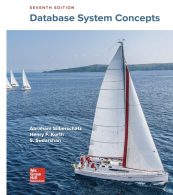


## 8.2: Map Reduce

- **Instructor:** Dr. GP Saggese - [gsaggese@umd.edu](mailto:gsaggese@umd.edu)
- **Resources**
  - Silberschatz: Chap 10
  - Ghemawat et al.: The Google File System, 2003
  - Dean et al.: MapReduce: Simplified Data Processing on Large Clusters, 2004



# MapReduce: Overview

- **MapReduce programming model**

- Inspired by functional programming (e.g., Lisp)
- Common pattern of parallel programming
- Basic algorithm
  - Process large number of records
  - Apply `map()` to each record
  - Group results by key
  - Apply `reduce()` to results of `map()`

- **Example**

- **Goal:** Sum length of all tuples in a document
  - E.g.,  
[() (a,) (a, b) (a, b, c)]
- **map(function, set of values)**
  - Apply function to each value (e.g., `len`)  
`map(len, [(), (a), (a, b), (a, b, c)]) -> [0, 1, 2, 3]`
- **reduce(function, set of values)**
  - Combine values using a binary function (e.g., `add`)  
`reduce(add, [0, 1, 2, 3]) -> 6`

# MapReduce: Overview

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- Structure of computation
  - **Read input**
    - Sequentially or in parallel
  - **Map**
    - Extract / compute from records
  - **Group by key**
    - Sort and shuffle
  - **Reduce**
    - Aggregate, summarize, filter, transform
  - **Write the result**
- MapReduce framework (e.g., Hadoop, Spark) implements algorithm
- User specifies `map()` and `reduce()` functions to solve problem

# MapReduce: Word Count

- **Word Count**
  - “Hello world” of MapReduce
  - Huge text file (can't fit in memory)
  - Count occurrences of each distinct word
- **Sample application**
  - Analyze web server logs for popular URLs
- **Linux solution**

```
> more doc.txt
One a penny, two a penny, hot cross buns.
> words doc.txt | sort | uniq -c
a 2
buns 1
cross 1
...
```

Hot cross buns!  
Hot cross buns!  
One a penny, two a penny,  
Hot cross buns!

If you have no daughters,  
Give them to your sons.  
One a penny, two a penny,  
Hot cross buns!<sup>[1]</sup>

- words outputs words one per line
- Unix pipeline is parallelizable in MapReduce sense



# MapReduce: Word Count

## Action

Read input

Map:

- Invoke **map()** on each input record
- Emit 0 or more output data items

Group by key:

- Gather all outputs from **map()** stage
- Collect outputs by keys

Reduce:

- Combine the list of outputs with same keys

## Python code

```
values = read(file_name)
```

```
def map(values):  
    # values: words in document  
    for word in values:  
        emit(word, 1)
```

```
def reduce(key, values):  
    # key: a word  
    # value: a list of counts  
    result = 0  
    # result = sum(values)  
    for count in values:  
        result += count  
    emit(key, result)
```

## Example

```
"One a penny, two a penny,  
hot cross buns."
```

Map:

```
[("one", 1), ("a", 1),  
 ("penny", 1), ("two", 1),  
 ("a", 1), ("penny", 1),  
 ("hot", 1), ("cross", 1),  
 ("buns", 1)]
```

Group by key:

```
[("a", [1, 1]),  
 ("buns", [1]),  
 ("cross", [1]),  
 ("hot", [1]),  
 ("one", [1]),  
 ("penny", [1, 1]),  
 ("two", [1])]
```

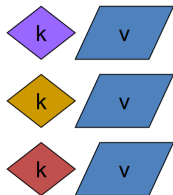
Reduce:

```
[("one", 1), ("a", 2),  
 ("penny", 2),  
 ("two", 1),  
 ("hot", 1),  
 ("cross", 1),  
 ("buns", 1)]
```

# MapReduce: Reduce Step

```
reduce(key: word, values: List[int]):  
    # key: a word  
    # value: an iterator over counts  
    result = 0  
    for count in values:  
        result += count  
    emit(key, result)
```

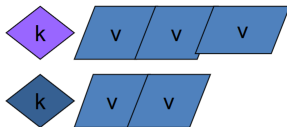
Intermediate  
key-value pairs



...



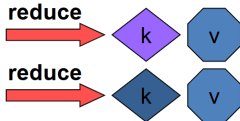
Key-value groups



...



Output  
key-value pairs



...



# MapReduce: Interfaces

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- Input: read key-value pairs **List[Tuple[k, v]]**
- Programmer specifies two methods map and reduce
- **Map(Tuple[k, v]) → List[Tuple[k, v]]**
  - Take a key-value pair and output a set of key-value pairs
    - E.g., key is a file, value is the number of occurrences
    - “One a penny” → [(“One”, 1), (“a”, 1), (“penny”, 1)]
  - There is one **Map** call for every  $(k, v)$  pair
- **GroupBy(List[Tuple[k, v]]) → List[Tuple[k, List[v]]]**
  - Group and optionally sort all the records with the reduce key
- **Reduce(Tuple[k, List[v]]) → Tuple[k, v]**
  - All values  $v'$  with same key  $k'$  are reduced together
  - There is one **Reduce** call per unique key  $k'$
- Output: write key-value pairs **List[Tuple[k, v]]**

# MapReduce: Log Processing

- Log file recording access to a website with format  
date, hour, filename
- **Goal:** find how many times each files is accessed during Feb 2013
- **Input**
  - Read the file and split into lines
- **Map**
  - Parse each line into the 3 fields
  - If the date is in the required interval  
emit(dir\_name, 1)
- **GroupBy**
  - The reduce key is the filename
  - Accumulate all the (key, value) with the same filename
- **Reduce**
  - Add the values for each list of (key, value) since they have the same filename
  - Output the number of access to each file
- **Output**
  - Write results on disk separated by

## After Input

```
...
2013/02/21 10:31:22.00EST /slide-dir/11.ppt
2013/02/21 10:43:12.00EST /slide-dir/12.ppt
2013/02/22 18:26:45.00EST /slide-dir/13.ppt
2013/02/22 18:26:48.00EST /exer-dir/2.pdf
2013/02/22 18:26:54.00EST /exer-dir/3.pdf
2013/02/22 20:53:29.00EST /slide-dir/12.ppt
...
```

## After Map

```
[`/slide-dir/11.ppt`, 1), ...]
```

## After GroupBy

```
[`/slide_dir/11.ppt`, 1), ...,
(`/slide-dir/12.ppt`, [1, 1]), ...]
```

## After Reduce

```
[`/slide_dir/11.ppt`, 1), ...,
(`/slide-dir/12.ppt`, 2), ...]
```

## Output

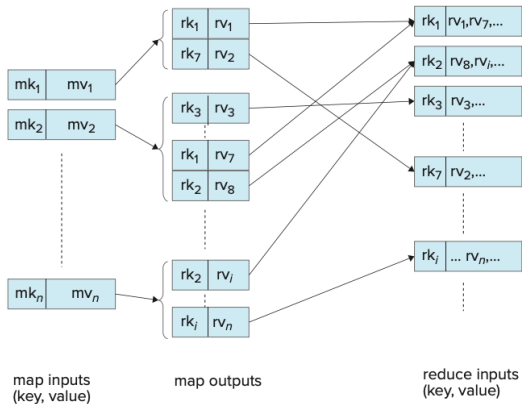
```
/slide_dir/11.ppt 1
...
/slide-dir/12.ppt 2
...
```





# MapReduce: Data Flow

Focusing on MapReduce functionality / flow of the data to expose the parallelism



- **Input**

- **Map**

- $mk_i$  = map keys
- $mvi$  = map input values

- **GroupBy**

- Shuffle / collect the data

- **Reduce**

- $rki$  = reduce keys
- $rvi^{**}$  = reduce input values
- Reduce outputs are not shown

Input

Map

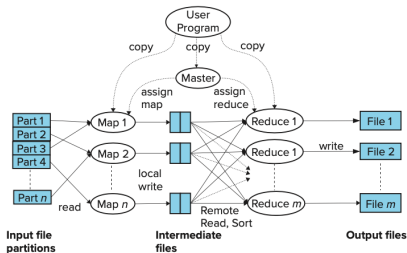
GroupBy

Reduce

.

# MapReduce: Parallel Data Flow

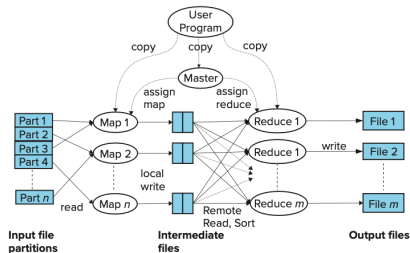
- **User program** specifies map/reduce code
- **Input data** is partitioned across multiple machines (HDFS)
- **Master** node sends copies of the code to all computing nodes
- **Map**
  - $n$  data chunks to process
  - Functions executed in parallel on multiple  $k$  machines
  - Output data from *Map* is saved on disk
- **GroupBy / Sort**
  - Output data from *Map* is sorted and partitioned based on reduce key
  - Different files are created for each *Reduce* task
- **Reduce**
  - Functions executed in parallel on multiple machines
  - Each work on some part of the data
  - Output data from *Reduce* is saved on disk



- All operations use HDFS as storage
- Machines are reused for multiple computations (Map, GroupBy, Reduce) at different times

# Master Node Responsibilities

- **Master node coordinates**
  - Task status: idle, in-progress, completed
  - Schedule idle tasks as workers become available
  - Map task completion sends location and sizes of intermediate files to Master
  - Master informs Reduce tasks
  - Schedule idle Reduce tasks
- **Master node pings workers to detect failures**



# Dealing with Failures

---

- **Map worker failure**
  - Reset failed map tasks to idle
  - Notify reduce workers when task is rescheduled
- **Reduce worker failure**
  - Reset in-progress tasks to idle
  - Restart reduce task
- **Master failure**
  - Abort MapReduce task
  - Notify client

# How many Map and Reduce jobs?

---

- **M** map tasks
- **R** reduce tasks
- **N** worker nodes
- Rules of thumb
  - $M \gg N$ 
    - Pros: Improve dynamic load balancing, Speed up recovery from worker failures
    - Cons: More communication between *Master* and *Worker Nodes*, Lots of smaller files
  - $R > N$ 
    - Usually  $R < M$ , Output is spread across fewer files

# Refinements: Backup Tasks

---

- **Problem**

- Slow workers significantly lengthen the job completion time
- Slow workers due to:
  - Older processor
  - Not enough RAM
  - Other jobs on the machine
  - Bad disks
  - OS thrashing / virtual memory hell

- **Solution**

- Near the end of Map / Reduce phase
  - Spawn backup copies of tasks
  - Whichever one finishes first “wins”

- **Result**

- Shorten job completion time

# Refinement: Combiners

## • Problem

- Often a *Map* task produces many pairs for the same key  $k$

$[(k1, v1), (k1, v2), \dots]$

- E.g., common words in the word count example
- Increase complexity of the GroupBy stage

## • Solution

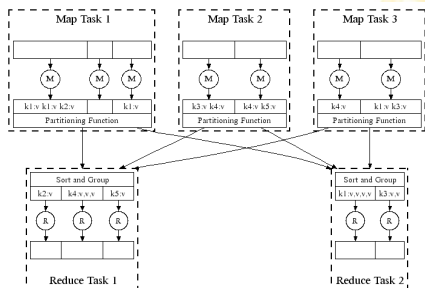
- Pre-aggregate values in the *Map* with a *Combine*

$[k1, (v1, v2, \dots), k2, ([\dots])]$

- Combine* is usually the same as the *Reduce* function
- Works only if *Reduce* function is commutative and associative

## • Result

- Better data locality
- Less shuffling and reordering
- Less network / disk traffic



# Refinement: Partition Function

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- **Problem**

- Sometimes users want to control how keys get partitioned
- Inputs to *Map* tasks are created by contiguous splits of input file
- MapReduce uses a default partition function  **$\text{hash}(\text{key}) \bmod R$**
- Reduce needs to ensure that records with the same intermediate key end up at the same worker

- **Solution**

- Sometimes useful to override the hash function:
- E.g.,  **$\text{hash}(\text{hostname}(\text{URL})) \bmod R$**  ensures URLs from a host end up in the same output file