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Assignment 04

Exercise 1.

1. (Cryptanalysis) Though encryption is primarily designed to preserve confidentiality and integrity of data, the mechanism itself is vulnerable to brute force (statistical analysis). In other words, the more we see the encrypted data, the easier we can hack it. In this exercise, you are asked to crack the following cipher text. Please provide the decrypted result and explain your strategy in decrypting this text.

Cipher text

PRCSOFQX FP QDR AFOPQ CZSPR LA JFPALOQSKR. QDFP FP ZK LIU BROJZK MOLTROE.

```
In [44]: cipher_text = "PRCSOFQX FP QDR AFOPQ CZSPR LA JFPALOQSKR. QDFP FP ZK LIU BROJZK MOLTROE."
```

a. Count the frequency of letters. List the top three most frequent characters.

```
In [45]: d = {
         c: cipher_text.count(c) for c in set(cipher_text) if c.isalpha()
}
d = dict(sorted(d.items(), key=lambda item: item[1], reverse=True)[:3])
print(d)
{'P': 7, '0': 6, 'F': 6}
```

Answer (a.): The top three most frequent characters are 'F' (6), 'P' (5), and 'R' (4)

b. Knowing that this is English, what are commonly used three-letter words and two-letter words. Does the knowledge give you a hint on cracking the given text?

```
In [46]: two_letter_words = [
                'IS',
                'TO'
                'AS'
                'AT'
                'AN',
                'AM',
                'BE',
               'HE',
               'SO',
               'HI'
          ]
          three_letter_words = [
               'AND',
               'THE',
                'Y0U'
               'ARE',
               'BUT',
               'FOR',
               'NOT',
                'NOW',
                'ONE',
                'OUT',
```

```
'SHE',
'HIM',
'HIS',
'CAN',
'HAD',
'HAS'
'HOW'
'WAS'
'WHO',
'ALL',
'ANY',
'NEW',
'OLD',
'BIG',
'BAD'
'DAY',
'GET',
'NEW',
'MAN'
'CAT',
'DOG',
'SUN'
'HAT'
```

Answer (b.): Commonly used

```
two-letter words: 'IS', 'TO', 'OF', 'IN', 'IF', 'ON', 'UP', 'NO', 'DO', 'OR', 'AS', 'AT', 'AN', 'AM', 'BE', 'HE', 'SO', 'HI'

three-letter words: 'AND', 'THE', 'YOU', 'ARE', 'BUT', 'FOR', 'NOT', 'NOW', 'ONE', 'OUT', 'SHE', 'HIM', 'HIS', 'CAN', 'HAD', 'HAS', 'HOW', 'WAS', 'WHO', 'ALL', 'ANY', 'NEW', 'OLD', 'BIG', 'BAD', 'DAY', 'GET', 'NEW', 'MAN', 'CAT', 'DOG', 'SUN', 'HAT'
```

Yes, the knowledge gives a hint on cracking the given text.

c. Cracking the given text. Measure the time that you have taken to crack this message.

Answer (c.): SECURITY IS THE FIRST CAUSE OF MISFORTUNE. THIS IS AN OLD GERMAN PROVERB.

This take me about 11 hours to crack the code.

d. Create a simple python program for cracking the Caesar cipher text using brute force attack. Explain the design and demonstrate your software. (You may use an English dictionary for validating results.)

Answer (d.): The following Python Code Below. Explain in comments.

```
In [47]: # Extract set of two-letter and three-letter tokens from the cipher text
         two_letter_tokens = set([w for w in cipher_text.split() if len(w) == 2])
         three_letter_tokens = set([w for w in cipher_text.split() if len(w) == 3])
         print(two_letter_tokens)
         print(three_letter_tokens)
        {'LA', 'ZK', 'FP'}
        {'QDR', 'LIU'}
In [48]: # Create an underscore version of the cipher text for pattern matching
         underscore_string = ''.join(['_' if c.isalpha() else c for c in cipher_text])
         print(underscore_string)
         # function to change the mapping dictionary
         def change_string(d, old_string, new_string):
             if len(set(old_string)) != len(set(new_string)):
             new_dict = {o: n for o, n in zip(old_string, new_string)}
             for k in new_dict:
                 if k in d and d[k] != new_dict[k]:
                     return None
             for k in new dict:
                 d[k] = new_dict[k]
             return d
```

.

Because "FP" appears in the second position of both sentences, it's likely to be "IS" or "AM".

Since "AM" can only be used with "I", "FP" is likely to be "IS".

```
In [49]: # Try to map "FP" -> "IS" and "QDR" -> each three-letter word
         file data = ''
         cnt = 0
         for i in range(len(three_letter_words)):
             text = underscore_string
             d = \{\}
             d = change_string(d, 'FP', 'IS')
             if not d:
                 continue
             d = change_string(d, 'QDR', three_letter_words[i])
             if not d:
                 continue
             for it in range(len(text)):
                 if cipher_text[it] in d:
                     text = text[:it] + d[cipher_text[it]] + text[it+1:]
             cnt += 1
             file_data += text + '\n'
         with open('output.txt', 'w') as f:
             lines = file data.splitlines()
             for i, line in enumerate(lines[:5]):
                 print(f"{i}\t{line}")
             print('...')
             for i, line in enumerate(lines[-5:]):
                 print(f''\{i + (len(lines) - 5) + 1\}\setminus t\{line\}'')
             f.write(file_data)
               SD__IA_ IS AND _I_SA ___SD __ _IS___A__D. ANIS IS __ __ _D_____
               SE__IT_ IS THE _I_ST ___SE __ _IS__T_E. THIS IS __ ___E___
        1
               SU__IY_ IS YOU _I_SY __SU __ IS__Y_U. YOIS IS __ __U
        2
        3
                SE__IA_ IS ARE _I_SA ___SE __ _IS___A__E. ARIS IS __ ____E_
        4
               ST__IB_ IS BUT _I_SB ___ST __ _IS___B__T. BUIS IS __ _T__T
        . . .
        28
                SN__IM_ IS MAN _I_SM ___SN __ _IS___M__N. MAIS IS __ ___ _N_
        29
                ST__IC_ IS CAT _I_SC ___ST __ _IS___C__T. CAIS IS __ ___ _T____
               SG__ID_ IS DOG _I_SD ___SG __ _IS___D_G. DOIS IS __ ___ _G___ ___G__.
        30
                                      ___SN ___IS__S_N. SUIS IS __ ___N_
                SN IS_ IS SUN _I_SS _
        31
               ST___IH_ IS HAT _I_SH ___ST __ _IS___H__T. HAIS IS __ ___ _T_
In [50]: | # Filter out grammatically incorrect results
         cipher_results = [
             # "SD___IA_ IS AND _I_SA ___SD __ _IS___A__D. ANIS IS __ ___ _D_____
             "SE__IT_ IS THE _I_ST ___SE __ _IS__T_E. THIS IS __ ___E___E_.",
                                                                   __ __U_
                                     ___SU __ _IS___Y__U. YOIS IS
             # "SU___IY_ IS YOU _I_SY _
             # "SE__IA_ IS ARE _I_SA __SE __ _IS__A_E. ARIS IS __ __E_
             # "ST___IB_ IS BUT _I_SB ___ST __ _IS___B__T. BUIS IS __ _T_
                   __IM_ IS MAN _I_SM ___SN __ _IS___M__N. MAIS IS __ ___ _N________N__."
             # "ST___IC_ IS CAT _I_SC ___ST __ _IS___C__T. CAIS IS __ ___ _T___
             # "SG__ID_ IS DOG _I_SD ___SG __ _IS___D_G. DOIS IS __ ___ _G___."
             # "SN__IS_ IS_SUN _I_SS ___SN __ _IS___S__N. SUIS IS __ ___ _N_
             # "ST__IH_ IS HAT _I_SH __ST __ IS__H__T. HAIS IS __ __T
```

Try the other 2 possible two_letter_words and 1 possible three_letter_word.

```
'BE',
     'HE',
     'SO',
     'HI'
 two_letter_words_after_IS = [
     'TO',
     # 'OF',
     'IN',
     # 'IF',
     'ON',
     'UP',
     'NO',
     # 'DO'
     # 'OR',
     # 'AS',
     'AT',
     'AN',
     # 'AM',
     # 'BE',
     # 'HE',
     'SO',
     # 'HI'
 file_data = ''
 cnt = 0
 for i in range(len(two_letter_words)):
     for j in range(len(three_letter_words)):
         for k in range(len(two_letter_words_after_IS)):
             if two_letter_words[i] == two_letter_words_after_IS[k]:
                 continue
             text = underscore_string
             for last_result in cipher_results:
                 d = \{\}
                 for it in range(len(cipher_text)):
                      if last_result[it].isalpha():
                          d[cipher_text[it]] = last_result[it]
                 d = change_string(d, 'LA', two_letter_words[i])
                 if not d:
                      continue
                 d = change_string(d, 'LIU', three_letter_words[j])
                 if not d:
                      continue
                 d = change_string(d, 'ZK', two_letter_words_after_IS[k])
                 if not d:
                      continue
                 for it in range(len(text)):
                     if cipher_text[it] in d:
                          text = text[:it] + d[cipher_text[it]] + text[it+1:]
                 cnt += 1
                 file_data += text + '\n'
 with open('output.txt', 'w') as f:
     lines = file_data.splitlines()
     for i, line in enumerate(lines[:5]):
         print(f"{i}\t{line}")
     print('...')
     for i, line in enumerate(lines[-5:]):
         print(f''\{i + (len(lines) - 5) + 1\}\setminus t\{line\}'')
     f.write(file_data)
            _IT_ IS THE OI_ST _I_SE TO _ISOT_T_NE. THIS IS IN THE _E__IN __T_E__.
0
          IT IS THE OI ST O SE TO ISOT T NE. THIS IS ON THE E ON TE.
2
        SE___IT_ IS THE OI_ST _U_SE TO _ISOT_T_PE. THIS IS UP THE _E__UP __T_E__.
3
        SE___IT_ IS THE OI_ST _N_SE TO _ISOT_T_OE. THIS IS NO THE _E__NO _
        SE___IT_ IS THE OI_ST _A_SE TO _ISOT_T_TE. THIS IS AT THE _E__AT __T_E__.
4
340
        SE___IT_ IS THE II_ST _U_SE HI _ISIH_T_PE. THIS IS UP HAT _E__UP __H_E__.
        SE___IT_ IS THE II_ST _N_SE HI _ISIH_T_OE. THIS IS NO HAT _E__NO __H_E__.
341
342
        SE___IT_ IS THE II_ST _A_SE HI _ISIH_T_TE. THIS IS AT HAT _E__AT __H_E__.
343
        SE___IT_ IS THE II_ST _A_SE HI _ISIH_T_NE. THIS IS AN HAT _E__AN __H_E__.
        SE___IT_ IS THE II_ST _S_SE HI _ISIH_T_OE. THIS IS SO HAT _E__SO __H_E__.
344
```

Filter out the impossible fourth word.

```
In [52]: with open('output.txt', 'r') as f:
    lines = f.readlines()
    s = set()
    for line in lines:
        line = line.split()
        s.add(line[3])
```

```
for line in s:
                  print(line)
             print(len(s))
        SI_ST
        TI_ST
        0I_ST
        NI_ST
        MI_ST
        EI_ST
        FI_ST
        RI_ST
        II_ST
In [53]: |possible_fourth_word = [
              'TI_ST', # TIEST
              'MI_ST', # MIDST
              'FI_ST', # FIRST
         with open('output.txt', 'r') as f:
             lines = f.readlines()
             final_results = []
             for line in lines:
                  line = line.split()
                  if line[3] in possible_fourth_word:
                      final_results.append(' '.join(line))
             # for line in final_results:
                   print(line)
             # print(len(final_results))
             with open('output2.txt', 'w') as f2:
                  lines = final_results
                  for i, line in enumerate(lines[:5]):
                      print(f"{i}\t{line}")
                  print('...')
                  for i, line in enumerate(lines[-5:]):
                      print(f''\{i + (len(lines) - 5) + 1\}\setminus t\{line\}'')
                  f2.write('\n'.join(final_results))
        0
                SE___IT_ IS THE FI_ST _T_SE OF _ISFO_T_OE. THIS IS TO ONE _E__TO __O_E__.
        1
                SE___IT_ IS THE FI_ST _I_SE OF _ISFO_T_NE. THIS IS IN ONE _E__IN __O_E__.
                SE__IT_ IS THE FI_ST _0_SE OF _ISFO_T_NE. THIS IS ON ONE _E__ON __O_E__.
        2
        3
                 SE__IT_ IS THE FI_ST _U_SE OF _ISFO_T_PE. THIS IS UP ONE _E__UP __O_E__.
        4
                 SE___IT_ IS THE FI_ST _N_SE OF _ISFO_T_OE. THIS IS NO ONE _E__NO __O_E__.
                 SE___IT_ IS THE MI_ST _U_SE AM _ISMA_T_PE. THIS IS UP ANY _E__UP __A_E__.
        65
        66
                 SE___IT_ IS THE MI_ST _N_SE AM _ISMA_T_OE. THIS IS NO ANY _E__NO __A_E__.
                SE___IT_ IS THE MI_ST _A_SE AM _ISMA_T_TE. THIS IS AT ANY _E__AT __A_E__.
        67
                 SE___IT_ IS THE MI_ST _A_SE AM _ISMA_T_NE. THIS IS AN ANY _E__AN __A_E__.
        68
                 SE___IT_ IS THE MI_ST _S_SE AM _ISMA_T_OE. THIS IS SO ANY _E__SO __A_E__.
In [54]: with open('output2.txt', 'r') as f:
             lines = f.readlines()
             s = set()
             for line in lines:
                  line = line.split()
                  s.add(' '.join(line[7:11]))
             # for line in s:
             #
                   print(line)
             # print(len(s))
             with open('output3.txt', 'w') as f2:
                  lines = list(s)
                  for i, line in enumerate(lines[:5]):
                      print(f"{i}\t{line}")
                  print('...')
                  for i, line in enumerate(lines[-5:]):
                      print(f''\{i + (len(lines) - 5) + 1\}\setminus t\{line\}'')
                  f2.write('\n'.join(s))
        0
                THIS IS IN OUT
        1
                THIS IS ON OLD
                THIS IS ON AND
        2
        3
                THIS IS UP ANY
        4
                THIS IS AN OLD
                THIS IS AN AND
        44
        45
                THIS IS UP ONE
                THIS IS UP OLD
        46
        47
                THIS IS AN ARE
        48
                THIS IS AT ARE
```

file:///Users/chayaninkongsareekul/Desktop/Comp_Security/act4/Comp_Sec_Assignment04.html

```
In [55]:
         possible_phrase = [
             # "THIS IS UP ANY",
             "THIS IS IN ONE",
             "THIS IS TO ONE",
             # "THIS IS UP OLD",
             "THIS IS AT OLD",
             # "THIS IS NO OUT",
             "THIS IS NO ONE",
             # "THIS IS NO ANY",
             "THIS IS AT ANY",
             # "THIS IS IN OLD",
             # "THIS IS IN ARE",
             # "THIS IS ON OUT"
             # "THIS IS ON AND"
              # "THIS IS AN OUT",
             # "THIS IS AT ARE",
             # "THIS IS ON ARE",
             # "THIS IS UP ONE",
             # "THIS IS SO AND",
             # "THIS IS SO ANY",
             "THIS IS ON ONE",
             # "THIS IS NO ARE",
             # "THIS IS ON OLD",
             # "THIS IS SO OUT",
             # "THIS IS AN ARE",
             "THIS IS IN ANY",
             "THIS IS AN OLD",
             "THIS IS UP AND",
             # "THIS IS TO AND",
             # "THIS IS IN OUT",
             # "THIS IS SO ONE",
             "THIS IS AT ONE",
             # "THIS IS TO OLD",
             # "THIS IS UP OUT",
             # "THIS IS NO OLD"
             # "THIS IS NO AND"
             # "THIS IS UP ARE",
             "THIS IS ON ANY",
             # "THIS IS AN ONE"
             # "THIS IS AN ANY",
             # "THIS IS AN AND",
             # "THIS IS TO OUT",
             # "THIS IS AT AND"
             "THIS IS TO ANY",
             # "THIS IS TO ARE",
             # "THIS IS IN AND",
             "THIS IS SO OLD",
             # "THIS IS AT OUT",
             # "THIS IS SO ARE"
```

Try search every 7th word in the dictionary, only MISFORTUNE is the only word that is possible.

```
In [56]: with open('output2.txt', 'r') as f:
             lines = f.readlines()
             final_results = []
             for line in lines:
                 line = line.split()
                 phrase = ' '.join(line[7:11])
                 if phrase in possible_phrase:
                     final_results.append(' '.join(line))
             for line in final_results:
                 print(line)
             print(len(final_results))
             with open('output4.txt', 'w') as f2:
                 f2.write('\n'.join(final_results))
             s2 = set()
             for line in final_results:
                 s2.add(line.split()[6])
             for line in s2:
                 print(line)
```

```
SE___IT_ IS THE FI_ST _T_SE OF _ISFO_T_OE. THIS IS TO ONE _E__TO __O_E__.
        SE___IT_ IS THE FI_ST _I_SE OF _ISFO_T_NE. THIS IS IN ONE _E__IN __O_E__.
        SE__IT_ IS THE FI_ST _0_SE OF _ISFO_T_NE. THIS IS ON ONE _E_ON __O_E_.
        SE___IT_ IS THE FI_ST _N_SE OF _ISFO_T_OE. THIS IS NO ONE _E__NO __O_E__.
        SE___IT_ IS THE FI_ST _A_SE OF _ISFO_T_TE. THIS IS AT ONE _E__AT __O_E__.
        SE___IT_ IS THE FI_ST _A_SE OF _ISFO_T_TE. THIS IS AT OLD _E__AT __O_E__.
        SE__IT_ IS THE FI_ST _A_SE OF _ISFO_T_NE. THIS IS AN OLD _E__AN __O_E_
        SE__IT_ IS THE FI_ST _S_SE OF _ISFO_T_OE. THIS IS SO OLD _E__SO __O_E__.
        SE__IT_ IS THE TI_ST_U_SE AT _ISTA_T_PE. THIS IS UP AND _E_UP __A_E__.
        SE___IT_ IS THE TI_ST _T_SE AT _ISTA_T_OE. THIS IS TO ANY _E__TO __A_E__.
        SE__IT_ IS THE TI_ST _I_SE AT _ISTA_T_NE. THIS IS IN ANY _E__IN __A_E__.
        SE___IT_ IS THE TI_ST _O_SE AT _ISTA_T_NE. THIS IS ON ANY _E__ON __A_E__.
        SE___IT_ IS THE MI_ST _U_SE AM _ISMA_T_PE. THIS IS UP AND _E__UP __A_E__.
        SE___IT_ IS THE MI_ST _T_SE AM _ISMA_T_OE. THIS IS TO ANY _E__TO __A_E__.
        SE__IT_ IS THE MI_ST _I_SE AM _ISMA_T_NE. THIS IS IN ANY _E__IN __A_E__.
        SE___IT_ IS THE MI_ST _O_SE AM _ISMA_T_NE. THIS IS ON ANY _E__ON __A_E__.
        SE__IT_ IS THE MI_ST _A_SE AM _ISMA_T_TE. THIS IS AT ANY _E_AT __A_E__.
        17
        _ISMA_T_NE.
        _ISTA_T_PE.
        _ISMA_T_PE.
        _ISFO_T_OE.
        _ISFO_T_TE.
        _ISFO_T_NE.
        _ISTA_T_0E.
        _ISTA_T_NE.
        _ISMA_T_TE.
        _ISMA_T_0E.
In [57]: possible_7th_word = "_ISF0_T_NE." #"MISFORTUNE."
         # "JFPALOOSKR."
         d = {
             'J': 'M',
             '0': 'R',
             'S': 'U',
         with open('output4.txt', 'r') as f:
             lines = f.readlines()
             final_results = []
             for line in lines:
                 line = line.split()
                 if line[6] == possible_7th_word:
                     string = ' '.join(line)
                     for i in range(len(cipher_text)):
                          if cipher_text[i] in d:
                              string = string[:i] + d[cipher_text[i]] + string[i+1:]
                     final_results.append(string)
             for line in final_results:
                 print(line)
             print(len(final_results))
             with open('output5.txt', 'w') as f2:
                 f2.write('\n'.join(final_results))
        SE_URIT_ IS THE FIRST _IUSE OF MISFORTUNE. THIS IS IN ONE _ERMIN _RO_ER_.
        SE_URIT_ IS THE FIRST _OUSE OF MISFORTUNE. THIS IS ON ONE _ERMON _RO_ER_.
        SE_URIT_ IS THE FIRST _AUSE OF MISFORTUNE. THIS IS AN OLD _ERMAN _RO_ER_.
In [58]: # SE_URIT_
         # SECURITY
         # PRCSOFQX
            ERMAN
         # GERMAN
         # BROJZK
         d = {
             'C': 'C',
             'X': 'Y',
             'B': 'G'
         }
         with open('output5.txt', 'r') as f:
             lines = f.readlines()
             final_results = []
             for line in lines:
                 line = line.split()
                 string = ' '.join(line)
                 for i in range(len(cipher_text)):
                     if cipher_text[i] in d:
                          string = string[:i] + d[cipher text[i]] + string[i+1:]
```

```
final_results.append(string)
for line in final_results:
    print(line)
print(len(final_results))
with open('final_output.txt', 'w') as f2:
    f2.write('\n'.join(final_results))

SECURITY IS THE FIRST CIUSE OF MISFORTUNE. THIS IS IN ONE GERMIN _RO_ER_.
SECURITY IS THE FIRST COUSE OF MISFORTUNE. THIS IS ON ONE GERMON _RO_ER_.
SECURITY IS THE FIRST CAUSE OF MISFORTUNE. THIS IS AN OLD GERMAN _RO_ER_.
3
```

Search Google



special-dictionary.com

https://www.special-dictionary.com > german_proverb

Security is the first cause of misfortune. - Special Dictionary

Famous Proverbs / German Proverbs. German Proverb: "Security is the first cause of misfortune." German Proverbs. Proverbs about: Cause, Misfortune, Security ...

Missing: OLD | Show results with: OLD

Ans: SECURITY IS THE FIRST CAUSE OF MISFORTUNE. THIS IS AN OLD GERMAN PROVERB.

Exercise 2. (Cryptanalysis on Symmetric Encryption) Vigenère is a complex version of the Caesar cipher. It is a polyalphabetic substitution.

a. Please review Kasiski examination Explain how it can be used to attack Vigenère.

Answer (a.):

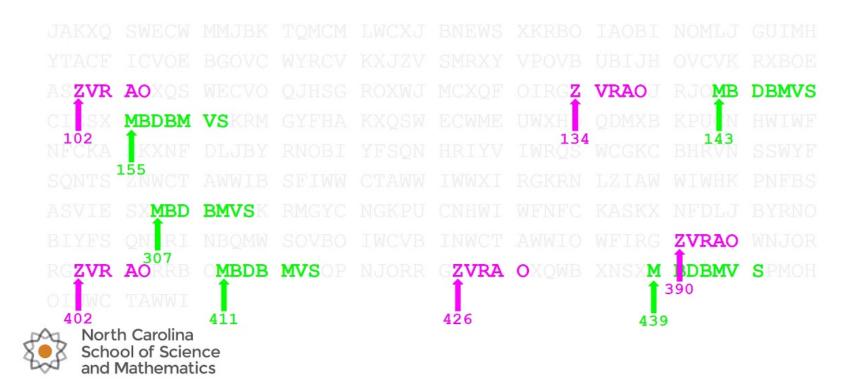
Kasiski Examination ใช้ค้นหาความยาวของ key ของรหัส Vigenère

• เมื่อรู้ความยาวของ key แล้ว เราจะสามารถแยกปัญหาออกเป็นหลาย ๆ Caesar cipher ได้

ขั้นตอนการใช้ Kasiski Examination

- 1. ค้นหาลำดับอักษรซ้ำใน ciphertext
 - เช่น "ABC" ปรากฏที่ตำแหน่ง 5, 15, 25
- 2. คำนวณระยะห่างระหว่างตำแหน่งที่ซ้ำ
 - เช่น 15-5=10, 25-15=10
- 3. หา **ตัวหารร่วมมาก (GCD)** ของระยะห่างเหล่านั้น
 - เช่น GCD(10, 10) = 10
- 4. แบ่ง ciphertext ออกเป็นกลุ่มตามความยาวของ key ที่คำนวณได้
 - เช่น ถ้า GCD = 10, แบ่ง ciphertext ออกเป็น 10 กลุ่ม
- 5. วิเคราะห์ความถี่ของแต่ละกลุ่ม (Frequency Analysis)
 - ใช้การวิเคราะห์ความถี่ (เหมือนโจมตี Caesar cipher) กับแต่ละกลุ่มเพื่อหาตัวอักษรกุญแจในแต่ละตำแหน่ง
 - รวมผลลัพธ์ทั้งหมด → ได้กุญแจเต็ม → ถอดรหัสข้อความ

Example



Exercise 3.

- 3. (Mode in Block Cipher) Block Cipher is designed to have more randomness in a block. However, an individual block still utilizes the same key. Thus, it is recommended to use a cipher mode with an initial vector, chaining or feedback between blocks. This exercise will show you the weakness of Electronic Code Book mode which does not include any initial vector, chaining or feedback.
- a. Find a bitmap image that is larger than 2000x2000 pixels. Note that you may resize any image. To simplify the pattern, we will change it to bitmap (1-bit per pixel) using the portable bitmap format (pbm). In this example, we will use imagemagick for the conversion.



b. c. d. e.

```
In [59]: # Block Cipher Mode Comparison
# Demonstrating the weakness of ECB mode vs other cipher modes

import matplotlib.pyplot as plt
import numpy as np
from PIL import Image
import os

# Use ImageMagick to convert and resize an image to PBM format
!magick image.jpg -resize 2000x2000 org.pbm

# Copy org.pbm to org.x for encryption (removing header)
!cp org.pbm org.x
```

```
# Create different cipher modes
 print("Creating encrypted files with different modes...")
 # ECB mode (Electronic Code Book)
 !openssl enc -aes-256-ecb -in org.x -nosalt -out enc_ecb.x -pass pass:mypassword
 # CBC mode (Cipher Block Chaining)
 !openssl enc -aes-256-cbc -in org.x -nosalt -out enc_cbc.x -pass pass:mypassword
 print("Encryption completed!")
 # Add headers back to create valid PBM files
 def add_pbm_header(input_file, output_file, width=2000, height=2000):
     """Add PBM header to encrypted file"""
     with open(input_file, 'rb') as f:
         data = f.read()
     header = f"P4\n{width} {height}\n".encode('ascii')
     with open(output_file, 'wb') as f:
         f.write(header)
         f.write(data)
 # Add headers to all encrypted files
 add_pbm_header('enc_ecb.x', 'enc_ecb.pbm')
 add_pbm_header('enc_cbc.x', 'enc_cbc.pbm')
 print("Headers added to encrypted files!")
 # Load and display images for comparison
 def load_pbm_image(filename):
     """Load PBM image and convert to numpy array for display"""
     try:
         img = Image.open(filename)
         return np.array(img)
     except Exception as e:
         print(f"Error loading {filename}: {e}")
         return None
 # Load all images
 original = load_pbm_image('org.pbm')
 ecb_encrypted = load_pbm_image('enc_ecb.pbm')
 cbc_encrypted = load_pbm_image('enc_cbc.pbm')
 # Create comparison plot
 fig, axes = plt.subplots(1, 3, figsize=(15, 5))
 fig.suptitle('Original vs ECB vs CBC', fontsize=16)
 # Plot original image if available
 if original is not None:
     axes[0].imshow(original, cmap='gray')
     axes[0].set_title('Original')
     axes[0].axis('off')
 # Plot ECB encrypted
 if ecb_encrypted is not None:
     axes[1].imshow(ecb_encrypted, cmap='gray')
     axes[1].set_title('AES-256-ECB')
     axes[1].axis('off')
 # Plot CBC encrypted
 if cbc encrypted is not None:
     axes[2].imshow(cbc_encrypted, cmap='gray')
     axes[2].set_title('AES-256-CBC')
     axes[2].axis('off')
Creating encrypted files with different modes...
*** WARNING : deprecated key derivation used.
Using -iter or -pbkdf2 would be better.
*** WARNING: deprecated key derivation used.
```

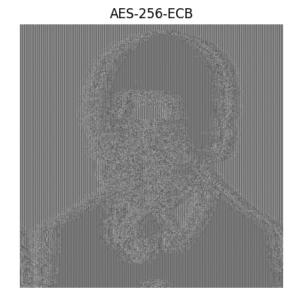
Headers added to encrypted files!

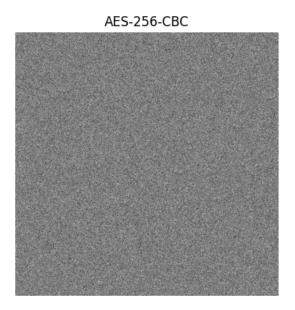
Encryption completed!

Using -iter or -pbkdf2 would be better.

Original vs ECB vs CBC







f. What does the result suggest about the mode of operation in block cipher? Please provide your analysis.

Answer (3.f.):

ECB

- encrypt(plaintext_block[i]) -> ciphertext_block[i]
- Identical plaintext blocks -> identical ciphertext blocks.
- In an image, the original pattern is still visible, like a mosaic.

CBC

- encrypt(plaintext_block[i] XOR ciphertext_block[i 1]) -> ciphertext_block[i]
- Identical plaintext blocks -> different ciphertext blocks (due to chaining).
- Requires an initialization vector (IV) for the first block.
- In an image, the original pattern is destroyed, looks random.

Analysis

ECB mode is insecure for data with structure

• Patterns in plaintext are visible in ciphertext.

CBC mode improves security by randomizing ciphertext

- Chaining ensures each block's encryption depends on all previous blocks. Removes visible patterns.
- Still vulnerable to certain attacks

Conclusion

- The underlying block cipher (AES-256) is strong, but the mode of operation (ECB) is weak for structured data.
- Use CBC or other modes for better security.

Exercise 4. (Encryption Protocol - Digital Signature)

a. Measure the performance of a hash function (sha1), RC4, Blowfish and DSA. Outline your experimental design. (Please use OpenSSL for your measurement)

Answer (4.a.): Use openssl speed command to measure the performance of each algorithm.

In [87]: !openssl speed -provider legacy -provider default sha1

```
Doing shal ops for 3s on 16 size blocks: 20847777 shal ops in 2.91s
        Doing shal ops for 3s on 64 size blocks: 18608664 shal ops in 2.93s
        Doing shal ops for 3s on 256 size blocks: 13066582 shal ops in 2.91s
        Doing shal ops for 3s on 1024 size blocks: 5572334 shal ops in 2.87s
        Doing shal ops for 3s on 8192 size blocks: 911671 shal ops in 2.97s
        Doing shal ops for 3s on 16384 size blocks: 461492 shal ops in 2.95s
        version: 3.5.2
        built on: Tue Aug 5 12:09:26 2025 UTC
        options: bn(64,64)
        compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
        L -DNDEBUG
        CPUINFO: OPENSSL_armcap=0x987d
        The 'numbers' are in 1000s of bytes per second processed.
                                                 256 bytes 1024 bytes 8192 bytes 16384 bytes
        type
                         16 bytes
                                     64 bytes
        sha1
                        114626.95k
                                    406469.11k 1149500.00k 1988177.71k 2514615.77k 2563079.64k
In [88]: !openssl speed -provider legacy -provider default rc4
        Doing rc4 ops for 3s on 16 size blocks: 179504546 rc4 ops in 2.97s
        Doing rc4 ops for 3s on 64 size blocks: 56846939 rc4 ops in 2.95s
        Doing rc4 ops for 3s on 256 size blocks: 14974308 rc4 ops in 2.92s
        Doing rc4 ops for 3s on 1024 size blocks: 3822397 rc4 ops in 2.92s
        Doing rc4 ops for 3s on 8192 size blocks: 494780 rc4 ops in 2.97s
        Doing rc4 ops for 3s on 16384 size blocks: 244863 rc4 ops in 2.94s
        version: 3.5.2
        built on: Tue Aug 5 12:09:26 2025 UTC
        options: bn(64,64)
        compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
        L -DNDEBUG
        CPUINFO: OPENSSL_armcap=0x987d
        The 'numbers' are in 1000s of bytes per second processed.
                                     64 bytes
                                                 256 bytes 1024 bytes 8192 bytes 16384 bytes
        type
                         16 bytes
                        967027.86k 1233289.52k 1312816.04k 1340457.03k 1364726.52k 1364569.86k
        rc4
In [89]: !openssl speed -provider legacy -provider default blowfish
        Doing blowfish ops for 3s on 16 size blocks: 26504696 blowfish ops in 2.93s
        Doing blowfish ops for 3s on 64 size blocks: 6841755 blowfish ops in 2.92s
        Doing blowfish ops for 3s on 256 size blocks: 1712506 blowfish ops in 2.91s
        Doing blowfish ops for 3s on 1024 size blocks: 438354 blowfish ops in 2.96s
        Doing blowfish ops for 3s on 8192 size blocks: 53665 blowfish ops in 2.92s
        Doing blowfish ops for 3s on 16384 size blocks: 27632 blowfish ops in 2.95s
        version: 3.5.2
        built on: Tue Aug 5 12:09:26 2025 UTC
        options: bn(64,64)
        compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
        L -DNDEBUG
        CPUINFO: OPENSSL_armcap=0x987d
        The 'numbers' are in 1000s of bytes per second processed.
        type
                        16 bytes
                                     64 bytes
                                                 256 bytes 1024 bytes
                                                                          8192 bytes 16384 bytes
                        144735.54k 149956.27k 150653.45k 151646.79k
        blowfish
                                                                          150556.05k 153465.32k
In [90]: !openssl speed -provider legacy -provider default dsa
        Doing 1024 bits sign dsa ops for 10s: 148139 1024 bits DSA sign ops in 9.76s
        Doing 1024 bits verify dsa ops for 10s: 192603 1024 bits DSA verify ops in 9.89s
        Doing 2048 bits sign dsa ops for 10s: 53381 2048 bits DSA sign ops in 9.87s
        Doing 2048 bits verify dsa ops for 10s: 58328 2048 bits DSA verify ops in 9.84s
        version: 3.5.2
        built on: Tue Aug 5 12:09:26 2025 UTC
        options: bn(64,64)
        compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
        L -DNDEBUG
        CPUINFO: OPENSSL armcap=0x987d
                        sign verify
                                           sign/s verify/s
        dsa 1024 bits 0.000066s 0.000051s 15178.2 19474.5
        dsa 2048 bits 0.000185s 0.000169s 5408.4 5927.6
In [91]: import subprocess
         import pandas as pd
         import re
         def run_openssl_speed(algorithm):
             """Run openssl speed and return raw stdout as string."""
             result = subprocess.run(
                 ["openssl", "speed", "-provider", "legacy", "-provider", "default", algorithm],
                 stdout=subprocess.PIPE,
                 stderr=subprocess.STDOUT,
                 text=True
             return result.stdout
         algorithms = {
```

```
"SHA1": "sha1",
             "RC4": "rc4",
             "Blowfish": "blowfish",
             "DSA": "dsa"
         output = dict()
         for algo in algorithms:
             output[algo] = run_openssl_speed(algorithms[algo])
In [92]: print(output['SHA1'])
        Doing sha1 ops for 3s on 16 size blocks: 21540457 sha1 ops in 2.94s
        Doing shal ops for 3s on 64 size blocks: 19056013 shal ops in 2.95s
        Doing shal ops for 3s on 256 size blocks: 13238546 shal ops in 2.92s
        Doing shal ops for 3s on 1024 size blocks: 5857653 shal ops in 2.95s
        Doing shal ops for 3s on 8192 size blocks: 914144 shal ops in 2.96s
        Doing shal ops for 3s on 16384 size blocks: 461701 shal ops in 2.93s
        version: 3.5.2
        built on: Tue Aug 5 12:09:26 2025 UTC
        options: bn(64,64)
        compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
        L -DNDEBUG
        CPUINFO: OPENSSL_armcap=0x987d
        The 'numbers' are in 1000s of bytes per second processed.
                        16 bytes
                                     64 bytes
                                                 256 bytes 1024 bytes 8192 bytes 16384 bytes
                        117226.98k 413418.59k 1160639.65k 2033300.57k 2529955.29k 2581743.75k
        sha1
In [93]: print(output['RC4'])
        Doing rc4 ops for 3s on 16 size blocks: 169798441 rc4 ops in 2.90s
        Doing rc4 ops for 3s on 64 size blocks: 56081978 rc4 ops in 2.94s
        Doing rc4 ops for 3s on 256 size blocks: 14191583 rc4 ops in 2.78s
        Doing rc4 ops for 3s on 1024 size blocks: 3813726 rc4 ops in 2.93s
        Doing rc4 ops for 3s on 8192 size blocks: 490755 rc4 ops in 2.96s
        Doing rc4 ops for 3s on 16384 size blocks: 245549 rc4 ops in 2.94s
        version: 3.5.2
        built on: Tue Aug 5 12:09:26 2025 UTC
        options: bn(64,64)
        compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
        L -DNDEBUG
        CPUINFO: OPENSSL_armcap=0x987d
        The 'numbers' are in 1000s of bytes per second processed.
                        16 bytes
                                     64 bytes 256 bytes 1024 bytes 8192 bytes 16384 bytes
        type
        rc4
                        936818.98k 1220832.17k 1306850.81k 1332851.68k 1358197.62k 1368392.79k
In [94]: print(output['Blowfish'])
        Doing blowfish ops for 3s on 16 size blocks: 26629382 blowfish ops in 2.95s
        Doing blowfish ops for 3s on 64 size blocks: 6873051 blowfish ops in 2.94s
        Doing blowfish ops for 3s on 256 size blocks: 1773275 blowfish ops in 2.98s
        Doing blowfish ops for 3s on 1024 size blocks: 437324 blowfish ops in 2.95s
        Doing blowfish ops for 3s on 8192 size blocks: 55663 blowfish ops in 2.97s
        Doing blowfish ops for 3s on 16384 size blocks: 27581 blowfish ops in 2.96s
        version: 3.5.2
        built on: Tue Aug 5 12:09:26 2025 UTC
        options: bn(64,64)
        compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
        L -DNDEBUG
        CPUINFO: OPENSSL_armcap=0x987d
        The 'numbers' are in 1000s of bytes per second processed.
        type
                        16 bytes
                                     64 bytes
                                                 256 bytes 1024 bytes 8192 bytes 16384 bytes
                                     149617.44k
                                                                           153532.42k 152664.56k
                        144430.55k
                                                 152335.03k 151803.31k
        blowfish
In [95]: print(output['DSA'])
```

```
Doing 1024 bits sign dsa ops for 10s: 147009 1024 bits DSA sign ops in 9.72s
Doing 1024 bits verify dsa ops for 10s: 190589 1024 bits DSA verify ops in 9.79s
Doing 2048 bits sign dsa ops for 10s: 52739 2048 bits DSA sign ops in 9.79s
Doing 2048 bits verify dsa ops for 10s: 58516 2048 bits DSA verify ops in 9.79s
version: 3.5.2
built on: Tue Aug 5 12:09:26 2025 UTC
options: bn(64,64)
compiler: clang -fPIC -arch arm64 -03 -Wall -DL_ENDIAN -DOPENSSL_PIC -D_REENTRANT -DOPENSSL_BUILDING_OPENSS
L -DNDEBUG
CPUINFO: OPENSSL_armcap=0x987d
                  sign
                          verify
                                    sign/s verify/s
dsa 1024 bits 0.000066s 0.000051s 15124.4 19467.7
dsa 2048 bits 0.000186s 0.000167s
                                    5387.0
```

b. Comparing performance and security provided by each method.

```
In [111... sha1_throughput = [float(e[:-1]) for e in output["SHA1"].splitlines()[-1].split()[1:]] # 1,000 KB/s
    rc4_throughput = [float(e[:-1]) for e in output["RC4"].splitlines()[-1].split()[1:]] # 1,000 KB/s
    blowfish_throughput = [float(e[:-1]) for e in output["Blowfish"].splitlines()[-1].split()[1:]] # 1,000 KB/s

sizes = [16, 64, 256, 1024, 8192, 16384] # bytes

plt.figure(figsize=(8,5))
    plt.plot(sizes, sha1_throughput, marker="o", label="SHA1")
    plt.plot(sizes, rc4_throughput, marker="o", label="RC4")
    plt.plot(sizes, blowfish_throughput, marker="o", label="Blowfish")
    plt.xscale("log")
    plt.xscale("log")
    plt.xlabel("Block Size (bytes)")
    plt.ylabel("Throughput (KB/s, log scale)")
    plt.title("Performance Comparison")
    plt.legend()
    plt.show()
```

Performance Comparison

SHA1 RC4 Blowfish

10²

```
In [122... dsa_1024_sign_per_s = [float(e) for e in output["DSA"].splitlines()[-2].split()[5:]] # sign/s
    dsa_2048_sign_per_s = [float(e) for e in output["DSA"].splitlines()[-1].split()[5:]] # sign/s
    sign = [dsa_1024_sign_per_s[0], dsa_2048_sign_per_s[0]]
    verify = [dsa_1024_sign_per_s[1], dsa_2048_sign_per_s[1]]

    key_sizes = ["1024-bit", "2048-bit"]

    print(sign)
    print(verify)

    x = range(len(key_sizes))

    plt.figure(figsize=(6,4))
    plt.bar(x, sign, width=0.4, label="Sign/s")
    plt.bar([i+0.4 for i in x], verify, width=0.4, label="Verify/s")
    plt.xticks([i+0.2 for i in x], key_sizes)
    plt.ylabel("Operations per second")
```

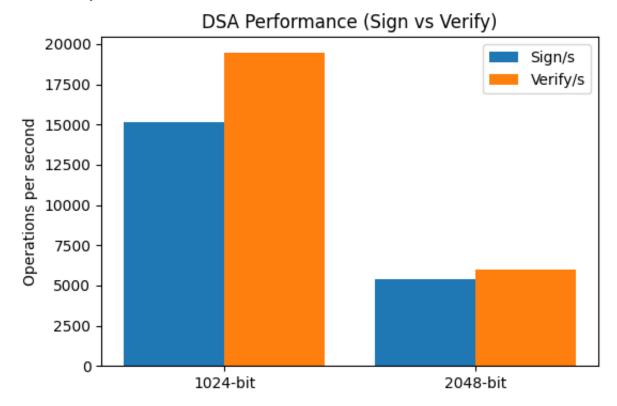
Block Size (bytes)

 10^{3}

 10^{4}

```
plt.title("DSA Performance (Sign vs Verify)")
plt.legend()
plt.show()
```

[15124.4, 5387.0] [19467.7, 5977.1]



Answer (4.b.):

SHA1

- ความเร็วสูงมาก โดยเฉพาะเมื่อ block size ใหญ่ (ประมาณ 2.5 GB/s สำหรับ block size 16384 bytes)
- ปัจจุบันถูกพบช่องโหว่ collision แล้ว → ไม่ปลอดภัยสำหรับการใช้งานที่ต้องการความมั่นคง

RC4

- ความเร็วสูงมาก สลับกันเป็นอันดับ 1 และ 2 กับ SHA1 ขึ้นอยู่กับ block size
- ไม่ปลอดภัยแล้ว (มี bias ใน keystream, ถูกห้ามใช้ใน TLS, Wi-Fi)

Blowfish

- ช้ากว่า SHA1/RC4 มาก (~150 MB/s)
- ปลอดภัยกว่า RC4 และ SHA1
- ข้อจำกัด: block size 64-bit ทำให้ไม่เหมาะกับข้อมูลขนาดใหญ่
- ปัจจุบันนิยมเปลี่ยนไปใช้ AES แทน

DSA

- วัดผลในรูปแบบจำนวนครั้งต่อวินาที (ops/s) ไม่ใช่ throughput
- เมื่อเปรียบเทียบกับ symmetric crypto (SHA1/RC4/Blowfish) ถือว่าช้ามาก
- 1024-bit: เลิกใช้แล้ว ไม่ปลอดภัย
- 048-bit: ยังปลอดภัย แต่ช้ากว่า และถูกแทนที่ด้วย ECDSA/EdDSA
- ใช้สำหรับลายเซ็นดิจิทัล ไม่ใช่สำหรับเข้ารหัส bulk data

สรุป

- RC4: เร็วที่สุด แต่ ไม่ปลอดภัย
- SHA1: เร็วมาก แต่ ถูกโจมตีได้แล้ว
- Blowfish: ช้ากว่า แต่ปลอดภัยกว่าสองตัวแรก อย่างไรก็ตามถือว่า **ล้าสมัย**
- DSA: ช้าที่สุด (หลายพัน ops/s เทียบกับ symmetric ที่วัดได้เป็น GB/s) แต่ใช้ในงาน digital signature ซึ่งมีวัตถุประสงค์ต่างออกไป

c. Explain the mechanism underlying Digital Signature. How does it combine the strength and weakness of each encryption scheme?

Answer (4.c.):

การสร้างลายเซ็น (Signing):

- ผู้ส่ง (Sender) มี public/private key pair
- ข้อความ (Message) ถูก Hash เพื่อสร้าง message digest
- message digest ถูกเข้ารหัสด้วย private key ของผู้ส่ง -> digital signature

การตรวจสอบลายเซ็น (Verification):

- ผู้รับ (Receiver) ได้รับข้อความและ digital signature
- ผู้รับทำ Hash ข้อความอีกครั้งด้วยฟังก์ชันเดียวกัน
- ผู้รับถอดรหัส digital signature ด้วย public key ของผู้ส่ง -> ได้ message digest ที่ผู้ส่งสร้าง
- เปรียบเทียบ message digest ที่ผู้รับสร้างกับที่ถอดรหัสได้
- ถ้าตรงกัน
 - แสดงว่าข้อความไม่ถูกแก้ไข (Integrity)
 - ยืนยันว่ามาจากผู้ส่งจริง (Authentication)
 - ผู้ส่งไม่สามารถปฏิเสธได้ว่าไม่ได้ส่งข้อความนั้น (Non-repudiation)

Strengths/Weaknesses:

- Symmetric crypto (เช่น AES, Blowfish)
 - Strength: เร็ว, เหมาะกับข้อมูลขนาดใหญ่
 - Weakness: ต้องมีการแลกเปลี่ยนกุญแจลับ (key) อย่างปลอดภัย
- Asymmetric crypto (เช่น DSA, RSA)
 - Strength: ไม่ต้องแลกเปลี่ยนกุญแจลับ ใช้ public key ได้เลย
 - Weakness: ช้ามาก, ไม่เหมาะกับข้อมูลขนาดใหญ่

Combination: - ไม่ได้เข้ารหัสข้อความทั้งฉบับ (ช้า) - เข้ารหัสแค่ Message Digest (เล็กมาก, เร็ว) - Message จริง ส่งตรง ๆ ได้ (เร็ว)

Conclusion: - Digital Signature = Hash (เร็ว) + Asymmetric Crypto (ปลอดภัย) - ได้ ความถูกต้องของข้อมูล (Integrity), การยืนยันตัว ตน (Authentication) และ การปฏิเสธไม่ได้ (Non-repudiation) - ดึงจุดแข็ง เลี่ยงจุดอ่อน

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