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## Unbounded Knapsack

You are given weight [ ]  $\Rightarrow$  weight of item  
values [ ]  $\Rightarrow$  value of item

$w$  = Maximum weight

Each item has infinite supplies (This is the diff btw 0/1 Knapsack & Unbound Knapsack).

Maximize the total value in Knapsack without exceeding capacity  $w$

Sol  
wt  $\rightarrow \{2, 4, 6\}$   
val  $\rightarrow \{5, 11, 13\}$

$$\boxed{w=10}$$

①  $\underset{13}{wt} - 6 + \underset{11}{wt} - 4 = 24$

②  $\underset{13}{wt} - 6 - \frac{(\underset{11}{wt} - 2) \times 2}{5 \times 2 = 10} = 23$

③  $(\underset{11 \times 2 = 22}{wt} \rightarrow 4) \times 2 + \underset{5}{wt} - 2 = 27$

④  $(\underset{5}{wt} - 2) 5 = 25$  Maximum possible solution.

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①  $f(ind, w) \Rightarrow$  till  $ind-1$  with baggage weight  $w$   
what will be the maximum value.

② Do possible stuff

$$\text{nottake} = 0 + f(ind-1, w);$$

$$\text{take} = \text{INT\_MIN};$$

$$\text{if } (\text{wt}[ind] \leq w) \{$$
  
$$\text{take} = \text{val}[ind] + f(ind, w - \text{wt}[ind])$$

return  $\max(\text{take}, \text{nottake})$  When there is infinite supply  
or can come to same ind twice  
don't increment ind twice,  
it's the same.

Base Cases

If there is only 1 element he should definitely steal it  
but the no. of times he can steal it matters.

$$\text{No. of times he can steal} = \frac{w}{\text{wt}[0]} \times \text{val}[0]$$

$\text{if } (ind == 0) \{$

$$\text{return } \left( \frac{w}{\text{wt}[0]} \right) \times \text{val}[0];$$

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## Recursion

```
public int f(int ind, int w, int[] val, int[] wt){  
    if(ind == 0){  
        return (w/wt[0]) * val[0];  
    }  
    int notTake = f(ind - 1, w, val, wt);  
    int take = 0;  
    if(wt[ind] <= w){  
        take = val[ind] + f(ind, w - wt[ind], val, wt);  
    }  
    return Math.max(take, notTake);  
}
```

~~public~~

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## Memoization

```
public int f(int ind, int w, int[] val, int[] wt, int[][] dp) {
    if (ind == 0) {
        return (w / wt[0]) * val[0];
    }
    if (dp[ind][w] != -1) {
        return dp[ind][w];
    }
    int notTake = f(ind - 1, w, val, wt, dp);
    int take = 0;
    if (wt[ind] <= w) {
        take = val[ind] + f(ind, w - wt[ind], val, wt, dp);
    }
    return dp[ind][w] = Math.max(take, notTake);
}
```

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### Tabulation

```
int [][] dp = new int [n] [(w+1)];
```

```
for (int w=0; w <= W; w++) {
```

```
    dp[0][w] = (w/wt[0]) * val[0];
```

```
}
```

```
for (int ind=1; ind < n; ind++) {
```

```
    for (int w=0; w <= W; w++) {
```

```
        int nottake = dp[ind-1][w];
```

```
        int take = 0;
```

```
        if (wt[ind] <= w) {
```

```
            take = val[ind] + dp[ind][w-wt[ind]];
```

```
}
```

```
        dp[ind][w] = Math.max(take, nottake);
```

```
}
```

```
}
```

```
return dp[n-1][w];
```

```
}
```

```
}
```

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## Space Optimization

```
int [] prev = new int[w+1]
int [] curr = new int[w+1]
for(int w=0; w<=w; w++) {
    prev[w] = (int)(w/wt[0]) * val[0];
}
for(int ind=1; ind<n; ind++) {
    for(int w=0; w<=w; w++) {
        int nottake = 0 + prev[w];
        int take = 0;
        if(wt[ind] <= w) {
            take = val[ind] + curr[w-wt[ind]];
        }
        curr[w] = max(take, nottake);
    }
    prev = curr;
}
return prev[w];
```