# Chapter 7 Programming Techniques with Additional Instructions

#### Register Pairs

- The 8085 is an 8-bit microprocessor. So, its general purpose registers are 8-bits wide.
  - But, it uses 16-bit addresses.
- To allow manipulation of the 16-bit addresses, the 8085 allows its registers to be <u>paired</u>.
  - In other words, it is possible to combine two 8-bit registers together into a 16-bit super register or a register pair.
- The 8085 recognizes only 3 register pairs: B and C, D and E and H and L.
  - It is not possible to pair B and E for example.

#### Identifying Register Pairs

- A register pair is identified with the name of the first register in the pair.
- Therefore, the three register pairs are called:
  - B for the B,C pair
  - D for the D,E pair
  - H for the H,L pair

#### Placing 16-bit Data into a Register Pair

 To place a 16-bit number into a register pair, we can place an 8-bit number in each of the two registers.

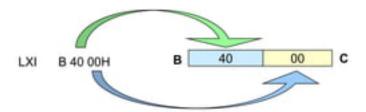
```
MVI L 00H
MVI H 32H
```

But, in what order?

```
MVI L 00H or MVI L 32H
MVI H 32H MVI H 00H
```

#### The LXI instruction

- The 8085 provides an instruction to place the 16-bit data into the register pair in one step.
  - LXI Rp, <16-bit address> (Load eXtended Immediate)
  - The instruction LXI B 4000H will place the 16-bit number 4000 into the register pair B, C.
    - The upper two digits are placed in the 1st register of the pair and the lower two digits in the 2nd.



- Most of the instructions of the 8085 can use a memory location in place of a register.
  - The memory location will become the "memory" register M.
    - MOV M, R
    - MOV M, B
      - copy the data from register B into a memory location.
  - Which memory location?
- The memory location is identified by the contents of the HL register pair.
  - The <u>16-bit contents</u> of the <u>HL register pair</u> are <u>treated</u> as a <u>16-bit</u> address and used to identify the <u>memory location</u>.

## **Direct Memory Access Operations**

Opcode	Operand	Bytes	M-Cycles	T-States	Hex Code	
LDA	16-bit address	3	4	13	3A	•throg0
a 3-byte in	struction; the	second by	tor. The content te specifies the			
	e high-order a flags are affe					
Flags No Example	flags are affe	cted.	on 2050H contr	ains byte F81	H. Load the as	ccumulator

#### **Direct Memory Access Operations**

STA: Store Accumulator Direct

Opcode	Operand	Bytes	M-Cycles	T-States	Hex Code
STA	16-bit	3	4	13	32

**Description** The contents of the accumulator are copied to a memory location specified by the operand. This is a 3-byte instruction; the second byte specifies the low-order address and the third byte specifies the high-order address.

Flags No flags are affected.

Example Assume the accumulator contains 9FH. Load the accumulator contents into memory location 2050H.

Instruction: STA 2050H Hex Code: 32 50 20

Register contents Memory contents before instruction after instruction

A 9F XX F 2050 9F

#### Indirect Addressing Mode

- Using data in memory directly (without loading first into a Microprocessor's register) is called Indirect Addressing.
- Indirect addressing uses the <u>data</u> in a <u>register pair</u> as a <u>16-bit</u> <u>address</u> to <u>identify</u> the <u>memory location</u> being accessed.
  - The HL register pair is <u>always</u> used in conjunction with the memory register "M".
  - The BC and DE register pairs can be used to load data into the <u>Accumultor</u> using indirect addressing.

#### Using the Other Register Pairs

- There is also an instruction for moving data from memory to the accumulator without disturbing the contents of the H and L register.
  - LDAX B/D (LoaD Accumulator extended)
    - Copy the 8-bit contents of the memory location identified by the Rp register pair into the Accumulator.
    - This instruction only uses the BC or DE pair.
    - It does not accept the HL pair.

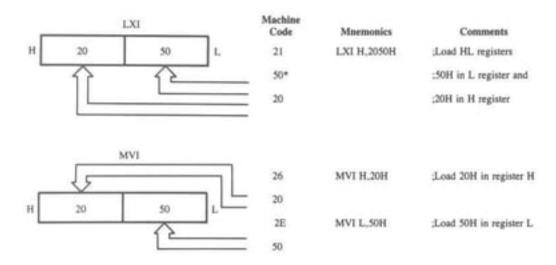
#### Indirect Memory Access Operations

- Use a register PAIR as an address pointer!
- We can define memory access operations using the memory location (16 bit address) stored in a register pair: BC, DE or HL.
- First, we have be able to load the register pairs.

```
LXI B, (16-bit address)
LXI D, (16-bit address)
LXI H, (16-bit address)
```

We can also increment / decrement register pairs.

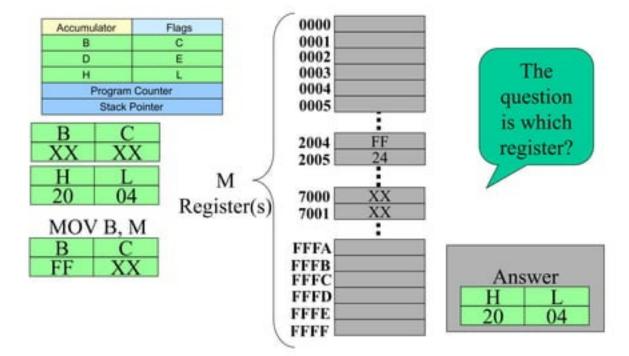
## **Loading Register Pairs**

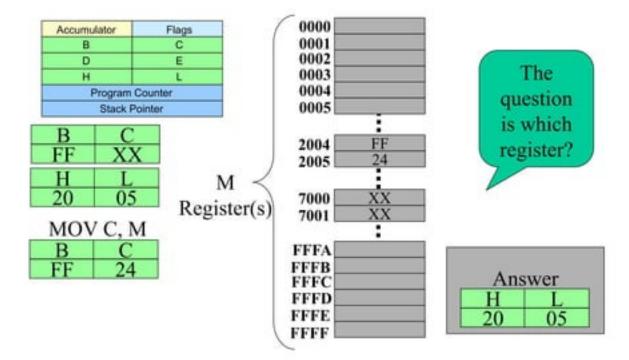


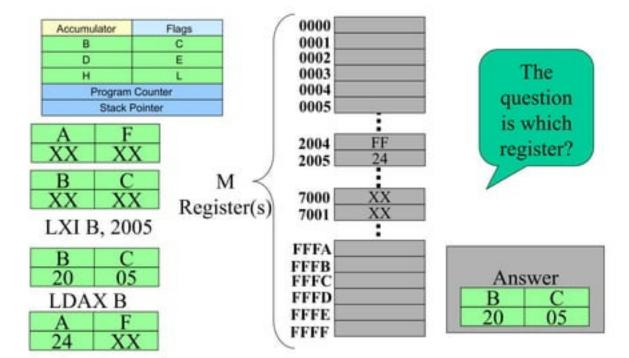
#### Data Transfer from Memory to Processor

 Once the register pairs are loaded with the memory address, we can use them as pointers.

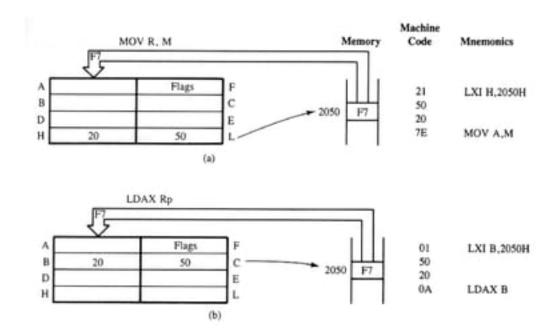
- MOV R,M
   Move the contents of the memory location stored in HL register pair into register (R).
- LDAX B / LDAX D
   Move the contents of the memory location stored in BC (or DE) register pair into the accumulator.







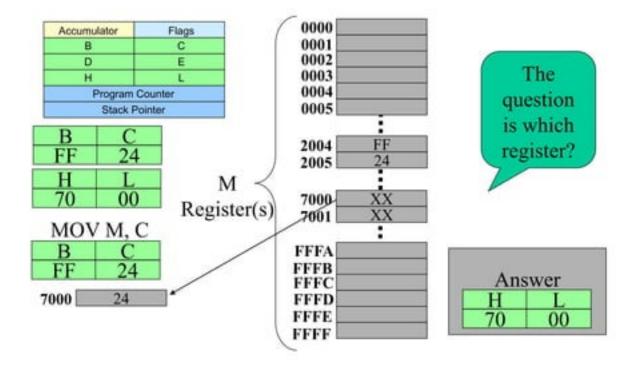
#### Data Transfer from Memory to Processor

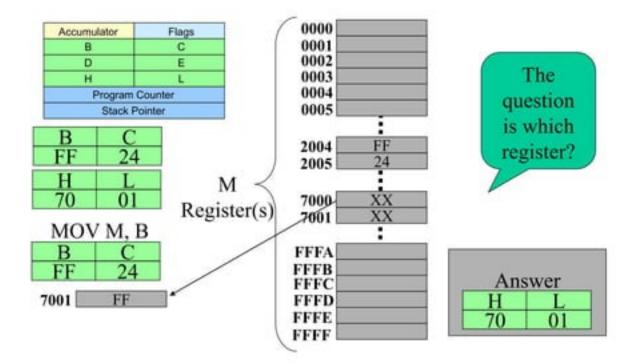


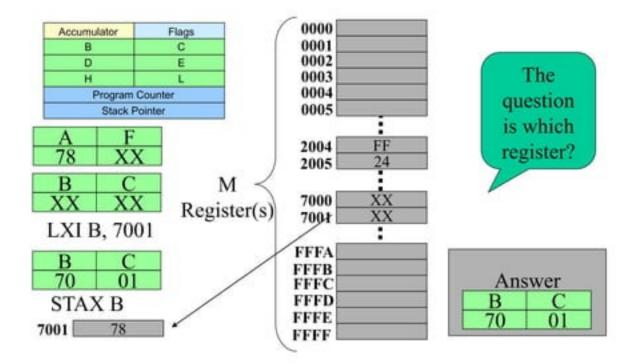
#### Data Transfer from Processor to Memory

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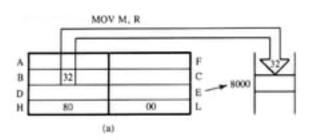
- MOV M,R
   Move the contents of the register (R) into memory location stored in HL register pair.
- STAX B / STAX D
   Move the contents of the accumulator into the memory location stored in BC (or DE) register pair.



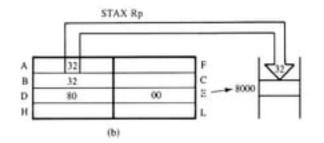




## **Data Transfer from Processor to Memory**



Machine Code	Mnemonics
21 00	LXI H,8000H
80 70	MOV M.B



This instruction copies the contents of the accumulator into memory. Therefore, it is necessary first to copy (B) into A.

11	LXI D,80008
00	
80	
78	MOV A.B
12	STAX D

#### Arithmetic Operations Related to 16 Bits or Register Pairs

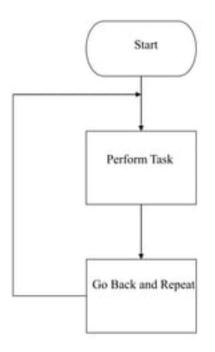
- Now that we have a 16-bit address in a register pair, how do we manipulate it?
  - It is possible to manipulate a 16-bit address stored in a register pair as one entity using some special instructions.

 The register pair is incremented or decremented as one entity. No need to worry about a carry from the lower 8-bits to the upper. It is taken care of automatically.

## Programing Tecniques: Looping, Counting, and Indexing

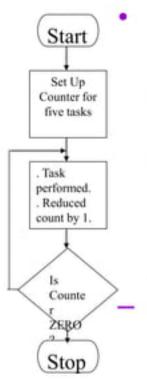
- The programming technique used to instruct the microprocessor to repeat tasks is called looping.
   A loop is set up by instructing the microprocessor to change the sequence of execution and perform the task again. This process is accomplished by using JUMP instructions. In addition, techniques such as counting and indexing are setting up a loop.
- Loops can be classified into two groups
  - Continuous loop: Repeats a task continuously
  - Conditional loop: Repeats a task until certain data conditions are met.

#### Continuos Loop



 A program with a continuous loop does not stop repeating the task until the system is reset.

#### Conditional Loop



The microprocessor needs a counter to repeat the task five times, and when the counting is completed, it needs a flag.

- Counter is set up by loading an appropriate count ain a register.
- Counting is performed by either incrementing or decrementing the counter.
- Loop is set up by a conditional JUMP instruction.
- End of counting is indicating by a flag.

It is easier to count down to zero than count up because the Z flag is set when the register becomes zero.

#### Programing Tecniques: Looping, Counting, and Indexing

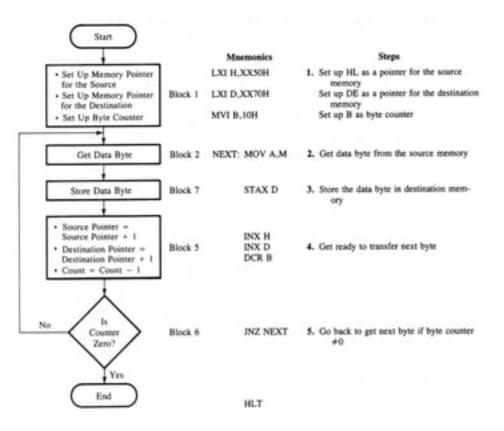
 Conditional loop, Counter, and Indexing Another type of loop includes indexing along with a counter. (Indexing means pointing or referencing objects with sequential numbers. In a library, books are arranged according to numbers, and they are referred to or stored by numbers. This is called indexing.) Similarly, data bytes are stored in memory locations, and those data bytes are referred to by their memory locations.

#### **Block Data Transfer Example**

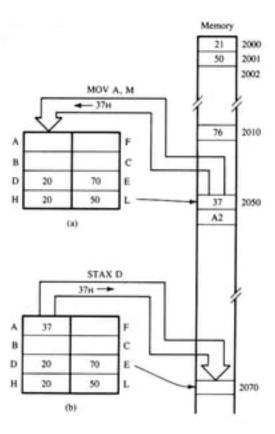
 16 bytes of data are stored in memory locations XX50 to XX5F (16 consecutive memory addresses).

 Transfer (copy) the entire block of data to new memory locations starting at XX70.

 Make use of memory address pointers and indirect addressing!



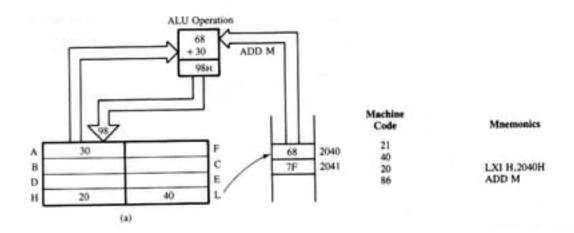
PROGRAM Memory	4				
Address	Hex		Instr	uctions	
HI-LO	Code	Label	Opcode	Operand	Comments
XX00	21	START:	LXI	H,XX50H	:Set up HL as a
0.1	50				; pointer for source
02	XX				; memory
03	11		LXI	D,XX70H	:Set up DE as
04	70				: a pointer for
05	XX				destination
06	06		MVI	B,10H	:Set up B to count
07	10				: 16 bytes
08	7E	NEXT:	MOV	A.M	:Get data byte from
					; source memory
09	12		STAX	D	;Store data byte at
					: destination
0A	23		INX	H	;Point HL to next
					; source location
0B	13		INX	D	:Point DE to
					; next destination
OC:	05		DCR	В	:One transfer is
					; complete,
					; decrement count
0D	C2		JNZ	NEXT	:If counter is not 0,
0E	08				; go back to transfer
06	XX				next byte
10	76		HLT		:End of program
XX50	37				;Data
XXSF	98				
AASE	98				



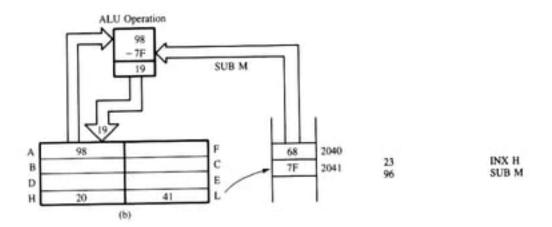
#### Arithmetic Operations Related to Memory

- These instructions perform an arithmetic operation using the contents of a memory location while they are still in memory.
  - ADD M
    - Add the contents of M to the Accumulator
  - SUB M
    - . Sub the contents of M from the Accumulator
  - INR M / DCR M
    - Increment/decrement the contents of the memory location in place.
  - All of these use the contents of the HL register pair to identify the memory location being used.

#### **Arithmetic Operations Using Memory**



## **Arithmetic Operations Using Memory**



#### Illustrative Program: Addition With Carry

- Six bytes of data are stored in memory locations starting at XX50H. Add all the data bytes. Use register B to save any carries generated, while adding the data bytes. Display the entire sum at two output ports, or store the sum at two consecutive memory locations XX70 H. And XX71 H.
  - Data(H): A2, FA, DF, E5, 98, 8B

Memory Address HI-LO	Machine Code	Label	Opcode	Instructions Operand	Comments
XX00	AF		XRA	Α	; Clear (A) to save sum
01	47		MOV	B,A	; Clear (B) to save carry
02	OE		MVI	C,06H	; Set up register C as counter
03	06		1		
04	21		LXI	н, жж50н	; Set up HL as memory pointer
05	50		52.07		
06	жx				
07	86	NXTBYTE:	ADD	м	; Add byte from memory
08	D2		JNC	NXTMEM	; If no carry , do not increment
09	0C				; carry register
OA.	xx				# 1818 7 F 181 E 1818
0В	04		INR	В	; If carry, save carry bit
oc.	23	NXTHEM:	INX	H	; Point to next memory location
00	0D		DCR	C	; One addition is completed
2000					; decrement counter
OE.	C2		JNZ	NXTBYTE	; If all bytes are not yet added,
OF	07				; go back to get next byte
10	xx				P. R. Harris & Bottom Co. Ph. C.
11	D3		OUT	PORT1	: Display low-order byte of the
12	PORT1				; sum at PORT1
13	78		MOV	A,B	: Transfer carry to accumulator
14	D3		OUT	PORT2	; Display Carry digits
15	PORT2				
16	76		HLT		: End of program

#### Illustrative Program: Addition With Carry

#### Storing in memory - Alternative to ouput display

Memory Address HI-LO	Machine Code	Label	Opcode	Instructions Operand	Comments
11	21		LXI	н, жж70н	; Point to the memory location
12	70				; to store answer
13	xx				F1 W 1 / 2010 1 10 10 00 00 00
14	77		MOV	M,A	; Store low-order byte at XX70H
15	23		INX	H	; Point to location XX71H
16	70		MOV	M,B	; Store carry bits
17	76		HLT		; End of the program

#### Additional Logic Operations

#### Rotate

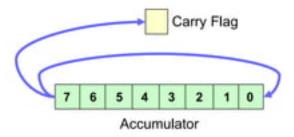
 Rotate the contents of the accumulator one position to the left or right.

```
    Rotate the accumulator left.
    Bit 7 goes to bit 0 AND the Carry flag.
```

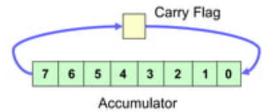
- RAL Rotate the accumulator left through the carry.
   Bit 7 goes to the carry and carry goes to bit 0.
- RRC Rotate the accumulator right.
   Bit 0 goes to bit 7 AND the Carry flag.
- RAR Rotate the accumulator right through the carry.
   Bit 0 goes to the carry and carry goes to bit 7.

#### RLC vs. RAL

RLC



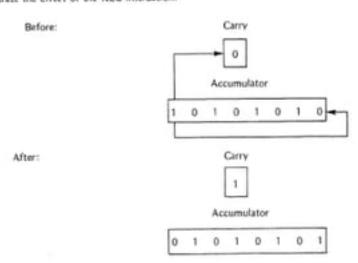
RAL



## RLC - Example

#### Example:

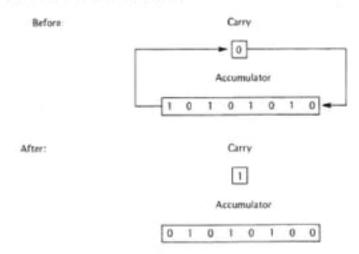
Assume that the accumulator contains the value OAAH and the carry flag is zero. The following diagrams illustrate the effect of the RLC imstruction.



## RAL - Example

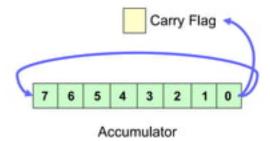
#### Example:

Assume that the accumulator contains the value OAAH and the carry flag is zero. The following diagrams illustrate the effect of the RAL instruction:

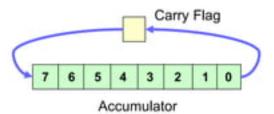


#### RRC vs. RAR

RRC



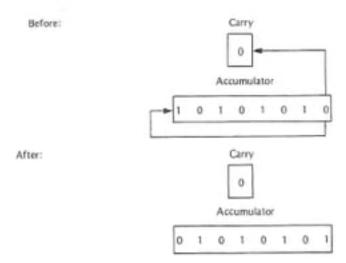
RAR



## RRC - Example

#### Example:

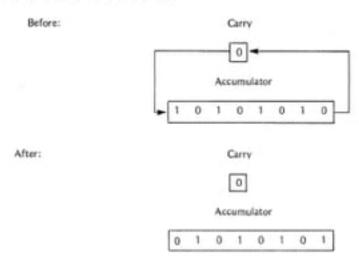
Assume that the accumulator contains the value 0AAH and the carry flag is zero. The following diagrams illustrate the effect of the RRC instruction:



#### RAR - Example

#### Example:

Assume that the accumulator contains the value 0AAH and the carry flag is zero. The following diagrams illustrate the effect of the RAR instruction:



#### **Logical Operations**

#### Compare

 Compare the contents of a register or memory location with the contents of the accumulator.

- CMP	R/M	Compare the contents of the register or memory location to the contents of the accumulator.
- CPI	#	Compare the 8-bit number to the contents of the accumulator.

- The compare instruction sets the flags (Z, Cy, and S).
- The compare is done using an internal subtraction that does not change the contents of the accumulator.

$$A - (R/M/#)$$

#### Logical Operations - Compare

- CMP R/M: This is a 1-byte instruction
  - It compares the data byte in register or memory with the contents of the accumulator.
  - If (A)<(R/M), the CY flag is set and the Z flag is reset</li>
  - If (A)=(R/M), the Z flag is set and the CY flag is reset
  - If (A)>(R/M), the CY and Z flags are reset
  - When memory is operand, its address is specified by (HL).
  - No contents are modified.
- CPI 8-bit: This is a 2 byte instruction
  - It compares the second byte with the contents of the accumulator.
  - If (A)<(8-bit), the CY flag is set and the Z flag is reset</li>
  - If (A)=(8-bit), the Z flag is set and the CY flag is reset
  - If (A)>(8-bit), the CY and Z flags are reset
  - No contents are modified.