

Algebra

A1 Let n be a positive integer and let a_1, \dots, a_{n-1} be arbitrary real numbers. Define the sequences u_0, \dots, u_n and v_0, \dots, v_n inductively by $u_0 = u_1 = v_0 = v_1 = 1$, and $u_{k+1} = u_k + a_k u_{k-1}$, $v_{k+1} = v_k + a_{n-k} v_{k-1}$ for $k = 1, \dots, n-1$.

Prove that $u_n = v_n$.

Solution. For each n , denote the set S_n as the set of subsets of $\{1, 2, \dots, n\}$ such that a set $W \in S_n$ iff W does not contain a pair of neighbouring numbers (that is, if $i \in W$

then $i+1 \notin W$). We claim that $u_n = \sum_{\{x_1, \dots, x_k\} \in S_{n-1}} \prod_{i=1}^k a_{x_i}$ (if $W = \emptyset$ then the term is the

constant term 1). We will go by induction. For $n = 2$ we have $u_2 = u_1 + a_1 u_0 = 1 + a_1$, and S_1 has elements $\{\}$ and $\{1\}$. For $n = 3$ we have $(1 + a_1) + a_2$ and it turned out that S_2 has elements $\{\}$, $\{1\}$ and $\{2\}$ ($\{1, 2\}$ is invalid). These settle the base cases. For inductive step, suppose for some n , we have u_n and u_{n-1} following the recurrence relations. This means we have:

$$u_{n+1} = u_n + a_n u_{n-1} = \sum_{\{x_1, \dots, x_k\} \in S_{n-1}} \prod_{i=1}^k a_{x_i} + a_n \cdot \sum_{\{y_1, \dots, y_m\} \in S_{n-2}} \prod_{i=1}^m a_{y_i}$$

We first notice that S_{n-1} is a subset of S_n that contains all subsets in S_n not containing n , thus $\{x_1, \dots, x_k\} \in S_{n-1}$ handles this. We also recognize that when $\{y_1, \dots, y_m\} \in S_{n-2}$, $\{y_1, \dots, y_m\} \cup \{n\} \in S_n$ since $\max\{y_i\} \leq n-2$. Conversely, this is also a necessary condition for $\{y_1, \dots, y_m\} \cup \{n\} \in S_n$ since $n-1$ cannot be in it. Combining these together, we see that S_{n-1} handles every element in S_n not containing n and S_{n-2} with n appended to each element handles every element in S_n containing n . This shows that u_{n+1} fulfills this property too.

Now that $u_n = \sum_{\{x_1, \dots, x_k\} \in S_{n-1}} \prod_{i=1}^k a_{x_i}$, we can deduce similarly that

$$v_n = \sum_{\{x_1, \dots, x_k\} \in S_{n-1}} \prod_{i=1}^k a_{n-x_i} = \sum_{\{n-x_1, \dots, n-x_k\} \in S_{n-1}} \prod_{i=1}^k a_{x_i}$$

But then $\{x_1, \dots, x_k\} \in S_{n-1}$ iff $\{n-x_1, \dots, n-x_k\} \in S_{n-1}$ by symmetry, so the coefficient of $a_{x_1} \cdots a_{x_k}$ of u_n and v_n are both equal regardless of the combinations of x_i 's. This proves that u_n and v_n .

A3 (IMO 5) Let $\mathbb{Q}_{>0}$ be the set of all positive rational numbers. Let $f : \mathbb{Q}_{>0} \rightarrow \mathbb{R}$ be a function satisfying the following three conditions:

- (i) for all $x, y \in \mathbb{Q}_{>0}$, we have $f(x)f(y) \geq f(xy)$;
- (ii) for all $x, y \in \mathbb{Q}_{>0}$, we have $f(x+y) \geq f(x) + f(y)$;
- (iii) there exists a rational number $a > 1$ such that $f(a) = a$.

Prove that $f(x) = x$ for all $x \in \mathbb{Q}_{>0}$.

Solution. Notice that (i) has the more general form:

$$f(a_1 a_2 \cdots a_n) \leq f(a_1) f(a_2 \cdots a_n) \leq f(a_1) f(a_2) f(a_3 \cdots a_n) \leq \cdots \leq f(a_1) f(a_2) \cdots f(a_n)$$

and same goes to (ii):

$$f(a_1 + \cdots + a_n) \geq f(a_1) + f(a_2 + \cdots + a_n) \geq \cdots \geq f(a_1) + \cdots + f(a_n)$$

We first show that f produces only positive values. We now have the following:

- Plugging $x := a, y := 1$ into (i) we get $f(a)f(1) \geq f(a)$, or $af(1) \geq a$. This means $f(1) \geq 1$ since $a > 0$.
- Next, by considering the general form of (ii), we let $a_1 = a_2 = \cdots a_1 = 1$ and get $f(n) \geq nf(1) \geq n$, so for each positive integer n , $f(n)$ is also positive.
- For any rational number p/q (with p and q integers) we have, from (i), $f(p/q)f(q) \geq f(p) \geq p$, so $f(p/q)f(q) > 0$. Since $f(q) > 0$, we have $f(p/q) > 0$.

Having this in mind, we notice from above that $f(n) \geq n$ so for any positive non-integer rational number x with $\lfloor x \rfloor = n$, we have $n - \lfloor x \rfloor > 0$, so $f(x) \geq f(n) + f(n - \lfloor x \rfloor) > f(n) \geq n$, so $f(x) \geq n$ if $x > n$. In other words, $f(x) \geq f(\lfloor x \rfloor)$ with equality holding only when x is an integer. More generally, for all $z > x$, there's $y > 0$ with $x + y = z$ (and if $z, x \in \mathbb{Q}$ then so is y) so $f(z) = f(x + y) \geq f(x) + f(y) > f(x)$, so f is also increasing. We therefore have $f(x) \geq f(\lfloor x \rfloor) \geq \lfloor x \rfloor > x - 1$, in particular.

With this, we consider the sequences $f(a^n), n = 1, 2, \dots$. From the generalized (i), we can plug $a_1 = \cdots a_n = a$ to get $a^n = f(a)^n \geq f(a^n)$. To claim that equality must hold, suppose otherwise, and $f(a^n) < a^n$ for some $n > 0$. Then by (i) we have the following:

$$a^m f(a^n) \geq f(a^n) f(a^m) \geq f(a^{n+m})$$

Let $\epsilon = a^n - f(a^n) > 0$, then $f(a^{n+m}) \leq a^m(a^n - \epsilon)$. Given that $a > 1$, for m sufficiently large we have $a^m \epsilon > 1$. This means that $f(a^{n+m}) \leq a^m(a^n - \epsilon) < a^{m+n} - 1$. But this contradicts what we had before: $f(x) > x - 1$. Therefore, $f(a^m) = a^m$ for all positive integers m .

Next, we show that $f(x) \leq x$ for all x . Suppose otherwise, then for each integer n the generalized form of (ii) gives $f(nx) = nf(x)$ by plugging $a_i = x, \forall i = 1, \dots, n$. If $\delta > 0$ is such that $f(x) = x + \delta$, then $f(nx) \geq n(x + \delta)$ and for sufficiently large n we have $f(nx) \geq nx + 1$. Choose one such n , for now. Now by (ii) again, for each positive integer m we have $f(nx + m) \geq f(nx) + m \geq nx + m + 1$. Since $a > 1$ again, there exists p such that $a^p > nx + 1$, and let m be the maximum integer with $nx + m \leq a^p$, i.e. $nx + m + 1 > a^p$. We have $nx + m \leq a^p$, but $f(nx + m) \geq nx + m + 1 > a^p = f(a^p)$. This contradicts our previous claim that f is increasing.

Therefore $f(n) \leq n$ for each positive integer n , and coupled with $f(n) \geq n$ from before we get $f(n) = n$ for each positive integer n . Finally for each rational number p/q we have $f(q)f(p/q) \geq f(p)$, so $qf(p/q) \geq p$, or $f(p/q) \geq p/q$. But we have shown that $f(p/q) \leq p/q$ must hold, too hence $f(p/q) = p/q$ for each rational number p/q .

A4 Let n be a positive integer, and consider a sequence a_1, a_2, \dots, a_n of positive integers. Extend it periodically to an infinite sequence a_1, a_2, \dots by defining $a_{n+i} = a_i$ for all $i \geq 1$. If

$$a_1 \leq a_2 \leq \cdots \leq a_n \leq a_1 + n$$

and

$$a_{a_i} \leq n + i - 1 \quad \text{for } i = 1, 2, \dots, n,$$

prove that

$$a_1 + \cdots + a_n \leq n^2.$$

Solution. The fact that $a_{a_1} \leq n$ means that there exists $i \in [1, n]$ with $a_i \leq n$ and since $a_1 \leq a_2 \leq \cdots a_n, a_1 \leq n$. Next, since $a_1, a_2, \dots, a_{a_1} \leq n$, we can consider the numbers $a_{a_1}, a_{a_2}, \dots, a_{a_{a_1}}$, which are all smaller than a_n . Also, each number is bounded by all the terms following itself, so if $x \leq a_k$ we have $a_x \leq a_{a_k} \leq n + k - 1$. We can now split the

numbers up into the following:

$$\begin{aligned}
\sum_{i=1}^n a_i &= \sum_{i=1}^{a_1} a_i + \sum_{i=a_1+1}^{a_2} a_i + \cdots + \sum_{i=a_{a_1}+1}^n a_i \\
&\leq \sum_{i=1}^{a_1} a_i + \sum_{i=a_1+1}^{a_2} (n+1) + \cdots + \sum_{i=a_{a_1}+1}^n (a_1+n) \\
&= \sum_{i=1}^{a_1} a_i + (n+1)(a_2 - a_1) + (n+2)(a_3 - a_2) + \cdots + (n+(a_1-1))(a_{a_1} - a_{a_1-1}) \\
&\quad + (n+a_1)(n - a_{a_1}) \\
&= \sum_{i=1}^{a_1} a_i + (n+1)(a_2 - a_1) + n \left(\sum_{i=2}^{a_1} a_i - a_{i-1} \right) + n(n - a_{a_1}) - a_1 + \sum_{i=2}^{a_1} a_i((i-1) - i) + na_1 \\
&= \sum_{i=1}^{a_1} a_i + n(n - a_1) - \sum_{i=1}^{a_1} a_i + na_1 \\
&= n^2
\end{aligned}$$

as desired.

- A5** Let $\mathbb{Z}_{\geq 0}$ be the set of all nonnegative integers. Find all the functions $f : \mathbb{Z}_{\geq 0} \rightarrow \mathbb{Z}_{\geq 0}$ satisfying the relation

$$f(f(f(n))) = f(n+1) + 1$$

for all $n \in \mathbb{Z}_{\geq 0}$.

Answer. There are two family of functions. One is $f(n) = n + 1$. The other is given by: $f(2) = 3, f(3) = 0, f(0) = 1, f(1) = 6$, and $f(n+4) = f(n) + 4$ for all $n \geq 0$.

Solution. The key idea is to analyze $f^k(n)$, i.e. f applied to n for k times. The condition gives $f^3(n) = f(n+1) + 1$.

We first notice the following:

$$f^4(n+1) = f^3(f(n+1)) = f(f(n+1) + 1) + 1 = f(f^3(n)) + 1 = f^4(n) + 1$$

so starting from $n = 0$ we can inductively deduce that $f^4(n) = f^4(0) + n$. This would also mean that for all integers $m \neq n$ we have $f^4(m) \neq f^4(n)$, which means f^4 is injective. It then follows that f must also be injective. If f were to be surjective, then so is f^3 . In particular, there's n_1 with $f(n_1) = 0$, n_2 with $f(n_2) = n_1$ and n_3 with $f(n_3) = n_2$. This gives $f^3(n_3) = 0$ and therefore $f(n_3 + 1) = -1$, contradicting that f only takes nonnegative values.

For convenience we denote $f^k(\mathbb{N}_{\geq 0})$ as $\{f^k(n) : n \in \mathbb{N}_{\geq 0}\}$. Since $f^k(\mathbb{N}_{\geq 0}) \subseteq f^\ell(\mathbb{N}_{\geq 0})$ for all $k \leq \ell$ and $f^4(\mathbb{N}_{\geq 0}) = \mathbb{N}_{\geq 0} \setminus \{0, 1, \dots, f^4(0) - 1\}$, we know that $\mathbb{N}_{\geq 0} \setminus f(\mathbb{N}_{\geq 0})$ is finite. Let $\mathbb{N}_{\geq 0} \setminus f(\mathbb{N}_{\geq 0}) = \{x_1, \dots, x_m\}$. We show by induction that $|\mathbb{N}_{\geq 0} \setminus f^k(\mathbb{N}_{\geq 0})| = mk$, with the base case given by $k = 1$. In fact, we will verify that for each $m \geq 1$, $\mathbb{N}_{\geq 0} \setminus f^{k+1}(\mathbb{N}_{\geq 0}) = (\mathbb{N}_{\geq 0} \setminus f^k(\mathbb{N}_{\geq 0})) \cup \{f^k(x_1), \dots, f^k(x_m)\}$. There are a few things to take care of:

- $f^k(x_1), \dots, f^k(x_m)$ are all distinct since f is injective (and so is f^k). They are also disjoint from $\mathbb{N}_{\geq 0} \setminus f^k(\mathbb{N}_{\geq 0})$ since these are values in f^k .
- Suppose for some x_i , there's some y with $f^{k+1}(y) = f^k(x_i)$. As mentioned, f^k is injective, so $f^k(f(y)) = f^k(x_i)$ means $f(y) = x_i$. This contradicts that x_i is not a value of f and so $\{f^k(x_1), \dots, f^k(x_m)\} \subseteq \mathbb{N}_{\geq 0} \setminus f^{k+1}(\mathbb{N}_{\geq 0})$.
- Finally if $y \neq x_i$ then $y = f(z)$ for some z and therefore $f^k(y) = f^{k+1}(z)$.

This justifies the claim of the set size.

Another realization is that, by the given identity, $f^3(\mathbb{N}_{\geq 0}) = f(\mathbb{N}_{\geq 0} + 1) + 1$. To determine $f(\mathbb{N}_{\geq 0} + 1) + 1$, we notice that:

- $f(\mathbb{N}_{\geq 0}) = \mathbb{N}_{\geq 0} \setminus \{x_1, \dots, x_k\}$.
- $f(\mathbb{N}_{\geq 0} + 1) = \mathbb{N}_{\geq 0} \setminus \{x_1, \dots, x_k, f(0)\}$.
- $(\mathbb{N}_{\geq 0} + 1) + 1 = \mathbb{N}_{\geq 0} \setminus \{0, x_1 + 1, \dots, x_k + 1, f(0) + 1\}$

and notice that 0 is different from $x_1 + 1, \dots, x_k + 1, f(0) + 1$ (since these numbers are all ≥ 1), and $f(0)$ is also different from x_1, \dots, x_k (x_i 's are not values of f but $f(0)$ is). Thus $|\mathbb{N}_{\geq 0} \setminus f(\mathbb{N}_{\geq 0} + 1) + 1| = k + 2$. Combined with $|\mathbb{N}_{\geq 0} \setminus f^3(\mathbb{N}_{\geq 0})| = 3k$ we have $3k = k + 2$, so $k = 1$.

Now that we have $k = 1$, it's rather easy so see how to proceed. Since $f^4(\mathbb{N}_{\geq 0}) = \{n \in \mathbb{N}_{\geq 0} : n \geq f^4(0)\}$, we have $f^4(0) = 4k = 4$, so $f^4(n) = n + 4$. This also means f cannot have any fixed point. That is, there's no n with $f(n) = n$, as we will see later. In an easier way, we have one value x_1 not a value of f , which also means $f(x_1), f^2(x_1), f^3(x_1)$ are not a value of f^4 . By the injectivity of f we also have $\{x_1, f(x_1), f^2(x_1), f^3(x_1)\} = \{0, 1, 2, 3\}$. Now we have a few cases:

- If $x_1 = 0$, then $f(1) + 1 = f^3(0) \leq 3$ so $f(1) \leq 2$. $f(1) = 0$ is impossible since here 0 is not a value of f . $f(1) = 1$ implies $f^k(1) = 1$ for all k , contradicting $f^4(1) = 1 + 4 = 5$. Thus we can only have $f(1) = 2$, i.e. $f^3(0) = 3$. which means the sequence $(x_1, f(x_1), f^2(x_1), f^3(x_1))$ must be in the form $(0, 1, 2, 3)$. We will deal with this later.
- If $x_1 = 1$, then by $f(2) + 1 = f^3(x_1) \leq 3$ we have $f(2) \leq 2$. Again by the similar logic above, $f(2) \neq 2$ so $f(2) = 0$ or $f(2) = 1$. Since 1 is not a value of f here, we can only have $f(2) = 0$. But this forces $f^3(1) = f(2) + 1 = 0 + 1 = 1$, contradicting that 1 is not a value of f (and hence not a value of f^3).
- If $x_1 = 2$, by the similar logic above we have by the same logic above we have $f(3) \leq 2$. Again 2 is not a value of f here so we have $f(3) = 0$ or $f(3) = 1$. In the first case, $f(3) = 0$ so $f(2) = 1$, which means the sequence $(x_1, f(x_1), f^2(x_1), f^3(x_1))$ has the form $(2, 3, 0, 1)$. We will come back and deal with this later. In the second case, $f(3) = 1$ so $f(2) = 2$, again contradicting f cannot have fixed point.
- If $x_1 = 3$ then we have $f(4) \leq 2$. Nevertheless, $f(4) = f^5(0) \in f^4(\mathbb{N}_{\geq 0})$ so $f(4) \geq 4$, which is a contradiction.

So our sequence of $(x_1, f(x_1), f^2(x_1), f^3(x_1))$ must go in the sequence $(0, 1, 2, 3)$ or $(2, 3, 0, 1)$. With $f^4(n) = f^4(0) + n = n + 4$, for each n and k we have $f^{4k}(n) = n + 4k$ (by repeatedly apply f^4 , which is k times addition of 4). This means that $f(n + 4k) = f(f^{4k}(n)) = f(n) + 4k$, too, so f is uniquely determined by $f(0), f(1), f(2), f(3)$. In the first case $(0, 1, 2, 3)$ we have $f(n) = n + 1$, which works easily. In the trickier second case, the function is extrapolated in the form $(2, 3, 0, 1, 6, 7, 4, 5, \dots)$. We now have $f^3(2) = 1 = 0 + 1 = f(3) + 1$, $f^3(3) = f^4(2) = 6 = 5 + 1 = f(4) + 1$, $f^3(0) = 7 = 6 + 1 = f(1) + 1$ and $f^3(1) = 4 = 3 + 1 = f(0) + 1$. Also if $f^3(n) = f(n + 1) + 1$ then $f^3(n + 4) = f^3(f^4(n)) = f^4(f(n) + 1) = f(n) + 1 + 4 = f(n) + 5 = f(n + 4) + 1$ (since $f(n) + 4 = f^4(f(n)) = f(f^4(n)) = f(n + 4)$) so the fact that our identity holds for $n = 0, 1, 2, 3$ means it holds for all n too. This also gives the following construction for f : $f(n) = n + 1$ for n even, $n - 3$ for $n \equiv 3 \pmod{4}$ and $n + 5$ for $n \equiv 1 \pmod{4}$, as claimed.

Combinatorics

- C1** Let n be an positive integer. Find the smallest integer k with the following property;
Given any real numbers a_1, \dots, a_d such that $a_1 + a_2 + \dots + a_d = n$ and $0 \leq a_i \leq 1$ for

$i = 1, 2, \dots, d$, it is possible to partition these numbers into k groups (some of which may be empty) such that the sum of the numbers in each group is at most 1.

Answer. $2n - 1$.

Solution. To see that $k = 2n - 1$ is necessary, consider $d = 2n - 1$ and $a_i = \frac{n}{2n-1}$ for all i . Since each $a_i = \frac{n}{2n-1} > \frac{n}{2n} = \frac{1}{2}$, each a_i must be in its own isolated group, hence $2n - 1$ groups are necessary.

To see that $2n - 1$ is sufficient, we sort the numbers into $a_1 \geq a_2 \geq \dots \geq a_d$ and place them into the $2n - 1$ groups in the following algorithmic manner: we let $a_1, a_2, \dots, a_{2n-1}$ go into their individual groups (assuming $d \geq 2n - 1$ otherwise the problem is trivial). Then for each $j \geq 2n$, we place a_j sequentially into the group that has sum of numbers at most $1 - a_j$.

To show that we can always find those groups, we first note that $\sum_{i=1}^j a_i \leq \sum_{i=1}^d a_i = n$ and since $a_i \geq a_j$ for all $i \leq j$, $n \geq \sum_{i=1}^j a_i \geq \sum_{i=1}^j a_j = ja_j \geq 2na_j$ (since $j \geq 2n$), and therefore $a_j \leq \frac{1}{2}$. The above identity also implies $\sum_{i=1}^{j-1} a_i \leq n - a_j$, and the first summation is the total size (i.e. sum of numbers) of groups by considering only a_1, \dots, a_{j-1} . The average group size under this consideration is currently $\frac{\sum_{i=1}^{j-1} a_i}{2n-1} \leq \frac{n-a_j}{2n-1}$. But also notice the following:

$$\frac{\sum_{i=1}^{j-1} a_i}{2n-1} + a_j \leq \frac{n-a_j}{2n-1} + a_j = \frac{n + (2n-2)a_j}{2n-1} \leq \frac{n + (2n-2) \cdot \frac{1}{2}}{2n-1} = 1$$

so the average of the size of groups are currently at most $1 - a_j$. This means, there must be a group of size at most $1 - a_j$, so we are done.

C2 (IMO 2) In the plane, 2013 red points and 2014 blue points are marked so that no three of the marked points are collinear. One needs to draw k lines not passing through the marked points and dividing the plane into several regions. The goal is to do it in such a way that no region contains points of both colors.

Find the minimal value of k such that the goal is attainable for every possible configuration of 4027 points.

Answer. 2013.

Solution. To show that 2013 is a lower bound, consider the regular polygon of 4027 vertices $A_i, i = 1, 2, \dots, 4027$. Moreover, let A_i blue if i odd, and A_i red if i even. Since each A_i, A_{i+1} are of different colour for $i = 1, \dots, 4026$, there must be at least a line passing through the segment $A_i A_{i+1}$. On the other hand, each line can only pass through two sides of a regular polygon, so at least $\frac{4026}{2}$ lines needed.

Now to show that 2013 is sufficient, we generalize the problem into any configuration of $2k + 1$ points (with no three collinear) with binary colours, and show that k lines suffice. To deal with $k = 1$ as base case, if all three points are the same colour we are easily done. Otherwise, there must be two points (say A_1 and A_2) of the same colour, and a point (say B) of the other colour. Now the segments BA_1 and BA_2 cannot be collinear because no three lines are collinear, so we can draw a line cutting through the internal (open) segments BA_1 and BA_2 .

For inductive step, consider an arbitrary configuration of $2k + 1$ points, and the convex hull \mathcal{P} of the configuration. Consider any line ℓ determined by any side of \mathcal{P} ; all other lines lie on the same side of ℓ . If any of the sides have both points of the same colour,

then we can draw a line to isolate the points from the others, and the rest follows from induction hypothesis: draw $k - 1$ lines to handle the case for the remaining $2k - 1$ points.

Otherwise, the points on \mathcal{P} must be alternating in colours, so we choose a side with a red vertex R and a blue vertex B randomly and consider the $k - 1$ lines that separate the remaining $2k - 1$ points into regions with no region containing points of both colours. If R and B are in the same region, then this region must contain either only red points and blue points other than R and B (we call this region “red region” or “blue region” in these respective cases). Since R and B are on the convex hull, we can draw a line to isolate R if this region is blue, or a line to isolate B if this region is red. Otherwise, R and B are in different regions, and we can now draw a line to isolate R and B from the rest of the $2k - 1$ points. This way, R and B are each isolated, and we are done.

C3 A crazy physicist discovered a new kind of particle which he called an imon, after some of them mysteriously appeared in his lab. Some pairs of imons in the lab can be entangled, and each imon can participate in many entanglement relations. The physicist has found a way to perform the following two kinds of operations with these particles, one operation at a time.

- (i) If some imon is entangled with an odd number of other imons in the lab, then the physicist can destroy it.
- (ii) At any moment, he may double the whole family of imons in the lab by creating a copy I' of each imon I . During this procedure, the two copies I' and J' become entangled if and only if the original imons I and J are entangled, and each copy I' becomes entangled with its original imon I ; no other entanglements occur or disappear at this moment.

Prove that the physicist may apply a sequence of much operations resulting in a family of imons, no two of which are entangled.

Solution. Throughout the solution we use the graph theory language: vertices and edges. We first apply (i) repeatedly until we can't do so (it terminates anyway since (i) decreases the number of vertices and there are only finitely many of them). Then, while the total number of edges is still positive, we do the following on the current graph $G = (V, E)$

- Starting with the configuration where the degree of all vertices are even (and total nonzero), perform operation (ii), and let the copied vertices to be V' with the copy of each $v \in V$ as v' . Now all vertices have odd degree.
- For each vertex $v \in V$ that had initial degree 0 before (ii) before, it now has exactly one neighbour v' . We can delete v' for all such v and the degree of v is now 0, while the degrees of other vertices are not impacted by the delete of v' .
- Now on the copied vertices V' (less the vertices that we deleted before), all of them with odd degree, we apply operation (i) repeatedly while both of the conditions below hold true:
 - At least two vertices in V' remain.
 - At least one vertex in V' has odd degree.
- Now we are in one of the following two situations: exactly one vertex v'_0 in V' remains (and has odd degree: its only neighbour now is v_0), or all vertices in V' have even degree.

In the first case, no neighbour of v_0 has been deleted after operation (ii) (all neighbours are either in V , or is v'_0), hence still has odd degree. Delete it, and we are left with $V \setminus \{v_0\} \cup \{v'_0\}$ with v'_0 isolated. By assumption, v_0 wasn't isolated before operation (ii), so the total number of edges decreases as compared to before operation (ii) was carried out.

In the second case, let $V = S \cup T$, with $v \in S$ if v' has been deleted, and $v \in T$ otherwise. Now, all vertices $v' \in T'$ (as assumed, we can't proceed with (i) anymore) and $v \in S$ (all its copies have been deleted) have even degree, while all vertices $v \in T$ have odd degree (none of its neighbours deleted). This means that all vertices $v \in T$ have oddly many neighbours in S , and T and T' are exact copies of each other (now none of them as been deleted, later on we will ensure that v is deleted iff v' is deleted). Assuming the invariant all vertices $v \in T$ have oddly many neighbours in S , and that T and T' are exact copies of each other, v and v' have different degree parity, so we can first delete the one with odd degree. The other one, having one neighbour deleted, had its degree changed from even to odd, so we can delete that as well. The invariant that all vertices $v \in T$ have oddly many neighbours in S , and that T and T' are exact copies of each other, still hold. Hence this operation can be continued indefinitely until T and T' are both empty. This means we just shrank V into S and by assumption, $V \neq \emptyset$ and all vertices in V have nonzero degree (even before operation(ii)). Thus the total number of edges decreases as compared to before operation (ii) was carried out.

- Finally we can still delete any vertex with odd degree repeatedly whenever possible.

Thus each pass above decreases the total number of edges available, and therefore repeating the passes above will make the total number of edges to drop to 0, eventually.

C4 Let n be a positive integer, and let A be a subset of $\{1, \dots, n\}$. An A -partition of n into k parts is a representation of n as a sum $n = a_1 + \dots + a_k$, where the parts a_1, \dots, a_k belong to A and are not necessarily distinct. The number of different parts in such a partition is the number of (distinct) elements in the set $\{a_1, a_2, \dots, a_k\}$. We say that an A -partition of n into k parts is optimal if there is no A -partition of n into r parts with $r < k$. Prove that any optimal A -partition of n contains at most $\sqrt[3]{6n}$ different parts.

Solution. Let $a_1 < a_2 < \dots < a_k$ be the different parts in an optimal partition of n , and let f_1, \dots, f_k be the frequency of the numbers in the representation (that is, $n = \sum_{i=1}^n a_i f_i$).

We first see that there cannot be $0 \leq i < j < k$ with $a_{j+1} - a_j = a_{i+1} - a_i$. Otherwise, we have $a_i + a_{j+1} = a_{i+1} + a_j$. Now suppose that $m = \min\{f_i, f_{i+1}, f_j, f_{j+1}\}$. Consider the following:

$$f_i a_i + f_{i+1} a_{i+1} + f_j a_j + f_{j+1} a_{j+1} = (f_i + m) a_i + (f_{i+1} - m) a_{i+1} + (f_j - m) a_j + (f_{j+1} + m) a_{j+1} \dots (1)$$

$$f_i a_i + f_{i+1} a_{i+1} + f_j a_j + f_{j+1} a_{j+1} = (f_i - m) a_i + (f_{i+1} + m) a_{i+1} + (f_j + m) a_j + (f_{j+1} - m) a_{j+1} \dots (2)$$

If $m = f_{i+1}$ or $m = f_j$, then in (1) one of these coefficients are zero, while the rest remain positive. This means that there is a representation of n with the original set of parts with either f_{i+1} or f_j removed. Similarly, if $m = f_{j+1}$ or $m = f_i$ then there is a representation of n with the original set of parts with either f_i or f_{j+1} removed. Thus the pairwise distance of the neighbouring elements must be different, which also implies the following:

$$a_j = a_1 + \sum_{i=1}^{j-1} a_{i+1} - a_i \geq a_1 + \sum_{i=1}^{j-1} i = a_1 + \frac{j(j-1)}{2}$$

and since $a_1 > 1$ (valid assumption, because otherwise we can have $f_1 = 1$ and this gives the representaiton of 1 part), $a_j \geq 2 + \frac{j(j-1)}{2}$. Since each $f_i \geq 1$, we have

$$n \geq \sum_{i=1}^k a_i \geq \sum_{i=1}^k \left(2 + \frac{j(j-1)}{2} \right) = 2k + \frac{k(k+1)(2k+1)}{12} - \frac{k(k-1)}{4} > \frac{k^3}{6}$$

and therefore $k < \sqrt[3]{6n}$.

C6 In some country several pairs of cities are connected by direct two-way flights. It is possible to go from any city to any other by a sequence of flights. The distance between two cities is defined to be the least possible numbers of flights required to go from one of them to the other. It is known that for any city there are at most 100 cities at distance exactly three from it. Prove that there is no city such that more than 2550 other cities have distance exactly four from it.

Solution. Let's ditch the city and flight notion and instead use the graph theory ones: vertices and edge. Let V and E be the set of vertices, in the said graph, respectively. Consider any node v_0 , and consider the set $S_i : \{v \in V : d(v, v_0) = i\}$ for $i = 1, 2, 3, 4$ (where $d(u, v)$ denoting distance of u and v). We note the following:

- The non-strict triangle inequality holds: if we can go from u to v in $d(u, v)$ steps and v to w in $d(v, w)$ steps, using the shortest route from u to v and then from v to w we can go from u to w in $d(u, w)$ steps. This means, $d(u, w) \leq d(u, v) + d(v, w)$.
- Consider any $v \in S_n$, and let $v_0, v_1, \dots, v_n = v$ be such a path of length n . For each $0 \leq i \leq n$, this path from v_0 to v_i has length i so for each i , $d(v_0, v_i) \leq i$. Conversely, $n = d(v_0, v_n) \leq d(v_0, v_i) + d(v_i, v_n) \leq d(v_0, v_i) + (n - i)$ since the path v_i, \dots, v_n takes $n - i$ steps. This means $d(v_0, v_i) \geq i$ and therefore $d(v_0, v_i) = i$ for all $i \leq n$.

Now let $S_4 = v_1, v_2, \dots, v_m$ and for each v_i consider its shortest path from v_0 : $v_0 = v_{0,i}, v_{1,i}, v_{2,i}, v_{3,i}, v_{4,i} = v_i$. Consider all such $v_{0,i}, v_{1,i}, v_{2,i}, v_{3,i}$. If, say, $v_{3,i} = v_{3,j}$ then $v_{0,i}, v_{1,i}, v_{2,i}, v_{3,j}, v_{4,j}$ is also a shortest path. Similarly $v_{2,i} = v_{2,j}$ implies that $v_{0,i}, v_{1,i}, v_{2,j}, v_{3,j}, v_{4,j}$ is a shortest path. This means we can assume that if $v_{1,i} \neq v_{1,j}$ then $v_{2,i} \neq v_{2,j}$ and $v_{3,i} \neq v_{3,j}$. Consider the set $T : \{v_{1,1}, \dots, v_{1,m}\}$, and suppose that it has n distinct elements. Consider, now, an arbitrary vertex $v_{1,i}$ and let $T_i = \{v_{4,j} : v_{1,i} = v_{1,j}\} \subseteq S_4$. From the definition of shortest path, for each $v_{4,j} \in T_i$ we have $d(v_{1,j}, v_{4,j}) = 4 - 1 = 3$. Consider, also, any $k \neq i$ with $v_{1,k} \neq v_{1,i}$, and the vertices $v_{2,k}, v_{3,k}$ and $v_{4,k}$. We claim that $v_{1,j}$ has a distance exactly three with at least one of those vertices. Now, $d(v_{1,j}, v_{4,k}) \geq d(v_0, v_{4,k}) - d(v_0, v_{1,j}) = 4 - 1 = 3$. If it's greater than 3 then we have $4 \leq d(v_{1,j}, v_{4,k}) \leq d(v_{1,j}, v_{3,k}) + d(v_{3,k}, v_{4,k}) = d(v_{1,j}, v_{3,k}) + 1$, so $d(v_{1,j}, v_{3,k}) \geq 3$. Again if $d(v_{1,j}, v_{3,k}) > 3$ we have $d(v_{1,j}, v_{2,k}) \geq 3$ by a similar logic, but then by triangle inequality again $d(v_{1,j}, v_{2,k}) \leq d(v_{1,j}, v_0) + d(v_0, v_{2,k}) = 1 + 2 = 3$ so the equality must hold.

To summarize, let $v_{1,a_1}, \dots, v_{1,a_n}$ be distinct elements in T . For each a_i , it has distance exactly 3 with each vertex in $|T_{a_i}|$, and with at least one member among v_{2,a_j}, v_{3,a_j} and v_{4,a_j} for any $j \neq i$. By our assumption, for each $j \neq i$ we have $v_{x,a_j} \neq v_{y,a_i}$ for $x, y \in \{1, 2, 3, 4\}$. This means that it has distance 3 with at least $|T_{a_i}| + (n - 1)$ elements, which is at most 100. This means $|T_{a_i}| \leq 100 - n + 1$. In addition, T_{a_1}, \dots, T_{a_n} together represent the partition of S_4 , we have

$$|S_4| = \sum_{i=1}^n |T_{a_i}| \leq \sum_{i=1}^n (100 - n + 1) = n(100 - n + 1)$$

The last expression obtains its maximum when $n = 101/2$; for n integer the maximum is when $n = 50$ or 51 , in which case $n(100 - n + 1) = 50 \times 51 = 2550$. Thus $|S_4| \leq 2550$ as desired.

Geometry

G1 (IMO 4) Let ABC be an acute triangle with orthocenter H , and let W be a point on the side BC , lying strictly between B and C . The points M and N are the feet of the altitudes from B and C , respectively. Denote by ω_1 is the circumcircle of BWN , and let X be the point on ω_1 such that WX is a diameter of ω_1 . Analogously, denote by ω_2 the

circumcircle of triangle CWM , and let Y be the point such that WY is a diameter of ω_2 . Prove that X, Y and H are collinear.

Solution. Consider, now, the point Z (possibly $Z = H$) which is the intersection of the circles BWN and CWM that is not W . As WX is a diameter of ω_1 , WZ and ZX are perpendicular to each other. Similarly, WZ and ZY are perpendicular to each other. This means that both X and Y are on the perpendicular to WZ passing through Z , say ℓ . In addition, from $\angle BYC = \angle BXC = 90^\circ$, $BYXC$ is cyclic, so A lies on the radical axis of the circles BWN and CWM , which is WZ . Therefore, A, W, Z are collinear.

By Miquel's theorem, Z also lies on the circumcircle of AMN , and since $\angle AMH = \angle ANH = 90^\circ$, we have AH a diameter of this circle, and therefore $AZH = 90^\circ$. We have shown that A, W, Z are collinear and X, Y on ℓ . Moreover, $AZ \perp ZH$ so H is also on ℓ . Thus X, Y, H are collinear.

- G2** Let ω be the circumcircle of a triangle ABC . Denote by M and N the midpoints of the sides AB and AC , respectively, and denote by T the midpoint of the arc BC of ω not containing A . The circumcircles of the triangles AMT and ANT intersect the perpendicular bisectors of AC and AB at points X and Y , respectively; assume that X and Y lie inside the triangle ABC . The lines MN and XY intersect at K . Prove that $KA = KT$.

Solution. Denote ℓ as the perpendicular bisector of segment AT . We first show that the perpendicular bisectors of AB and AC (namely ℓ_C, ℓ_B) are symmetric about ℓ . To see this, if O is the center of ω , we have $OA = OT$ so $O \in \ell$. Moreover, AT is the internal angle bisector of the angle $\angle BAC$, and thus:

$$\angle(\ell_C, AT) = \angle(\ell_C, AB) + \angle(AB, AT) = 90^\circ + \angle(AT, AC) = \angle(AC, \ell_B) + \angle(AT, AC) = \angle(AT, \ell_B)$$

and with the fact that the lines ℓ_C and ℓ_B intersect at O (which is on ℓ), ℓ_C and ℓ_B are indeed symmetric in ℓ .

Next, we show that M and X are also symmetric about ℓ . Let M' to be the reflection of M in ℓ . Then since M is on ℓ_C , M' is on ℓ_B and therefore $M'A = M'C$. Moreover, $AMM'T$ is an isosceles trapezoid, hence cyclic. It follows that M' is one of the intersections of the line ℓ_B and the circle AMT . Since $\angle AMX > \angle AMO = 90^\circ$, $\angle AMX$ and similarly $\angle ANX$ are obtuse. It follows that the other intersection of ℓ and circle AMT must be on the different side with X with respect to AC , so it reduces to consider the half-line of ℓ_B passing through O . We see that the reflection of this half line intersects M (for X to stay inside $\triangle ABC$), so M' is on the half-line of ℓ_B passing through O , and therefore $M = X'$. Similarly N and Y are symmetric in ℓ .

The above argument readily implies that $MXYN$ is an isosceles trapezoid with parallel sides MX and NY . It then follows that the intersection of MN and XY (i.e. K) lies on ℓ . Therefore $KA = KT$.

- G3** In a triangle ABC , let D and E be the feet of the angle bisectors of angles A and B , respectively. A rhombus is inscribed into the quadrilateral $AEDB$ (all vertices of the rhombus lie on different sides of $AEDB$). Let φ be the non-obtuse angle of the rhombus. Prove that $\varphi \leq \max\{\angle BAC, \angle ABC\}$.

Solution. Denote $PQRS$ as the rhombus, with P, Q, R, S lying on AB, BD, DE, EA respectively. Now let U and V be the intersections the diagonal QS and the lines AD and BE , respectively. Now suppose for a contradiction, that φ is greater than both $\angle BAC$ and $\angle ABC$. Since QS bisects angle $\angle RSP$, we also have $\angle QSP = \frac{\varphi}{2} > \frac{\angle BAC}{2} = \angle DAB$. Looking at the quadrilateral $USAP$, we have $\angle USP > \angle UAP$, which implies that S lies strictly inside the circumcircle of UAP , too. This has the following consequences:

- By a similar logic we immediately have $\angle ASP > \angle AUP$.

- The opposite angles, $\angle ASU$ and $\angle APU$, have sum greater than 180° .
- Point A is now outside the circumcircle of SUP , and therefore $\angle UPS > \angle UAS$. Similarly we have $\angle SAP > \angle SUA$.

(TL; DR we are tweaking properties of a cyclic quadrilateral on a non-cyclic quadrilateral). Similarly we have $\angle QPB > \angle QVB$. If we denote I and the incenter of ABC then we now have the following:

$$\begin{aligned}
\angle SRV &= \angle SPV \\
&= \angle SPQ - \angle VPQ \\
&= 180^\circ - \varphi - \angle VPQ \\
&< 180^\circ - \varphi - \angle VQB \\
&< 180^\circ - \varphi - \frac{\angle ABC}{2}
\end{aligned}$$

and bearing in mind that $\angle SEV = \angle AEB = 180^\circ - \angle BAC - \frac{\angle ABC}{2}$. Since $\varphi > \angle BAC$, this inequality also implies

$$\angle SRV < 180^\circ - \varphi - \frac{\angle ABC}{2} = 180^\circ - \angle BAC - \frac{\angle ABC}{2} = \angle SEV$$

and therefore by the similar logic $\angle RES > \angle RSV = \frac{\varphi}{2}$. Similarly $\angle RDU > \angle RQU = \frac{\varphi}{2}$ and therefore

$$\angle AED + \angle EDB = \angle AEB + \angle ADB + \angle RES + \angle RDU > \varphi + 360^\circ - \frac{3}{2}\angle CAB - \frac{3}{2}\angle CBA$$

which would mean that $\angle CER + \angle CDR < \frac{3}{2}\angle CAB + \frac{3}{2}\angle CBA - \varphi$. But then E, R, D are collinear so this angle $\angle CER + \angle CDR$ is supposed to be the same as $\angle CAB + \angle CBA$ so we have

$$\angle CAB + \angle CBA < \frac{3}{2}\angle CAB + \frac{3}{2}\angle CBA - \varphi$$

which means $\varphi < \frac{1}{2}\angle CAB + \frac{1}{2}\angle CBA$. This clearly contradicts that $\varphi > \angle CAB$ and $\varphi > \angle CBA$.

G4 Let ABC be a triangle with $\angle B > \angle C$. Let P and Q be two different points on line AC such that $\angle PBA = \angle QBA = \angle ACB$ and A is located between P and C . Suppose that there exists an interior point D of segment BQ for which $PD = PB$. Let the ray AD intersect the circle ABC at $R \neq A$. Prove that $QB = QR$.

Solution. From $\angle ACB = \angle PBA$ and $PD = PB$ we have triangles PBA and PCB similar, so $PD^2 = PB^2 = PA \cdot PC$. In a similar way we can deduce $AB^2 = AQ \cdot AC = \angle AD \cdot AR$ from the fact that $\angle QBA = \angle ACB = \angle ARB$ (since R is on circle ABC). Therefore, the quadrilateral $DQCR$ is cyclic. We can now compute the following:

$$\angle QBR = \angle QBC + \angle CBR = \angle ABC - \angle ABQ + \angle RAC = \angle ABC - \angle ACB + (\angle PDA + \angle DPQ)$$

and the fact $PD^2 = PA \cdot PC$ also implies $\angle PDA = \angle PDC$ and therefore

$$\angle QRB = \angle QRD + \angle DRB = \angle QCD + \angle ACB = \angle PDA + \angle ACB$$

but since

$$2\angle ACB = \angle PBA + \angle ABQ = \angle PBQ = \angle PDB = \angle PDQ + \angle PQB = \angle DPQ + \angle ABC$$

and therefore

$$\begin{aligned}
\angle QBR &= \angle ABC - \angle ACB + (\angle PDA + \angle DPQ) \\
&= \angle ABC - \angle ACB + \angle PDA + (2\angle ACB - \angle ABC) \\
&= \angle ACB + \angle PDA \\
&= \angle QRB
\end{aligned}$$

and therefore $QR = QB$.

- G5** Let $ABCDEF$ be a convex hexagon with $AB = DE$, $BC = EF$, $CD = FA$, and $\angle A - \angle D = \angle C - \angle F = \angle E - \angle B$. Prove that the diagonals AD , BE , and CF are concurrent.

Solution. The fact that their opposite sides have equal length motivates us to try the spiral similarity between them. To be precise, let X be the intersection of AD and BE and we consider the second intersection O_1 of circles ABX and DEX , with this intersection being equal to X only when the two circles are tangent at X . We first notice the following realizations by directed angles:

$$\angle(AO_1, BO_1) = \angle(AX, BX) = \angle(AD, BD) = \angle(DX, EX) = \angle(DO_1, EO_1)$$

$$\angle(AB, BO_1) = \angle(AX, XO_1) = \angle(DX, XO_1) = \angle(DE, EO_1)$$

$$\angle(AB, AO_1) = \angle(BX, XO_1) = \angle(EX, XO_1) = \angle(DE, DO_1)$$

and therefore triangles ABO_1 and DEO_1 are similar (it's not quite immediate from directed angles since the corresponding angles could either be equal or supplementary, but then the claim follows that the angles in a triangle are positive and add up to 180°). Since $AB = DE$, these triangles are indeed congruent, and therefore $AO_1 = DO_1$ and $BO_1 = EO_1$. Since the congruent triangles can be obtained by each other via rotation (technically reflection might be involved but that's just rotation by 180°) we have $\angle(AB, DE) = \angle(AO_1, DO_1) = \angle(BO_1, EO_1)$, which in turn motivates the following:

$$\angle(AB, DE) = \angle(AO_1, DO_1) = \angle(AO_1, AD) + \angle(AD, DO_1) = 2\angle(AO_1, AD)$$

and similarly $\angle(AB, DE) = 2\angle(BO_1, BE)$.

Now to properly investigate $\angle(AB, DE)$, we need the angle condition given above. In essence $\angle(AB, DE) = \angle(AB, BC) + \angle(BC, CD) + \angle(CD, DE) = \angle B + \angle C + \angle D$ (we may assume that our hexagon $ABCDEF$ is indeed oriented in this manner). Denoting the angle $\alpha := \angle A - \angle D = \angle C - \angle F = \angle E - \angle B$, and bearing in mind that the sum of the six angles $\angle A$ through $\angle F$ is 720° , we get $\angle B + \angle C + \angle D = 360^\circ - \frac{\alpha}{2}$. Therefore, denoting $\beta = \frac{\alpha}{2}$ (we want to avoid fractions in directed angles...if you get what I mean) we have $\angle(AB, DE) = 360^\circ - \beta = -\beta$, i.e. $2\angle(BO_1, BE) = -\beta$.

Consider, now, the intersection Y of the lines BE and CF , and O_2 as the second intersection of circles BCY and EFY . Then by the similar logic as above we get $2\angle(BO_2, BE) = \angle(BC, EF)$ and we also have $\angle(BC, EF) = \angle C + \angle D + \angle E$ which can be calculated as $360^\circ + \beta$, so we now have $2\angle(BO_2, BE) = \angle(BC, EF) = \beta = -2\angle(BO_1, BE)$. Thus considering the sides O_1 and O_2 belong to w.r.t. BE we have BE as the internal angle bisector of $\angle O_1BO_2$ and similarly O_1E is the internal angle bisector of $\angle O_1EO_2$. Therefore BE is the perpendicular bisector of O_1O_2 .

Finally, denote the intersection Z of the lines CF and AD , and O_3 as the second intersection of CDZ and FAZ . By the similar logic as above, again, AD is the perpendicular bisector of O_1O_3 and CF is the perpendicular bisector of O_2O_3 . Therefore AD, BE, CF concur at the circumcenter of $O_1O_2O_3$, i.e. the point $X = Y = Z$.

Number Theory

- N1** Let $\mathbb{Z}_{>0}$ be the set of positive integers. Find all functions $f : \mathbb{Z}_{>0} \rightarrow \mathbb{Z}_{>0}$ such that

$$m^2 + f(n) \mid mf(m) + n$$

for all positive integers m and n .

Answer. The only such function is the identity function $f(x) = x$, where both sides are equal to $m^2 + n$ so this function works.

Solution. By plugging $m = n = 2$ we have $4 + f(2) \mid 2f(2) + 2 = 2(4 + f(2)) - 6$, so $4 + f(2) \mid 6$. Since $f(2) > 0$, the only possibility is $f(2) = 2$. Then, plugging $m = 2$ gives $4 + f(n) \mid 2f(2) + n = 4 + n$. Given that $4 + f(n)$ and $4 + n$ are both positive, we have $4 + f(n) \leq 4 + n$, so $f(n) \leq n$. Finally, plugging $n = 2$ gives $m^2 + 2 = m^2 + f(n) \mid mf(m) + 2$, so $m^2 \leq mf(m)$, or $m \leq f(m)$. This means, $f(m) \geq m$ and $f(m) \leq m$ both hold for all m , hence $f(m) = m$ for all m .

N2 (IMO 1) Assume that k and n are two positive integers. Prove that there exist positive integers m_1, \dots, m_k such that

$$1 + \frac{2^k - 1}{n} = \left(1 + \frac{1}{m_1}\right) \cdots \left(1 + \frac{1}{m_k}\right).$$

Solution. A solution that fixes n and does induction on k is possible, as per my own solution to the IMO. But let me give a solution that's based on the following intuition: we want to find a sequence of $k + 1$ numbers $n = a_0 < a_1 < \dots < a_k = n + 2^k - 1$ such that for each $0 \leq i < k$, $a_{i+1} - a_i \mid a_i$. This way, we have:

$$1 + \frac{2^k - 1}{n} = \frac{n + 2^k - 1}{n} = \prod_{i=1}^k \frac{a_i}{a_{i-1}} = \prod_{i=1}^k \left(1 + \frac{1}{\frac{a_i - a_{i-1}}{a_{i-1}}}\right)$$

and notice that $\frac{a_i - a_{i-1}}{a_{i-1}}$ is an integer based on our construction. Hence this sequence of a_i 's gives rise of a valid construction.

To construct such sequence, we first notice that among $n, n + 1, \dots, n + 2^k - 1$ which is a consecutive sequence of 2^k numbers, exactly one of them is divisible by 2^k . Name this number n_0 . Next, we consider the binary representation of the two numbers:

$$n_0 - n = 2^{c_1} + 2^{c_2} + \dots + 2^{c_m} \quad (n + 2^k - 1) - n_0 = 2^{b_1} + 2^{b_2} + \dots + 2^{b_\ell}$$

with $c_1 < c_2 < \dots < c_m < k$ and $b_1 < b_2 < \dots < b_\ell < k$ (k and ℓ could be 0 in case n or $n + 2^k - 1$ coincides with n_0). We now construct the sequence based on the following: for $i \leq m$, we let $a_i = a_{i-1} + 2^{c_i}$ and for $i > m$, $a_i = a_{i-1} + 2^{b_{\ell-i+1}}$. To see it works, we need the two following observations:

- First, we need $a_i - a_{i-1} \mid a_i$. We notice that for each $j \leq m$, $a_j = n + \sum_{i=1}^j 2^{c_i} = n_0 - \sum_{i=j+1}^m 2^i$. Since n_0 is divisible by 2^k and $k > a_i$ for all i , we have $2^{\min\{c_i : i \geq j+1\}} = 2^{c_{j+1}}$

so in fact $a_j - a_{j-1} = 2^{c_j} \mid a_j$, as desired. As for $j > m$, $a_j = a_n + \sum_{i=0}^{j-m-1} 2^{b_{\ell-i}}$ so by similar logic, $2^{\ell-(j-m-1)} \mid a_j$ and thus $2^{\ell-(j-m)} \mid a_j$. But $2^{\ell-(j-m)} = a_j - a_{j-1}$, so the conclusion follows.

- We can see that $n_{\ell+m}$ is indeed equal to $n_0 + 2^{b_1} + 2^{b_2} + \dots + 2^{b_\ell} = n + 2^k - 1$, so it remains to show that $\ell + m = k$. This is the same as showing that $n_0 - n$ and $(n + 2^k - 1) - n_0$ have k ones in total in their binary representation. We however notice that both numbers are at most $2^k - 1$, so they can be written with k digits. Moreover, since they add up to $2^k - 1$, all digits within the first k digits are complementary of each other. This means, half of the $2k$ digits are 1 and half are zero, and we have a total of k ones, as desired.

The conclusion follows from the combination of the two claims.

N3 Prove that there exist infinitely many positive integers n such that the largest prime divisor of $n^4 + n^2 + 1$ is equal to the largest prime divisor of $(n+1)^4 + (n+1)^2 + 1$.

Solution. Denote by $f(n) = n^2 + n + 1$. Then $n^4 + n^2 + 1 = (n^2 + n + 1)(n^2 - n + 1) = f(n)f(n-1)$. Denote also the largest prime divisor of $f(n)$ as $p(n)$. We first show that $\gcd(f(n), f(n-1)) = 1$. If k is a common divisor of $f(n) = n^2 + n + 1$ and $f(n-1) = n^2 - n + 1$, then $k \mid 2n$ too. We have $\gcd(f(n), n) = 1$ (since) and both $f(n)$ and $f(n-1)$ are odd, and therefore $k = \pm 1$ must hold. Next, if $p(n) \geq p(n-1)$ and $p(n) \geq p(n+1)$, considering that the largest prime divisor of $n^4 + n^2 + 1$ is $\max\{p(n), p(n-1)\} = p(n)$, we have the largest prime divisor of $n^4 + n^2 + 1$ and $(n+1)^4 + (n+1)^2 + 1$ being equal, and therefore this n would work.

It remains to show that there are infinitely many n such that $p(n) > p(n-1)$ and $p(n) > p(n+1)$. Since the $\gcd\{f(n), f(n-1)\}$, $p(n)$ cannot be equal to $p(n-1)$. Suppose our desired condition doesn't hold, then either $p(n) < p(n-1)$ for all sufficiently large n , or $p(n) > p(n-1)$ for sufficiently large n . The first case is impossible since all prime numbers are greater than 0, and thus for each n there's only finitely many prime numbers smaller than $p(n)$. For the second case to hold, let n_0 be that $p(n+1) > p(n)$ for all $n \geq n_0$. By taking a tail of the sequence $\{n : n \geq n_0\}$, we may also assume that $n_0 \geq 2$. Now $(n_0+1)^2 > n_0+1$ for $n_0 \geq 2$, so it also follows that $p((n_0+1)^2) > p(n_0+1) > p(n_0)$. However, we have $f((n_0+1)^2) = f(n_0)f(n_0+1)$ so $p((n_0+1)^2) = \max\{p(n_0+1), p(n_0)\}$, contradicting our initial assumption that $p((n_0+1)^2) > p(n_0+1) > p(n_0)$. The conclusion therefore follows.

N5 Fix an integer $k > 2$. Two players, called Ana and Banana, play the following game of numbers. Initially, some integer $n \geq k$ gets written on the blackboard. Then they take moves in turn, with Ana beginning. A player making a move erases the number m just written on the blackboard and replaces it by some number m' with $k \leq m' < m$ that is coprime to m . The first player who cannot move anymore loses.

An integer $n \geq k$ is called good if Banana has a winning strategy when the initial number is n , and bad otherwise.

Consider two integers $n, n' \geq k$ with the property that each prime number $p \leq k$ divides n if and only if it divides n' . Prove that either both n and n' are good or both are bad.

Solution. Alternatively, a number is good if the second player has the winning strategy and bad if otherwise. We have the following observation:

- The number k is good.
- For each number $n > k$, if there's a good number n' with $k \leq n' < n$ and $\gcd(n', n) = 1$, then n is bad because the player handling n' is guaranteed to lose, and therefore player handling n can replace n with n' .
- Otherwise, n is not coprime to any good number n' , so whoever handling n will either lose a move or choose the number n_1 that is bad. This means their opponent who handles n_1 has a winning strategy, so the player handling n themselves lose. This means n is a good number.

This also means that k is good, any two good numbers cannot be coprime to each other, and any bad number n is coprime to at least one good number m with $k \leq m < n$. In particular, if n is coprime to k (and so is n' by the problem condition), then n is bad (and so is n'). So we now assume that $\gcd(n, k)$ and $\gcd(n, k')$ are both greater than 1.

To solve the problem, we first consider any two numbers $n \geq k$ and n_0 such that they share the same set of prime divisors that are at most k , and n_0 has no divisors greater than k . We show that such n_0 with $k \leq n_0 \leq n$ exists. If n itself has no divisor greater than k we are done. Otherwise, let $n = n_1 n_2$, with n_1 having only prime divisors at most

p and n_2 having only prime divisors greater than p . Since n_2 has at least one prime divisor (by assumption) and greater than k , we have $n = n_1 n_2 > k n_1$. If $n_1 \geq k$, then n_1 has the same set of prime divisors at most k as n , and $k \leq n_1 \leq n_1 n_2 = n$, so $n_0 = n_1$ would work. Otherwise, consider the sequence $\{n_1^m : m \geq 1\}$, where each member in the sequence also has prime divisors satisfying our desired property. Since $\gcd(n, k) > 1$, we have $n_1 > 1$ and so we can pick m that is the smallest integer such that $n_1^m \geq k$. Since $n_1 < k$, we have $m \geq 2$. This also means that $n_1^{m-1} < k$ and therefore $n_1^m < k n_1 < n$. But we have just assumed that $n_1^m \geq k$, so this $n_0 := n_1^m$ would work.

Now if n_0 is good, then for each prime p we have $p \mid n_0 \rightarrow p \mid n$ and so for any good number m , we have $p \mid m$ and $p \mid n_0$ simultaneously for some prime p , and therefore for this prime we have $p \mid \gcd(m, n)$. Considering this for all good numbers m we know that n is also good. Conversely, suppose that for some n , n is good but n_0 with the said property as above is bad. By our argument above, we may assume that n is the smallest such example. This means there's a good number $m < n_0$ such that $\gcd(m, n_0) = 1$ but since n is good, m and n_0 have a common prime divisor, say p_0 . Since the set of prime divisors of n_0 and n only differs on those that are greater than k , we have $p_0 > k$. Now, let m_0 have the same set of prime divisors $\leq k$ as m , but has no prime divisor $> k$. By our previous paragraph, we may assume that $k \leq m_0 < m$, and by the minimality of n , since m is good, m_0 must also be good. Since m_0 and n are both good, there must be a prime divisor p_1 dividing both numbers, and by the property of m_0 , $p_1 \leq k$. This means $p_1 \mid n_0$ and $p_1 \mid m$ too, contradicting $\gcd(m, n_0) = 1$.

Finally, consider any n and n' that share the same set of prime divisors $\leq k$. Consider n_0 that has this same property without any prime divisor $> k$ (and smaller than both n and n'). If n_0 is good then both n and n' are good; otherwise, n_0 is bad so n and n' are both bad.

N6 Determine all functions $f : \mathbb{Q} \rightarrow \mathbb{Z}$ satisfying

$$f\left(\frac{f(x) + a}{b}\right) = f\left(\frac{x + a}{b}\right)$$

for all $x \in \mathbb{Q}$, $a \in \mathbb{Z}$, and $b \in \mathbb{Z}_{>0}$. (Here, $\mathbb{Z}_{>0}$ denotes the set of positive integers.)

Answer. There are three classes of functions: the constant function $f(x) \equiv a$ for any $a \in \mathbb{Z}$, the floor function $f(x) \equiv \lfloor x \rfloor$ and the ceiling function $f(x) \equiv \lceil x \rceil$.

Solution. It's easy to verify that constant function fulfills the condition (both sides are wrapped by $f(\cdot)$). For the rest two, observe that

$$\lfloor \frac{x+a}{b} \rfloor \leq \frac{x+a}{b} < \lfloor \frac{x+a}{b} \rfloor + 1$$

I.e.

$$b \lfloor \frac{x+a}{b} \rfloor \leq x+a < b \lfloor \frac{x+a}{b} \rfloor + b$$

Notice that $\lfloor x \rfloor + a \leq x + a$ so $\lfloor x \rfloor + a < b \lfloor \frac{x+a}{b} \rfloor + b$ (note: we implicitly used the fact b is positive integer here). Given that a is also an integer, $\lfloor x \rfloor + a = \lfloor x + a \rfloor$ and therefore

$$b \lfloor \frac{x+a}{b} \rfloor \leq \lfloor x + a \rfloor = \lfloor x \rfloor + a$$

using the fact that $a \lfloor x \rfloor \leq \lfloor ax \rfloor$ for any positive integer a . Therefore, we also have

$$b \lfloor \frac{x+a}{b} \rfloor \leq \lfloor x \rfloor + a < b \lfloor \frac{x+a}{b} \rfloor + b$$

and so

$$\lfloor \frac{x+a}{b} \rfloor = \lfloor \frac{\lfloor x \rfloor + a}{b} \rfloor$$

a similar method can be used to verify for the ceiling function.

To show that these are the only functions, we first consider the case where for some integer x , $f(x) \neq x$. Then $b = |x - f(x)|$ and plugging a to make $f(x) - x \mid x + a$ yields $f(b) = f(b + 1)$ for all integers b , so f is constant on integers. Now let x be any rational number, we have $f(x) = f(f(x))$ but $f(x)$ is an integer so f is, in fact, constant.

Thus for f to be nonconstant, f is identity on the scope of integers. Let's focus on this case now. Letting $b = 1$ gives

$$f(x) + a = f(f(x) + a) = f(x + a)$$

given that $f(x)$ and a are both integers. This means for each integer n , since $f(\frac{1}{n})$ is an integer, there's an a_0 such that $f(a_0 + \frac{1}{n}) = 0$. Plugging $x = a_0 + \frac{1}{n}$ and $b = |na_0 + 1|$ gives

$$f\left(\frac{f(x)}{b}\right) = f(0) = 0 = f\left(\frac{x}{b}\right)$$

and given that $\frac{x}{b}$ is either $\frac{1}{n}$ or $-\frac{1}{n}$, we have, for each n , either $f(\frac{1}{n}) = 0$ or $f(-\frac{1}{n}) = 0$.

Assume that $f(\frac{1}{2}) = 0$, and we show that $f \equiv \lfloor \cdot \rfloor$. We first show this property for all the numbers in terms of $x = \frac{b}{2^k}$ for b integers.

We will proceed into the following lemma: let $b > 1$, and suppose that $f(\frac{x}{b}) = \lfloor \frac{x}{b} \rfloor$ for any integer x . Then $f(\frac{x}{b^k}) = \lfloor \frac{x}{b^k} \rfloor$ for any integer x and any integer k .

Proof: Let's do induction on k . Now supposed that the said property holds for some integer $k \geq 1$ and consider $y = \frac{x}{2^{k+1}}$. Then plugging $x := \frac{x}{2^k}$, $a = 0$ we get

$$f(y) = f\left(\frac{x}{b^k}/b\right) = f\left(f\left(\frac{x}{b^k}\right)/b\right) = f\left(\lfloor \frac{x}{b^k} \rfloor / b\right) = \lfloor \lfloor \frac{x}{b^k} \rfloor / b \rfloor$$

and using the fact that $\lfloor \lfloor \frac{x}{b} \rfloor \rfloor = \lfloor \frac{x}{b} \rfloor$,

$$f\left(\frac{x}{b^{k+1}}\right) = f(y) = \lfloor \lfloor \frac{x}{b^k} \rfloor / b \rfloor = \lfloor \frac{x}{b^k} / b \rfloor = \lfloor \frac{x}{b^{k+1}} \rfloor$$

Now that $f(\frac{1}{2}) = 0$, f identity in integers and $f(x + a) = f(x) + a$, we have $f(\frac{x}{2}) = \lfloor \frac{x}{2} \rfloor$ for all integers x and consequently, $f(\frac{x}{2^k}) = \lfloor \frac{x}{2^k} \rfloor$ for all integers x and k .

Next, with the 'base case' above, let $p \geq 3$ be some prime number and suppose that the claim $f(\frac{b}{q}) = \lfloor \frac{b}{q} \rfloor$ holds for any q with all prime divisors $< p$. Let's extend this claim to 'any q with all prime divisors $\leq p$ '. In other words, we want to induct on prime numbers $\leq p$. Given that $p + 1$ is composite for $p \geq 3$, it has all prime divisors less than p . Now, plug $b := p + 1$, $a := 0$ and $x := \frac{a(p+1)}{p}$ for some positive integer p and we have

$$f\left(\frac{a}{p}\right) = f\left(\frac{f\left(\frac{a(p+1)}{p}\right)}{p+1}\right) = \left\lfloor \frac{f\left(\frac{a(p+1)}{p}\right)}{p+1} \right\rfloor = \left\lfloor \frac{f\left(\frac{a}{p} + a\right)}{p+1} \right\rfloor = \left\lfloor \frac{f\left(\frac{a}{p}\right) + a}{p+1} \right\rfloor$$

which then becomes

$$0 \leq \frac{f\left(\frac{a}{p}\right) + a}{p+1} - f\left(\frac{a}{p}\right) = \frac{p}{p+1} \left(\frac{a}{p} - f\left(\frac{a}{p}\right) \right) < 1$$

which means that $0 \leq \frac{a}{p} - f\left(\frac{a}{p}\right) < 1 + \frac{1}{p}$. Considering that f is identity on integers, and $f(\frac{a}{p})$ is also an integer, the only way for this to be fulfilled is for $f(\frac{a}{p}) = \lfloor \frac{a}{p} \rfloor$, and by the lemma above, $f(\frac{a}{p^k}) = \lfloor \frac{a}{p^k} \rfloor$. Finally, for any number q with prime divisors $\leq p$, we can write $q = rp^k$ where r has all divisors $< p$. Now

$$f\left(\frac{a}{q}\right) = f\left(\frac{a}{rp^k}\right) = f\left(\frac{a/p^k}{r}\right) = f\left(\frac{f(a/p^k)}{r}\right) = \left\lfloor \frac{\lfloor a/p^k \rfloor}{r} \right\rfloor = \left\lfloor \frac{a/p^k}{r} \right\rfloor = \left\lfloor \frac{a}{q} \right\rfloor$$

as desired. Using this induction on primes p , we have $f \equiv$ the floor function. Similarly, if we assume earlier on that $f(-\frac{1}{2}) = 0$ then f would be the ceiling function.