



**Sunbeam Institute of Information Technology
Pune and Karad**

Algorithms and Data structures

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Bubble sort

1. Compare all pairs of consecutive elements of the array one by one
2. If left element is greater than right element, then swap both
3. Repeat above steps until array is sorted ($N-1$)

No. of elements = n
 No. of passes = $n-1$

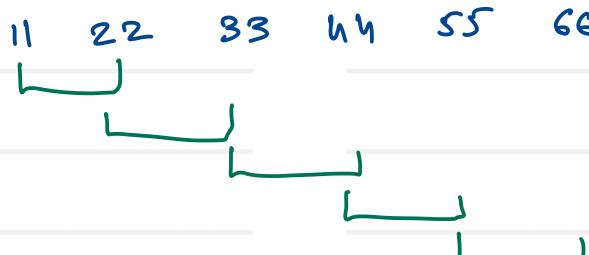
pass	comps
1	$n-1$
2	$n-2$
3	\vdots
i	2
5	1

$$n=6$$

$$j < N-1 \quad j < N-i$$

i	j	j
1	$0 \rightarrow 4$	$0 \rightarrow 4$ (5)
2	$0 \rightarrow 4$	$0 \rightarrow 3$ (4)
3	$0 \rightarrow 4$	$0 \rightarrow 2$ (3)
4	$0 \rightarrow 4$	$0 \rightarrow 1$ (2)
5	$0 \rightarrow 4$	$0 \rightarrow 0$ (1)

Best case :



$$T(n) = O(n)$$

$$\begin{aligned} \text{Total comps} &= 1+2+3+\dots+n \\ &= \frac{n(n+1)}{2} \end{aligned}$$

$$\text{Time} \propto \frac{1}{2}(n^2 + n)$$

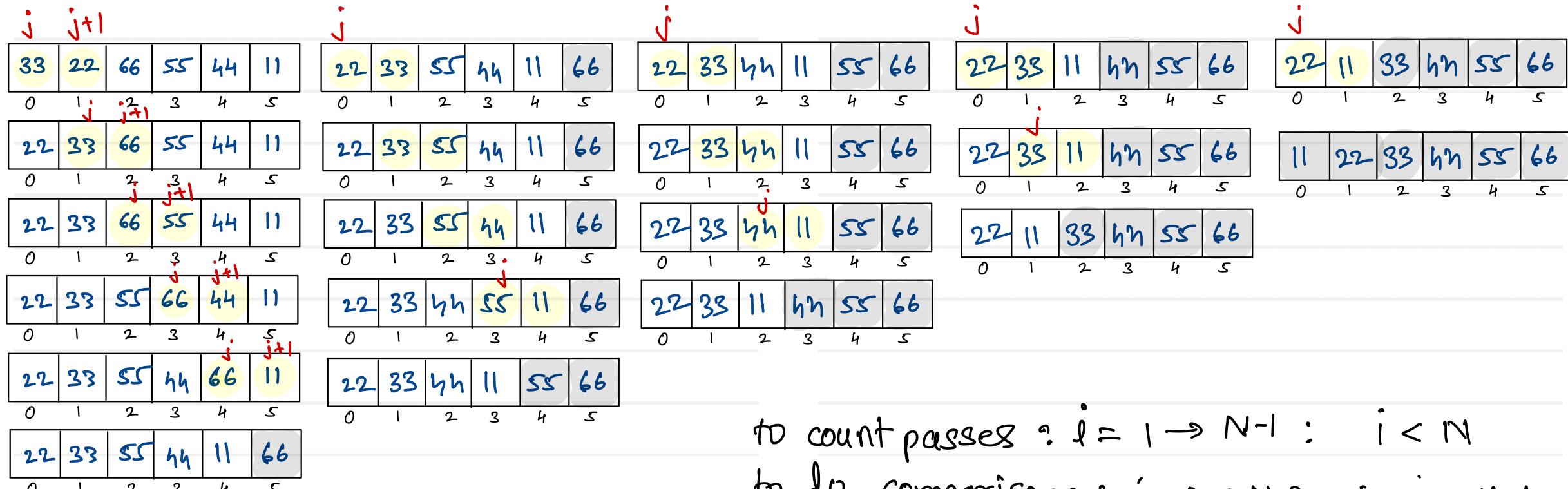
$$T(n) = O(n^2)$$

Avg
Worst



Bubble sort

33	22	66	55	44	11
0	1	2	3	4	5



to count passes : $i = 1 \rightarrow N-1$: $i < N$

to do comparisons : $j = 0 \rightarrow N-2$: $j < N-i$

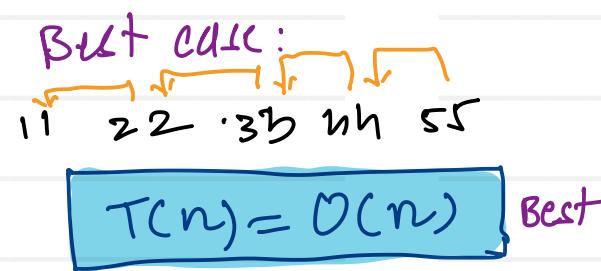




Insertion sort

1. Pick one element of the array (start from 2nd index)
2. Compare picked element with all its left neighbours one by one
3. If left neighbour is greater, move it one position ahead
4. Insert picked element at its appropriate position
5. Repeat above steps until array is sorted ($N-1$)

	passes	comps
No. of elements = n	1	1
No. of passes = $n-1$	2	2
	3	3
	\vdots	\vdots
	$n-1$	$n-1$
	n	<u>n</u>



$$\text{Total comps} = 1 + 2 + 3 + \dots + n$$

$$= \frac{n(n+1)}{2}$$

$$\text{Time} \propto \frac{1}{2}(n^2 + n)$$

$$T(n) = O(n^2)$$

worst
Avg





Insertion sort

55	44	22	66	11	33
0	1	2	3	4	5

44
j
temp

22
j
temp

66
temp
j

11
j
temp

33
j
temp

55	44	22	66	11	33
0	1	2	3	4	5

44	55	22	66	11	33
0	1	2	3	4	5

22	44	55	66	11	33
0	1	2	3	4	5

22	44	55	66	11	33
0	1	2	3	4	5

11	22	44	55	66	33
0	1	2	3	4	5

55	55	22	66	11	33
0	1	2	3	4	5

44	55	55	66	11	33
0	1	2	3	4	5

22	44	55	66	11	33
0	1	2	3	4	5

22	44	55	66	66	33
0	1	2	3	4	5

11	22	44	55	66	66
0	1	2	3	4	5

44	55	22	66	11	33
0	1	2	3	4	5

44	44	55	66	11	33
0	1	2	3	4	5

22	44	55	66	11	33
0	1	2	3	4	5

22	44	55	55	66	33
0	1	2	3	4	5

11	22	44	55	55	66
0	1	2	3	4	5

22	44	55	55	66	11
0	1	2	3	4	5

to pick element : $i = 1 \rightarrow N-1$ ($i < N$)

to compare with left : $j = i-1 \rightarrow 0$ ($j \geq 0$)



Insertion sort

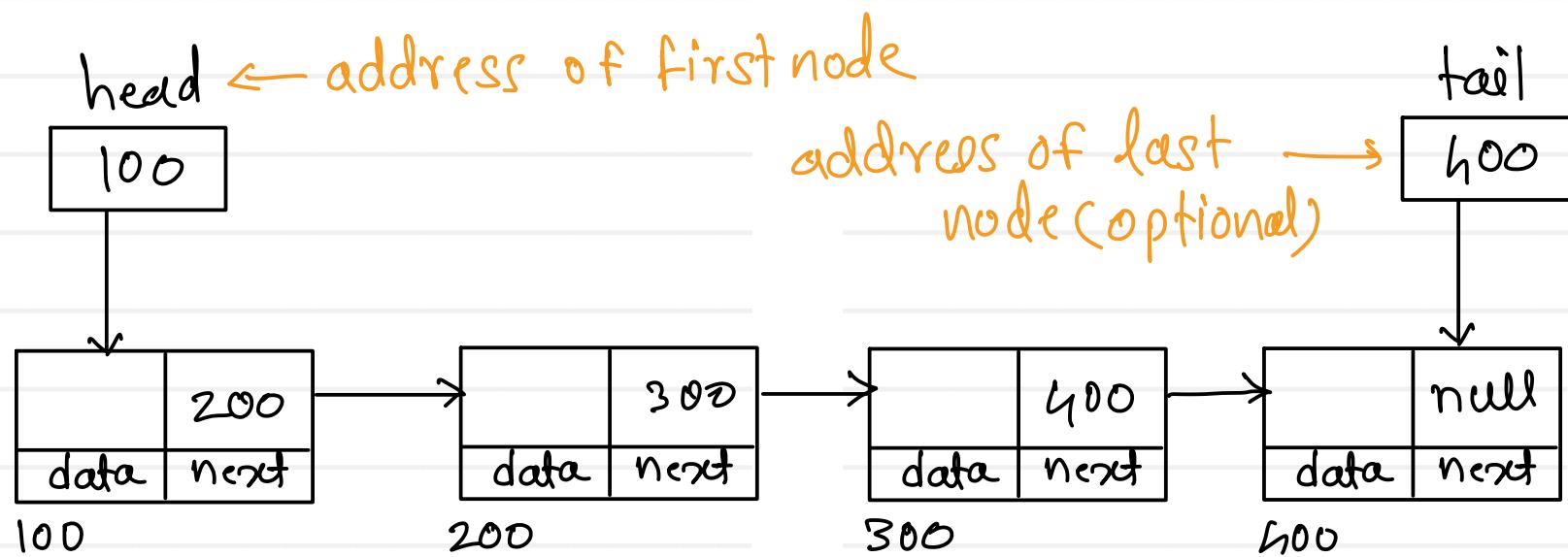
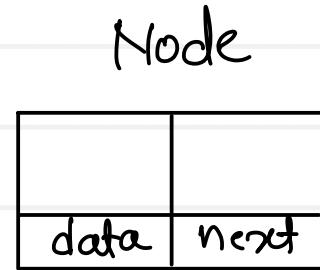
```
for(i=1; i<N; i++) {  
    temp = arr[i];  
    for(j=i-1; j>=0; j--) {  
        if(arr[j] > temp)  
            arr[j+1] = arr[j];  
        else  
            break;  
    }  
    arr[j+1] = temp;  
}
```

11	22	33	44	55	66
0	1	2	3	4	5

i	i<6	temp	j
1	T	44	0,-1
2	T	22	1,0,-1
3	T	66	2
4	T	11	3,2,1,0,-1
5	T	33	4,3,2,1

Linked List

- linked list is a linear data structure.
- link/address of next data is kept with its previous data.
- every element of linked list has two parts (data, link/addr) and referred as node





Linked List

Operations

1. Add first
 2. Add last
 3. Add position (insert)
-
1. Delete first
 2. Delete last
 3. Delete position
-
1. Display (traverse) (forward/ backward)
-
1. Search
 2. Sort
 3. Reverse

Types

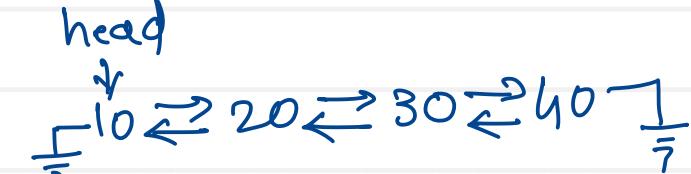
1. Singly linear linked list



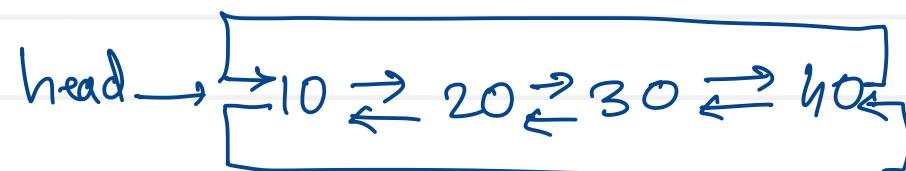
2. Singly circular linked list



3. Doubly linear linked list



4. Doubly circular linked list





Linked List

Node :

data → int, char, double, class
next → address / reference of next node

class Node { ← self referential class

int data ;

Node next ;

}

why inner class ?

↳ to access private fields of Node class directly into List class.

why static ?

↳ to restrict access of private fields of List class into Node class

↳ there is no dependency of outer class to create object of inner class .

class List {

static class Node {

int data ;

Node next ;

}

Node head ;

Node tail ; (optional)

int size ;

public List() { ... }

public addNode() { ... }

public deleteNode() { ... }

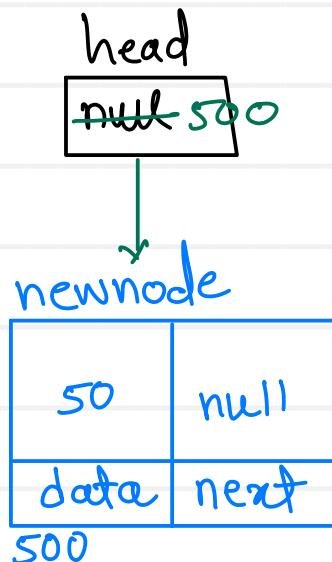
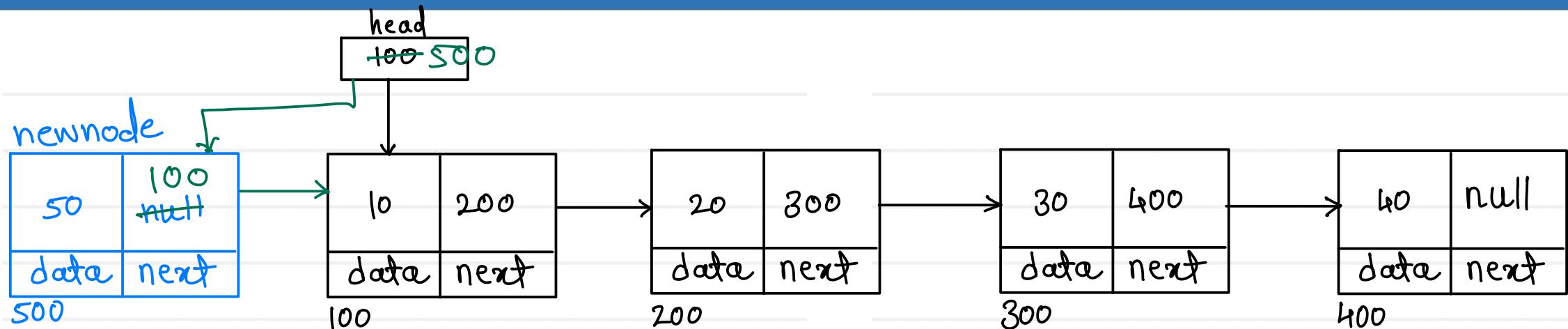
public display() { ... }

public reverse() { ... }

}



Singly linear Linked List - Add first

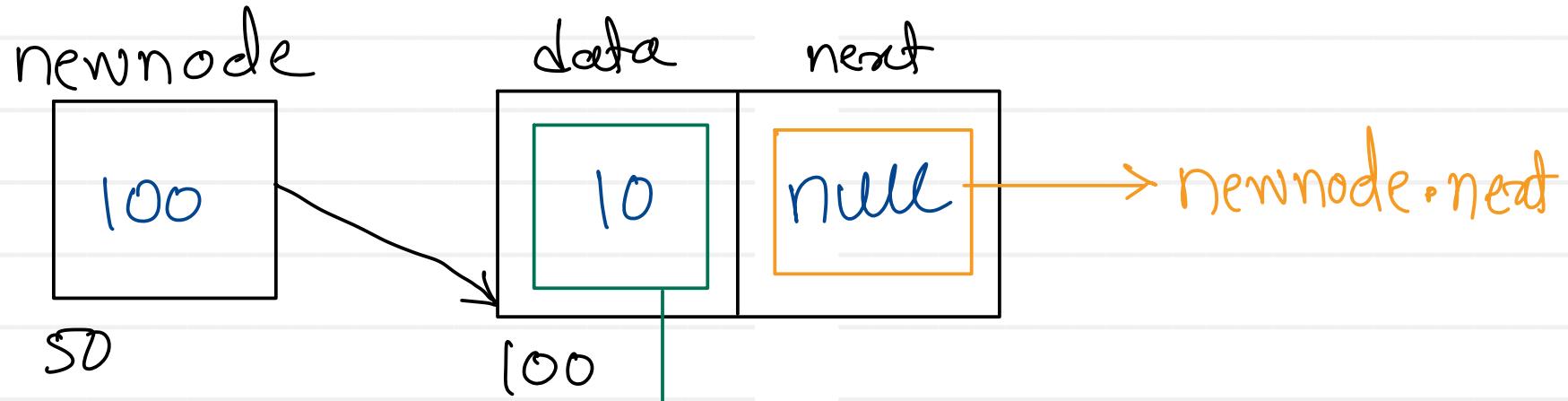


newnode.next = head;
head = newnode;

1. Create node with given data
2. add first node into next of newnode
3. move head on newnode

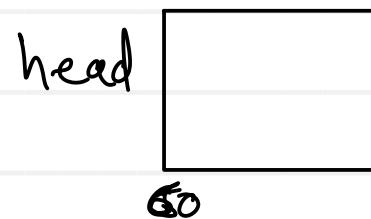
$$T(n) = O(1)$$

Node newnode = new Node(10);

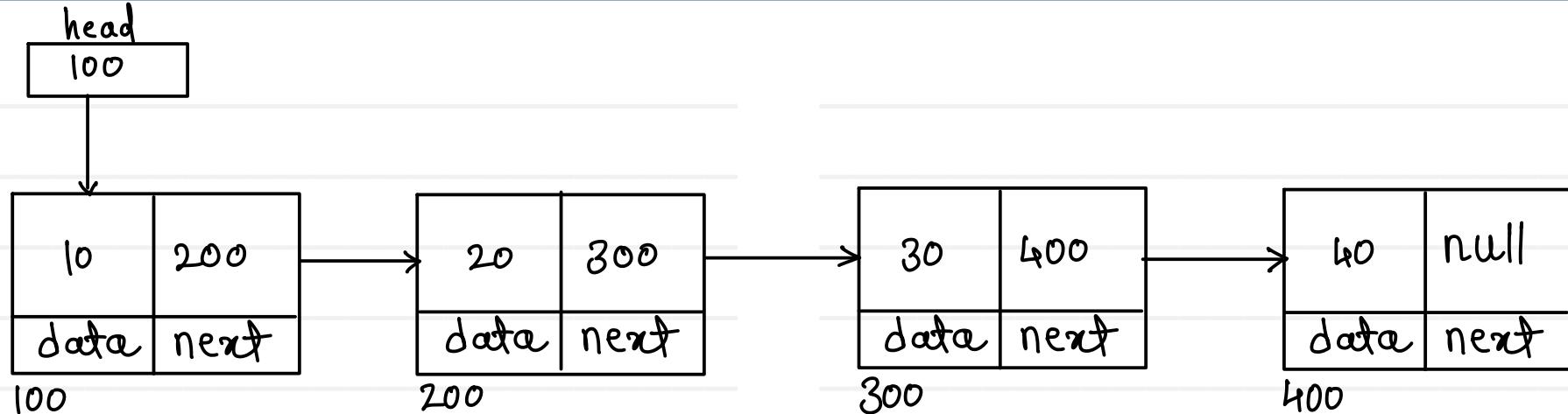


`newnode.data = 2; ← writing`
`2 = newnode.data ← reading`

Node head ;



Singly linear Linked List - Display



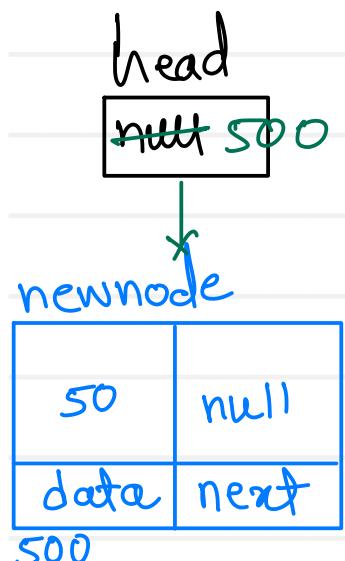
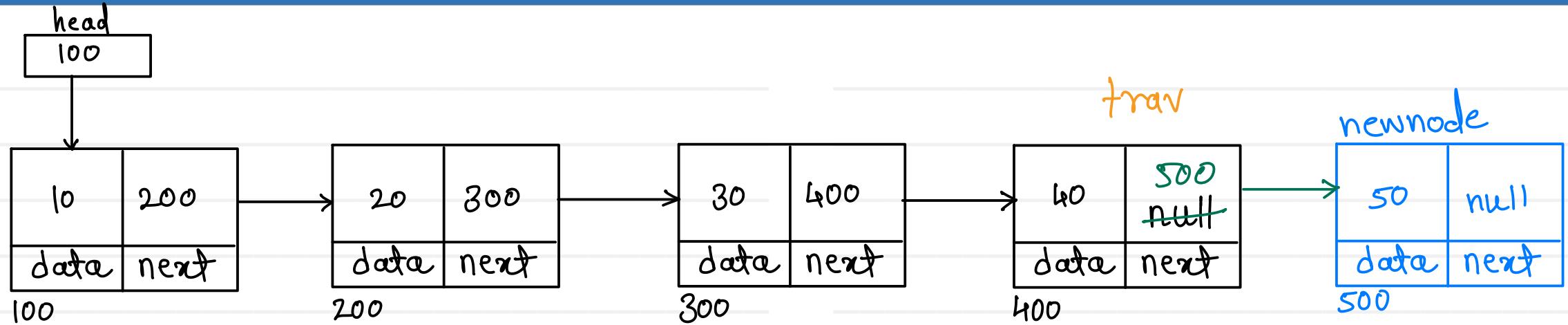
```
Node trav = head;
while(trav != null) {
    System.out.println(trav.data);
    trav = trav.next;
}
```

	trav	trav!=null	trav.data
100	T	T	10
200	T	T	20
300	T	T	30
400	T	T	40
null	F		

1. create `trav` & start at `head`
2. print/visit current node (`trav.data`)
3. go on next node (`trav.next`)
4. repeat step 2 & 3 till last node

$$T(n) = O(n)$$

Singly linear Linked List - Add last

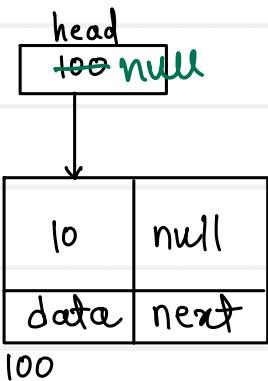
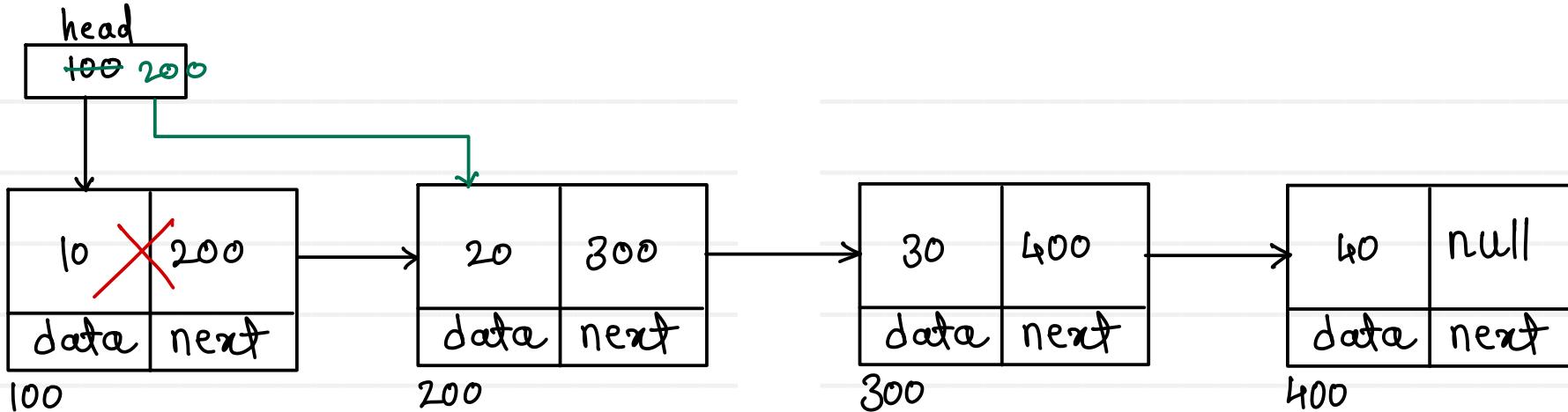


1. create node
2. if list is empty
add newnode into head
3. if list is not empty
 - a. traverse till last node
 - b. add newnode into next of last node.

$$T(n) = O(n)$$

```
Node trav = head;
while (trav.next != null) {
    trav = trav.next;
}
trav.next = newnode;
```

Singly linear Linked List - Delete first

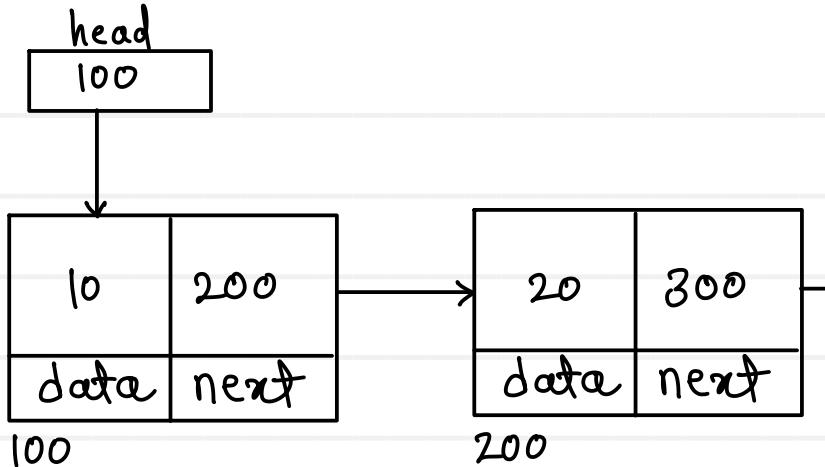


1. if list is empty
do nothing
2. if list is not empty
move head on second node

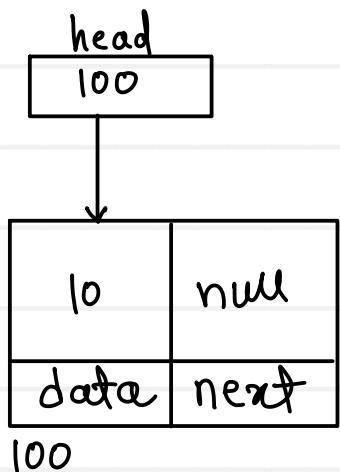
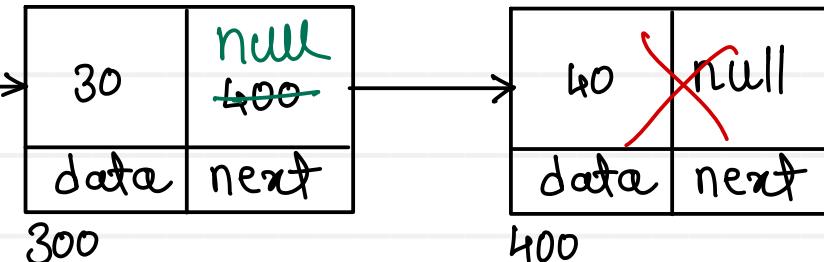


$$T(n) = O(1)$$

Singly linear Linked List - Delete last



trav



1. if list empty
return
2. if list has single node
head = null
3. if list has multiple nodes
 - a. traverse till second last node
 - b. add null into next of second last node

$$T(n) = O(n)$$

`while(trav != null) → trav = null`

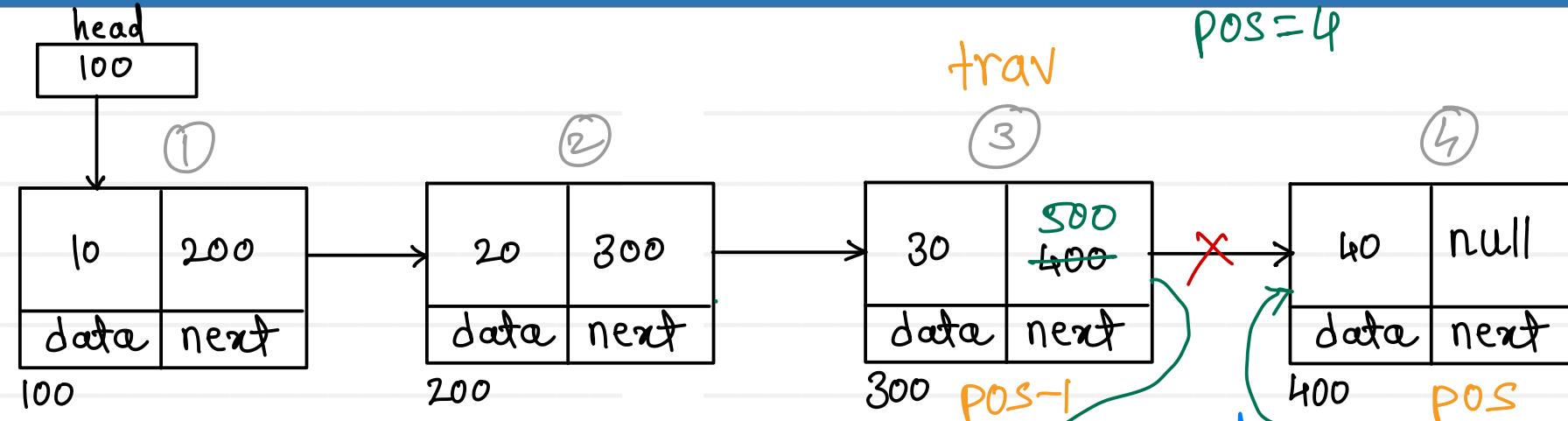
`while(trav.next != null) → trav = last node`

`while(trav.next.next != null) → trav = second last node`



Singly linear Linked List - Add position

make before break



1. create a newnode
if list is empty
 add newnode into head
2. if list is not empty
 2. traverse till pos-1 node
3. add pos node into next of newnode
4. add newnode into next of pos-1 node

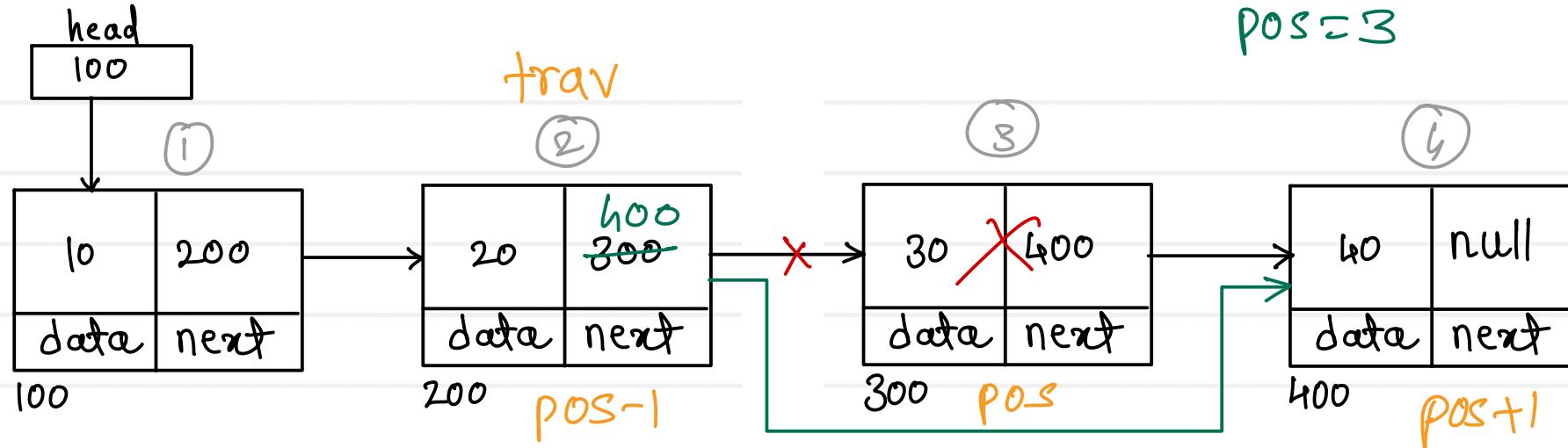
$$T(n) = O(n)$$

```
Node trav = head;
for( i=1; i< pos-1 ; i++)
    trav = trav.next;
```

trav	i	i < 3	pos = 4	trav	i	i < 4	pos = 5
100	1	T		100	1	T	
200	2	T		200	2	T	
300	3	F		300	3	T	
				400	4	F	



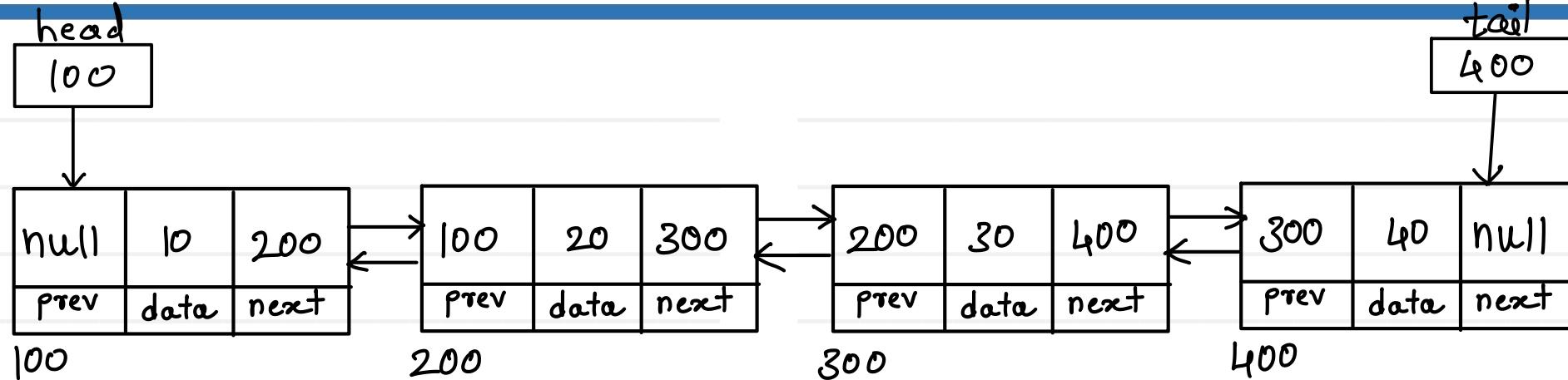
Singly linear Linked List - Delete position



1. if list is empty
return
2. if list is not empty
 - a. traverse till pos-1 node
 - b. add pos+1 node into next of pos-1 node

$$T(n) = O(n)$$

Doubly Linear Linked List - Display



Forward traversal

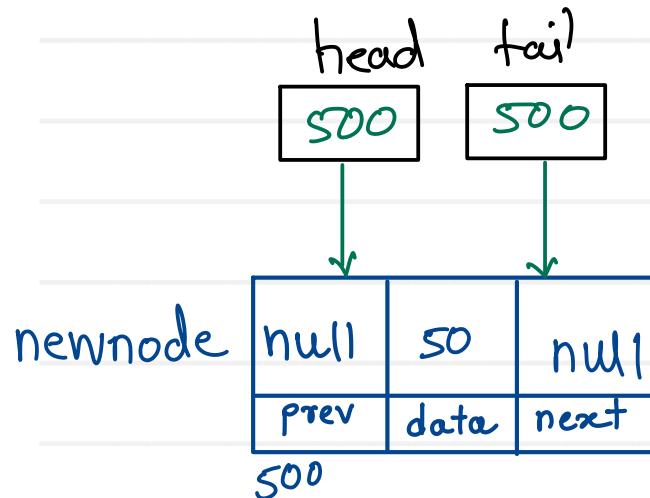
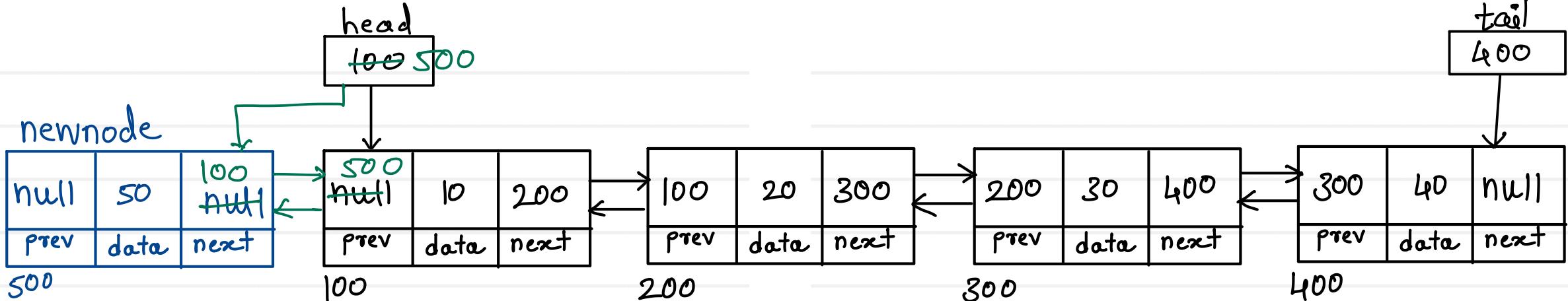
1. create trav & start at head
2. print current node data (trav.data)
3. go on next node (trav.next)
4. repeat above two steps till last node

Backward traversal

1. create trav & start at tail
2. print current node data (trav.data)
3. go on prev node (trav.prev)
4. repeat above two steps till first node

$$T(n) = O(n)$$

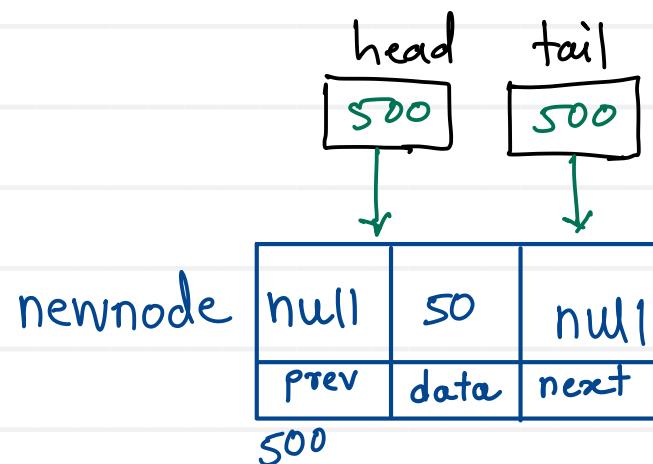
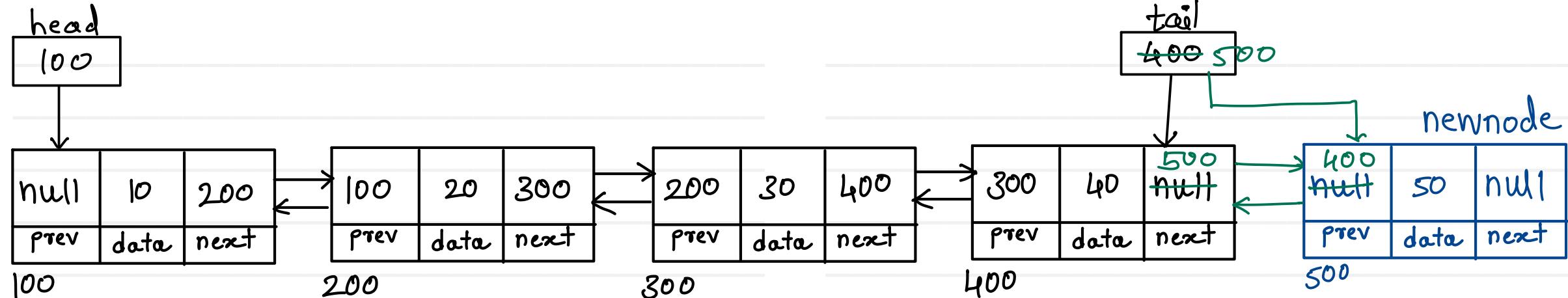
Doubly Linear Linked List - Add first



1. create node
2. if list is empty,
add newnode into head & tail
3. if list is not empty
 - a. add first node into next of newnode
 - b. add newnode into prev of first node
 - c. move head on newnode

$$T(n) = O(1)$$

Doubly Linear Linked List - Add last



1. create node
2. if list is empty,
add newnode into head & tail
3. if list is not empty
 - a. add last node into prev of newnode
 - b. add newnode into next of last node
 - c. move tail on newnode

$$T(n) = O(1)$$



Thank you!!!

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