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# Microprocessor Systems and Interfacing

EEE 342

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Nesruminallah

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# About the Instructor

- MSc, Communication Engineering, 2014
  - University of Portsmouth, United Kingdom, .
- BSc, Telecommunication Engineering, 2012
  - UET Peshawar , Pakistan.

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# Course Introduction

## ■ Textbook

- ❑ The Intel Microprocessors by Barry B. Brey, 8th Edition, Pearson
- ❑ Assembly Language Programming and Organization of the IBM PC by Ytha Yu and Charles Marut, 1992, 1st Edition, McGraw-Hill
- ❑ Embedded Systems: Introduction to Arm® Cortex™-M Microcontrollers, by Jonathan W Valvano, Vol 1, 5th Edition, 2019, CreateSpace

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# Grading Policy

- Assignments 10%
  - Minimum 4
- Quizzes (scheduled/surprised) 15%
  - Minimum 4
- Midterm 25%
- Final exam 50%

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# Academic Honesty

- Your work and participation in the course **must** be your own
- If students are found to have collaborated excessively or to have cheated (e.g. by copying or sharing answers in assignments or during an examination), all involved will at a minimum receive grades of 0 for the first infraction
- Further infractions may result in failure in the course

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# Lectures

- Lecture notes given by instructor.
- Please be courteous in class
  - Arrive on time
  - Turn off cell phones / sound on laptop
  - Keep quiet ...
  - Drinking or eating is strictly prohibited
- Attendance is important
  - There are just things that you cannot learn from reading notes
  - 80% is must to appear in final exam and pass the course

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# Few Recommendations

- “Eighty percent of success is showing up.”
  - Come to lectures, discussions, lab
  
- If you are not sure: ask!
  - Talk to us after class, send email, come to office hours
  - Early communication solves problems easiest
  - Don't wait until it's too late
  
- Email protocol
  - Write your full name and registration ID.
    - We need to know who you are.
  - Do not forget to write subject of email
  - Be professional

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# Couse CLOs and PLOs

- **Theory CLOs:**

- CLO-1: To write the Intel-assembly code using the knowledge of programmer model, addressing mode and assembly language programming concepts. (PLO3-C5)
- CLO-2: To integrate the memory, timer, I/O and PPI with microprocessor using address decoding techniques. (PLO3-C5)
- CLO-3: To design digital system based on microprocessor using the knowledge of architecture, memory, timer, I/O and PPI interfacing. (PLO3-C5)

- **Lab CLOs:**

- CLO4: To explain and reproduce the Intel-assembly and STM32F407VG C-Programming codes using software and hardware platforms. (PLO5-P3)
- CLO5: To design digital system using the knowledge of STM32F407VG C-Programming and peripherals. (PLO3-C5)
- CLO6: To write effective report(s) of the assigned project. (PLO10-A2).
- CLO7: To describe the impact of digital system on our society and environment using existing industrial standards (PLO7-C6).
- CLO8: To justify the significance of designed project to the society using existing engineering practices (PLO6-C6).



# Course Contents

- Introduction to microprocessor and microcontroller and
- Basic concepts and definitions of computer architecture and organizations
- Introduction to Programmers model of 8086/88
- Assembly Language Programming for 8086/88 Architecture
- Interfacing of RAM/ROM with 8088 microprocessors
- Introduction to Microcontroller (STM32F407VG)
- Interfacing of RAM/ROM with 8086 microprocessors
- Stack programming and memory mapping
- 8254 timer/counter interfacing with 8088 microprocessors
- I/O interfacing (isolated and memory-mapped) with 8088 microprocessors
- 8255 PPI interfacing with 8088 microprocessors
- A/D and D/A Conversion
- Hardware Interrupts
- Interfacing output devices with 8088 microprocessors using PPI

# Number System and Conversions

CLO	Bloom Taxonomy	Specific Outcome
CLO1	C2	Comprehend the theoretical knowledge of number systems such as Binary, Octal, Decimal and Hexadecimal numbers using standard conversion methods.

## ■ Outline

- ❑ Binary numbers
- ❑ Number-Base conversions
- ❑ Octal and hexadecimal numbers
- ❑ Complements and signed binary numbers

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# Digital Computer Systems

- Digital systems consider *discrete* amounts of data.
- Examples
  - 26 letters in the alphabet
  - 10 decimal digits
- Larger quantities can be built from discrete values
  - Words made of letters
  - Numbers made of decimal digits (e.g. 239875.32)

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# Digital Computer Systems

- Questions to ask
  - How the numbers are represented in digital systems?
  - How computer performs basic arithmetic operations?
- Computers operate on *binary* values (0 and 1)
- Easy to represent binary values electrically
  - Voltages and currents
  - Advantages
    - Can be implemented using circuits
    - Create the building blocks of modern computers

# Understanding Decimal Numbers

- Decimal numbers are made of decimal digits:  
(0,1,2,3,4,5,6,7,8,9)
- But how many items does a decimal number represent?
  - $8653 = 8 \times 10^3 + 6 \times 10^2 + 5 \times 10^1 + 3 \times 10^0$
- What about fractions?
  - $97654.35 = 9 \times 10^4 + 7 \times 10^3 + 6 \times 10^2 + 5 \times 10^1 + 4 \times 10^0 + 3 \times 10^{-1} + 5 \times 10^{-2}$
  - In formal notation  $\rightarrow (97654.35)_{10}$
- Why do we use 10 digits, anyway?



# Understanding Binary Numbers

- Binary numbers are made of binary digits (bits)
  - 0 and 1
- How many items does a binary number represent?
  - $(1011)_2 = 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 = (11)_{10}$
- What about fractions?
  - $(110.10)_2 = 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0 + 1 \times 2^{-1} + 0 \times 2^{-2}$
- Groups of eight bits are called a *byte*
  - $(11001001)_2$
- Groups of four bits are called a *nibble*
  - $(1101)_2$

# Understanding Octal Numbers

- Octal numbers are made of octal digits
  - 0,1,2,3,4,5,6,7
- How many items does an octal number represent?
  - $(4536)_8 = 4 \times 8^3 + 5 \times 8^2 + 3 \times 8^1 + 6 \times 8^0 = (1362)_{10}$
- What about fractions?
  - $(465.27)_8 = 4 \times 8^2 + 6 \times 8^1 + 5 \times 8^0 + 2 \times 8^{-1} + 7 \times 8^{-2}$
- Octal numbers don't use digits 8 or 9
- Why would someone use octal number, anyway?

# Understanding Hexadecimal Numbers

- Hexadecimal numbers are made of 16 digits
  - (0,1,2,3,4,5,6,7,8,9,A, B, C, D, E, F)
- How many items does an hex number represent?
  - $(3A9F)_{16} = 3 \times 16^3 + 10 \times 16^2 + 9 \times 16^1 + 15 \times 16^0 = 14999_{10}$
- What about fractions?
  - $(2D3.5)_{16} = 2 \times 16^2 + 13 \times 16^1 + 3 \times 16^0 + 5 \times 16^{-1} = 723.3125_{10}$
- Note that *each* hexadecimal digit can be represented with four bits.
  - $(1110)_2 = (E)_{16}$



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## Exercise (Convert to decimal)

- $(1011.101)_2$
- Answer = 11.625

- $(24.6)_8$
- Answer = 20.75

- $(IBC2)_{16}$
- Answer = 7106

# Putting It All Together

Decimal	Binary	Octal	Hexadecimal
0	0	0	0
1	1	1	1
2	10	2	2
3	11	3	3
4	100	4	4
5	101	5	5
6	110	6	6
7	111	7	7
8	1000	10	8
9	1001	11	9
10	1010	12	A
11	1011	13	B
12	1100	14	C
13	1101	15	D
14	1110	16	E
15	1111	17	F

# Convert an Integer from Decimal to Another Base

- For each digit position
  - Divide decimal number by the base (e.g. 2)
  - The *remainder* is the lowest-order digit
  - Repeat first two steps until no *divisor* remains

Example for  $(13)_{10}$

	Integer Quotient		Remainder	Coefficient
$13/2 =$	6	+	$\frac{1}{2}$	$a_0 = 1$
$6/2 =$	3	+	0	$a_1 = 0$
$3/2 =$	1	+	$\frac{1}{2}$	$a_2 = 1$
$1/2 =$	0	+	$\frac{1}{2}$	$a_3 = 1$

Answer  $(13)_{10} = (a_3 a_2 a_1 a_0)_2 = (1101)_2$

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## Exercise (Convert decimal to other bases)

- 291 to binary
- Answer =  $(100100011)_2$
- 291 to octal
- Answer =  $(443)_8$
- 291 to hexadecimal
- Answer =  $(123)_{16}$

# Convert a Fraction from Decimal to Another Base

- For each digit position
  - Multiply decimal number by the base (e.g. 2)
  - The integer is the highest-order digit
  - Repeat first two steps until fraction becomes zero

Example for  $(0.625)_{10}$

	Integer	Fraction	Coefficient
$0.625 \times 2 =$	1	+	0.25 $a_{-1} = 1$
$0.250 \times 2 =$	0	+	0.50 $a_{-2} = 0$
$0.500 \times 2 =$	1	+	0 $a_{-3} = 1$

Answer  $(0.625)_{10} = (0.a_{-1} a_{-2} a_{-3})_2 = (0.101)_2$

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## Exercise (Convert fraction to decimal)

- $(0.513)_{10}$  to octal
- Answer =  $(0.406517\dots)_8$
- $0.513 \times 8 = 4.104$
- $0.104 \times 8 = 0.832$
- $0.832 \times 8 = 6.656$
- $0.656 \times 8 = 5.248$
- $0.248 \times 8 = 1.984$
- $0.984 \times 8 = 7.872$

# Binary Addition

- Binary addition is very simple.
- An example of adding two binary numbers

$$\begin{array}{rcccccc} & 1 & 1 & 1 & 1 & 1 & 1 & \leftarrow \text{carries} \\ & & 1 & 1 & 1 & 1 & 0 & 1 \\ + & & & 1 & 0 & 1 & 1 & 1 \\ \hline 1 & 0 & 1 & 0 & 1 & 0 & 0 \end{array}$$

# Binary Subtraction

- We can also perform subtraction (with borrows in place of carries)
- Example: subtract  $(10111)_2$  from  $(1001101)_2$

$$\begin{array}{r}
 \begin{array}{ccccccc}
 & & 1 & & 10 & & \\
 0 & \cancel{1}0 & 10 & 0 & \cancel{0}10 & & 
 \end{array}
 \begin{array}{c} \longleftarrow \text{borrows} \end{array} \\
 \begin{array}{ccccccc}
 \cancel{1} & \cancel{0} & \cancel{0} & \cancel{1} & \cancel{1} & \cancel{0} & 1 \\
 - & & & 1 & 0 & 1 & 1 & 1 \\
 \hline
 & & 1 & 1 & 0 & 1 & 1 & 0
 \end{array}
 \end{array}$$



# Binary Multiplication

- Binary multiplication is much the same as decimal multiplication
  - The multiplication operations are much simpler

$$\begin{array}{r} \phantom{X} \phantom{00000} 1 \phantom{00} 0 \phantom{00} 1 \phantom{00} 1 \phantom{00} 1 \\ X \phantom{00000} \phantom{00000} 1 \phantom{00} 0 \phantom{00} 1 \phantom{00} 0 \\ \hline \phantom{00000} \phantom{00000} 0 \phantom{00} 0 \phantom{00} 0 \phantom{00} 0 \phantom{00} 0 \\ \phantom{00000} \phantom{0000} 1 \phantom{00} 0 \phantom{00} 1 \phantom{00} 1 \phantom{00} 1 \\ \phantom{00000} 0 \phantom{00} 0 \phantom{00} 0 \phantom{00} 0 \phantom{00} 0 \\ 1 \phantom{00} 0 \phantom{00} 1 \phantom{00} 1 \phantom{00} 1 \\ \hline 1 \phantom{00} 1 \phantom{00} 1 \phantom{00} 0 \phantom{00} 0 \phantom{00} 1 \phantom{00} 1 \phantom{00} 0 \end{array}$$

# Converting Between Base 16 and Base 2

$$3A9F_{16} = \begin{array}{cccc} \underline{0011} & \underline{1010} & \underline{1001} & \underline{1111}_2 \\ 3 & A & 9 & F \end{array}$$

- Conversion is easy
  - Determine 4-bit value for each hex digit
- Note that there are  $2^4 = 16$  different values of four bits
- Easier to read and write in hexadecimal
- Representations are equivalent

## Converting Between Base 16 and Base 8

Diagram illustrating the conversion of hexadecimal 3A9F<sub>16</sub> to octal 35237<sub>8</sub> using 3-bit groups:

Hexadecimal: 3A9F<sub>16</sub> = 0011 1010 1001 1111<sub>2</sub>

Octal: 35237<sub>8</sub> = 011 101 010 011 111<sub>2</sub>

The mapping shows that each hexadecimal digit (4 bits) is converted to two octal digits (3 bits each) by grouping the bits into 3-bit groups from the right.

1. Convert from Base 16 to Base 2
2. Regroup bits into groups of three starting from right
3. Ignore leading zeros
4. Each group of three bits forms an octal digit

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# Exercise

- Convert base 16 to base 8 without intermediate stage of base 10
- $(B98D)_{16}$
- Answer =  $(134615)_8$

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# Converting Between Base 8 and Base 16

$$\begin{aligned} \blacksquare (673.124)_8 &= (110 \ 111 \ 011 \ . \ 001 \ 010 \ 100)_2 \\ &\quad \downarrow \\ &= (\underline{1} \ \underline{1011} \ \underline{1011} \ . \ \underline{0010} \ \underline{1010} \ 0) \\ &= (\underline{1} \ \underline{B} \ \underline{B} \ . \ \underline{2} \ \underline{A}) \end{aligned}$$

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# Complements

- Used in computers to simplify the subtraction operation
  - For each base- $r$  system
    - Diminished radix complement or  $r-1$ 's complement
    - Radix complement or  $r$ 's complement
  - For example for base-2 system
    - 1's complement
    - 2's complement
  - For base-10 system
    - 9's complement
    - 10's complement
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## Diminished Radix Complement ( $r-1$ 's complement)

- Given a number 'N' in base 'r' with 'n' digits,  $r-1$ 's complement is defined as  $(r^n - 1) - N$
- For example if  $N = (546700)_{10}$  then  $r = 10$  and  $n = 6$ 
  - $(r^n - 1) = 10^6 - 1 = 999999$
  - 9's complement of  $N = (r^n - 1) - N = 999999 - 546700 = 453299$
  - Similarly for  $N = (012398)_{10}$ , 9's complement of N is  $999999 - 012398 = 987601$

## Diminished Radix Complement ( $r-1$ 's complement)

- For a binary number,  $r = 2$  and  $r-1$  or 1's complement can be found just like base-10 numbers
- For example if  $N = (1010)_2$  then  $r = 2$  and  $n = 4$ 
  - $(r^n - 1) = 2^4 - 1 = (15)_{10} = (1111)_2$
  - 1's complement of  $N = (r^n - 1) - N = 1111 - 1010 = (0101)_2$
- Shortcut: Invert all the bits of  $N$  in order to take its 1's complement
  - 1's complement of 1011000  $\rightarrow$  0100111
  - 1's complement of 0101101  $\rightarrow$  1010010



# Radix Complement (r's complement)

- Given a number 'N' in base 'r' with 'n' digits, r's complement is defined as  $r^n - N$  for  $N \neq 0$  and 0 for  $N = 0$ 
  - r's complement can also be obtained by adding 1 to r-1's complement
- For example if  $N = (546700)_{10}$ 
  - 9's complement of  $N = (r^n - 1) - N = 999999 - 546700 = 453299$
  - 10's complement of  $N = 9's \text{ complement} + 1 = r^n - N = 453300$

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## Radix Complement (r's complement)

- For a binary number,  $r = 2$  and  $r$  or 2's complement can be found just like base-10 numbers
- For example if  $N = (1010)_2$  then  $r = 2$  and  $n = 4$ 
  - 1's complement of  $N = (r^n - 1) - N = 1111 - 1010 = (0101)_2$
  - 2's complement of  $N = 1$ 's complement of  $N + 1 = 0110$

# 2's Complement Shortcuts

## ■ Algorithm 1

- Complement each bit and then add 1 to the result
- Example: Find the 2's complement of  $(01100101)_2$  and of its 2's complement

$$\begin{array}{r} N = 01100101 \\ \quad 10011010 \\ + \quad \quad 1 \\ \hline 10011011 \end{array}$$

$$\begin{array}{r} [N] = 10011011 \\ \quad 01100100 \\ + \quad \quad 1 \\ \hline 01100101 \end{array}$$

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# 2's Complement Shortcuts

## ■ Algorithm 2

- Starting with the least significant bit, copy all of the bits up to and including the first 1 bit and then complementing the remaining bits
- Example: Find the 2's complement of  $(01100101)_2$

$$N = 0\ 1\ 1\ 0\ 0\ 1\ 0\ 1$$

$$[N] = 1\ 0\ 0\ 1\ 1\ 0\ 1\ 1$$

# Signed Numbers and their representation

- Plus and minus sign used for decimal numbers: 25 (or +25), -16, etc.
- For computers, desirable to represent everything as *bits*
- Three types of signed binary number representations
  - signed magnitude, 1's complement, 2's complement
- In each case: left-most bit indicates sign: positive (0) or negative (1)

## **signed magnitude**

$$\begin{array}{c} \text{00001100}_2 = 12_{10} \\ \swarrow \quad \nwarrow \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

$$\begin{array}{c} \text{10001100}_2 = -12_{10} \\ \swarrow \quad \nwarrow \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

# 1's Complement Representation

- Invert all bits
  - $00110011 \rightarrow 11001100$
  - $10101010 \rightarrow 01010101$
- For an  $n$  bit number  $N$ , the 1's complement is  $(2^n - 1) - N$
- To find negative of 1's complement number take the 1's complement

$$\begin{array}{c} \text{00001100}_2 = 12_{10} \\ \swarrow \quad \nwarrow \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

$$\begin{array}{c} \text{11110011}_2 = -12_{10} \\ \swarrow \quad \nwarrow \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

# 2's Complement Representation

- Invert all bits and add 1
  - $00110011 \rightarrow 11001101$
  - $10101010 \rightarrow 01010110$
- For an  $n$  bit number  $N$  the 2's complement is  $(2^n - 1) - N + 1$
- To find negative of 2's complement number take the 2's complement

$$\begin{array}{c} \text{00001100}_2 = 12_{10} \\ \swarrow \quad \nwarrow \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

$$\begin{array}{c} \text{11110100}_2 = -12_{10} \\ \swarrow \quad \nwarrow \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

# 1's Complement Addition

- Add  $+(1100)_2$  and  $+(0001)_2$ 
  - $(12)_{10} = +(1100)_2 = 01100_2$
  - $(1)_{10} = +(0001)_2 = 00001_2$

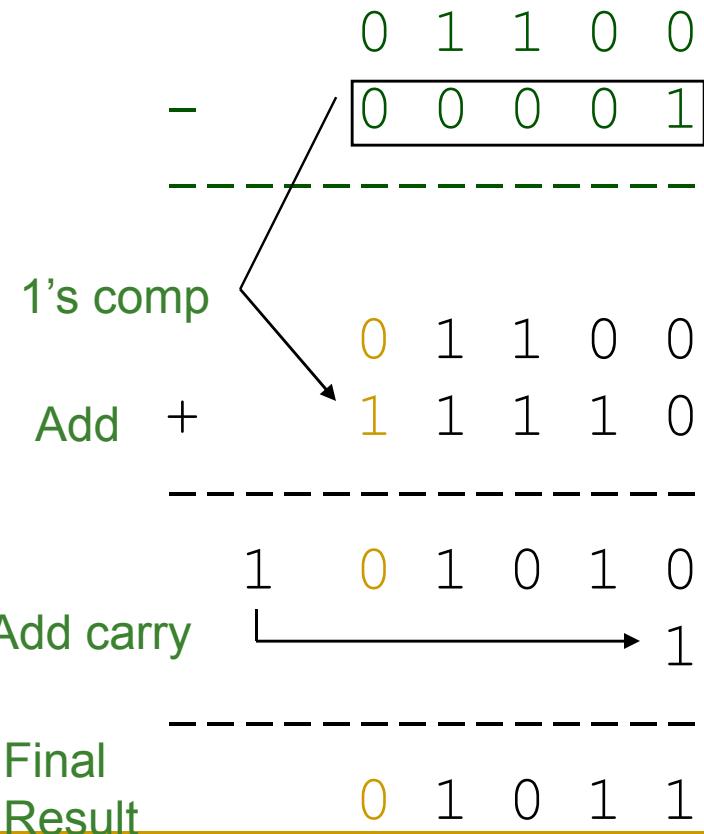
Step 1: Add binary numbers  
Step 2: Add carry to low-order bit

		0	1	1	0	0
Add	+	0	0	0	0	1
		-----				
		0	0	1	1	0
Add carry						1
						0
		-----				
Final Result		0	1	1	0	1



## 1's Complement Subtraction

- **Subtract**  $+(0001)_2$  from  $+(1100)_2$ 
  - $(12)_{10} = +(1100)_2 = 01100_2$
  - $(-1)_{10} = -(0001)_2 = 11110_2$



## Step 1: Take 1's complement of 2<sup>nd</sup> operand

## Step 2: Add binary numbers

### Step 3: Add carry to low order bit

# 2's Complement Addition

- Add  $+(1100)_2$  and  $+(0001)_2$ .

- $(12)_{10} = +(1100)_2 = 01100_2$

- $(1)_{10} = +(0001)_2 = 00001_2$

Step 1: Add binary numbers

Step 2: Ignore carry bit

Add

$$\begin{array}{r} + \quad \begin{array}{ccccc} 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{array} \\ \hline \end{array}$$

Final Result

0

0 1 1 0 1

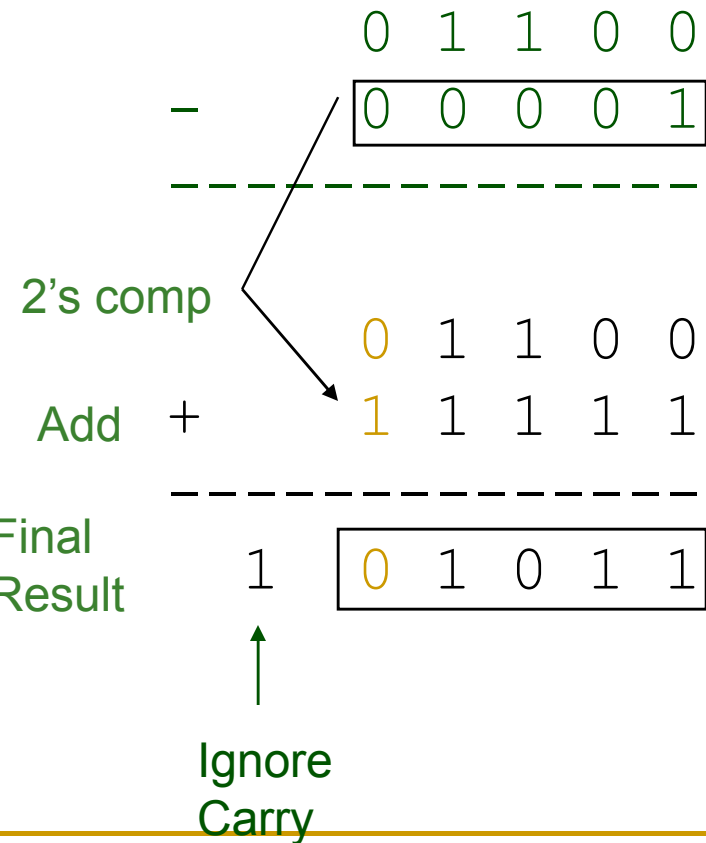
↑  
Ignore

# 2's Complement Subtraction

- **Subtract**  $+(0001)_2$  from  $+(1100)_2$

- $(12)_{10} = +(1100)_2 = 01100_2$

- $(-1)_{10} = -(0001)_2 = 11111_2$



Step 1: Take 2's complement of 2<sup>nd</sup> operand

Step 2: Add binary numbers

Step 3: Ignore carry bit

## 2's Complement Subtraction: Example 2

- Let's compute  $(13)_{10} - (5)_{10}$ 
  - $(13)_{10} = +(1101)_2 = (01101)_2$
  - $(-5)_{10} = -(0101)_2 = (11011)_2$
- Adding these two 5-bit codes

$$\begin{array}{rcccccc} & & & 0 & 1 & 1 & 0 & 1 \\ & & + & 1 & 1 & 0 & 1 & 1 \\ & & \hline \text{carry} \swarrow & 1 & 0 & 1 & 0 & 0 & 0 \end{array}$$

- Discarding the carry bit, the sign bit is seen to be zero, indicating a correct result. Indeed,  
 $(01000)_2 = +(1000)_2 = +(8)_{10}$

## 2's Complement Subtraction: Example 3

- Let's compute  $(5)_{10} - (12)_{10}$ 
  - $(-12)_{10} = -(1100)_2 = (10100)_2$
  - $(5)_{10} = +(0101)_2 = (00101)_2$
- Adding these two 5-bit codes

$$\begin{array}{r} \phantom{+} \phantom{0} \phantom{0} \phantom{1} \phantom{0} \phantom{1} \\ \phantom{+} \phantom{0} \phantom{0} \phantom{1} \phantom{0} \phantom{1} \\ + \phantom{0} \phantom{0} \phantom{1} \phantom{0} \phantom{1} \phantom{0} \phantom{0} \\ \hline \phantom{0} \phantom{0} \phantom{1} \phantom{0} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \\ \phantom{0} \phantom{0} \phantom{1} \phantom{0} \phantom{1} \phantom{0} \phantom{0} \phantom{0} \end{array}$$

- Here, there is no carry bit and the sign bit is 1
  - This indicates a negative result, which is what we expect.  $(11001)_2 = -(7)_{10}$

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# Subtraction using $r$ 's complements

- Subtraction of two  $n$ -digit unsigned numbers  $M - N$  in base  $r$  is performed as follows
  - Add  $M$  to  $r$ 's complement of  $N$
  - If  $M \geq N$  sum will produce an end carry which can be discarded
  - If  $M < N$  sum does not produce end carry. Take  $r$ 's complement of the result to know exact result

# Subtraction using r's complements

- Let  $M = (52532)_{10}$ ,  $N = (3250)_{10}$ ,  $M - N = ?$

$M =$  52532

$N =$  03250

10's complement of  $N =$  96750

$M - N =$  52532  
+96750

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149282

Discard end carry and the result is

Answer =  $(49282)_{10}$

# Subtraction using r's complements

- Let  $M = (3250)_{10}$ ,  $N = (72532)_{10}$ ,  $M - N = ?$

$M =$  03250

$N =$  72532

10's complement of  $N =$  27468

$M - N =$  03250  
+27468

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30718

No end carry in this case

Answer = - (10's complement of  $M - N$ ) = - 69282

Same rules can be applied to base-2 numbers



# Subtraction using $r-1$ 's complements

- Subtraction of two  $n$ -digit unsigned numbers  $M-N$  in base  $r$  is performed as follows
  - Add  $M$  to  $r-1$ 's complement of  $N$
  - If  $M \geq N$  sum will produce an end carry which will be added to least significant digit of the sum
  - If  $M < N$  sum does not produce end carry. Take  $r-1$ 's complement of the result to know exact answer

# Subtraction using r-1's complements

- Let  $M = (1010100)_2$ ,  $N = (1000011)_2$ ,  $M - N = ?$

$M =$  1010100

$N =$  1000011

1's complement of  $N =$  0111100

$M - N =$  1010100  
+0111100

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10010000

└──────────┘↑

Add end carry around

Answer =  $(0010001)_2$

# Subtraction using r-1's complements

- Let  $M = (1000011)_2$ ,  $N = (1010100)_2$ ,  $M - N = ?$

$M =$  1000011

$N =$  1010100

1's complement of  $N =$  0101011

$M - N =$  1000011  
+0101011

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1101110

No end carry in this case

Answer = - (1's complement of  $M - N$ ) = -(0010001)