Chapter 17
Background Debug Module (BDM) Block Description

17.1 Introduction

This section describes the functionality of the background debug module (BDM) sub-block of the HCS12 core platform.

A block diagram of the BDM is shown in Figure 17-1.
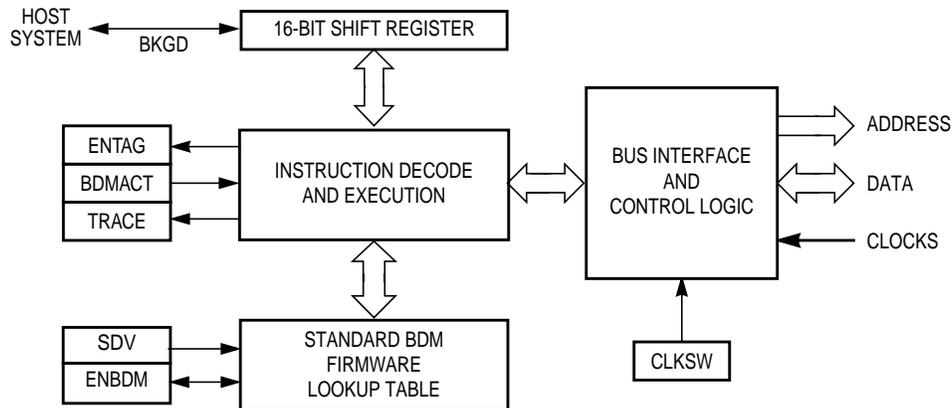
Figure 17-1. BDM Block Diagram

The background debug module (BDM) sub-block is a single-wire, background debug system implemented in on-chip hardware for minimal CPU intervention. All interfacing with the BDM is done via the BKGD pin.

BDMV4 has enhanced capability for maintaining synchronization between the target and host while allowing more flexibility in clock rates. This includes a sync signal to show the clock rate and a handshake signal to indicate when an operation is complete. The system is backwards compatible with older external interfaces.

17.1.1 Features

- Single-wire communication with host development system
- BDMV4 (and BDM2): Enhanced capability for allowing more flexibility in clock rates
- BDMV4: SYNC command to determine communication rate
- BDMV4: GO_UNTIL command
- BDMV4: Hardware handshake protocol to increase the performance of the serial communication
- Active out of reset in special single-chip mode
- Nine hardware commands using free cycles, if available, for minimal CPU intervention
- Hardware commands not requiring active BDM
- 15 firmware commands execute from the standard BDM firmware lookup table

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- Instruction tagging capability
- Software control of BDM operation during wait mode
- Software selectable clocks
- When secured, hardware commands are allowed to access the register space in Special Single-Chip mode, if the FLASH and EEPROM erase tests fail.

17.1.2   Modes of Operation

BDM is available in all operating modes but must be enabled before firmware commands are executed. Some system peripherals may have a control bit which allows suspending the peripheral function during background debug mode.

17.1.2.1   Regular Run Modes

All of these operations refer to the part in run mode. The BDM does not provide controls to conserve power during run mode.

- Normal operation

  General operation of the BDM is available and operates the same in all normal modes.
- Special single-chip mode

  In special single-chip mode, background operation is enabled and active out of reset. This allows programming a system with blank memory.
- Special peripheral mode

  BDM is enabled and active immediately out of reset. BDM can be disabled by clearing the BDMACT bit in the BDM status (BDMSTS) register. The BDM serial system should not be used in special peripheral mode.
- Emulation modes

  General operation of the BDM is available and operates the same as in normal modes.

17.1.2.2   Secure Mode Operation

If the part is in secure mode, the operation of the BDM is reduced to a small subset of it's regular run mode operation. Secure operation prevents access to FLASH or EEPROM other than allowing erasure.

17.1.2.3   Low-Power Modes

- Wait mode

  The BDM cannot be used in wait mode if the system disables the clocks to the BDM.

  There is a clearing mechanism associated with the WAIT instruction when the clocks to the BDM (CPU core platform) are disabled. As the clocks restart from wait mode, the BDM receives a soft reset (clearing any command in progress) and the ACK function will be disabled. This is a change from previous BDM modules.
- Stop mode

The BDM is completely shutdown in stop mode.

There is a clearing mechanism associated with the STOP instruction. STOP must be enabled and the part must go into stop mode for this to occur. As the clocks restart from stop mode, the BDM receives a soft reset (clearing any command in progress) and the ACK function will be disabled. This is a change from previous BDM modules.

17.2   External Signal Description

A single-wire interface pin is used to communicate with the BDM system. Two additional pins are used for instruction tagging. These pins are part of the multiplexed external bus interface (MEBI) sub-block and all interfacing between the MEBI and BDM is done within the core interface boundary. Functional descriptions of the pins are provided below for completeness.

- BKGD — Background interface pin
- $\overline{\text{TAGHI}}$ — High byte instruction tagging pin
- $\overline{\text{TAGLO}}$ — Low byte instruction tagging pin
- BKGD and $\overline{\text{TAGHI}}$ share the same pin.
- $\overline{\text{TAGLO}}$ and $\overline{\text{LSTRB}}$ share the same pin.

NOTE

Generally these pins are shared as described, but it is best to check the device overview chapter to make certain. All MCUs at the time of this writing have followed this pin sharing scheme.

17.2.1   Background Interface Pin (BKGD)

Debugging control logic communicates with external devices serially via the single-wire background interface pin (BKGD). During reset, this pin is a mode select input which selects between normal and special modes of operation. After reset, this pin becomes the dedicated serial interface pin for the background debug mode.

17.2.2   High Byte Instruction Tagging Pin ($\overline{\text{TAGHI}}$)

This pin is used to tag the high byte of an instruction. When instruction tagging is on, a logic 0 at the falling edge of the external clock (ECLK) tags the high half of the instruction word being read into the instruction queue.

17.2.3   Low Byte Instruction Tagging Pin ($\overline{\text{TAGLO}}$)

This pin is used to tag the low byte of an instruction. When instruction tagging is on and low strobe is enabled, a logic 0 at the falling edge of the external clock (ECLK) tags the low half of the instruction word being read into the instruction queue.

17.3   Memory Map and Register Descriptions

A summary of the registers associated with the BDM is shown in Table 17-1. Registers are accessed by host-driven communications to the BDM hardware using READ_BD and WRITE_BD commands. Detailed descriptions of the registers and associated bits are given in the subsections that follow..

Table 17-1. BDM Register Map Summary

Address	Register Name		Bit 7	6	5	4	3	2	1	Bit 0
$FF00	Reserved	Read:	X	X	X	X	X	X	0	0
		Write:								
$FF01	BDMSTS	Read:	ENBDM	BDMACT	ENTAG	SDV	TRACE	CLKSW	UNSEC	0
		Write:	ENBDM		ENTAG			CLKSW		
$FF02	Reserved	Read:	X	X	X	X	X	X	X	X
		Write:								
$FF03	Reserved	Read:	X	X	X	X	X	X	X	X
		Write:								
$FF04	Reserved	Read:	X	X	X	X	X	X	X	X
		Write:								
$FF05	Reserved	Read:	X	X	X	X	X	X	X	X
		Write:								
$FF06	BDMCCR	Read:	CCR7	CCR6	CCR5	CCR4	CCR3	CCR2	CCR1	CCR0
		Write:								
$FF07	BDMINR	Read:	0	REG14	REG13	REG12	REG11	0	0	0
		Write:								
$FF08	Reserved	Read:	0	0	0	0	0	0	0	0
		Write:								
$FF09	Reserved	Read:	0	0	0	0	0	0	0	0
		Write:								

▨	= Unimplemented, Reserved	▮	= Implemented (do not alter)
X	= Indeterminate	0	= Always read zero

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Table 17-1. BDM Register Map Summary

Address	Register Name		Bit 7	6	5	4	3	2	1	Bit 0
$FF0A	Reserved	Read:	X	X	X	X	X	X	X	X
		Write:								
$FF0B	Reserved	Read:	X	X	X	X	X	X	X	X
		Write:								

	= Unimplemented, Reserved		= Implemented (do not alter)
X	= Indeterminate	0	= Always read zero

17.3.1 BDM Status Register (BDMSTS)

Register Address: $FF01

	7	6	5	4	3	2	1	0
R	ENBDM	BDMACT	ENTAG	SDV	TRACE	CLKSW	UNSEC	0
W								
Reset:								
Special single-chip mode:	1	1	0	0	0	0	0	0
Special peripheral mode:	0	1	0	0	0	0	0	0
All other modes:	0	0	0	0	0	0	0	0

	= Unimplemented or Reserved		= Implemented (do not alter)

Figure 17-2. BDM Status Register (BDMSTS)

Read: All modes through BDM operation

Write: All modes but subject to the following:

— BDMACT can only be set by BDM hardware upon entry into BDM. It can only be cleared by the standard BDM firmware lookup table upon exit from BDM active mode.

— CLKSW can only be written via BDM hardware or standard BDM firmware write commands.

— All other bits, while writable via BDM hardware or standard BDM firmware write commands, should only be altered by the BDM hardware or standard firmware lookup table as part of BDM command execution.

— ENBDM should only be set via a BDM hardware command if the BDM firmware commands are needed. (This does not apply in special single-chip mode).

ENBDM — Enable BDM

This bit controls whether the BDM is enabled or disabled. When enabled, BDM can be made active to allow firmware commands to be executed. When disabled, BDM cannot be made active but BDM hardware commands are still allowed.

    1 = BDM enabled
    0 = BDM disabled

NOTE

ENBDM is set by the firmware immediately out of reset in special single-chip mode. In secure mode, this bit will not be set by the firmware until after the EEPROM and FLASH erase verify tests are complete.

BDMACT — BDM active status

This bit becomes set upon entering BDM. The standard BDM firmware lookup table is then enabled and put into the memory map. BDMACT is cleared by a carefully timed store instruction in the standard BDM firmware as part of the exit sequence to return to user code and remove the BDM memory from the map.

    1 = BDM active
    0 = BDM not active

ENTAG — Tagging enable

This bit indicates whether instruction tagging in enabled or disabled. It is set when the TAGGO command is executed and cleared when BDM is entered. The serial system is disabled and the tag function enabled 16 cycles after this bit is written. BDM cannot process serial commands while tagging is active.

    1 = Tagging enabled
    0 = Tagging not enabled or BDM active

SDV — Shift data valid

This bit is set and cleared by the BDM hardware. It is set after data has been transmitted as part of a firmware read command or after data has been received as part of a firmware write command. It is cleared when the next BDM command has been received or BDM is exited. SDV is used by the standard BDM firmware to control program flow execution.

    1 = Data phase of command is complete
    0 = Data phase of command not complete

TRACE — TRACE1 BDM firmware command is being executed

This bit gets set when a BDM TRACE1 firmware command is first recognized. It will stay set as long as continuous back-to-back TRACE1 commands are executed. This bit will get cleared when the next command that is not a TRACE1 command is recognized.

    1 = TRACE1 command is being executed
    0 = TRACE1 command is not being executed

CLKSW — Clock switch

The CLKSW bit controls which clock the BDM operates with. It is only writable from a hardware BDM command. A 150 cycle delay at the clock speed that is active during the data portion of the command will occur before the new clock source is guaranteed to be active. The start of the next BDM command uses the new clock for timing subsequent BDM communications.

Table 17-2 shows the resulting BDM clock source based on the CLKSW and the PLLSEL (Pll select from the clock and reset generator) bits.

Table 17-2. BDM Clock Sources

PLLSEL	CLKSW	BDMCLK
0	0	Bus clock
0	1	Bus clock
1	0	Alternate clock (refer to the device overview chapter to determine the alternate clock source)
1	1	Bus clock dependent on the PLL

NOTE

The BDM alternate clock source can only be selected when CLKSW = 0 and PLLSEL = 1. The BDM serial interface is now fully synchronized to the alternate clock source, when enabled. This eliminates frequency restriction on the alternate clock which was required on previous versions. Refer to the device overview section to determine which clock connects to the alternate clock source input.

NOTE

If the acknowledge function is turned on, changing the CLKSW bit will cause the ACK to be at the new rate for the write command which changes it.

UNSEC — Unsecure

This bit is only writable in special single-chip mode from the BDM secure firmware and always gets reset to zero. It is in a zero state as secure mode is entered so that the secure BDM firmware lookup table is enabled and put into the memory map along with the standard BDM firmware lookup table.

The secure BDM firmware lookup table verifies that the on-chip EEPROM and FLASH EEPROM are erased. This being the case, the UNSEC bit is set and the BDM program jumps to the start of the standard BDM firmware lookup table and the secure BDM firmware lookup table is turned off. If the erase test fails, the UNSEC bit will not be asserted.

    1 = System is in a unsecured mode
    0 = System is in a secured mode

NOTE

When UNSEC is set, security is off and the user can change the state of the secure bits in the on-chip FLASH EEPROM. Note that if the user does not change the state of the bits to "unsecured" mode, the system will be secured again when it is next taken out of reset.

17.3.2   BDM CCR Holding Register (BDMCCR)

Register Address: $FF06

	7	6	5	4	3	2	1	0
R W	CCR7	CCR6	CCR5	CCR4	CCR3	CCR2	CCR1	CCR0
Reset:	0	0	0	0	0	0	0	0

Figure 17-3. BDM CCR Holding Register (BDMCCR)

Read: All modes

Write: All modes

NOTE

When BDM is made active, the CPU stores the value of the CCR register in the BDMCCR register. However, out of special single-chip reset, the BDMCCR is set to $D8 and not $D0 which is the reset value of the CCR register.

When entering background debug mode, the BDM CCR holding register is used to save the contents of the condition code register of the user's program. It is also used for temporary storage in the standard BDM firmware mode. The BDM CCR holding register can be written to modify the CCR value.

17.3.3   BDM Internal Register Position Register (BDMINR)

Register Address: $FF07

	7	6	5	4	3	2	1	0
R	0	REG14	REG13	REG12	REG11	0	0	0
W								
Reset:	0	0	0	0	0	0	0	0

        = Unimplemented or Reserved

Figure 17-4. BDM Internal Register Position (BDMINR)

Read: All modes

Write: Never

REG14–REG11 — Internal register map position

These four bits show the state of the upper five bits of the base address for the system's relocatable register block. BDMINR is a shadow of the INITRG register which maps the register block to any 2K byte space within the first 32K bytes of the 64K byte address space.

17.4 Functional Description

The BDM receives and executes commands from a host via a single wire serial interface. There are two types of BDM commands, namely, hardware commands and firmware commands.

Hardware commands are used to read and write target system memory locations and to enter active background debug mode, see Section 17.4.3, "BDM Hardware Commands." Target system memory includes all memory that is accessible by the CPU.

Firmware commands are used to read and write CPU resources and to exit from active background debug mode, see Section 17.4.4, "Standard BDM Firmware Commands." The CPU resources referred to are the accumulator (D), X index register (X), Y index register (Y), stack pointer (SP), and program counter (PC).

Hardware commands can be executed at any time and in any mode excluding a few exceptions as highlighted, see Section 17.4.3, "BDM Hardware Commands." Firmware commands can only be executed when the system is in active background debug mode (BDM).

17.4.1 Security

If the user resets into special single-chip mode with the system secured, a secured mode BDM firmware lookup table is brought into the map overlapping a portion of the standard BDM firmware lookup table. The secure BDM firmware verifies that the on-chip EEPROM and FLASH EEPROM are erased. This being the case, the UNSEC bit will get set. The BDM program jumps to the start of the standard BDM firmware and the secured mode BDM firmware is turned off and all BDM commands are allowed. If the EEPROM or FLASH do not verify as erased, the BDM firmware sets the ENBDM bit, without asserting UNSEC, and the firmware enters a loop. This causes the BDM hardware commands to become enabled, but does not enable the firmware commands. This allows the BDM hardware to be used to erase the EEPROM and FLASH. After execution of the secure firmware, regardless of the results of the erase tests, the CPU registers, INITEE and PPAGE, will no longer be in their reset state.

17.4.2 Enabling and Activating BDM

The system must be in active BDM to execute standard BDM firmware commands. BDM can be activated only after being enabled. BDM is enabled by setting the ENBDM bit in the BDM status (BDMSTS) register. The ENBDM bit is set by writing to the BDM status (BDMSTS) register, via the single-wire interface, using a hardware command such as WRITE_BD_BYTE.

After being enabled, BDM is activated by one of the following[1]:

- Hardware BACKGROUND command
- BDM external instruction tagging mechanism

1. BDM is enabled and active immediately out of special single-chip reset.

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- CPU BGND instruction
- Breakpoint sub-block's force or tag mechanism[1]

When BDM is activated, the CPU finishes executing the current instruction and then begins executing the firmware in the standard BDM firmware lookup table. When BDM is activated by the breakpoint sub-block, the type of breakpoint used determines if BDM becomes active before or after execution of the next instruction.

NOTE

If an attempt is made to activate BDM before being enabled, the CPU resumes normal instruction execution after a brief delay. If BDM is not enabled, any hardware BACKGROUND commands issued are ignored by the BDM and the CPU is not delayed.

In active BDM, the BDM registers and standard BDM firmware lookup table are mapped to addresses $FF00 to $FFFF. BDM registers are mapped to addresses $FF00 to $FF07. The BDM uses these registers which are readable anytime by the BDM. However, these registers are not readable by user programs.

17.4.3 BDM Hardware Commands

Hardware commands are used to read and write target system memory locations and to enter active background debug mode. Target system memory includes all memory that is accessible by the CPU such as on-chip RAM, EEPROM, FLASH EEPROM, I/O and control registers, and all external memory.

Hardware commands are executed with minimal or no CPU intervention and do not require the system to be in active BDM for execution, although, they can still be executed in this mode. When executing a hardware command, the BDM sub-block waits for a free CPU bus cycle so that the background access does not disturb the running application program. If a free cycle is not found within 128 clock cycles, the CPU is momentarily frozen so that the BDM can steal a cycle. When the BDM finds a free cycle, the operation does not intrude on normal CPU operation provided that it can be completed in a single cycle. However, if an operation requires multiple cycles the CPU is frozen until the operation is complete, even though the BDM found a free cycle.

The BDM hardware commands are listed in Table 17-3.

Table 17-3. Hardware Commands

Command	Opcode (hex)	Data	Description
BACKGROUND	90	None	Enter background mode if firmware is enabled. If enabled, an ACK will be issued when the part enters active background mode.
ACK_ENABLE	D5	None	Enable handshake. Issues an ACK pulse after the command is executed.
ACK_DISABLE	D6	None	Disable handshake. This command does not issue an ACK pulse.
READ_BD_BYTE	E4	16-bit address 16-bit data out	Read from memory with standard BDM firmware lookup table in map. Odd address data on low byte; even address data on high byte.

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1. This method is only available on systems that have a a breakpoint or a debug sub-block.

Table 17-3. Hardware Commands (Continued)

Command	Opcode (hex)	Data	Description
READ_BD_WORD	EC	16-bit address 16-bit data out	Read from memory with standard BDM firmware lookup table in map. Must be aligned access.
READ_BYTE	E0	16-bit address 16-bit data out	Read from memory with standard BDM firmware lookup table out of map. Odd address data on low byte; even address data on high byte.
READ_WORD	E8	16-bit address 16-bit data out	Read from memory with standard BDM firmware lookup table out of map. Must be aligned access.
WRITE_BD_BYTE	C4	16-bit address 16-bit data in	Write to memory with standard BDM firmware lookup table in map. Odd address data on low byte; even address data on high byte.
WRITE_BD_WORD	CC	16-bit address 16-bit data in	Write to memory with standard BDM firmware lookup table in map. Must be aligned access.
WRITE_BYTE	C0	16-bit address 16-bit data in	Write to memory with standard BDM firmware lookup table out of map. Odd address data on low byte; even address data on high byte.
WRITE_WORD	C8	16-bit address 16-bit data in	Write to memory with standard BDM firmware lookup table out of map. Must be aligned access.

NOTE:
> If enabled, ACK will occur when data is ready for transmission for all BDM READ commands and will occur after the write is complete for all BDM WRITE commands.

The READ_BD and WRITE_BD commands allow access to the BDM register locations. These locations are not normally in the system memory map but share addresses with the application in memory. To distinguish between physical memory locations that share the same address, BDM memory resources are enabled just for the READ_BD and WRITE_BD access cycle. This allows the BDM to access BDM locations unobtrusively, even if the addresses conflict with the application memory map.

17.4.4   Standard BDM Firmware Commands

Firmware commands are used to access and manipulate CPU resources. The system must be in active BDM to execute standard BDM firmware commands, see Section 17.4.2, "Enabling and Activating BDM." Normal instruction execution is suspended while the CPU executes the firmware located in the standard BDM firmware lookup table. The hardware command BACKGROUND is the usual way to activate BDM.

As the system enters active BDM, the standard BDM firmware lookup table and BDM registers become visible in the on-chip memory map at $FF00–$FFFF, and the CPU begins executing the standard BDM firmware. The standard BDM firmware watches for serial commands and executes them as they are received.

The firmware commands are shown in Table 17-4.

Table 17-4. Firmware Commands

Command[1]	Opcode (hex)	Data	Description
READ_NEXT	62	16-bit data out	Increment X by 2 (X = X + 2), then read word X points to.
READ_PC	63	16-bit data out	Read program counter.
READ_D	64	16-bit data out	Read D accumulator.
READ_X	65	16-bit data out	Read X index register.
READ_Y	66	16-bit data out	Read Y index register.
READ_SP	67	16-bit data out	Read stack pointer.
WRITE_NEXT	42	16-bit data in	Increment X by 2 (X = X + 2), then write word to location pointed to by X.
WRITE_PC	43	16-bit data in	Write program counter.
WRITE_D	44	16-bit data in	Write D accumulator.
WRITE_X	45	16-bit data in	Write X index register.
WRITE_Y	46	16-bit data in	Write Y index register.
WRITE_SP	47	16-bit data in	Write stack pointer.
GO	08	none	Go to user program. If enabled, ACK will occur when leaving active background mode.
GO_UNTIL[2]	0C	none	Go to user program. If enabled, ACK will occur upon returning to active background mode.
TRACE1	10	none	Execute one user instruction then return to active BDM. If enabled, ACK will occur upon returning to active background mode.
TAGGO	18	none	Enable tagging and go to user program. There is no ACK pulse related to this command.

[1] If enabled, ACK will occur when data is ready for transmission for all BDM READ commands and will occur after the write is complete for all BDM WRITE commands.

[2] Both WAIT (with clocks to the S12 CPU core disabled) and STOP disable the ACK function. The GO_UNTIL command will not get an Acknowledge if one of these two CPU instructions occurs before the "UNTIL" instruction. This can be a problem for any instruction that uses ACK, but GO_UNTIL is a lot more difficult for the development tool to time-out.

17.4.5   BDM Command Structure

Hardware and firmware BDM commands start with an 8-bit opcode followed by a 16-bit address and/or a 16-bit data word depending on the command. All the read commands return 16 bits of data despite the byte or word implication in the command name.

NOTE

8-bit reads return 16-bits of data, of which, only one byte will contain valid data. If reading an even address, the valid data will appear in the MSB. If reading an odd address, the valid data will appear in the LSB.

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NOTE

16-bit misaligned reads and writes are not allowed. If attempted, the BDM will ignore the least significant bit of the address and will assume an even address from the remaining bits.

For hardware data read commands, the external host must wait 150 bus clock cycles after sending the address before attempting to obtain the read data. This is to be certain that valid data is available in the BDM shift register, ready to be shifted out. For hardware write commands, the external host must wait 150 bus clock cycles after sending the data to be written before attempting to send a new command. This is to avoid disturbing the BDM shift register before the write has been completed. The 150 bus clock cycle delay in both cases includes the maximum 128 cycle delay that can be incurred as the BDM waits for a free cycle before stealing a cycle.

For firmware read commands, the external host should wait 44 bus clock cycles after sending the command opcode and before attempting to obtain the read data. This includes the potential of an extra 7 cycles when the access is external with a narrow bus access (+1 cycle) and / or a stretch (+1, 2, or 3 cycles), (7 cycles could be needed if both occur). The 44 cycle wait allows enough time for the requested data to be made available in the BDM shift register, ready to be shifted out.

NOTE

This timing has increased from previous BDM modules due to the new capability in which the BDM serial interface can potentially run faster than the bus. On previous BDM modules this extra time could be hidden within the serial time.

For firmware write commands, the external host must wait 32 bus clock cycles after sending the data to be written before attempting to send a new command. This is to avoid disturbing the BDM shift register before the write has been completed.

The external host should wait 64 bus clock cycles after a TRACE1 or GO command before starting any new serial command. This is to allow the CPU to exit gracefully from the standard BDM firmware lookup table and resume execution of the user code. Disturbing the BDM shift register prematurely may adversely affect the exit from the standard BDM firmware lookup table.

NOTE

If the bus rate of the target processor is unknown or could be changing, it is recommended that the ACK (acknowledge function) be used to indicate when an operation is complete. When using ACK, the delay times are automated.

Figure 17-5 represents the BDM command structure. The command blocks illustrate a series of eight bit times starting with a falling edge. The bar across the top of the blocks indicates that the BKGD line idles in the high state. The time for an 8-bit command is 8×16 target clock cycles.[1]
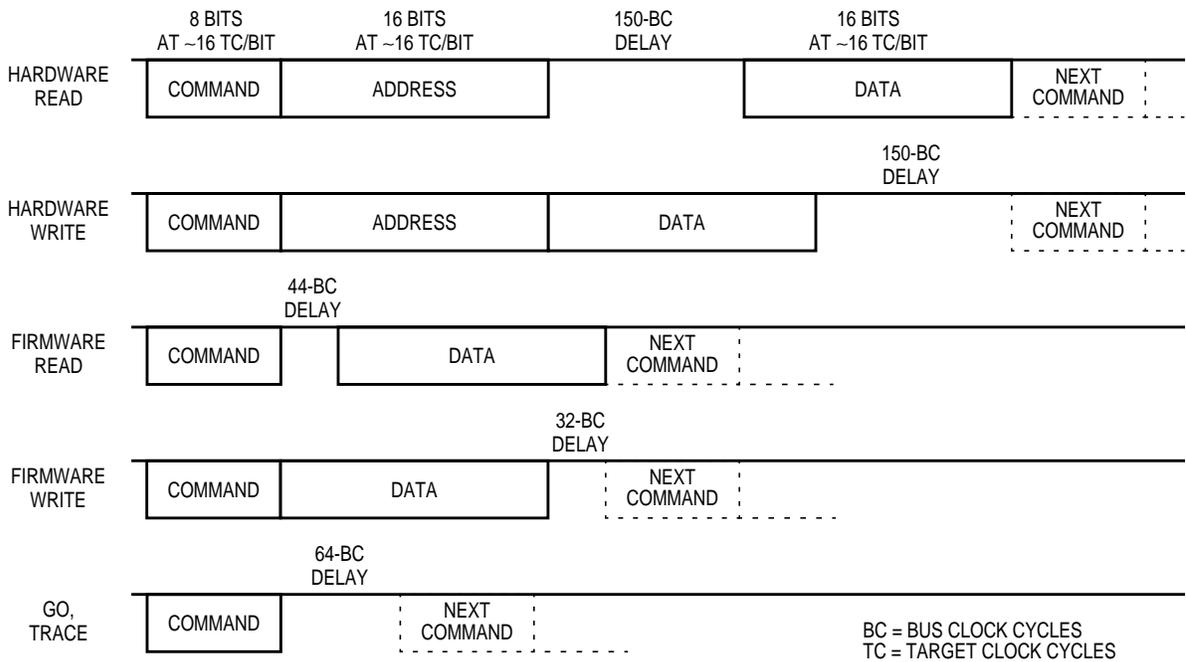

Figure 17-5. BDM Command Structure

17.4.6 BDM Serial Interface

The BDM communicates with external devices serially via the BKGD pin. During reset, this pin is a mode select input which selects between normal and special modes of operation. After reset, this pin becomes the dedicated serial interface pin for the BDM.

The BDM serial interface is timed using the clock selected by the CLKSW bit in the status register see Section 17.3.1, "BDM Status Register (BDMSTS)." This clock will be referred to as the target clock in the following explanation.

The BDM serial interface uses a clocking scheme in which the external host generates a falling edge on the BKGD pin to indicate the start of each bit time. This falling edge is sent for every bit whether data is transmitted or received. Data is transferred most significant bit (MSB) first at 16 target clock cycles per bit. The interface times out if 512 clock cycles occur between falling edges from the host.

The BKGD pin is a pseudo open-drain pin and has an weak on-chip active pull-up that is enabled at all times. It is assumed that there is an external pull-up and that drivers connected to BKGD do not typically drive the high level. Because R-C rise time could be unacceptably long, the target system and host provide brief driven-high (speedup) pulses to drive BKGD to a logic 1. The source of this speedup pulse is the host for transmit cases and the target for receive cases.

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1. Target clock cycles are cycles measured using the target MCU's serial clock rate. See Section 17.4.6, "BDM Serial Interface," and Section 17.3.1, "BDM Status Register (BDMSTS)," for information on how serial clock rate is selected.

The timing for host-to-target is shown in Figure 17-6 and that of target-to-host in Figure 17-7 and Figure 17-8. All four cases begin when the host drives the BKGD pin low to generate a falling edge. Because the host and target are operating from separate clocks, it can take the target system up to one full clock cycle to recognize this edge. The target measures delays from this perceived start of the bit time while the host measures delays from the point it actually drove BKGD low to start the bit up to one target clock cycle earlier. Synchronization between the host and target is established in this manner at the start of every bit time.

Figure 17-6 shows an external host transmitting a logic 1 and transmitting a logic 0 to the BKGD pin of a target system. The host is asynchronous to the target, so there is up to a one clock-cycle delay from the host-generated falling edge to where the target recognizes this edge as the beginning of the bit time. Ten target clock cycles later, the target senses the bit level on the BKGD pin. Internal glitch detect logic requires the pin be driven high no later that eight target clock cycles after the falling edge for a logic 1 transmission.

Because the host drives the high speedup pulses in these two cases, the rising edges look like digitally driven signals.
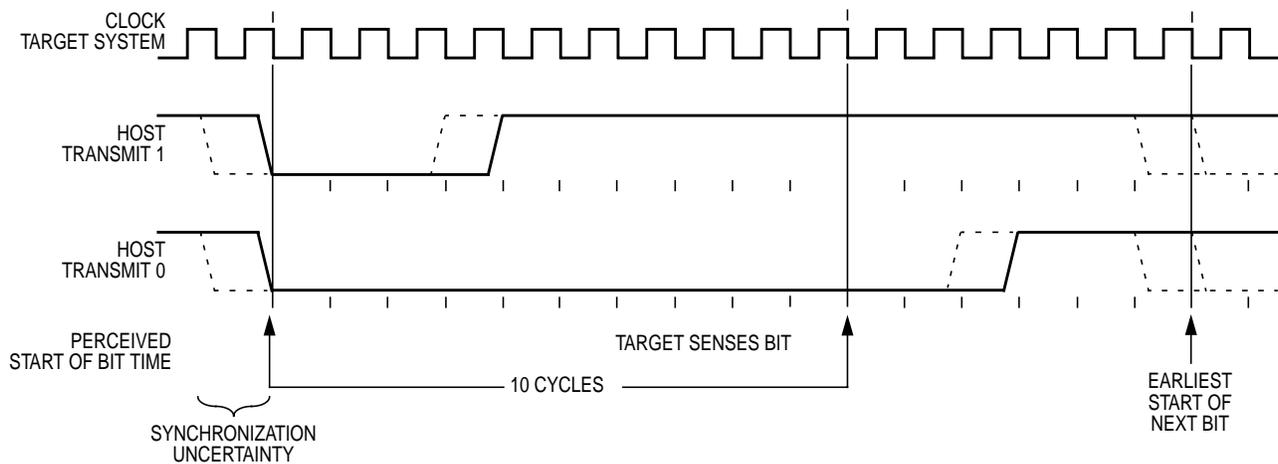
Figure 17-6. BDM Host-to-Target Serial Bit Timing

The receive cases are more complicated. Figure 17-7 shows the host receiving a logic 1 from the target system. Because the host is asynchronous to the target, there is up to one clock-cycle delay from the host-generated falling edge on BKGD to the perceived start of the bit time in the target. The host holds the BKGD pin low long enough for the target to recognize it (at least two target clock cycles). The host must release the low drive before the target drives a brief high speedup pulse seven target clock cycles after the perceived start of the bit time. The host should sample the bit level about 10 target clock cycles after it started the bit time.
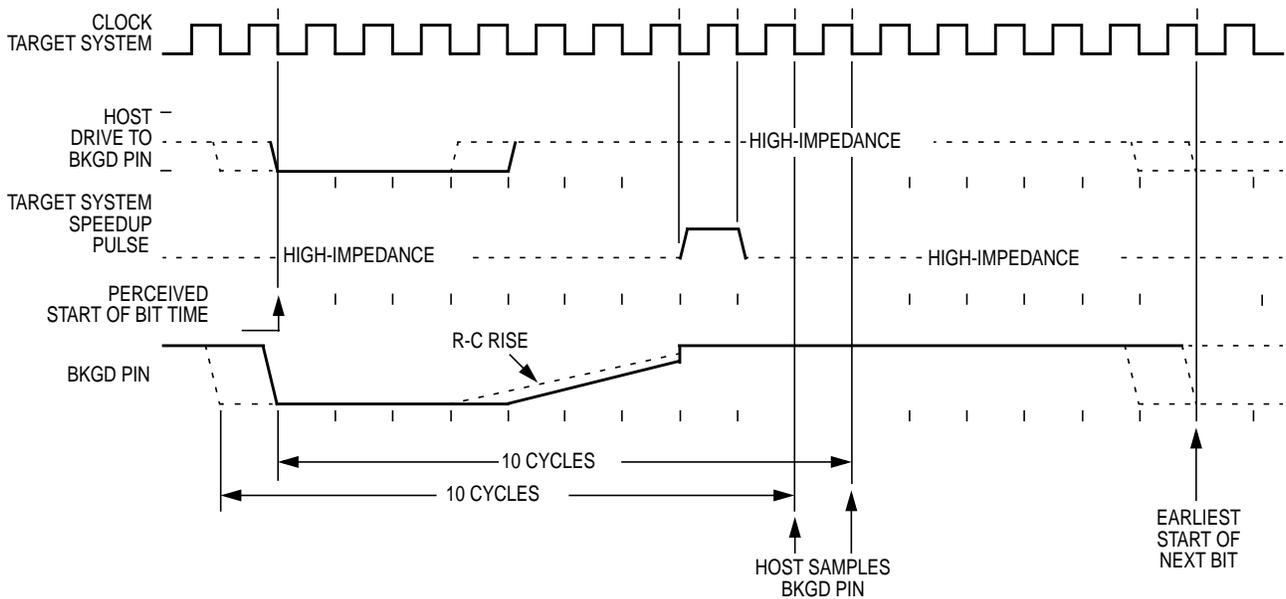
Figure 17-7. BDM Target-to-Host Serial Bit Timing (Logic 1)

Figure 17-8 shows the host receiving a logic 0 from the target. Because the host is asynchronous to the target, there is up to a one clock-cycle delay from the host-generated falling edge on BKGD to the start of the bit time as perceived by the target. The host initiates the bit time but the target finishes it. Because the target wants the host to receive a logic 0, it drives the BKGD pin low for 13 target clock cycles then briefly drives it high to speed up the rising edge. The host samples the bit level about 10 target clock cycles after starting the bit time.
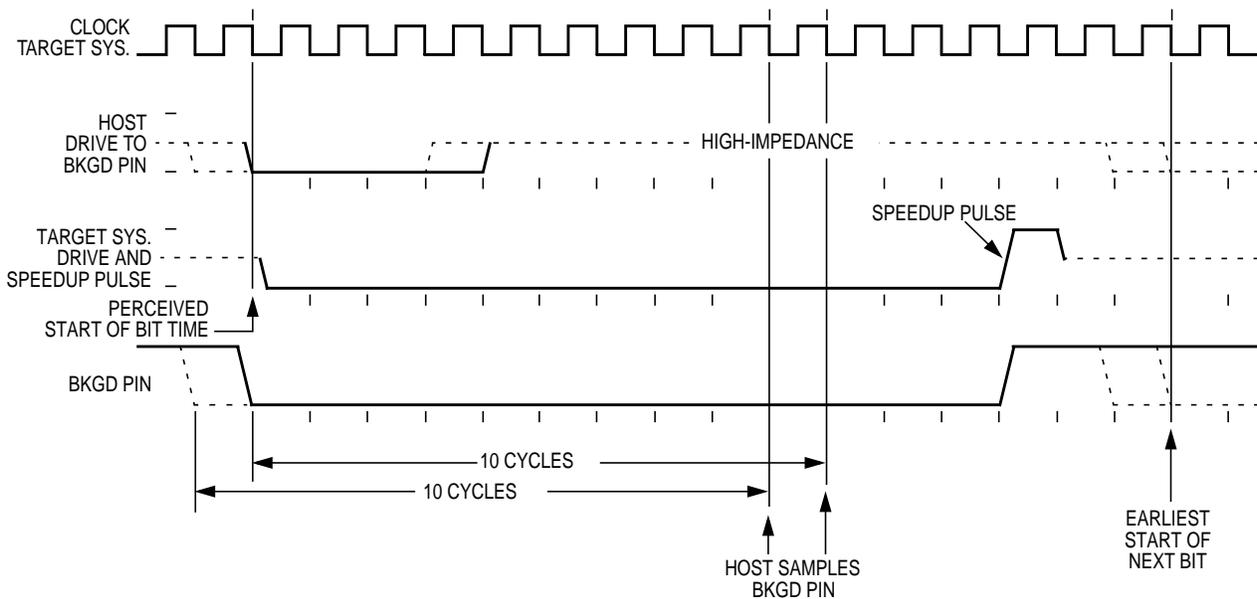
Figure 17-8. BDM Target-to-Host Serial Bit Timing (Logic 0)

17.4.7   Serial Interface Hardware Handshake Protocol

BDM commands that require CPU execution are ultimately treated at the MCU bus rate. Because the BDM clock source can be asynchronously related to the bus frequency, when CLKSW = 0, it is very helpful to provide a handshake protocol in which the host could determine when an issued command is executed by the CPU. The alternative is to always wait the amount of time equal to the appropriate number of cycles at the slowest possible rate the clock could be running. This sub-section will describe the hardware handshake protocol.

The hardware handshake protocol signals to the host controller when an issued command was successfully executed by the target. This protocol is implemented by a 16 serial clock cycle low pulse followed by a brief speedup pulse in the BKGD pin. This pulse is generated by the target MCU when a command, issued by the host, has been successfully executed (see Figure 17-9). This pulse is referred to as the ACK pulse. After the ACK pulse has finished: the host can start the bit retrieval if the last issued command was a read command, or start a new command if the last command was a write command or a control command (BACKGROUND, GO, GO_UNTIL, or TRACE1). The ACK pulse is not issued earlier than 32 serial clock cycles after the BDM command was issued. The end of the BDM command is assumed to be the 16th tick of the last bit. This minimum delay assures enough time for the host to perceive the ACK pulse. Note also that, there is no upper limit for the delay between the command and the related ACK pulse, because the command execution depends upon the CPU bus frequency, which in some cases could be very slow compared to the serial communication rate. This protocol allows a great flexibility for the POD designers, because it does not rely on any accurate time measurement or short response time to any event in the serial communication.
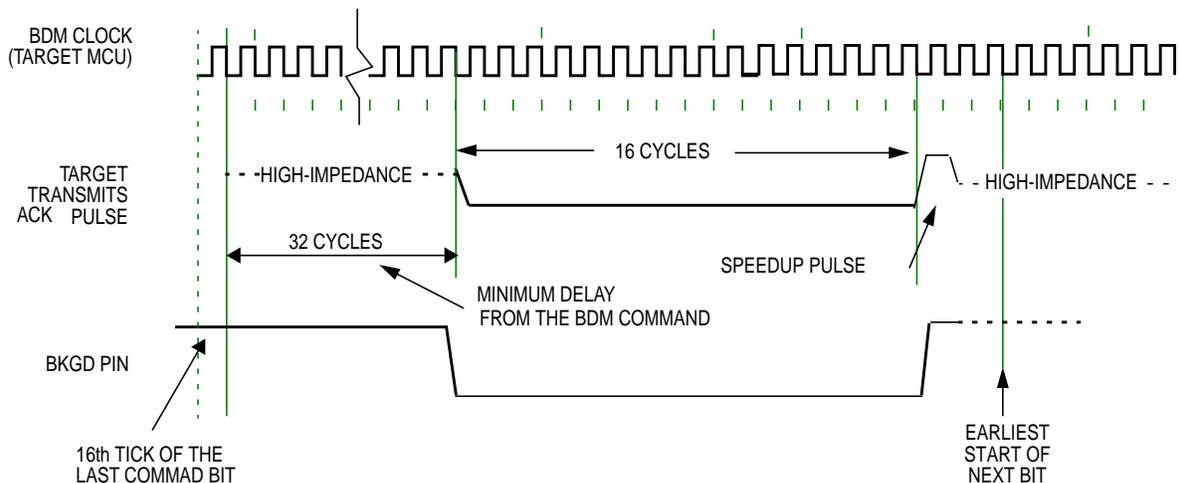

Figure 17-9. Target Acknowledge Pulse (ACK)

NOTE

If the ACK pulse was issued by the target, the host assumes the previous command was executed. If the CPU enters WAIT or STOP prior to executing a hardware command, the ACK pulse will not be issued meaning that the BDM command was not executed. After entering wait or stop mode, the BDM command is no longer pending.

Figure 17-10 shows the ACK handshake protocol in a command level timing diagram. The READ_BYTE instruction is used as an example. First, the 8-bit instruction opcode is sent by the host, followed by the address of the memory location to be read. The target BDM decodes the instruction. A bus cycle is grabbed (free or stolen) by the BDM and it executes the READ_BYTE operation. Having retrieved the data, the BDM issues an ACK pulse to the host controller, indicating that the addressed byte is ready to be retrieved. After detecting the ACK pulse, the host initiates the byte retrieval process. Note that data is sent in the form of a word and the host needs to determine which is the appropriate byte based on whether the address was odd or even.
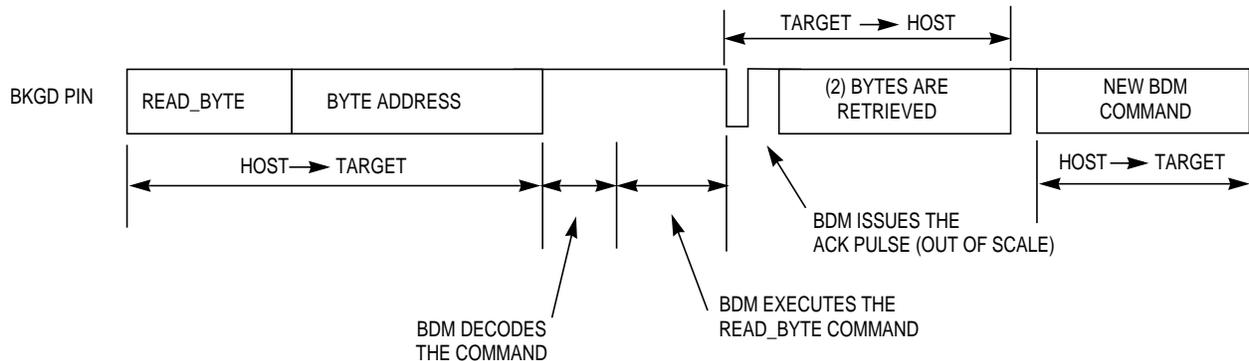

Figure 17-10. Handshake Protocol at Command Level

Differently from the normal bit transfer (where the host initiates the transmission), the serial interface ACK handshake pulse is initiated by the target MCU by issuing a falling edge in the BKGD pin. The hardware handshake protocol in Figure 17-9 specifies the timing when the BKGD pin is being driven, so the host should follow this timing constraint in order to avoid the risk of an electrical conflict in the BKGD pin.

NOTE

The only place the BKGD pin can have an electrical conflict is when one side is driving low and the other side is issuing a speedup pulse (high). Other "highs" are pulled rather than driven. However, at low rates the time of the speedup pulse can become lengthy and so the potential conflict time becomes longer as well.

The ACK handshake protocol does not support nested ACK pulses. If a BDM command is not acknowledge by an ACK pulse, the host needs to abort the pending command first in order to be able to issue a new BDM command. When the CPU enters WAIT or STOP while the host issues a command that requires CPU execution (e.g., WRITE_BYTE), the target discards the incoming command due to the WAIT or STOP being detected. Therefore, the command is not acknowledged by the target, which means that the ACK pulse will not be issued in this case. After a certain time the host should decide to abort the ACK sequence in order to be free to issue a new command. Therefore, the protocol should provide a mechanism in which a command, and therefore a pending ACK, could be aborted.

NOTE

> Differently from a regular BDM command, the ACK pulse does not provide
> a time out. This means that in the case of a WAIT or STOP instruction being
> executed, the ACK would be prevented from being issued. If not aborted, the
> ACK would remain pending indefinitely. See the handshake abort procedure
> described in Section 17.4.8, "Hardware Handshake Abort Procedure."

17.4.8   Hardware Handshake Abort Procedure

The abort procedure is based on the SYNC command. In order to abort a command, which had not issued the corresponding ACK pulse, the host controller should generate a low pulse in the BKGD pin by driving it low for at least 128 serial clock cycles and then driving it high for one serial clock cycle, providing a speedup pulse. By detecting this long low pulse in the BKGD pin, the target executes the SYNC protocol, see Section 17.4.9, "SYNC — Request Timed Reference Pulse," and assumes that the pending command and therefore the related ACK pulse, are being aborted. Therefore, after the SYNC protocol has been completed the host is free to issue new BDM commands.

Although it is not recommended, the host could abort a pending BDM command by issuing a low pulse in the BKGD pin shorter than 128 serial clock cycles, which will not be interpreted as the SYNC command. The ACK is actually aborted when a falling edge is perceived by the target in the BKGD pin. The short abort pulse should have at least 4 clock cycles keeping the BKGD pin low, in order to allow the falling edge to be detected by the target. In this case, the target will not execute the SYNC protocol but the pending command will be aborted along with the ACK pulse. The potential problem with this abort procedure is when there is a conflict between the ACK pulse and the short abort pulse. In this case, the target may not perceive the abort pulse. The worst case is when the pending command is a read command (i.e., READ_BYTE). If the abort pulse is not perceived by the target the host will attempt to send a new command after the abort pulse was issued, while the target expects the host to retrieve the accessed memory byte. In this case, host and target will run out of synchronism. However, if the command to be aborted is not a read command the short abort pulse could be used. After a command is aborted the target assumes the next falling edge, after the abort pulse, is the first bit of a new BDM command.

NOTE

> The details about the short abort pulse are being provided only as a reference
> for the reader to better understand the BDM internal behavior. It is not
> recommended that this procedure be used in a real application.

Because the host knows the target serial clock frequency, the SYNC command (used to abort a command) does not need to consider the lower possible target frequency. In this case, the host could issue a SYNC very close to the 128 serial clock cycles length. Providing a small overhead on the pulse length in order to assure the SYNC pulse will not be misinterpreted by the target. See Section 17.4.9, "SYNC — Request Timed Reference Pulse."

Figure 17-11 shows a SYNC command being issued after a READ_BYTE, which aborts the READ_BYTE command. Note that, after the command is aborted a new command could be issued by the host computer.


MC9S12NE64 Data Sheet, Rev 1.0
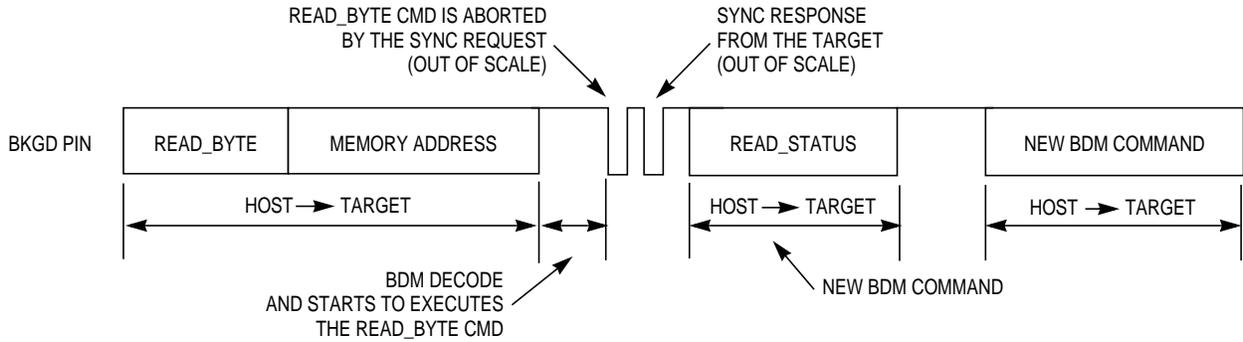
Figure 17-11. ACK Abort Procedure at the Command Level

NOTE

Figure 17-11 does not represent the signals in a true timing scale

Figure 17-12 shows a conflict between the ACK pulse and the SYNC request pulse. This conflict could occur if a POD device is connected to the target BKGD pin and the target is already in debug active mode. Consider that the target CPU is executing a pending BDM command at the exact moment the POD is being connected to the BKGD pin. In this case, an ACK pulse is issued along with the SYNC command. In this case, there is an electrical conflict between the ACK speedup pulse and the SYNC pulse. Because this is not a probable situation, the protocol does not prevent this conflict from happening.
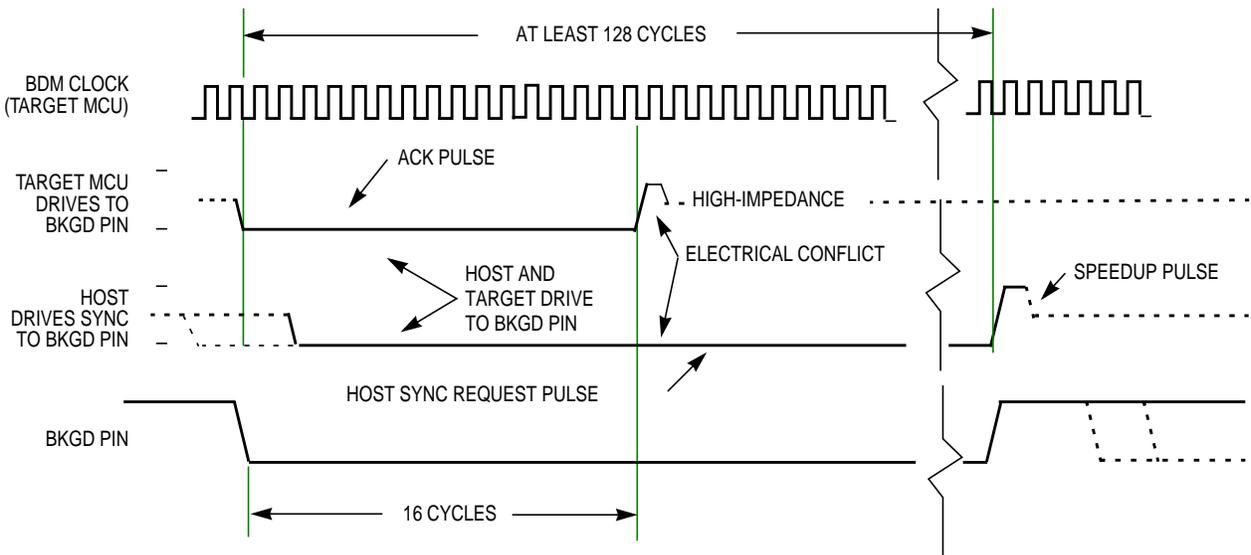

Figure 17-12. ACK Pulse and SYNC Request Conflict

NOTE

This information is being provided so that the MCU integrator will be aware that such a conflict could eventually occur.

The hardware handshake protocol is enabled by the ACK_ENABLE and disabled by the ACK_DISABLE BDM commands. This provides backwards compatibility with the existing POD devices which are not able to execute the hardware handshake protocol. It also allows for new POD devices, that support the

hardware handshake protocol, to freely communicate with the target device. If desired, without the need for waiting for the ACK pulse.

The commands are described as follows:

- ACK_ENABLE — enables the hardware handshake protocol. The target will issue the ACK pulse when a CPU command is executed by the CPU. The ACK_ENABLE command itself also has the ACK pulse as a response.
- ACK_DISABLE — disables the ACK pulse protocol. In this case, the host needs to use the worst case delay time at the appropriate places in the protocol.

The default state of the BDM after reset is hardware handshake protocol disabled.

All the read commands will ACK (if enabled) when the data bus cycle has completed and the data is then ready for reading out by the BKGD serial pin. All the write commands will ACK (if enabled) after the data has been received by the BDM through the BKGD serial pin and when the data bus cycle is complete. See Section 17.4.3, "BDM Hardware Commands," and Section 17.4.4, "Standard BDM Firmware Commands," for more information on the BDM commands.

The ACK_ENABLE sends an ACK pulse when the command has been completed. This feature could be used by the host to evaluate if the target supports the hardware handshake protocol. If an ACK pulse is issued in response to this command, the host knows that the target supports the hardware handshake protocol. If the target does not support the hardware handshake protocol the ACK pulse is not issued. In this case, the ACK_ENABLE command is ignored by the target since it is not recognized as a valid command.

The BACKGROUND command will issue an ACK pulse when the CPU changes from normal to background mode. The ACK pulse related to this command could be aborted using the SYNC command.

The GO command will issue an ACK pulse when the CPU exits from background mode. The ACK pulse related to this command could be aborted using the SYNC command.

The GO_UNTIL command is equivalent to a GO command with exception that the ACK pulse, in this case, is issued when the CPU enters into background mode. This command is an alternative to the GO command and should be used when the host wants to trace if a breakpoint match occurs and causes the CPU to enter active background mode. Note that the ACK is issued whenever the CPU enters BDM, which could be caused by a breakpoint match or by a BGND instruction being executed. The ACK pulse related to this command could be aborted using the SYNC command.

The TRACE1 command has the related ACK pulse issued when the CPU enters background active mode after one instruction of the application program is executed. The ACK pulse related to this command could be aborted using the SYNC command.

The TAGGO command will not issue an ACK pulse because this would interfere with the tagging function shared on the same pin.

17.4.9   SYNC — Request Timed Reference Pulse

The SYNC command is unlike other BDM commands because the host does not necessarily know the correct communication speed to use for BDM communications until after it has analyzed the response to the SYNC command. To issue a SYNC command, the host should perform the following steps:

1. Drive the BKGD pin low for at least 128 cycles at the lowest possible BDM serial communication frequency (the lowest serial communication frequency is determined by the crystal oscillator or the clock chosen by CLKSW.)

2. Drive BKGD high for a brief speedup pulse to get a fast rise time (this speedup pulse is typically one cycle of the host clock.)

3. Remove all drive to the BKGD pin so it reverts to high impedance.

4. Listen to the BKGD pin for the sync response pulse.

Upon detecting the SYNC request from the host, the target performs the following steps:

1. Discards any incomplete command received or bit retrieved.

2. Waits for BKGD to return to a logic 1.

3. Delays 16 cycles to allow the host to stop driving the high speedup pulse.

4. Drives BKGD low for 128 cycles at the current BDM serial communication frequency.

5. Drives a one-cycle high speedup pulse to force a fast rise time on BKGD.

6. Removes all drive to the BKGD pin so it reverts to high impedance.

The host measures the low time of this 128 cycle SYNC response pulse and determines the correct speed for subsequent BDM communications. Typically, the host can determine the correct communication speed within a few percent of the actual target speed and the communication protocol can easily tolerate speed errors of several percent.

As soon as the SYNC request is detected by the target, any partially received command or bit retrieved is discarded. This is referred to as a soft-reset, equivalent to a time-out in the serial communication. After the SYNC response, the target will consider the next falling edge (issued by the host) as the start of a new BDM command or the start of new SYNC request.

Another use of the SYNC command pulse is to abort a pending ACK pulse. The behavior is exactly the same as in a regular SYNC command. Note that one of the possible causes for a command to not be acknowledged by the target is a host-target synchronization problem. In this case, the command may not have been understood by the target and so an ACK response pulse will not be issued.

17.4.10 Instruction Tracing

When a TRACE1 command is issued to the BDM in active BDM, the CPU exits the standard BDM firmware and executes a single instruction in the user code. As soon as this has occurred, the CPU is forced to return to the standard BDM firmware and the BDM is active and ready to receive a new command. If the TRACE1 command is issued again, the next user instruction will be executed. This facilitates stepping or tracing through the user code one instruction at a time.

If an interrupt is pending when a TRACE1 command is issued, the interrupt stacking operation occurs but no user instruction is executed. Upon return to standard BDM firmware execution, the program counter points to the first instruction in the interrupt service routine.

17.4.11 Instruction Tagging

The instruction queue and cycle-by-cycle CPU activity are reconstructible in real time or from trace history that is captured by a logic analyzer. However, the reconstructed queue cannot be used to stop the CPU at a specific instruction. This is because execution already has begun by the time an operation is visible outside the system. A separate instruction tagging mechanism is provided for this purpose.

The tag follows program information as it advances through the instruction queue. When a tagged instruction reaches the head of the queue, the CPU enters active BDM rather than executing the instruction.

NOTE

Tagging is disabled when BDM becomes active and BDM serial commands are not processed while tagging is active.

Executing the BDM TAGGO command configures two system pins for tagging. The $\overline{\text{TAGLO}}$ signal shares a pin with the $\overline{\text{LSTRB}}$ signal, and the $\overline{\text{TAGHI}}$ signal shares a pin with the BKGD signal.

Table 17-5 shows the functions of the two tagging pins. The pins operate independently, that is the state of one pin does not affect the function of the other. The presence of logic level 0 on either pin at the fall of the external clock (ECLK) performs the indicated function. High tagging is allowed in all modes. Low tagging is allowed only when low strobe is enabled (LSTRB is allowed only in wide expanded modes and emulation expanded narrow mode).

Table 17-5. Tag Pin Function

$\overline{\text{TAGHI}}$	$\overline{\text{TAGLO}}$	Tag
1	1	No tag
1	0	Low byte
0	1	High byte
0	0	Both bytes

17.4.12 Serial Communication Time-Out

The host initiates a host-to-target serial transmission by generating a falling edge on the BKGD pin. If BKGD is kept low for more than 128 target clock cycles, the target understands that a SYNC command was issued. In this case, the target will keep waiting for a rising edge on BKGD in order to answer the SYNC request pulse. If the rising edge is not detected, the target will keep waiting forever without any time-out limit.

Consider now the case where the host returns BKGD to logic one before 128 cycles. This is interpreted as a valid bit transmission, and not as a SYNC request. The target will keep waiting for another falling edge marking the start of a new bit. If, however, a new falling edge is not detected by the target within 512 clock

cycles since the last falling edge, a time-out occurs and the current command is discarded without affecting memory or the operating mode of the MCU. This is referred to as a soft-reset.

If a read command is issued but the data is not retrieved within 512 serial clock cycles, a soft-reset will occur causing the command to be disregarded. The data is not available for retrieval after the time-out has occurred. This is the expected behavior if the handshake protocol is not enabled. However, consider the behavior where the BDC is running in a frequency much greater than the CPU frequency. In this case, the command could time out before the data is ready to be retrieved. In order to allow the data to be retrieved even with a large clock frequency mismatch (between BDC and CPU) when the hardware handshake protocol is enabled, the time out between a read command and the data retrieval is disabled. Therefore, the host could wait for more then 512 serial clock cycles and still be able to retrieve the data from an issued read command. However, as soon as the handshake pulse (ACK pulse) is issued, the time-out feature is re-activated, meaning that the target will time out after 512 clock cycles. Therefore, the host needs to retrieve the data within a 512 serial clock cycles time frame after the ACK pulse had been issued. After that period, the read command is discarded and the data is no longer available for retrieval. Any falling edge of the BKGD pin after the time-out period is considered to be a new command or a SYNC request.

Note that whenever a partially issued command, or partially retrieved data, has occurred the time out in the serial communication is active. This means that if a time frame higher than 512 serial clock cycles is observed between two consecutive negative edges and the command being issued or data being retrieved is not complete, a soft-reset will occur causing the partially received command or data retrieved to be disregarded. The next falling edge of the BKGD pin, after a soft-reset has occurred, is considered by the target as the start of a new BDM command, or the start of a SYNC request pulse.