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How to identify if it is a DP Problem.

$[0, 0, -1]$
 $[0, 1, 1]$

1. Count total no. of ways
 2. There are multiple ways of doing it & you are asked to find Minimum/Maximum
- } Tend to do it Using Recursion

Try all possible ways

- Count
- Best Way

} This is when you try to apply recursion

(Shortcut to ~~DP~~ Recursion)

- 1) Try to represent the problem in terms of index
- 2) Do all possible stuffs on that index according to the problem statement
- 3) If Count All Ways \Rightarrow Sum up all stuffs
- If Find Min \Rightarrow Minimum of all stuffs
- If Find Max \Rightarrow Maximum of all stuffs

How to Convert Recursion into DP?

Memoization

- 1) Look at the parameters changing)
- 2) Declare an array with max size of parameter & initialize it to -1
- 3) Use the array & store the already computed recursive values.

Tabulation

In recursion you start from the top $f(5)$ and go to the bottom hence it is a Top-Down Approach, Tabulation is the exact opposite i.e., we calculate & Bottom-Up

Steps

- 1) Initialize the dp with $dp[n+1] = -1$
- 2) Check for Base Cases & write in the form of dp
In Frog Jump if $(ind == 0)$ return 0 $\Rightarrow dp[0] = 0;$

In Memoization/Recursion $jump1, jump2$ operations are being performed from $f(5) \rightarrow f(0)$ (since it is Top Down Approach)

Space Optimization

wherever you see $dp[ind-1]$, $dp[ind-2]$, you can perform space optimization. In Space Optimization instead of using array you store the values in variables.

FROG JUMP

0 $dp[i-2]$ $dp[i-1]$ $dp[i]$ 4 ----- $n-1$

↓ next step

0 1 $dp[i-2]$ $dp[i-1]$ $dp[i]$ 4 ----- $n-1$

At any position you only need $dp[i-1]$ & $dp[i-2]$ & you don't need anyone else. Therefore you store $dp[i-1]$ & $dp[i-2]$ in variables as $prev1$, $prev2$ instead of using array

As next steps you update $prev1 = curr$
 $prev2 = prev1$; ;