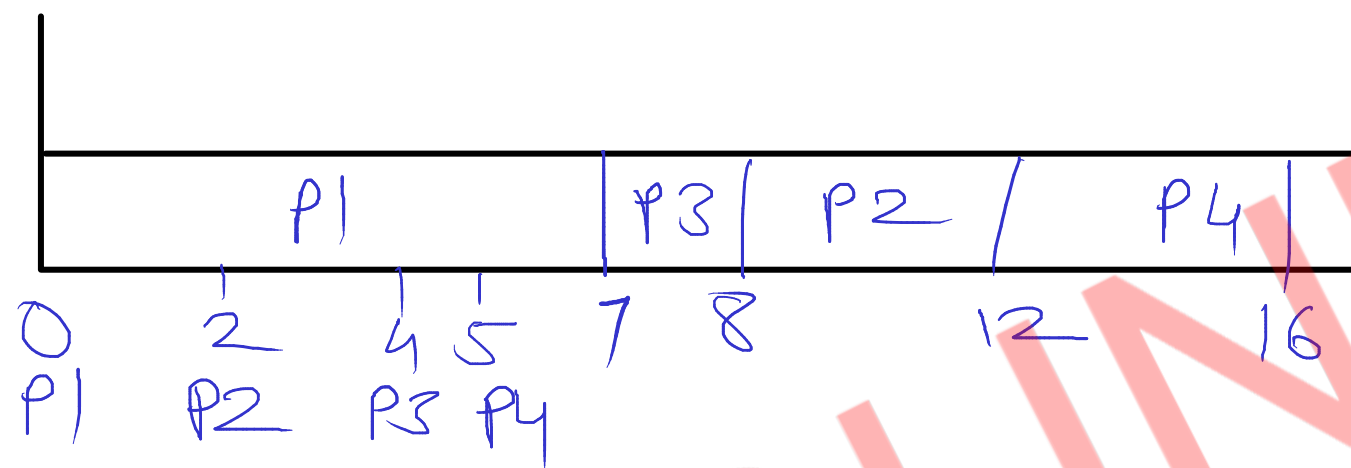
SJF (Shortest Job First)
 (Non-preemptive) (preemptive)
 (Shortest Remaining Time First)

Process	Arrival	CPU Burst
P1	0	7
P2	2	4
P3	4	1
P4	5	4

WT RT TAT
 0 0 7
 6 6 10
 3 3 4
 7 7 11

Process	Arrival	CPU Burst
P1	0	7
P2	2	4
P3	4	1
P4	5	4

remain WT RT TAT
 7.5 X 9 0 16
 4.2 X 1 0 5
 1.0 X 0 0 1
 4 X 2 2 6
 0.5



Starvation:

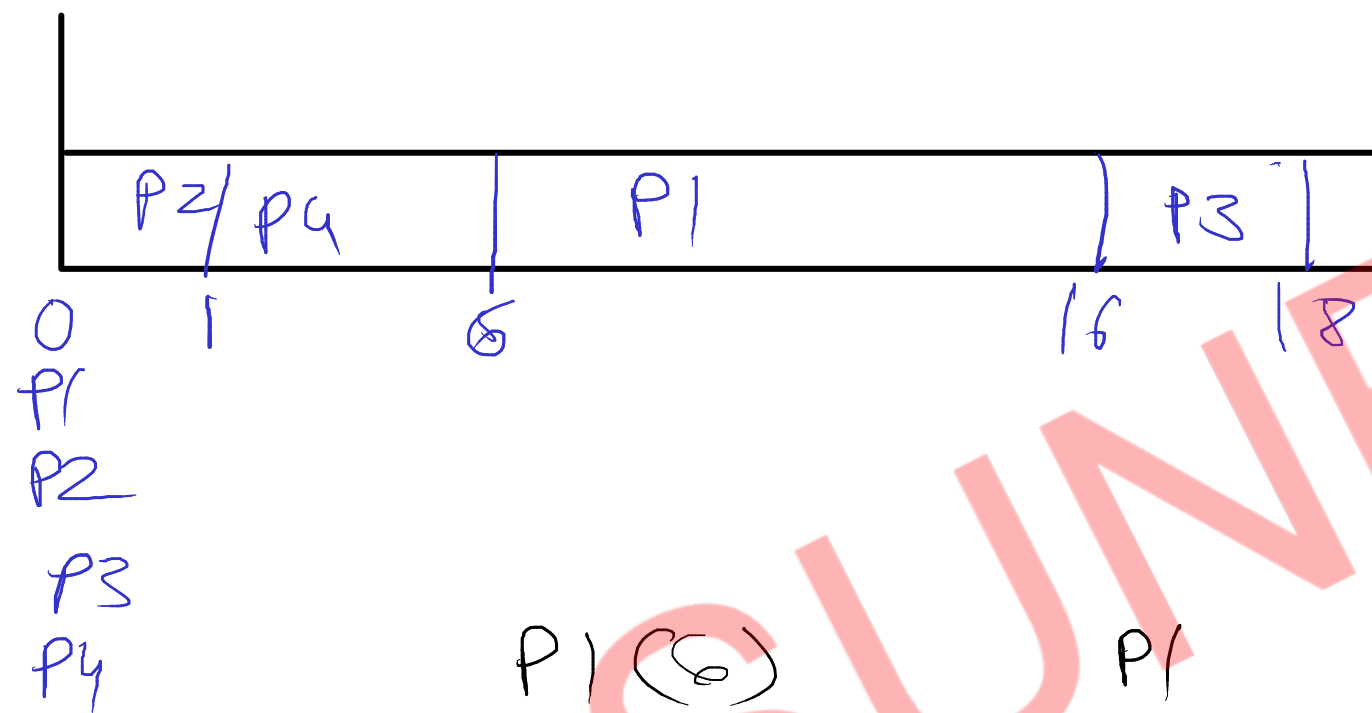
due to longer CPU burst, process is not getting scheduled for longer time.

Priority

(Non-preemptive)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	0	1	1 (H)
P3	0	2	4 (L)
P4	0	5	2

WT RT TAT
 6 6 16
 0 0 1
 16 16 18
 1 1 6



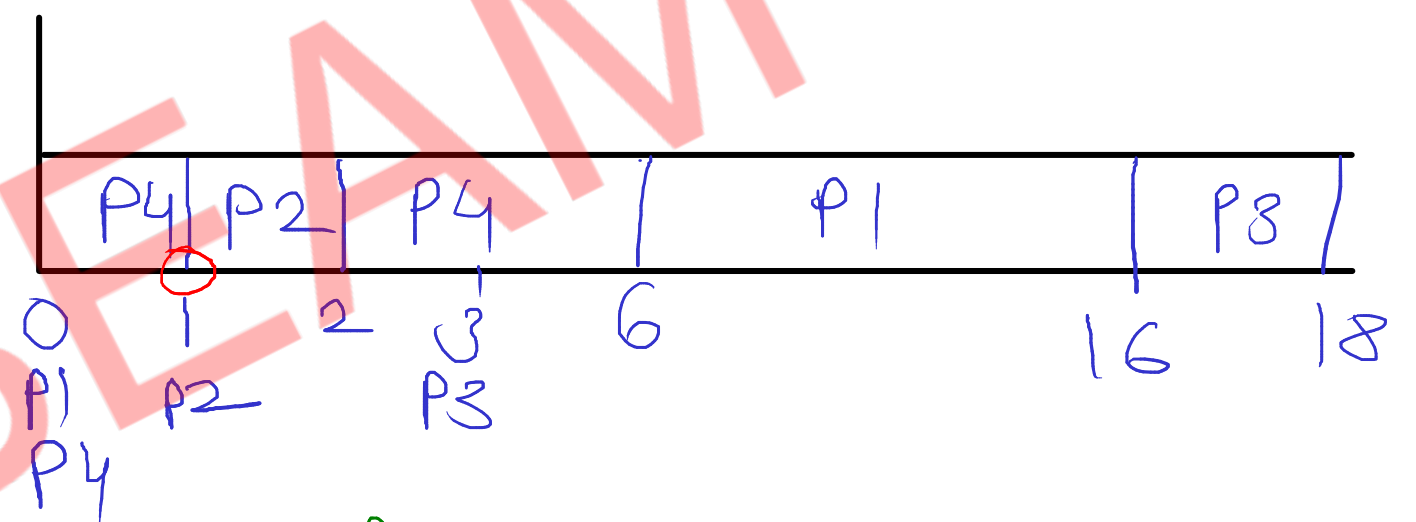
P1 (6)
 P2 (9)
 P3 (7)
 P4 (5)
 P5 (6)
 P6 (6)
 P7 (5)

P1
 P4
 P5
 P7
 P6
 P3
 P2

(preemptive)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	1	1	1
P3	3	2	4
P4	0	5	2

WT RT TAT
 6 6
 0 0
 13 13
 1 0



Starvation:

due to low priority, process has to wait for longer time to get CPU time.

Aging:

increase priority of process gradually

RR (pre-emptive)

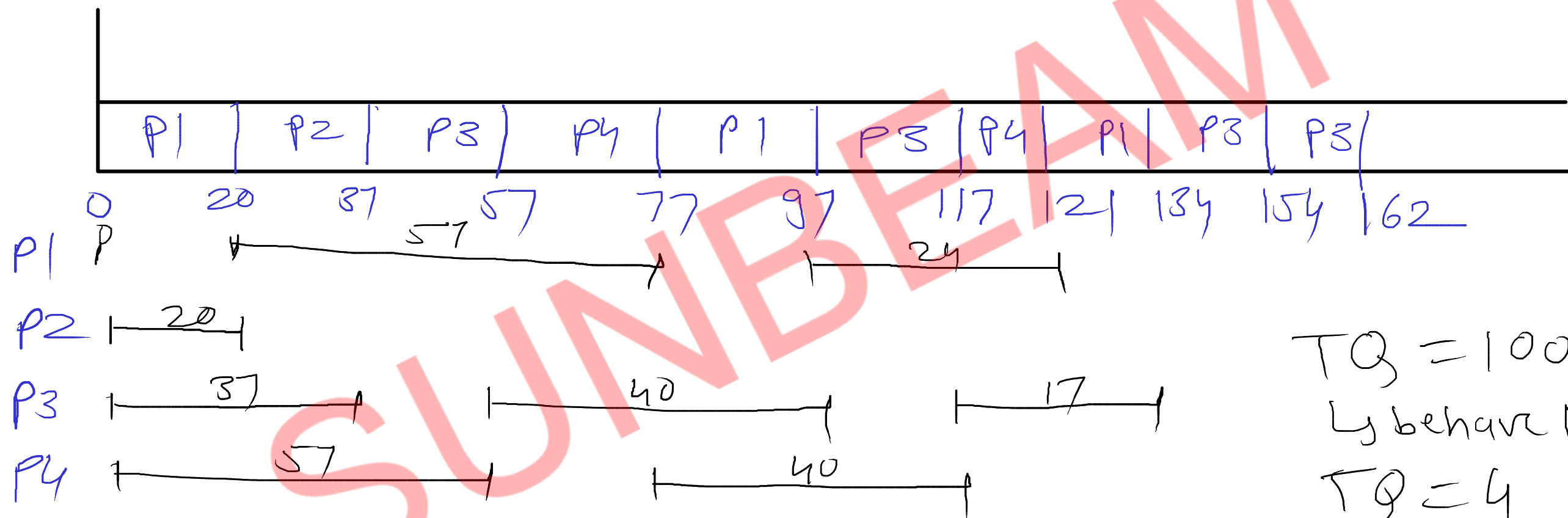
Time Quantum = 20

Process	CPU Burst
P1	53
P2	17
P3	68
P4	24

remain
53, 13X
X
48, 28, 8
4 X

WT
0 + 57 + 24
20
37 + 40 + 7
57 + 40

RT
0
20
37
57



TQ = 100

↳ behave like FCFS

TQ = 4

↳ CPU overhead will increase

Fair Share

- CPU time is divided into time slices (epoch)
- some share of each epoch is given to the processes which are in ready queue.
- share is given to the process on the basis of their priority
- priority of every process is decided by its nice value
- nice values range ---> -20 to +19 (40 values)
 - * -20 - highest priority
 - * +19 - lowest priority

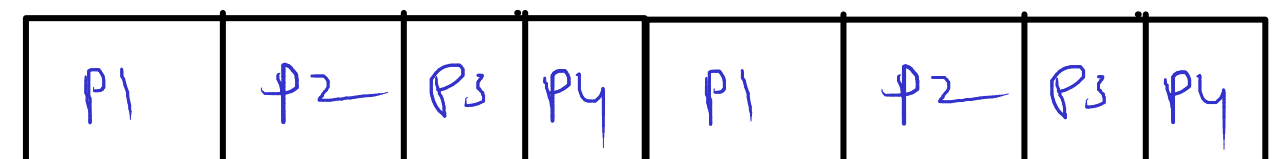
Process	Nice Value
P1	10
P2	10
P3	10
P4	10

Epoch - 100

Process	Nice Value
P1	5
P2	5
P3	10
P4	10



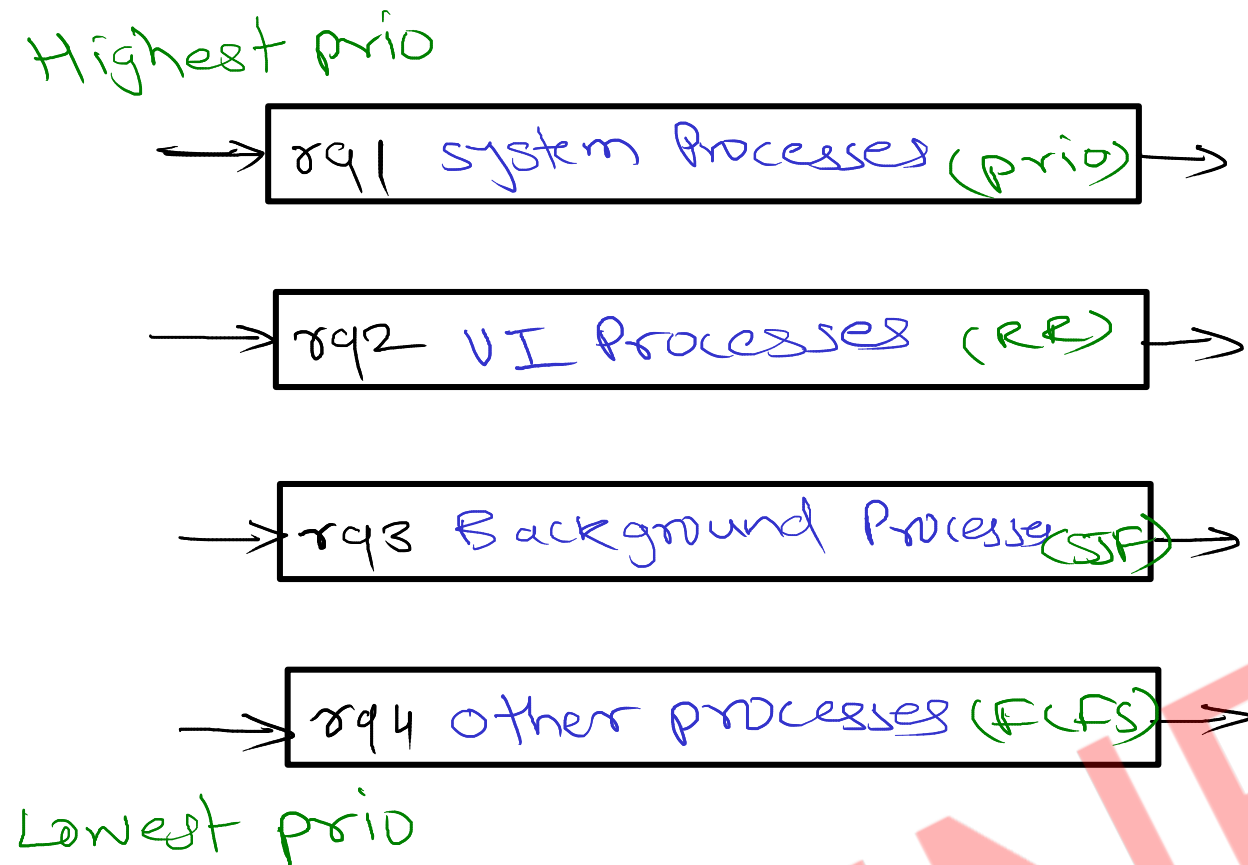
0 100 200



0 100 200

Completely Fair Scheduler (CFS)

Multi Level Ready Queue

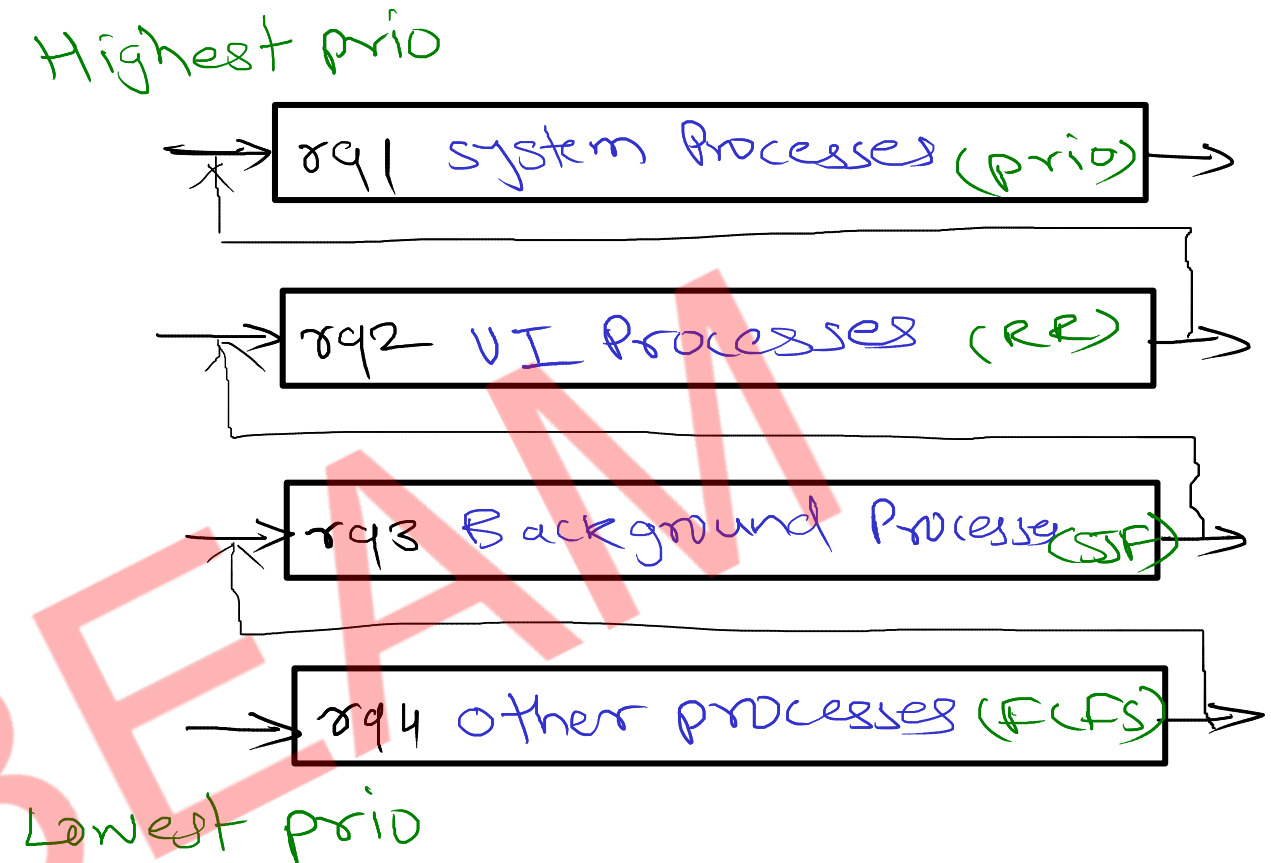


1. SCHED_FIFO
2. SCHED_RR

3. SCHED_OTHERS
4. SCHED_BATCH
5. SCHED_IDLE

CFS

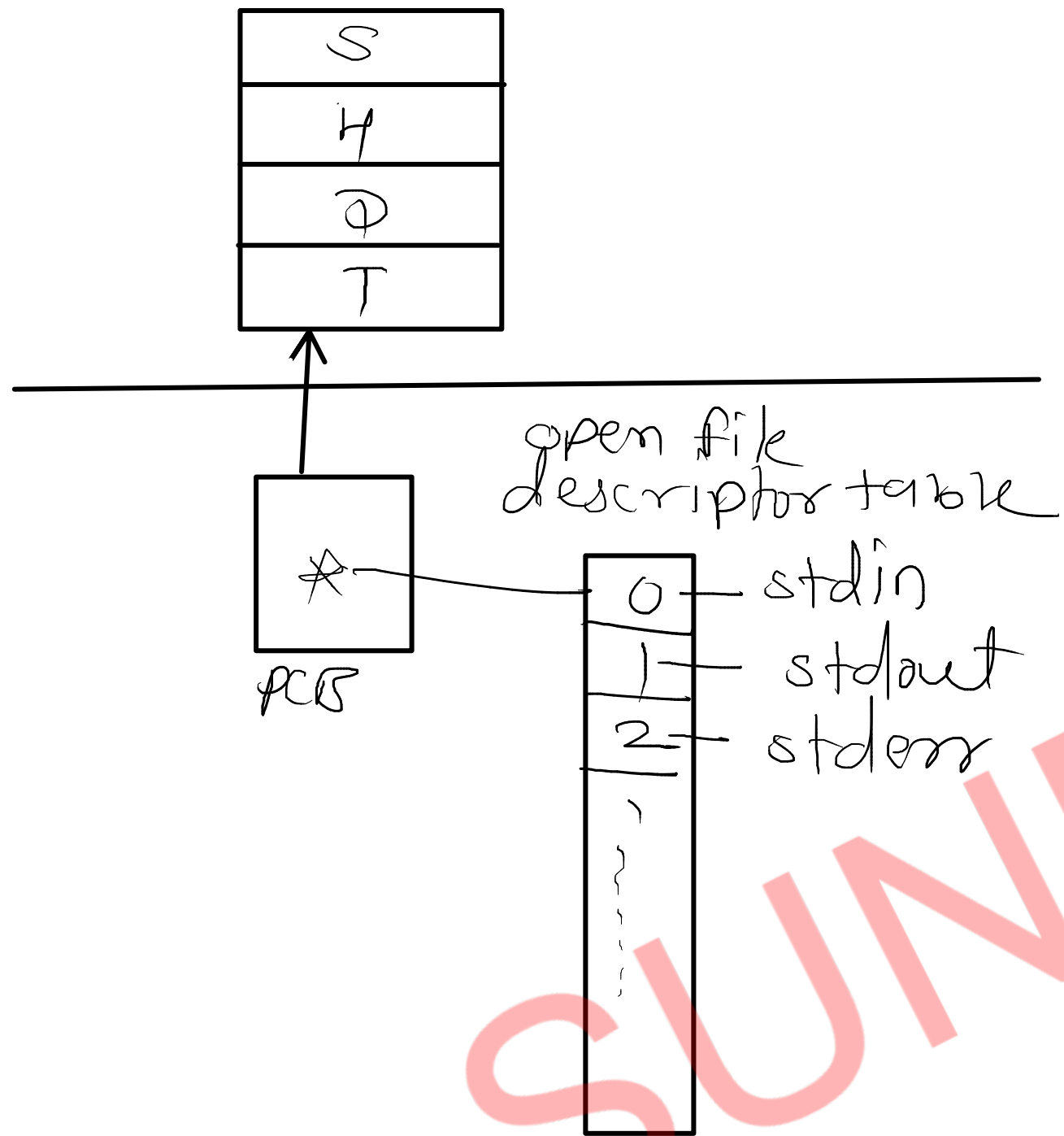
Multi Level Feedback Ready Queue



} real-time classes/policies

} non real time classes/policies

Redirection



Pipe

