

Heap Structures

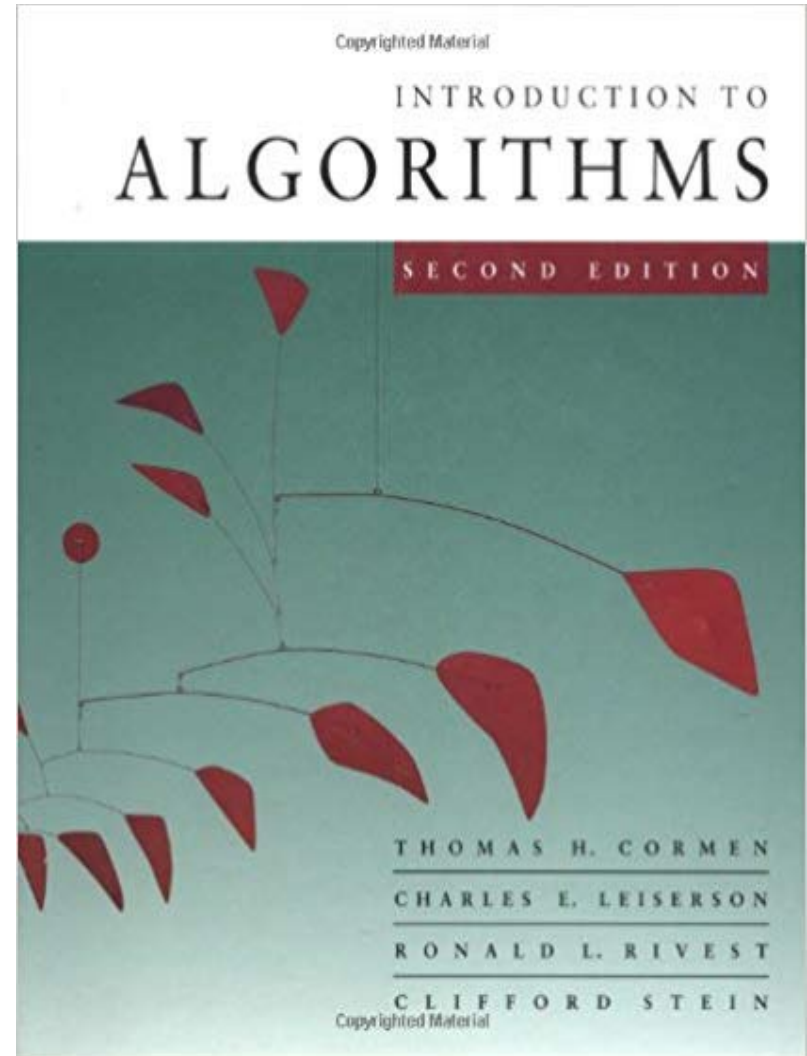
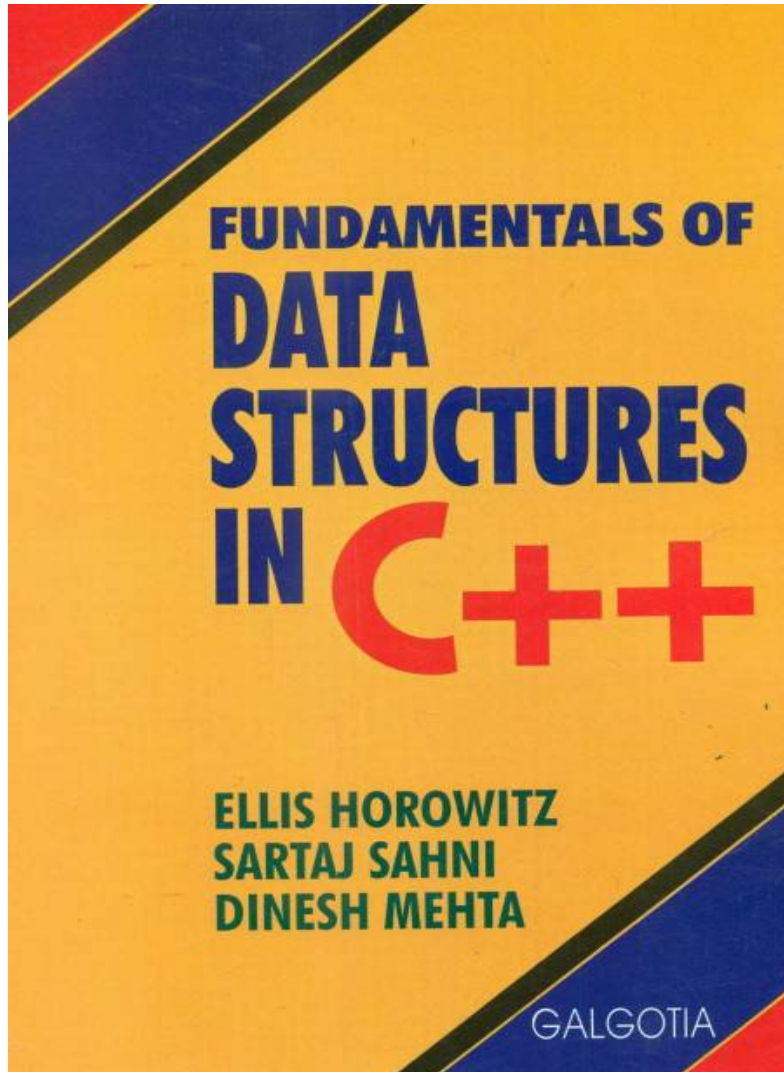
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Courses of Heap Data Structures

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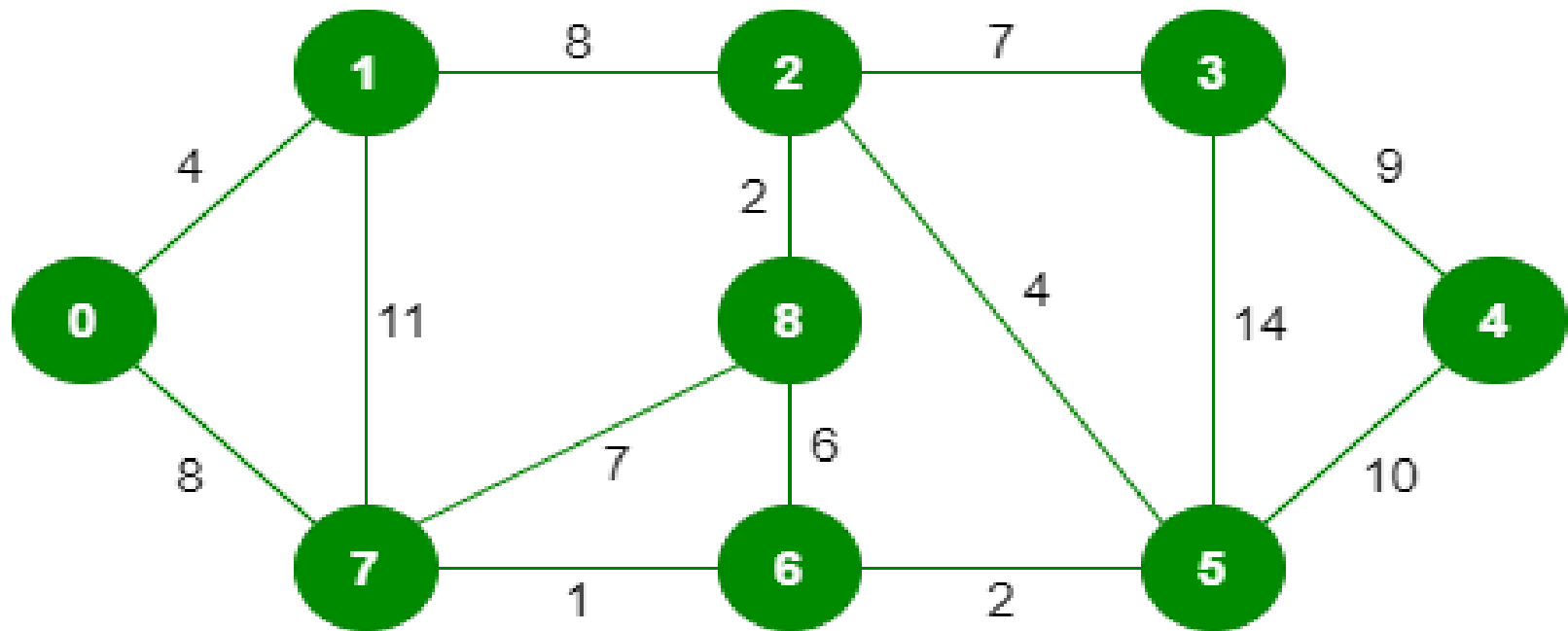
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Heap

Minimum Spanning Tree (MST) problem

- The **Minimum Spanning Tree (MST)** problem is a fundamental problem in graph theory, where the goal is to find a subset of edges in a graph that connects all vertices with the minimum possible total edge weight, forming a tree without cycles. Different types of heaps can be used to efficiently solve the MST problem, primarily through **Prim's Algorithm**.



1. Binary Heap

- Complexity: $O((V + E) \log V)$
- Explanation:
 - Prim's algorithm starts with an arbitrary vertex and grows the MST by selecting the minimum-weight edge connecting a vertex inside the MST to one outside it.
 - In each step, the algorithm uses a **priority queue** (which can be implemented using a binary heap) to select the smallest edge.
 - The binary heap supports **insertion**, **deletion**, and **decrease-key** operations, all in $O(\log V)$ time. Since there are V vertices and up to E edges, the complexity becomes $O((V + E) \log V)$.

Pros: Simple and efficient for most graphs.

Cons: Decrease-key operation is somewhat costly compared to more advanced heaps.

Example:

- Start with a graph, initialize all vertices' keys to infinity, and the starting vertex's key to 0.
- Use the binary heap to extract the vertex with the smallest key, update its adjacent vertices using the decrease-key operation, and repeat until the MST is complete.

2. Fibonacci Heap

- Complexity: $O(E + V \log V)$
- Explanation:
 - A Fibonacci Heap improves the complexity of Prim's algorithm, especially when dealing with dense graphs.
 - The Fibonacci Heap provides **deletion** and **extract-min** operations in $O(\log V)$ time, but crucially, it supports the **decrease-key** operation in **amortized $O(1)$** time. This is beneficial because decrease-key is frequently used in Prim's algorithm to update the keys of adjacent vertices.
 - The overall complexity becomes $O(E + V \log V)$ since the number of decrease-key operations is proportional to the number of edges (E), and there are V extract-min operations.

Pros: Faster decrease-key operation, which significantly reduces running time for dense graphs.

Cons: Fibonacci heaps are more complex to implement and involve higher constant factors, which may not be ideal for smaller or sparse graphs.

Example:

- Start the algorithm with an empty Fibonacci Heap.
- As you add vertices to the MST, update the keys of the neighboring vertices using the decrease-key operation in $O(1)$ amortized time.
- This approach optimizes performance for graphs with a large number of edges (dense graphs).

3. Pairing Heap

- Complexity: $O(E \log V)$
- Explanation:
 - A **Pairing Heap** is another alternative to a binary or Fibonacci heap that is simple to implement and performs well in practice. Like Fibonacci heaps, pairing heaps provide good amortized performance for Prim's algorithm.
 - The **decrease-key** operation in pairing heaps takes **amortized $O(\log V)$** time, which is slower than Fibonacci heaps but faster than binary heaps in many practical cases.
 - The **extract-min** operation also runs in $O(\log V)$ time.

Pros: Easier to implement than Fibonacci heaps and performs well in practice, making it a good compromise between simplicity and performance.

Cons: While faster in practice, the theoretical complexity for decrease-key is worse than Fibonacci heaps.

Example:

- Use the pairing heap to manage the priority queue for Prim's algorithm. Decrease the keys of vertices adjacent to the currently selected vertex. The heap's structure ensures logarithmic time operations for insertion and deletion.

4. d-ary Heap

- Complexity: $O((E \log_d V) + V \log_d V)$
- Explanation:
 - A **d-ary Heap** is a generalization of the binary heap where each node has d children instead of 2. The idea is to reduce the number of levels in the heap, which reduces the number of decrease-key operations needed to reach the root.
 - The **extract-min** operation is $O(\log_d V)$, while **decrease-key** is $O(d \log_d V)$. If d is chosen carefully (typically around the number of edges divided by vertices), this can improve performance, especially in dense graphs.

Pros: Flexible in tuning the value of d to balance the trade-offs between depth and branching factor.

Cons: Decrease-key is slower than Fibonacci heap but can outperform binary heaps depending on d .

Example:

- Prim's algorithm can utilize the d-ary heap by adjusting the branching factor d based on graph density. For denser graphs, a higher d reduces the heap height, improving the performance of extract-min and decrease-key operations.

5. Bucket Heap

- Complexity: $O(E + V)$ (for certain edge weights)
- Explanation:
 - A **Bucket Heap** works well when edge weights are small integers. It divides edges into buckets based on their weights.
 - This can lead to **linear time** performance for Prim's algorithm, i.e., $O(E + V)$, since extracting the minimum edge and performing updates becomes efficient when edge weights are bounded.

Pros: Extremely fast for graphs with small integer weights.

Cons: Not suitable for graphs with arbitrary or large edge weights.

Example:

- If the graph's edge weights are bounded integers, create buckets corresponding to the edge weights.
- Use Prim's algorithm by extracting the smallest bucket and updating the neighboring vertices accordingly.

Summary of Heap Performance in Prim's Algorithm:

Heap Type	Insert	Extract Min	Decrease Key	Overall Complexity (Prim's Algorithm)
Binary Heap	$O(\log V)$	$O(\log V)$	$O(\log V)$	$O((V + E) \log V)$
Fibonacci Heap	$O(1)$	$O(\log V)$	$O(1)$ (amortized)	$O(E + V \log V)$
Pairing Heap	$O(1)$	$O(\log V)$	$O(\log V)$	$O(E \log V)$
d-ary Heap	$O(\log_d V)$	$O(\log_d V)$	$O(d \log_d V)$	$O((E \log_d V) + V \log_d V)$
Bucket Heap	$O(1)$	$O(1)$	$O(1)$	$O(E + V)$

Summary of Heap Performance in Prim's Algorithm:

Heap Type	Insert	Extract Min	Decrease Key	Overall Complexity (Prim's Algorithm)
Binary Heap	$O(\log V)$	$O(\log V)$	$O(\log V)$	$O((V + E) \log V)$
Fibonacci Heap	$O(1)$	$O(\log V)$	$O(1)$ (amortized)	$O(E + V \log V)$
Pairing Heap	$O(1)$	$O(\log V)$	$O(\log V)$	$O(E \log V)$
d-ary Heap	$O(\log_d V)$	$O(\log_d V)$	$O(d \log_d V)$	$O((E \log_d V) + V \log_d V)$
Bucket Heap	$O(1)$	$O(1)$	$O(1)$	$O(E + V)$

Conclusion:

- For **general graphs**, the **Fibonacci heap** provides the best theoretical performance for Prim's algorithm with $O(E + V \log V)$ complexity, but it is complex to implement.
- **Binary heaps** and **pairing heaps** are simpler and more practical for many real-world applications, with pairing heaps being easier to implement than Fibonacci heaps.
- **Bucket heaps** offer excellent performance for graphs with small integer weights, achieving linear time complexity.



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Exercises

Exercises

1. How is a heap used to implement a priority queue?
2. How can a heap be utilized in Dijkstra's algorithm for shortest path finding?
3. What role does a heap play in the heapsort algorithm?
4. How does a heap support the efficient implementation of a median-finding algorithm?
5. How are heaps used in the process of merging k sorted arrays or lists?
6. What is the difference between using a min-heap and a max-heap for job scheduling problems?
7. How can a heap be used to find the k largest (or smallest) elements in an unsorted array?
8. How does a heap help in maintaining the top k elements in a stream of data?
9. How can a heap be used to solve the problem of finding the k-th smallest element in a matrix sorted by rows and columns?
10. How can heaps be used in graph algorithms like Prim's Minimum Spanning Tree algorithm?

Exercises

11. How can a heap be applied to solve the Huffman coding problem in data compression?
12. What is the application of heaps in the construction of interval trees?
13. How can heaps help optimize a real-time event simulation system?
14. How does a heap improve the performance of cache management systems?
15. How can heaps be used to maintain a dynamic set of median values as elements are inserted or removed?
16. How does a heap assist in memory management, particularly in garbage collection algorithms?
17. How can a heap be used in network packet scheduling?
18. How can heaps solve the problem of scheduling tasks with varying deadlines and priorities?
19. How can a heap be used to solve the "meeting rooms" problem in interval scheduling?
20. How do heaps assist in the implementation of external sorting algorithms for large data sets that do not fit in memory?

Exercises

21. How does a Fibonacci heap improve the efficiency of graph algorithms compared to binary heaps?
22. In what scenarios would a binomial heap be more advantageous over a binary heap?
23. How is a heap used in the implementation of an event-driven simulation, such as in discrete event simulation systems?
24. How can a heap be applied to solve the skyline problem in computational geometry?
25. How does a heap help solve the problem of finding the closest points to a given point in a 2D plane?
26. How can heaps be used in the context of paging systems for virtual memory management?
27. How is a heap used in online algorithms, such as online sorting or online interval scheduling?
28. How can heaps be utilized to implement a dynamic array that supports fast retrieval of the smallest or largest element?
29. How does a heap assist in optimizing search operations in best-first search algorithms, like A* search?
30. How can heaps help in solving the range minimum query (RMQ) problem in arrays?

Exercises

31. How can a heap be applied to manage a pool of active threads in a multi-threaded system?
32. How does a heap assist in handling buffer overflow attacks in memory allocation?
33. How can heaps be applied to load balancing in distributed systems?
34. How are heaps used in sensor networks to track the top k measurements in real-time?
35. How does a heap optimize task assignment in a cloud computing environment?
36. How can heaps assist in managing time-stamped data, such as in a real-time data streaming application?
37. How are heaps applied in search engines for ranking top search results efficiently?
38. How can heaps improve the performance of packet routing algorithms in a network?
39. How can a heap be used to solve the interval cover problem in computational geometry?
40. How does a heap aid in implementing efficient algorithms for decoding topological graphs?

Exercises

41. What are the time complexities of insertion and removal of elements on a heap?
42. In the description of reheapification downward, we specified that the out-of-place element must be swapped with the larger of its two children. Can we swap with the smaller child instead?
43. Start with an empty heap and enter 10 elements with priorities 1 through 10. Draw the resulting heap (max-heap).
44. Remove three elements from the heap you created in the above heap. Draw the resulting heap.
45. How to create a max-min heap in linear time?
46. Find the Top K Frequent: Elements Given the array [1, 1, 1, 3, 2, 2, 4], find the 2 most frequent elements using a Max-Heap.
47. Kth Largest Element in a Stream of Data: Given a stream of numbers arriving in this order: 10, 7, 4, 3, 20, 15, find the 3rd largest element after each insertion using a Min-Heap.
48. Minimum Cost to Connect Ropes: Given ropes of lengths [4, 3, 2, 6], find the minimum cost to connect them into one rope. Use a Min-Heap to combine the ropes step by step, and show the total cost.
49. Sort Nearly Sorted (K-Sorted) Array: Given the array [6, 5, 3, 2, 8, 10, 9], where each element is at most 3 positions away from its correct place, use a Min-Heap to sort it efficiently. Show the heap after each step.

Exercises

50. Give an example of a min-heap such that there exist an element x in level i , y in level j , $i < j$ and $x > y$ in the min-heap.
51. Is it possible to construct an n -element Deap in $O(n)$ time?
52. Justify true or false. For a Deap, there exist an element x in min-heap at level i , y in max heap at level j , $i > j$ and $x > y$.
53. Construct a Deap having l levels (root is at level 0) such that when extract max operation is initiated, after replacing the max element with the last leaf, $2\lfloor l/2 \rfloor - 1$ exchanges are required to set right the deap. Does there exist a deap with more than the $2\lfloor l/2 \rfloor - 1$ exchanges? Give a proper justification.
54. Find the 5th largest element in an array of 1 million elements. Show how a Min-Heap can be used to maintain the top 5 elements efficiently.
55. You have tasks arriving dynamically with priorities: (Task1, 8), (Task2, 4), (Task3, 10), (Task4, 5). Use a Max-Heap to schedule the tasks, ensuring the highest-priority task is always executed next. Show the heap after each task insertion and extraction.