

Problem 1

Step 1: Array of distances

	h	u	a	m	p	f	e	s
h	0	3	4	inf	inf	inf	inf	inf
u	inf	0	6	1	6	inf	inf	inf
a	inf	inf	0	2	7	inf	inf	inf
m	inf	inf	inf	0	2	5	5	inf
p	inf	inf	inf	inf	0	3	4	inf
f	inf	inf	inf	inf	inf	0	inf	2
e	inf	inf	inf	inf	inf	inf	0	3
s	inf	inf	inf	inf	inf	inf	inf	0

Priority Queue:

V	H	U	A	M	P	F	E	S
D[v]	0	inf	inf	inf	inf	inf	inf	inf
Prev[v]	null	null	null	null	null	null	null	null
Visited	false	false	false	false	false	false	false	false

Step 2:

Update adjacent edges to the initial vertex and updating H as being visited

V	H	U	A	M	P	F	E	S
D[v]	0	3	4	inf	inf	inf	inf	inf
Prev[v]	null	H	H	null	null	null	null	null
Visited	true	false	false	false	false	false	false	false

U is minimum now

Step 3: Update adjacent edges to the initial vertex and updating U as being visited

V	H	U	A	M	P	F	E	S
D[v]	0	3	4	4	9	inf	inf	inf
Prev[v]	null	H	H	U	U	null	null	null
Visited	True	True	false	false	false	false	false	false

Step 4: Update adjacent edges to the initial vertex and updating A as being visited

V	H	U	A	M	P	F	E	S
D[v]	0	3	4	4	9	inf	inf	inf
Prev[v]	null	H	H	U	U	null	null	null
Visited	True	True	True	false	false	false	false	false

Step 5: Update adjacent edges to the initial vertex and updating M as being visited

V	H	U	A	M	P	F	E	S
D[v]	0	3	4	4	9	9	9	inf
Prev[v]	null	H	H	U	M	M	M	null
Visited	True	True	True	True	false	false	false	false

Step 6: Update adjacent edges to the initial vertex and updating P as being visited

V	H	U	A	M	P	F	E	S
D[v]	0	3	4	4	9	9	9	inf
Prev[v]	null	H	H	U	M	M	M	null
Visited	True	True	True	True	True	false	false	false

Step 7: Update adjacent edges to the initial vertex and updating F as being visited

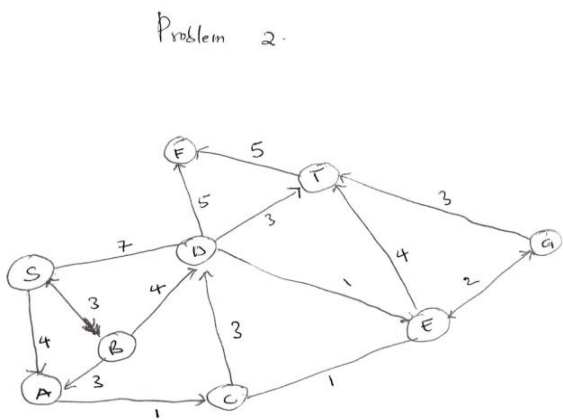
V	H	U	A	M	P	F	E	S
D[v]	0	3	4	4	9	9	9	11
Prev[v]	null	H	H	U	M	M	M	F
Visited	True	True	True	True	True	True	false	false

Step 8: Update adjacent edges to the initial vertex and updating E as being visited

V	H	U	A	M	P	F	E	S
D[v]	0	3	4	4	9	9	9	11
Prev[v]	null	H	H	U	M	M	M	F
Visited	True	True	True	True	True	True	True	false

Minimum Path: F-M-U-H

Problem 2



Note: Not drawn to Scale.

Problem 3

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	-1	7	-1	-1	-1	-1
Prev[v]	null	null	null	null	null	null	null	null	null
Visited	False	False	False	False	False	False	False	False	False

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	-1	7	-1	-1	-1	-1
Prev[v]	null	S	S	null	S	null	null	null	null
Visited	True	False	False	False	False	False	False	False	False

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	-1	7	-1	-1	-1	-1
Prev[v]	null	S	S	null	B	null	null	null	null
Visited	True	False	True	False	False	False	False	False	False

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	5	7	-1	-1	-1	-1
Prev[v]	null	S	S	A	B	null	null	null	null
Visited	True	True	True	False	False	False	False	False	False

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	5	7	10	-1	-1	-1
Prev[v]	null	S	S	A	B	C	null	null	null
Visited	True	True	True	True	False	False	False	False	False

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	5	7	10	15	-1	13
Prev[v]	null	S	S	A	B	C	D	null	D
Visited	True	True	True	True	True	False	False	False	False

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	5	7	10	15	21	13
Prev[v]	null	S	S	A	B	C	D	E	D
Visited	True	True	True	True	True	True	False	False	False

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	5	7	10	15	21	13
Prev[v]	null	S	S	A	B	C	D	E	D
Visited	True	True	True	True	True	True	False	False	True

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	5	7	10	15	21	13
Prev[v]	null	S	S	A	B	C	D	E	D
Visited	True	True	True	True	True	True	True	False	True

V	S	A	B	C	D	E	F	G	T
D[v]	0	4	3	5	7	10	15	21	13
Prev[v]	null	S	S	A	B	C	D	E	D
Visited	True	True	True	True	True	True	True	True	True

Final minimum path using Dijkstra's algorithm is S-B-D-T

Problem 4

Matrix when path is undirected

	S	A	B	C	D	E	F	G	T
S	0	4	3	-1	7	-1	-1	-1	-1
A	4	0	-1	1	-1	-1	-1	-1	-1
B	3	-1	0	-1	4	-1	-1	-1	-1
C	-1	1	-1	0	3	1	-1	-1	-1
D	7	-1	4	3	0	1	5	-1	3
E	-1	-1	-1	-1	-1	0	-1	2	4
F	-1	-1	-1	-1	5	-1	0	-1	-1
G	-1	-1	-1	-1	-1	2	-1	0	3
T	-1	-1	-1	-1	3	4	-1	3	0

Problem 5

It is impossible to work with negative cost in Dijkstra's algorithm because with negative costs, it is likely the path with the minimum cost will be infinity.

Problem 6

To find the largest-cost path, we would have to negate all the cost values on each path. Doing so will make it impossible to use the Dijkstra's algorithm since it is not applicable to negative costs.