



Microprocessor Based Systems

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CHAPTER 8

THE STACK AND INTRODUCTION TO PROCEDURES

The Stack

- The stack segment of a program is used for temporary storage of data and addresses
- A stack is a one-dimensional data structure
- Items are added to and removed from one end of the structure using a "Last In - First Out" technique (LIFO)
- The top of the stack is the last addition to the stack
- The statement `.STACK 100H` in your program sets aside a block of 256 bytes of memory to hold the stack
- The SS (Stack Segment Register) contains the segment number of the stack segment

The Stack (cont'd)

- The complete segment:offset address to access the stack is SS:SP
- Initially before any data or addresses have been placed on the stack, the **SP** contains the offset address of the memory location immediately following the stack segment

Empty Stack

Offset

0000

0002

0004

0006

•

•

•

•

0100



SP:

0100

(Beyond the end of the stack)

PUSH Instruction

- **PUSH** instruction adds a new word to the stack
- SYNTAX: **PUSH source**
where source is a 16-bit register or memory word
- **PUSH** instruction causes
 - the stack pointer (**SP**) to be decreased by 2.
 - Then a copy of the value in the source field is placed in the address specified by **SS:SP**.

- Initially SP points to a location immediately following the stack. The first push decreases SP by 2, making it point to the last word in the stack
- Because **each PUSH decreases the SP**, the stack is filled a word at a time backwards from the last available word in the stack toward the beginning of the stack.

How Words Are Added To Stack

Offset

0000

•
•
•

00F4

00F6

00F8

00FA

00FC

00FE

0100



SP:

It stack is empty, SP has a value of 100h; otherwise it has a value between 0000-00FEh

(Beyond the end of the stack)

How Words Are Added To Stack

**EMPTY
STACK**

Offset

0000

•
•
•

00F4

00F6

00F8

00FA

00FC

00FE

0100

SP: 0100

AX: 1234

BX: 5678

← SP

(Beyond the end of the stack)

How Words Are Added To Stack

**AFTER
PUSH AX**

Offset

0000

⋮

00F4

00F6

00F8

00FA

00FC

00FE

0100



SP: 00FE

AX: 1234

BX: 5678

← SP

(Beyond the end of the stack)

How Words Are Added To Stack

**AFTER
PUSH BX**

Offset

0000

⋮

00F4

00F6

00F8

00FA

00FC

00FE

0100



SP: 00FC

AX: 1234

BX: 5678

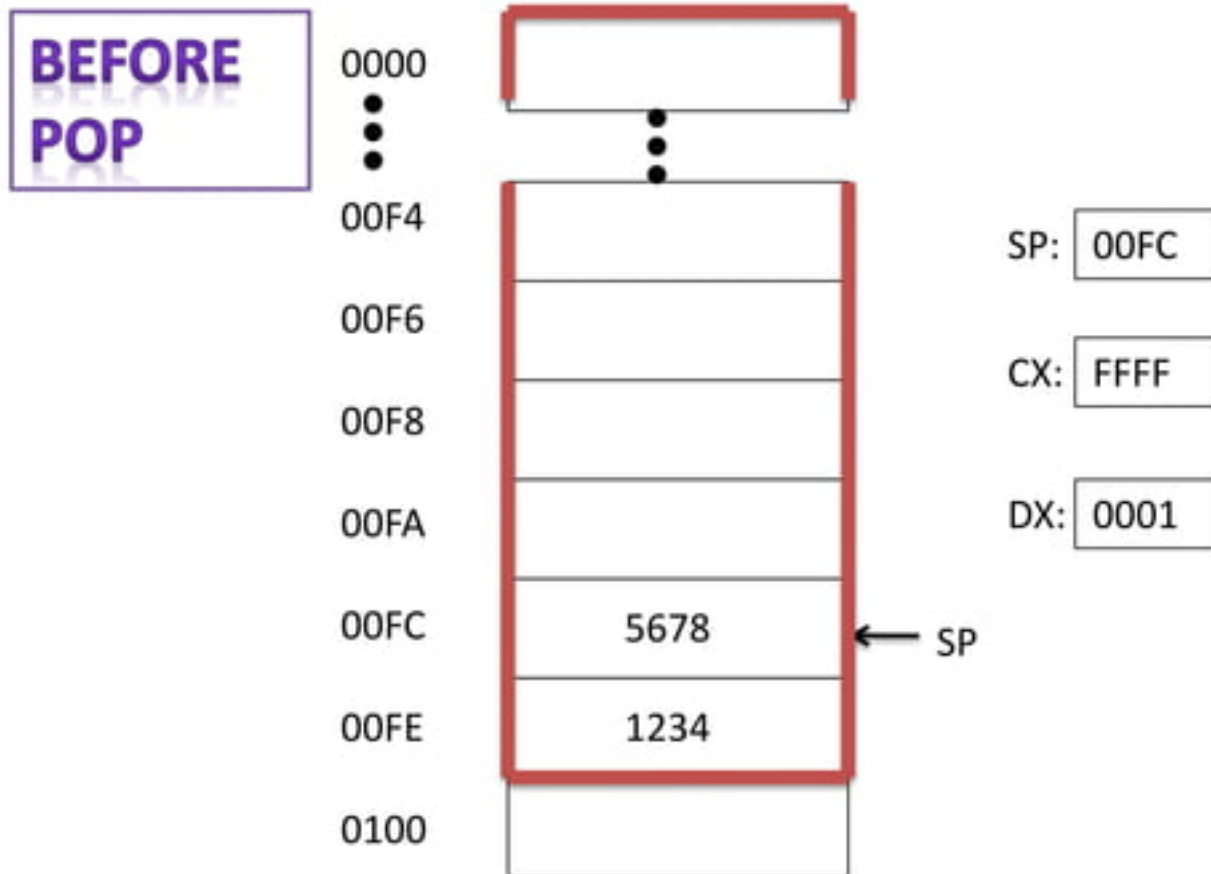
← SP

(Beyond the end of the stack)

POP Instruction

- POP instruction removes the last word placed on the stack
- SYNTAX: POP destination
 - where source is a 16-bit register or memory word
- POP instruction causes
 - the contents of SS:SP to be moved to the destination field
 - It increases the stack pointer (SP) by 2
- *Restrictions:*
 1. PUSH and POP work only with words
 2. Byte and immediate data operands are illegal

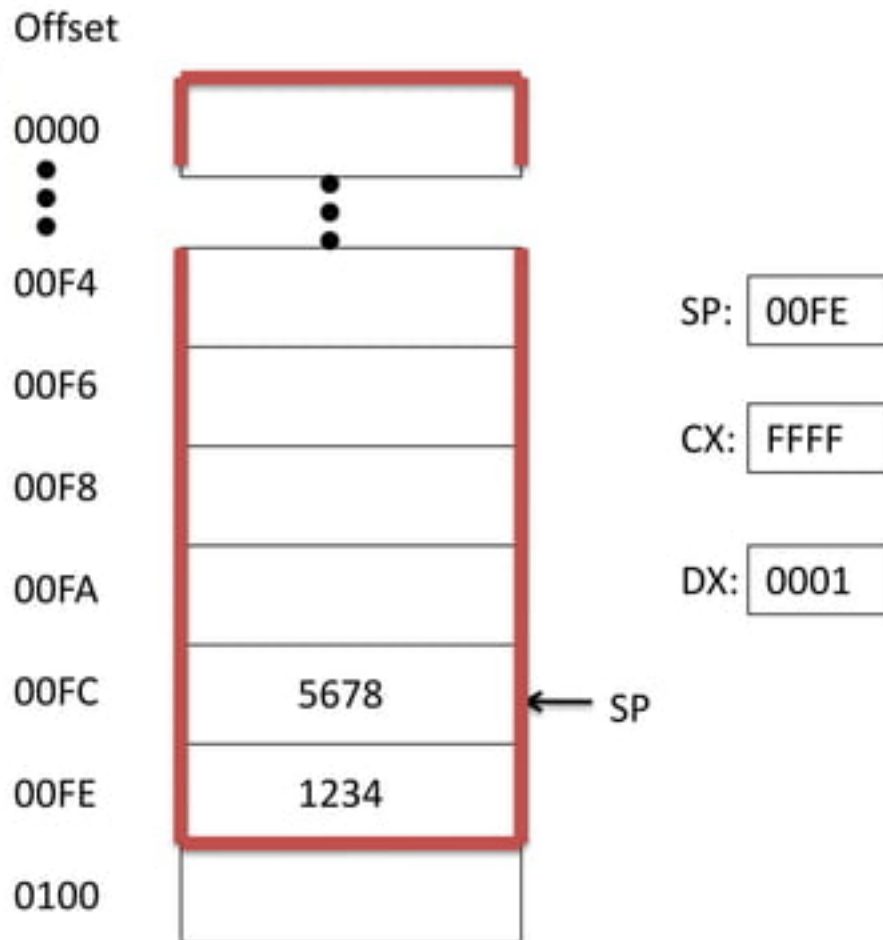
How Words Are Added To Stack



(Beyond the end of the stack)

How Words Are Added To Stack

**AFTER
POP CX**



(Beyond the end of the stack)

How Words Are Added To Stack

**AFTER
POP DX**

Offset

0000

⋮

00F4

00F6

00F8

00FA

00FC

00FE

0100



SP: 0100

CX: FFFF

DX: 0001

← SP

(Beyond the end of the stack)

FLAGS Register and Stack

- **PUSHF**
 - pushes (copies) the contents of the **FLAGS** register onto the stack. It has no operands
- **POPF**
 - pops (copies) the contents of the top word in the stack to the **FLAGS** register. It has no operands
- **NOTES:**
 - *PUSH, POP, and PUSHF do not affect the flags !!*
 - *POPF could theoretically change all the flags because it resets the **FLAGS REGISTER** to some original value that you have previously saved with the **PUSHF** instruction*

Example: Fill up the trace table given below.

[illegible]

Example:

- AX = 3245H
- BX = 1234H
- CX = ABCDH
- SP = FEH

PUSH AX

PUSH CX

POP BX

AX = ?

BX = ?

CX = ?

SP = ?

- AX = 3245H
- BX = 1234H
- CX = ABCDH
- SP = FEH

PUSH BX

PUSH CX

POP BX

POP AX

PUSH CX

PUSH BX

POP CX

PUSH AX

POP BX

AX = ?

BX = ?

CX = ?

SP = ?

- AX = 3245H
- BX = 1234H
- CX = ABCDH
- SP = FEH

PUSH BX

PUSHF

POPF

PUSH CX

POP BX

POP AX

PUSH CX

PUSH BX

POP CX

PUSH AX

POP BX

AX = ?

BX = ?

CX = ?

SP = ?

Important Notes

- *Not only can the programmer use the stack but DOS can and also does use the stack*
- *In fact **DOS** uses the stack every time the user executes an **INT 21h** function*
- *Because of the "last-in first-out" nature of the stack, the order that items are removed from the stack is the reverse of the order in which they are placed on the stack*

Example Program

- The following code allows a user to input a string consisting of 10 characters and then displays the 10 characters in reverse order on the screen

```
TITLE  DISPLAY THE 10 CHARACTERS IN REVERSE ORDER
.MODEL  SMALL
.STACK  100H
.DATA
CR      EQU    0DH
LF      EQU    0AH
MESSAGE DB    CR,LF,'PLEASE TYPE ANY 10 '
        DB    ' CHARACTERS',CR,LF,'$'
REVERSE DB    CR,LF,'THE CHARACTERS IN REVERSE'
        DB    ' ARE:',CR,LF,'$'
```

.CODE

MAIN PROC

;-----INITIALIZE DATA SEGMENT REGISTER

MOV AX,@DATA

MOV DS,AX

;----- SOUND BELL AND PRINT A MESSAGE FOR

INPUT

MOV AH,2

MOV DL,07H

INT 21H

MOV AH,9

LEA DX,MESSAGE

INT 21H

;-----ACCEPT CHARACTERS

MOV CX,10

MOV AH,1

**Example
Program
(cont'd)**

**Example
Program
(cont'd)**

READ:

INT 21H

PUSH AX ;CAN'T PUSH AL SO PUSH AX!

LOOP READ

;-----PRINT REVERSE MESSAGE

MOV AH,9

LEA DX,REVERSE

INT 21H

;-----PREPARE TO PRINT IN REVERSE

MOV CX,10

MOV AH,2

**Example
Program
(cont'd)**

DISP:

POP DX

INT 21H

LOOP DISP

;-----RETURN TO DOS

MOV DL,CR

INT 21h

MOV DL,LF

INT 21h

MOV AH,4CH

INT 21H

MAIN ENDP

END MAIN

Terminology of Procedures

- Top-down program design
 - Decompose the original problem into a series of **subproblems** that are easier to solve than the original problem
- **Subproblems** in assembler language can be structured as a collection of **procedures**
- **Main procedure** contains the entry point to the program and can call one of the other procedures using a **CALL** statement
- It is possible for a called **sub-procedure** to call other **procedures**
- In **AL**, it is also possible for a called sub-procedure to call itself (**recursion**)!

Terminology of Procedures (cont'd)

- When a procedure calls another procedure, control transfers to the called procedure
- When the instructions in a called procedure have been executed, the called procedure usually returns control to the calling procedure at the next sequential instruction after the CALL statement
- In high level languages, mechanism for call and return are hidden from programmer

Procedure Declaration

User defined
name

```
name PROC type  
; body of procedure  
RET  
name ENDP
```

Optional
operand

Causes the
control to
transfer back to
the calling
procedure

Every procedure should
have a **RET** coded
somewhere within the
procedure - usually the last
instruction in a procedure

Near

Statement calls
the procedure in
the same segment

Far

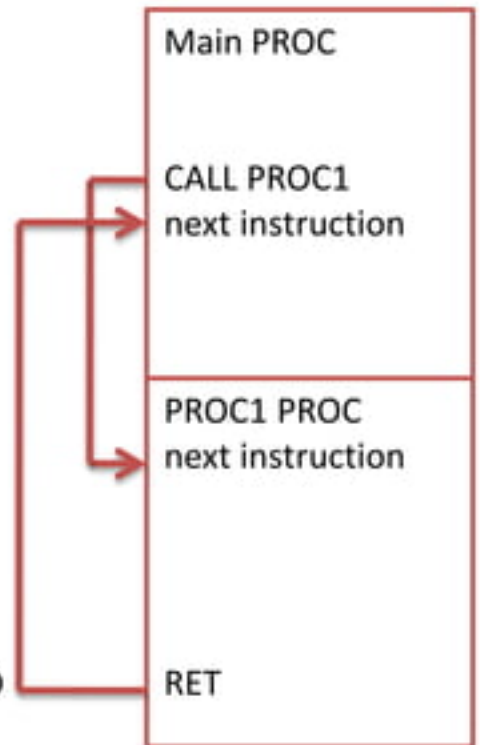
Statement calls the
procedure in the
different segment

PROC Instruction

- **PROC** instruction establishes a procedure
- Procedure declaration syntax:

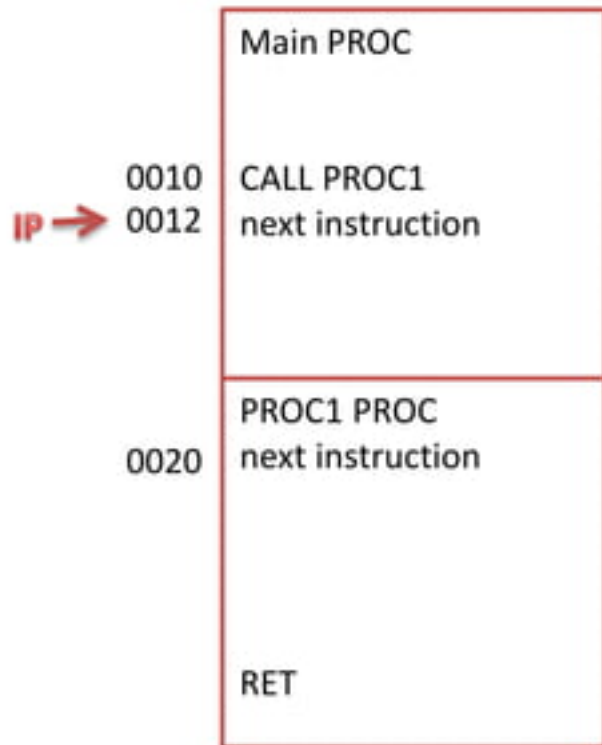
```
name          PROC
; body of the procedure
              RET
name  ENDP
```

- **name** is a user-defined variable.
- **RET** instruction causes control to transfer back to the calling Procedure.
- Every procedure should have a **RET** coded somewhere within the procedure - usually the last instruction in a procedure

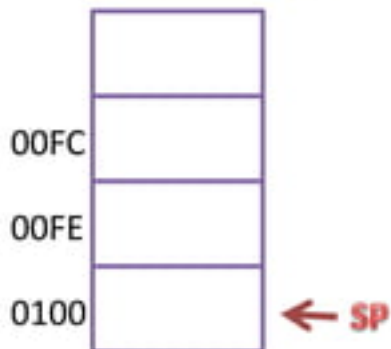


Before Call

CODE SEGMENT



STACK SEGMENT



COMMUNICATION BETWEEN PROCEDURES

- Programmers must devise a way to communicate between procedures –
- there are no parameter lists !!!
- Typically in assembler language, procedures often pass data to each other through registers

Procedures Documentation

- Procedures should be well-documented
 - Describe what the procedure does
 - Indicate how it receives its input from the calling program
 - Indicate how it delivers the results to the calling program
 - Indicate the names of any other procedures that this procedure calls

```
; Describe what the procedure does  
; input:  Indicate how it receives its input  
          from the calling program  
  
; output: Indicate how it delivers the results  
          to the calling program  
  
; uses:   Indicate the names of any other  
          procedures that this procedure calls
```

Procedures (cont'd)

- A procedure usually begins by **PUSHing** (saving) the current contents of all of the registers on the stack.
- A procedure usually ends by **POPing** the stack contents back into the registers before returning to the **CALLing** procedure
- When writing a procedure, **do NOT PUSH or POP any registers in which you intend to return output!!**

CALL Instruction

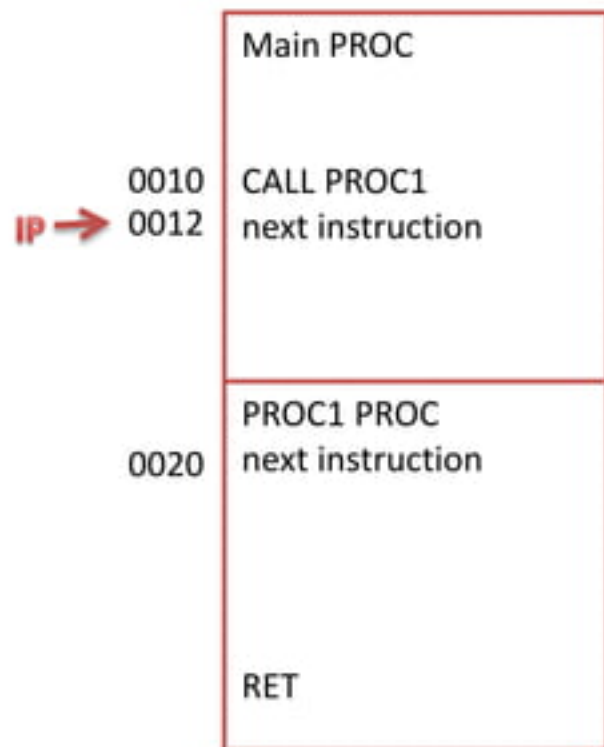
- A **CALL** instruction invokes a procedure
- SYNTAX: **CALL name** (direct CALL)
where **name** is the name of a procedure.
- Executing a **CALL** instruction causes the following to happen:
 - The return address of the **CALLing** program which is in the **IP** register is pushed (saved) on the **STACK**. This saved address is the offset of the next sequential instruction after the **CALL** statement (**CS:IP**)
 - The **IP** then gets the offset address of the first instruction in the procedure

RET Instruction

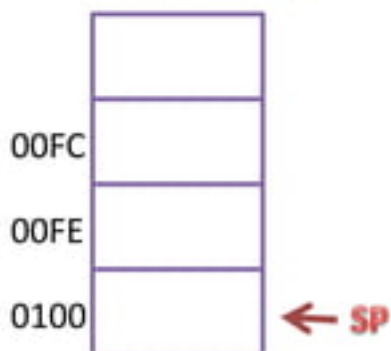
- **RET** statement cause the stack to be popped into **IP**. Procedures typically end with a **RET** statement.
- Syntax: **RET**
- Once the **RET** is executed, **CS:IP** now contains the segment offset of the return address and control returns to the calling program
- In order for the return address to be accessible, each procedure must ensure that the return address is at the top of the stack when the **RET** instruction is executed.

Before Call

CODE SEGMENT

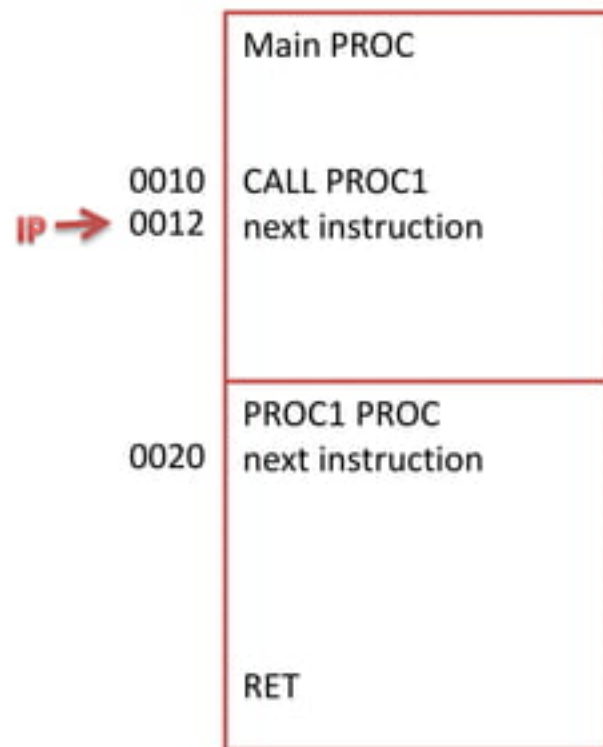


STACK SEGMENT



AFTER Call

CODE SEGMENT

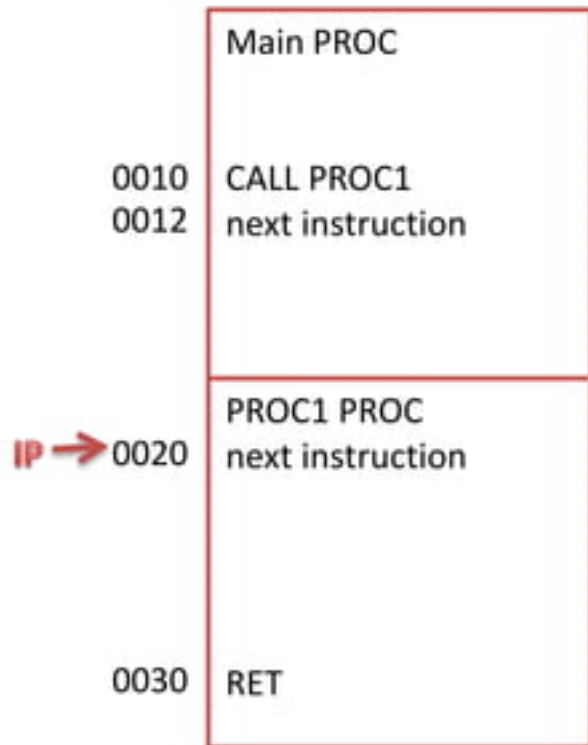


STACK SEGMENT



Before RET

CODE SEGMENT

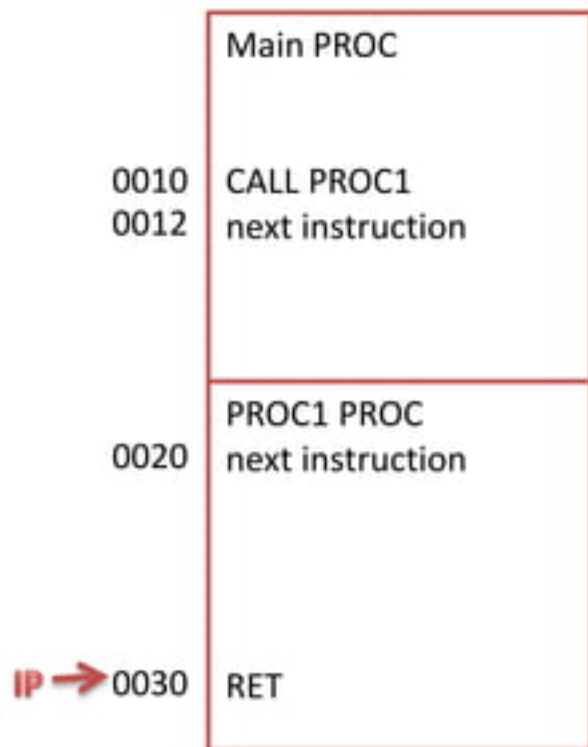


STACK SEGMENT



After RET

CODE SEGMENT



STACK SEGMENT



Typical Layout of a Program Containing Procedures

```
TITLE    A PROGRAM THAT CONTAINS SEVERAL PROCEDURES
.MODEL   SMALL
.STACK   100H
.DATA
;*** define data elements here      ****
.CODE
MAIN     PROC
; INITIALIZE DATA SEGMENT REGISTER
        MOV     AX,@DATA
        MOV     DS,AX
;*** code any necessary statements here      ***
;*** get ready to call procedure ABC by      ***
;*** first moving input values to appropriate registers ***
        CALL    ABC
```

Typical Layout (cont'd)

```
*** code any necessary statements here ***
*** get ready to call procedure DEF by ***
*** first moving input values to appropriate registers ***
        CALL    DEF
*** code any necessary statements here ***
*** get ready to call procedure GHI by ***
*** first moving input values to appropriate registers ***
        CALL    GHI
*** code any necessary statements here ***
; RETURN TO DOS
        MOV     AH,4CH
        INT     21H
MAIN    ENDP
```

Typical Layout (cont'd)

```
ABC      PROC
        PUSH      .....    ; as many PUSHes as you need

        POP       .....    ; POPs in reverse order of PUSHes
        RET
ABC      ENDP
;
DEF      PROC
        PUSH      .....    ; as many PUSHes as you need

        POP       .....    ; POPs in reverse order of PUSHes
        RET
DEF      ENDP
;
```


Typical Layout (cont'd)

```
GHI      PROC
        PUSH     .....    ; as many PUSHes as you need

        POP      .....    ; POPs in reverse order of PUSHes
        RET
GHI      ENDP
        END      MAIN
```

Example Program

- Now let's study the Multiplication Procedure

Product = 0

Repeat

IF lsb of B is 1

then

product = product + A

end_if

shift left A

shift right B

Until B = 0

Example Program (cont'd)

```
TITLE  MULTIPLICATION BY ADDING AND SHIFTING (8 BITS BY 8 BITS)
.MODEL  SMALL
.STACK  100H
.CODE
MAIN  PROC
;-----> INITIALIZE AX AND BX
    MOV  AX,13      ;SOME ARBITRARY VALUE
    MOV  BX,10      ;SOME ARBITRARY VALUE
;-----> INVOKE PROCEDURE
    CALL     MULTIPLY
;-----> DX NOW CONTAINS PRODUCT
; RETURN TO DOS
    MOV  AH,4CH
    INT  21H
MAIN  ENDP
```

Example (cont'd)

MULTIPLY PROC

```
;-----> THIS PROCEDURE MULTIPLIES THE  
;-----> VALUE IN AX BY THE VALUE IN BX  
;-----> RETURNING THE PRODUCT IN DX.  
;-----> VALUES IN AX AND BX ARE LIMITED ;-----> TO 00 - FFh.  
;-----> IT USES SHIFTING AND ADDITION  
;-----> TO ACCOMPLISH THE MULTIPLICATION
```

```
    PUSH  AX      ; DON'T DESTROY AX  
    PUSH  BX      ; DON'T DESTROY BX  
    XOR   DX,DX    ; CLEAR DX WHERE PRODUCT WILL BE
```

REPEAT:

```
    TEST  BX,1     ; IS LSB = 0?  
    JZ    END_IF  
    ADD   DX,AX     ; PRODUCT = PRODUCT + A
```

END_IF:

```
    SHL   AX,1  
    SHR   BX,1  
    JNZ   REPEAT  
    POP   BX  
    POP   AX  
    RET
```

MULTIPLY ENDP

```
    END    MAIN
```

Important Notes on Stack

- PUSHES AND POPS are often used to save data temporarily on the program stack. They are also used implicitly each time a CALL and a RETURN sequence is executed.
- Remember that the SP is decremented **BEFORE** placing a word on the stack at PUSH time but it is incremented **AFTER** removing a word from the stack at POP time.
- If, for some reason, you want a copy of the FLAGS register in BX, you can accomplish this by:
PUSHF
POP BX
- Stack allows you to save the contents of a register, use the register for something else temporarily, and then restore the register to its original value.

Notes on Stack (cont'd)

- Pushing and popping the contents of registers is preferable to storing their contents as variables in the DATA segment
- Reasons:
 - Using the stack is more economical. Instead of allocating data space, by pushing and popping data into the stack, you use space as you need it and release it when you no longer need it.
 - Since the 8088/8086 allows recursion, if a routine called itself and saved the contents of a register to a data location each time it was invoked, it would be overwriting the previous contents of that location with each recursion!
 - Using a stack instead of a data location makes code more portable. Once you have written a good routine, you may choose to incorporate that routine into several different programs. Or if you are working with a programming team piecing smaller subroutines into a one larger main routine, subroutines that do their work without referencing particular data locations are more easily patched into main programs than subroutines that do reference particular data locations. Therefore, should not refer to ANY variable data names in ANY procedure that you write!!!'

Notes on Stack (cont'd)

- Care must be taken not to corrupt the STACK because not only does it save values for the programmer but it also saves values for the CPU. These values get interwoven on the stack. If the SP becomes confused, the CPU could get lost throwing your computer into a system error!!
- Always check to see that the PUSHES and POPS in a program are paired --- or --- at least that each of them is balanced by program code that restores the stack pointer to its proper value.
- If you find yourself in the middle of a system error, more than likely you look for a problem in the way you implemented the stack.

Example Procedure on Safe Use of Stack

- A procedure to display a carriage return and a line feed:

```
CRLF  PROC
      PUSH        AX        ; Save AX
      PUSH        DX        ; Save DX
      MOV         AH,2      ; Display a Carriage Return
      MOV         DL,0Dh
      INT         21h
      MOV         DL,0Ah    ; Display a Line Feed
      INT         21h
      POP         DX        ; Restore DX
      POP         AX        ; Restore AX
      RET
CRLF  ENDP
```


Example (cont'd)

- This procedure can be called by the programmer at any time regardless of what is in his/her AX or DX registers. As far as the programmer is concerned, all you know is that this procedure issues the CR/LF sequence to the console and all of your registers will be unchanged when the procedure has finished executing!