

1 データ構造

Union-Find

```
1 struct UnionFind{
2     vector<int> dat;
3
4     UnionFind(int sz){
5         dat.assign(sz, -1);
6     }
7
8     bool unite(int x, int y){
9         x = root(x), y = root(y);
10        if(x == y) return(false);
11        if(dat[x] > dat[y]) swap(x, y);
12        dat[x] += dat[y];
13        dat[y] = x;
14        return(true);
15    }
16
17    int root(int k){
18        if(dat[k] < 0) return(k);
19        return(dat[k] = root(dat[k]));
20    }
21
22    int size(int k){
23        return(-dat[root(k)]);
24    }
25 };
```

Binary Indexed Tree

```
1 template< class T >
2 struct BinaryIndexedTree {
3     vector< T > dat;
4
5     BinaryIndexedTree(int sz) {
6         dat.assign(++sz, 0);
7     }
8
9     T sum(int k){
10        T ret = 0;
11        for(++k; k > 0; k -= k & -k) ret += data[k];
12        return ret;
13    }
14
15    void add(int k, T x){
16        for(++k; k < data.size(); k += k & -k) data[k] += x;
17    }
18 };
```

Segment Tree

```
1  /*
2  !) 0-indexed
3  !) [a, b) に対する演算
4  !) merge, updateNodeを書く
5  */
6  template<typename Monoid>
7  struct Segtree{
8      int n;
9      vector<Monoid> dat;
10     Monoid m0; // データの初期化値
11     Segtree(int sz, Monoid m0) : m0(m0){
12         n = 1;
13         while(n < sz) n *= 2;
14         dat.assign(2*n-1, m0);
15     }
16
17     Monoid merge(Monoid a, Monoid b) // 区間をマージする二項演算
18     void updateNode(int k, Monoid x) // 区間を操作する二項演算
19
20     void update(int k, Monoid x){
21         k += n-1;
22         updateNode(k, x);
23         while(k > 0) {
24             k = (k-1)/2;
25             dat[k] = merge(dat[k*2+1], dat[k*2+2]);
26         }
27     }
28     Monoid query(int a, int b){
29         Monoid L = m0, R = m0;
30         int A = a+n-1;
31         int B = b+n-1;
32         while(A < B) {
33             if((A&1) == 0) L = merge(L, dat[A++]);
34             if((B&1) == 0) R = merge(dat[--B], R);
35             A >>= 1;
36             B >>= 1;
37         }
38         return merge(L, R);
39     }
40     Monoid operator[](const int &k) const { return dat[k+n-1]; }
41 };
42
43 //##### 例 #####
44 // Range min (AOJ DSL_2_A)
45 Monoid merge(Monoid a, Monoid b){ return min(a,b); }
46 void updateNode(int k, Monoid x){ dat[k] = x; }
47 Segtree<LL> a(n, (1LL<<31)-1);
48
49 // Range add (AOJ DSL_2_B)
50 Monoid merge(Monoid a, Monoid b){ return a + b; }
51 void updateNode(int k, Monoid x){ dat[k] += x; }
52 Segtree<LL> a(n, 0);
```

遅延伝搬 Segment Tree

```
1  /*
2  !) 0-indexed 遅延伝搬セグメント木
3  !) [a, b)に対する演算
4  !) M0, L0, merge, updateNode, propagateを書く
5  */
6  template <typename Monoid>
7  struct LazySegtree {
8      int n;
9      vector<Monoid> dat, lazy;
10
11      Monoid M0, L0 // データと遅延配列の初期化値 queryに合わせて選択する
12
13      LazySegtree(int sz, Monoid dat_init){
14          n = 1;
15          while (n < sz) n *= 2;
16          dat.assign(2*n-1, dat_init);
17          lazy.assign(2*n-1, L0);
18      }
19
20      Monoid merge(Monoid a, Monoid b) // 区間をマージする二項演算
21      void updateNode(int k, Monoid x) // 区間を操作する二項演算
22      void propagate(int k, int l, int r) // 遅延配列の伝搬のさせ方
23
24      void eval(int k, int l, int r) {
25          if(lazy[k] == L0) return;
26          propagate(k, l, r);
27          if(r-l > 1) {
28              updateNode(2*k+1, lazy[k]);
29              updateNode(2*k+2, lazy[k]);
30          }
31          lazy[k] = L0;
32      }
33
34      void update(int a, int b, Monoid x, int k, int l, int r) {
35          eval(k, l, r);
36          if (r <= a || b <= l) return;
37          if (a <= l && r <= b) {
38              updateNode(k, x);
39              eval(k, l, r);
40          }else{
41              update(a, b, x, k*2+1, l, (l+r)/2);
42              update(a, b, x, k*2+2, (l+r)/2, r);
43              dat[k] = merge(dat[2*k+1], dat[2*k+2]);
44          }
45      }
46
47      void update(int a, int b, Monoid x) {
48          update(a, b, x, 0, 0, n);
49      }
50
51      Monoid query(int a, int b, int k, int l, int r) {
52          eval(k, l, r);
53          if (r <= a || b <= l) return M0;
54          if (a <= l && r <= b) return dat[k];
55          Monoid L = query(a, b, k*2+1, l, (l+r)/2);
```

```

56         Monoid R = query(a, b, k*2+2, (l+r)/2, r);
57         return merge(L, R);
58     }
59
60     Monoid query(int a, int b){
61         return query(a, b, 0, 0, n);
62     }
63 };
64
65 //##### 例 #####
66 // Range update - min (AOJ DSL_2_F)
67 Monoid M0 = LLINF, L0 = LLINF;
68 Monoid merge(Monoid a, Monoid b){ return min(a, b); }
69 void updateNode(int k, Monoid x){ lazy[k] = x; }
70 void propagate(int k, int l, int r){ dat[k] = lazy[k]; }
71 LazySegtree<LL> seg(n+1, (1LL<<31)-1);
72
73 // Range update - sum (AOJ DSL_2_I)
74 Monoid M0 = 0, L0 = LLINF;
75 Monoid merge(Monoid a, Monoid b){ return a + b; }
76 void updateNode(int k, Monoid x){ lazy[k] = x; }
77 void propagate(int k, int l, int r){ dat[k] = lazy[k]*(r-l); }
78 LazySegtree<LL> seg(n+1, 0);
79
80 // Range add - min (AOJ DSL_2_H)
81 Monoid M0 = LLINF, L0 = 0;
82 Monoid merge(Monoid a, Monoid b){ return min(a, b); }
83 void updateNode(int k, Monoid x){ lazy[k] += x; }
84 void propagate(int k, int l, int r){ dat[k] += lazy[k]; }
85 LazySegtree<LL> seg(n+1, 0);
86
87 // Range add - sum (AOJ DSL_2_G)
88 Monoid M0 = 0, L0 = 0;
89 Monoid merge(Monoid a, Monoid b){ return a + b; }
90 void updateNode(int k, Monoid x){ lazy[k] += x; }
91 void propagate(int k, int l, int r){ dat[k] += lazy[k]*(r-l); }
92 LazySegtree<LL> seg(n+1, 0);

```

2 グラフ

Grid 上での BFS

```
1  int W, H;
2  vector<vector<char>> s;
3  vector<vector<int>> cost;
4  int bfs(){
5      int dx[] = {0, 1, 0, -1}, dy[] = {1, 0, -1, 0};
6      queue<pair<int, int>> que;
7      que.push(make_pair(0, 0));
8      cost[0][0] = 0;
9
10     while(!que.empty()) {
11         pair<int, int> p = que.front();
12         que.pop();
13         if(p == make_pair(H-1, W-1)){
14             // ゴールに到達
15             return cost[p.first][p.second];
16         }
17         for(int i = 0; i < 4; i++) {
18             int ny = p.first + dy[i], nx = p.second + dx[i];
19             if(nx < 0 || ny < 0 || nx >= W || ny >= H) continue;
20             if(s[ny][nx] == '#') continue;
21             if(cost[ny][nx] != -1) continue;
22
23             cost[ny][nx] = cost[p.first][p.second] + 1;
24             que.push(make_pair(ny, nx));
25         }
26     }
27     return -1;
28 }
```

Dijkstra 法 (単一始点最短経路)

```
1  struct edge{
2      int to, cost;
3  };
4  using WeightedGraph = vector<vector<edge>>;
5
6  vector<int> dijkstra(WeightedGraph &G, int st){
7      vector<int> dist(G.size(), INF);
8      using pi = pair<int, int>;
9      priority_queue<pi, vector<pi>, greater<pi>> que;
10     dist[st] = 0;
11     que.push(mp(dist[st], st));
12     while(!que.empty()){
13         int cost, idx;
14         tie(cost, idx) = que.top(); que.pop();
15         if(dist[idx] < cost) continue;
16         for(auto &e: G[idx]){
17             if(dist[e.to] <= cost+e.cost) continue;
18             dist[e.to] = cost+e.cost;
19             que.push(mp(dist[e.to], e.to));
20         }
21     }
```

```

21     }
22     return dist;
23 }

```

Bellman-Ford 法 (負路あり単一始点最短経路)

```

1  struct edge{
2      int src, to, cost;
3  };
4  using Edges = vector<edge>;
5
6  vector<int> bellman_ford(Edges &E, int V, int st){
7      vector<int> dist(V, INF);
8      dist[st] = 0;
9      rep(i, V-1){
10         for(auto &e: E){
11             if(dist[e.src] == INF) continue;
12             dist[e.to] = min(dist[e.to], dist[e.src]+e.cost);
13         }
14     }
15     for(auto &e: E){
16         if(dist[e.src] == INF) continue;
17         if(dist[e.to] > dist[e.src]+e.cost){
18             // 負閉路が存在
19             return vector<int>();
20         }
21     }
22     return dist;
23 }

```

Warshall-Floyd 法 (全点对間最短経路)

```

1  using Graph = vector<vector<int>>;
2
3  void warshall_floyd(Graph &G){
4      int V = G.size();
5      rep(k, V)rep(i, V)rep(j, V){
6          if(G[i][k] == INF || G[k][j] == INF) continue;
7          G[i][j] = min(G[i][j], G[i][k]+G[k][j]);
8      }
9      // G[i][i] < 0が存在 <=> 負閉路が存在
10 }

```

Kruskal 法 (最小全域木)

```

1  // UnionFindが必要
2
3  struct edge{
4      int src, to, cost;
5  };
6  using Edges = vector<edge>;
7
8  int kruskal(Edges &E, int V)
9  {
10     sort(all(E), [](const edge &a, const edge &b)
11         {

```

```

12         return (a.cost < b.cost);
13     });
14     UnionFind tree(V);
15     int res = 0;
16     for(auto &e : E) {
17         if(tree.unite(e.src, e.to)) res += e.cost;
18     }
19     return (res);
20 }

```

トポロジカルソート

```

1  struct edge{
2      int to, cost;
3  };
4  using WeightedGraph = vector<vector<edge>>;
5
6  vector<int> tsort(WeightedGraph &G){
7      vector<int> tsorted;
8      vector<int> used(G.size(), 0);
9      bool f = false;
10     function<void(int)> dfs = [&](int u){
11         if(used[u] > 0){
12             if(used[u] == 1) f = true;
13             return;
14         }
15         used[u] = 1;
16         for(auto &e : G[u]) dfs(e.to);
17         used[u] = 2;
18         tsorted.pb(u);
19     };
20     rep(i, G.size()) dfs(i);
21     if(f){
22         // 閉路が存在
23         return vector<int>();
24     }
25     reverse(all(tsorted));
26     return tsorted;
27 }

```

Dinic 法 (最大流)

```

1  template<typename flow_t>
2  struct Dinic{
3      const flow_t INF_flow_t = INF;    // WRITE HERE
4
5      struct edge{
6          int to;
7          flow_t cap;
8          int rev;
9      };
10     using WeightedGraph = vector<vector<edge>>;
11     int V;
12     WeightedGraph G;
13     vector<int> itr, level;
14

```

```

15     Dinic(int V) : V(V) { G.assign(V, vector<edge>()); }
16
17     void add_edge(int from, int to, int cap) {
18         G[from].push_back((edge){to, cap, (int)G[to].size()});
19         G[to].push_back((edge){from, 0, (int)G[from].size()-1});
20     }
21
22     void bfs(int s) {
23         level.assign(V, -1);
24         queue<int> que;
25         level[s] = 0;
26         que.push(s);
27         while (!que.empty()) {
28             int v = que.front(); que.pop();
29             for(auto &&e: G[v]){
30                 if (e.cap > 0 && level[e.to] < 0) {
31                     level[e.to] = level[v] + 1;
32                     que.push(e.to);
33                 }
34             }
35         }
36     }
37
38     flow_t dfs(int v, int t, flow_t f) {
39         if(v == t) return f;
40         for(int &i = itr[v]; i < (int)G[v].size(); i++) {
41             edge &e = G[v][i];
42             if (e.cap > 0 && level[v] < level[e.to]) {
43                 flow_t d = dfs(e.to, t, min(f, e.cap));
44                 if (d > 0) {
45                     e.cap -= d;
46                     G[e.to][e.rev].cap += d;
47                     return d;
48                 }
49             }
50         }
51         return 0;
52     }
53
54     flow_t max_flow(int s, int t) {
55         flow_t res = 0, f;
56         while(bfs(s), level[t] >= 0) {
57             itr.assign(V, 0);
58             while((f = dfs(s, t, INF_flow_t)) > 0) res += f;
59         }
60         return res;
61     }
62 };
63
64 // 最小流量制限付き最大流
65 // 各辺に[lb, ub]の容量の辺を張る
66 template<typename flow_t>
67 struct DinicWithLowerBound{
68     Dinic<flow_t> flow;
69     int S, T;
70     flow_t sum_lb;
71

```



```

72     DinicWithLowerBound(int V) : flow(V+2), S(V), T(V+1), sum_lb(0) {}
73
74     void add_edge(int from, int to, flow_t lb, flow_t ub) {
75         flow.add_edge(from, to, ub-lb);
76         flow.add_edge(S, to, lb);
77         flow.add_edge(from, T, lb);
78         sum_lb += lb;
79     }
80
81     flow_t max_flow(int s, int t) {
82         auto a = flow.max_flow(S, T);
83         auto b = flow.max_flow(s, T);
84         auto c = flow.max_flow(S, t);
85         auto d = flow.max_flow(s, t);
86         return (b == c && a + b == sum_lb) ? b+d : -1;
87     }
88 };

```

二部マッチング

```

1  struct BipartiteMatching {
2      using Graph = vector<vector<int>>>;
3      Graph G;
4      vector<int> match, alive, used;
5      int timestamp;
6
7      BipartiteMatching(int n) : G(n), alive(n, 1),
8                               used(n, 0), match(n, -1), timestamp(0) {}
9
10     void add_edge(int u, int v) {
11         G[u].push_back(v);
12         G[v].push_back(u);
13     }
14
15     int dfs(int idx) {
16         used[idx] = timestamp;
17         for(auto &&to : G[idx]) {
18             int w = match[to];
19             if(alive[to] == 0) continue;
20             if(w < 0 || (used[w] != timestamp && dfs(w))) {
21                 match[idx] = to;
22                 match[to] = idx;
23                 return 1;
24             }
25         }
26         return 0;
27     }
28
29     int bipartite_matching() {
30         int res = 0;
31         for(int i = 0; i < G.size(); i++) {
32             if(alive[i] == 0) continue;
33             if(match[i] == -1) {
34                 ++timestamp;
35                 res += dfs(i);
36             }
37         }

```

```

38         return res;
39     }
40
41     void output() {
42         for(int i = 0; i < G.size(); i++) {
43             if(i < match[i]) {
44                 cout << i << "-" << match[i] << endl;
45             }
46         }
47     }
48 };

```

3 木

木の直径

```

1  struct edge{
2      int to, cost;
3  };
4  using WeightedGraph = vector<vector<edge>>;
5  using pi = pair<int, int>;
6
7  pi dfs(WeightedGraph &G, int idx, int src){
8      pi res(0, idx);
9      for(auto &e : G[idx]) {
10         if(e.to == src) continue;
11         pi cost = dfs(G, e.to, idx);
12         cost.first += e.cost;
13         res = max(res, cost);
14     }
15     return res;
16 }
17
18 int tree_diameter(WeightedGraph &G)
19 {
20     auto far = dfs(G, 0, -1);
21     auto res = dfs(G, far.second, -1);
22     return (res.first);
23 }

```

4 数学

GCD・LCM

```
1 LL gcd(LL a, LL b){
2     if(a < b) swap(a, b);
3     if(b == 0) return a;
4     return gcd(b, a%b);
5 }
6
7 LL lcm(LL a, LL b){
8     return a*b/gcd(a,b);
9 }
```

5 文字列

KMP 法

文字列 $S[0, i-1]$ の接頭辞と接尾辞が最大何文字一致しているかを記録した配列を $O(|S|)$ で構築する

```
1 vector<int> A(s.size()+1);
2 A[0] = -1;
3 int j = -1;
4 for (int i = 0; i < s.size(); i++) {
5     while (j >= 0 && s[i] != s[j]) j = A[j];
6     j++;
7     A[i+1] = j;
8 }
```

6 テクニック

座標圧縮

```
1 vector<int> unzip = a;
2 map<int, int> zip;
3 sort(all(unzip));
4 unzip.erase(unique(all(unzip)), unzip.end());
5 for(int i=0; i<unzip.size(); i++) zip[unzip[i]] = i;
```

スライド最小値 $[i-k, i]$ の最小値を格納した vector を返す

```
1 vector<int> slide_min(vector<int> &a, int k){
2     deque<int> deq;
3     vector<int> b;
4     rep(i, a.size()){
5         // maxはこの不等号の向きを変える
6         while(!deq.empty() && a[deq.back()] >= a[i]) deq.pop_back();
7         deq.push_back(i);
8         b.push_back(a[deq.front()]);
9         if(i-k+1 >= 0 && deq.front() == i-k+1) deq.pop_front();
10    }
11    return b;
12 }
```