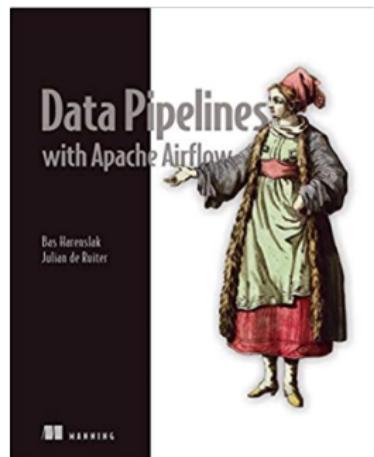


8.2: Map Reduce

- **Instructor:** Dr. GP Saggese -
gsaggese@umd.edu
- Resources
 - Silberschatz: Chap 10
 - Seminal papers
 - Sanjay Ghemawat, Howard Gobioff, and Shun-Tak Leung: The Google File System, 2003
 - Jeffrey Dean and Sanjay Ghemawat: MapReduce: Simplified Data Processing on Large Clusters, 2004



MapReduce: Overview

- **MapReduce programming model**

- Inspired by functional programming (e.g., Lisp)
- Common pattern of parallel programming
- Basic algorithm
 - Process large number of records
 - Apply `map()` to each record
 - Group results by key
 - Apply `reduce()` to results of `map()`

- **Example**

- **Goal:** Sum length of all tuples in a document
 - E.g.,
[(), (a,), (a, b), (a, b, c)]
- **map(function, set of values)**
 - Apply function to each value (e.g., `len`)
`map(len, [(), (a,), (a, b), (a, b, c)])` -> [0, 1, 2, 3]
- **reduce(function, set of values)**
 - Combine values using a binary function (e.g., `add`)
`reduce(add, [0, 1, 2, 3])` -> 6

MapReduce: Overview

- Structure of computation
 - **Read input**
 - Sequentially or in parallel
 - **Map**
 - Extract / compute from records
 - **Group by key**
 - Sort and shuffle
 - **Reduce**
 - Aggregate, summarize, filter, transform
 - **Write the result**
- MapReduce framework (e.g., Hadoop, Spark) implements algorithm
- User specifies `map()` and `reduce()` functions to solve problem

MapReduce: Word Count

- **Word Count**

- “Hello world” of MapReduce
- Huge text file (can’t fit in memory)
- Count occurrences of each distinct word

- **Sample application**

- Analyze web server logs for popular URLs

- **Linux solution** :::: columns :::: {.column width=70%}

```
> more doc.txt
```

```
One a penny, two a penny, hot cross buns.
```

```
> words doc.txt | sort | uniq -c
```

```
a 2
```

```
buns 1
```

```
cross 1
```

```
...
```

```
:::: :::: {.column width=30%}
```

Hot cross buns!



MapReduce: Word Count

Action

Read input

Map:

- Invoke `map()` on each input record
- Emit 0 or more output data items

Group by key:

- Gather all outputs from `map()` stage
- Collect outputs by keys

Reduce:

- Combine the list of outputs with same keys

Python code

```
values = read(file_name)

def map(values):
    # values: words in document
    for word in values:
        emit(word, 1)

def reduce(key, values):
    # key: a word
    # value: a list of counts
    result = 0
    # result = sum(values)
    for count in values:
        result += count
    emit(key, result)
```

Example

"One a penny, two a penny,
hot cross buns."

Map:

```
[("one", 1), ("a", 1),
("penny", 1), ("two", 1),
("a", 1), ("penny", 1),
("hot", 1), ("cross", 1),
("buns", 1)]
```

Group by key:

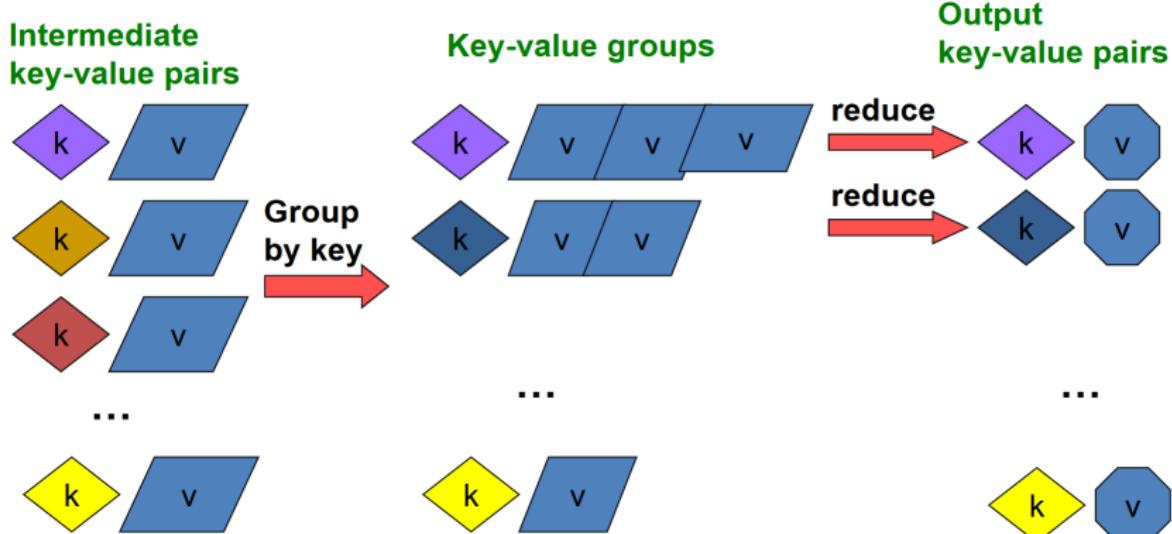
```
[("a", [1, 1]),
("buns", [1]),
("cross", [1]),
("hot", [1]),
("one", [1]),
("penny", [1, 1]),
("two", [1])]
```

Reduce:

```
[("one", 1), ("a", 2),
("penny", 2),
("two", 1),
("hot", 1),
("cross", 1),
("buns", 1)]
```

MapReduce: Reduce Step

```
reduce(key: word, values: List[int]):  
    # key: a word  
    # value: an iterator over counts  
    result = 0  
    for count in values:  
        result += count  
    emit(key, result)
```



MapReduce: Interfaces

- Input: read key-value pairs $\text{List}[\text{Tuple}[k, v]]$
- Programmer specifies two methods map and reduce
- $\text{Map}(\text{Tuple}[k, v]) \rightarrow \text{List}[\text{Tuple}[k, v]]$
 - Take a key-value pair and output a set of key-value pairs
 - E.g., key is a file, value is the number of occurrences
 - “One a penny” $\rightarrow [(\text{“One”, 1}), (\text{“a”, 1}), (\text{“penny”, 1})]$
 - There is one **Map** call for every (k, v) pair
- $\text{GroupBy}(\text{List}[\text{Tuple}[k, v]]) \rightarrow \text{List}[\text{Tuple}[k, \text{List}[v]]]]$
 - Group and optionally sort all the records with the reduce key
- $\text{Reduce}(\text{Tuple}[k, \text{List}[v]]) \rightarrow \text{Tuple}[k, v]$
 - All values v' with same key k' are reduced together
 - There is one **Reduce** call per unique key k'
- Output: write key-value pairs $\text{List}[\text{Tuple}[k, v]]$

MapReduce: Log Processing

- Log file recording access to a website with format date, hour, filename
- **Goal:** find how many times each files is accessed during Feb 2013
- **Input**
 - Read the file and split into lines
- **Map**
 - Parse each line into the 3 fields
 - If the date is in the required interval emit(dir_name, 1)
- **GroupBy**
 - The reduce key is the filename
 - Accumulate all the (key, value) with the same filename
- **Reduce**
 - Add the values for each list of (key, value) since they have the same filename
 - Output the number of access to each file
- **Output**
 - Write results on disk separated by newline

After Input

...
2013/02/21 10:31:22.00EST /slide-dir/11.ppt
2013/02/21 10:43:12.00EST /slide-dir/12.ppt
2013/02/22 18:26:45.00EST /slide-dir/13.ppt
2013/02/22 18:26:48.00EST /exer-dir/2.pdf
2013/02/22 18:26:54.00EST /exer-dir/3.pdf
2013/02/22 20:53:29.00EST /slide-dir/12.ppt
...

After Map

[`/slide-dir/11.ppt`, 1), ...]

After GroupBy

[`/slide_dir/11.ppt`, 1), ..., (`/slide-dir/12.ppt`, [1, 1]), ...]

After Reduce

[`/slide_dir/11.ppt`, 1), ..., (`/slide-dir/12.ppt`, 2), ...]

Output

/slide_dir/11.ppt 1

...

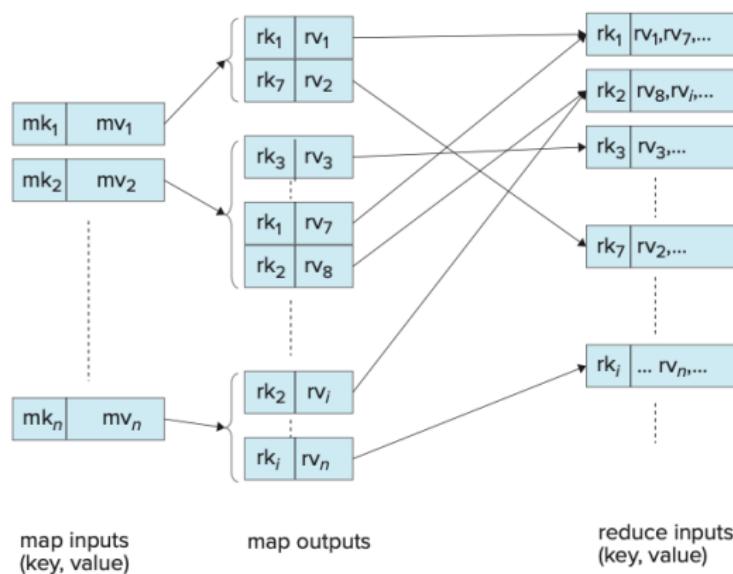
/slide-dir/12.ppt 2

...



MapReduce: Data Flow

Focusing on MapReduce functionality / flow of the data to expose the parallelism



Input Map

GroupBy Reduce .

- **Input**
- **Map**

- mki = map keys
- mvi = map input values

- **GroupBy**

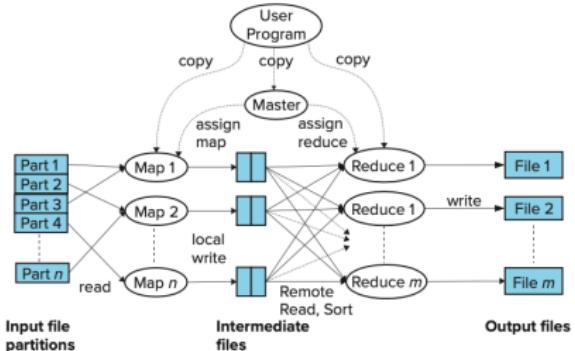
- Shuffle / collect the data

- **Reduce**

- rki = reduce keys
- rvi^{**} = reduce input values
- Reduce outputs are not shown

MapReduce: Parallel Data Flow

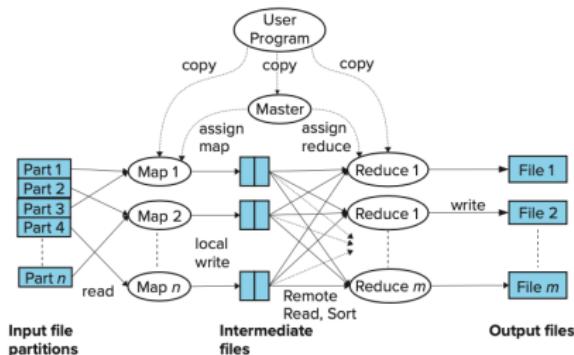
- User program specifies map/reduce code
- Input data is partitioned across multiple machines (HDFS)
- Master node sends copies of the code to all computing nodes
- Map
 - n data chunks to process
 - Functions executed in parallel on multiple k machines
 - Output data from Map is saved on disk
- GroupBy / Sort
 - Output data from Map is sorted and partitioned based on reduce key
 - Different files are created for each Reduce task
- Reduce
 - Functions executed in parallel on multiple machines
 - Each work on some part of the data
 - Output data from Reduce is saved on disk



- All operations use HDFS as storage
- Machines are reused for multiple computations (Map, GroupBy, Reduce) at different times

Master Node Responsibilities

- **Master node coordinates**
 - Task status: idle, in-progress, completed
 - Schedule idle tasks as workers become available
 - Map task completion sends location and sizes of intermediate files to Master
 - Master informs Reduce tasks
 - Schedule idle Reduce tasks
- **Master node pings workers to detect failures**



Dealing with Failures

- **Map worker failure**
 - Reset failed map tasks to idle
 - Notify reduce workers when task is rescheduled
- **Reduce worker failure**
 - Reset in-progress tasks to idle
 - Restart reduce task
- **Master failure**
 - Abort MapReduce task
 - Notify client

How many Map and Reduce jobs?

- M map tasks
- R reduce tasks
- N worker nodes
- Rules of thumb
 - $M \gg N$
 - Pros: Improve dynamic load balancing, Speed up recovery from worker failures
 - Cons: More communication between *Master* and *Worker Nodes*, Lots of smaller files
 - $R > N$
 - Usually $R < M$, Output is spread across fewer files

Refinements: Backup Tasks

- **Problem**

- Slow workers significantly lengthen the job completion time
- Slow workers due to:
 - Older processor
 - Not enough RAM
 - Other jobs on the machine
 - Bad disks
 - OS thrashing / virtual memory hell

- **Solution**

- Near the end of Map / Reduce phase
 - Spawn backup copies of tasks
 - Whichever one finishes first “wins”

- **Result**

- Shorten job completion time

Refinement: Combiners

- **Problem**

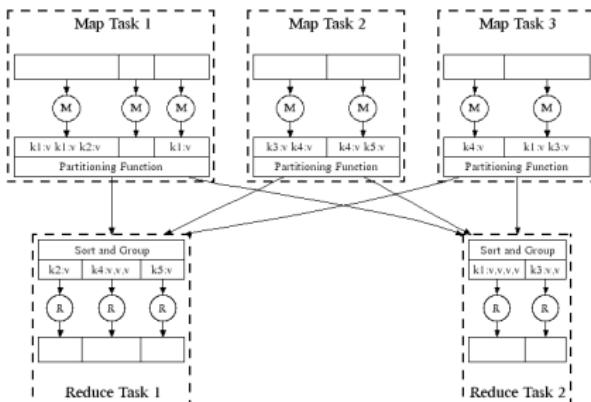
- Often a *Map* task produces many pairs for the same key k
[(k_1, v_1), (k_1, v_2), ...]
- E.g., common words in the word count example
- Increase complexity of the *GroupBy* stage

- **Solution**

- Pre-aggregate values in the *Map* with a *Combine*
[$k_1, (v_1, v_2, \dots)$, $k_2, ([\dots])$]
- *Combine* is usually the same as the *Reduce* function
- Works only if *Reduce* function is commutative and associative

- **Result**

- Better data locality
- Less shuffling and reordering
- Less network / disk traffic



Refinement: Partition Function

- **Problem**

- Sometimes users want to control how keys get partitioned
- Inputs to *Map* tasks are created by contiguous splits of input file
- MapReduce uses a default partition function **hash(key) mod R**
- Reduce needs to ensure that records with the same intermediate key end up at the same worker

- **Solution**

- Sometimes useful to override the hash function:
- E.g., **hash(hostname(URL)) mod R** ensures URLs from a host end up in the same output file