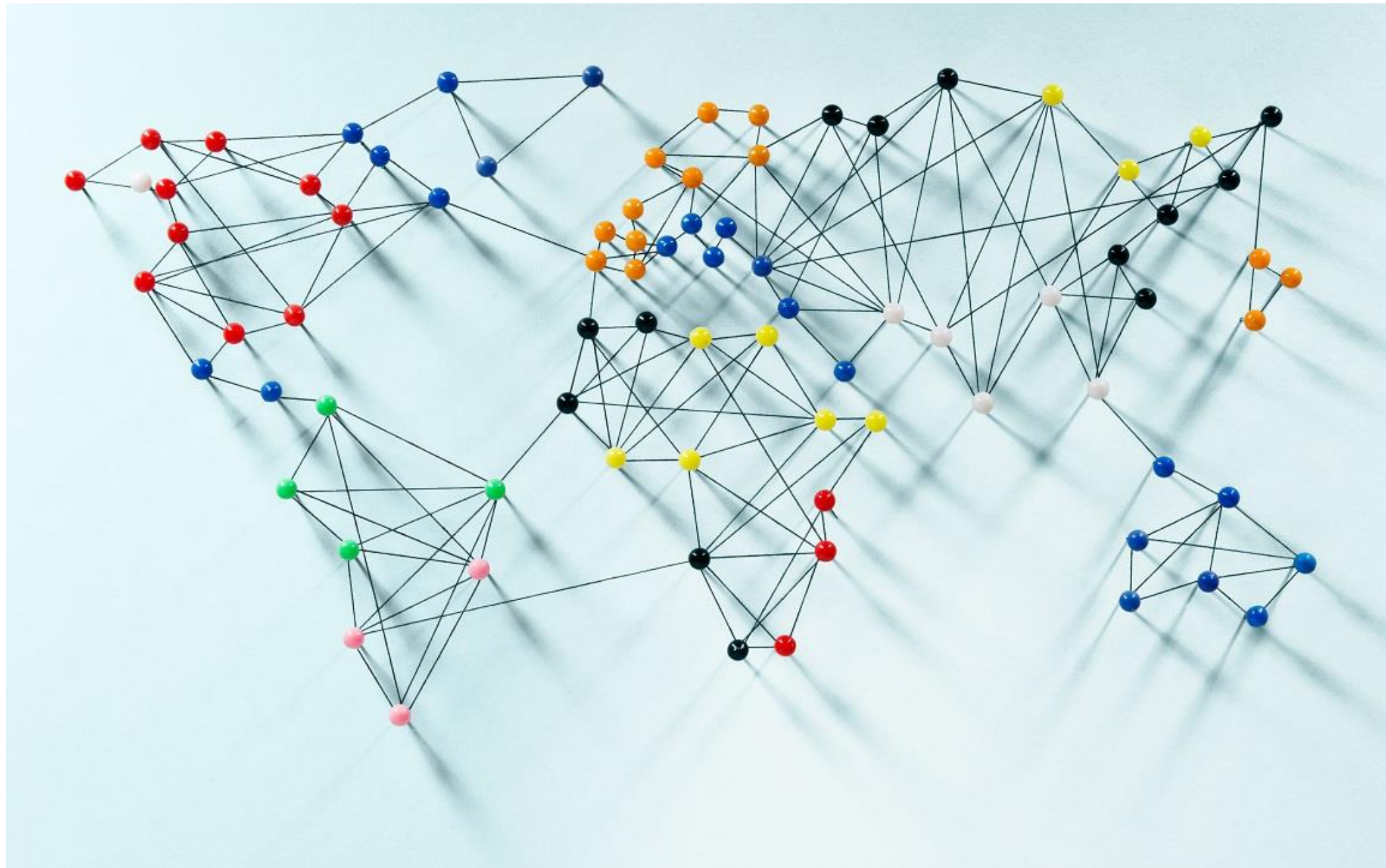


# Cut-Sets & Cut-Vertices

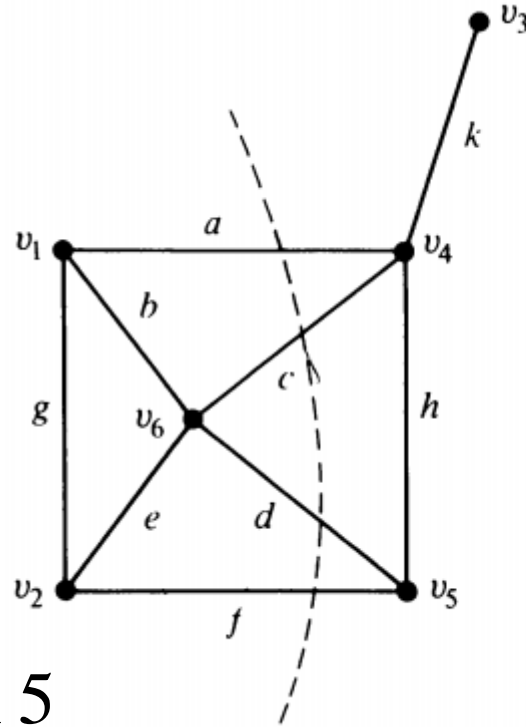
A.B.M. Ashikur Rahman



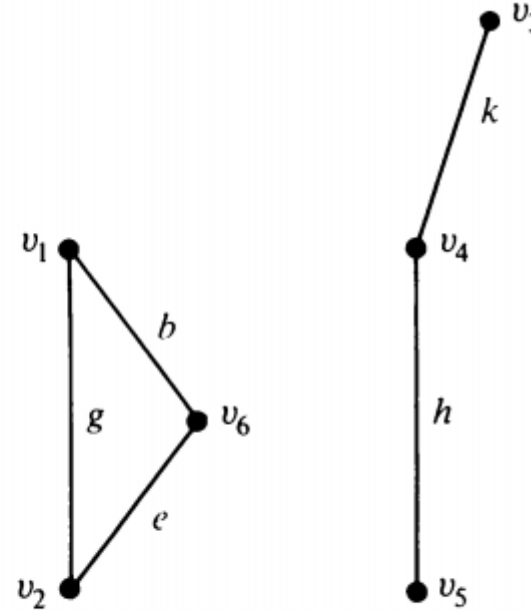
# Cut-Sets

- In a connected graph  $G$ , a *cut-set* is a set of edges whose removal from  $G$  leaves  $G$  disconnected, provided removal of no proper subset of these edges disconnects  $G$ .
- Minimal cut-set/Proper cut-set/simple cut-set/cocycle
- Cut-set always cuts the graph in two.
- Removal of cut-set reduces the rank of graph by one.

# Cut-Sets



Rank 5



Rank 4

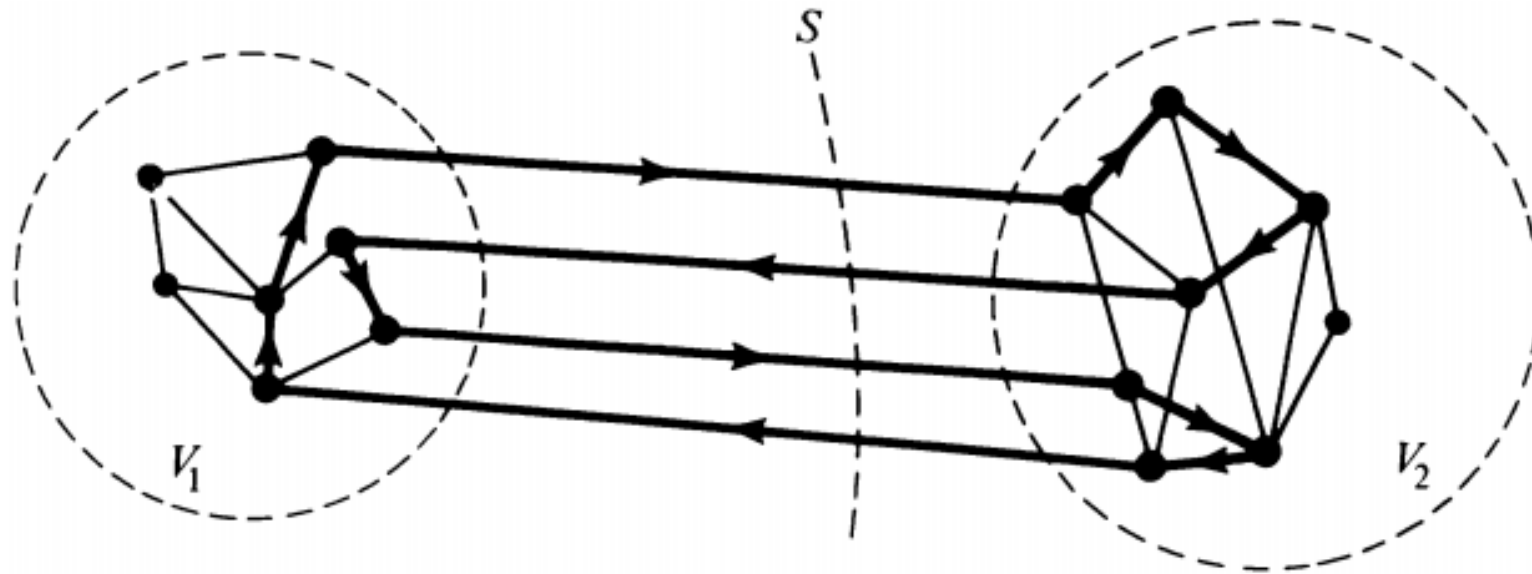
- For tree every edge is a cut-set

# Properties of a Cut-Set

- Theorem 4.1 – Every cut-set in a connected graph  $G$  must contain at least one branch of every spanning tree of  $G$ .
- Theorem 4.2 – In a connected graph  $G$ , any minimal set of edges containing at least one branch of every spanning tree of  $G$  is a cut-set.

# Properties of a Cut-Set

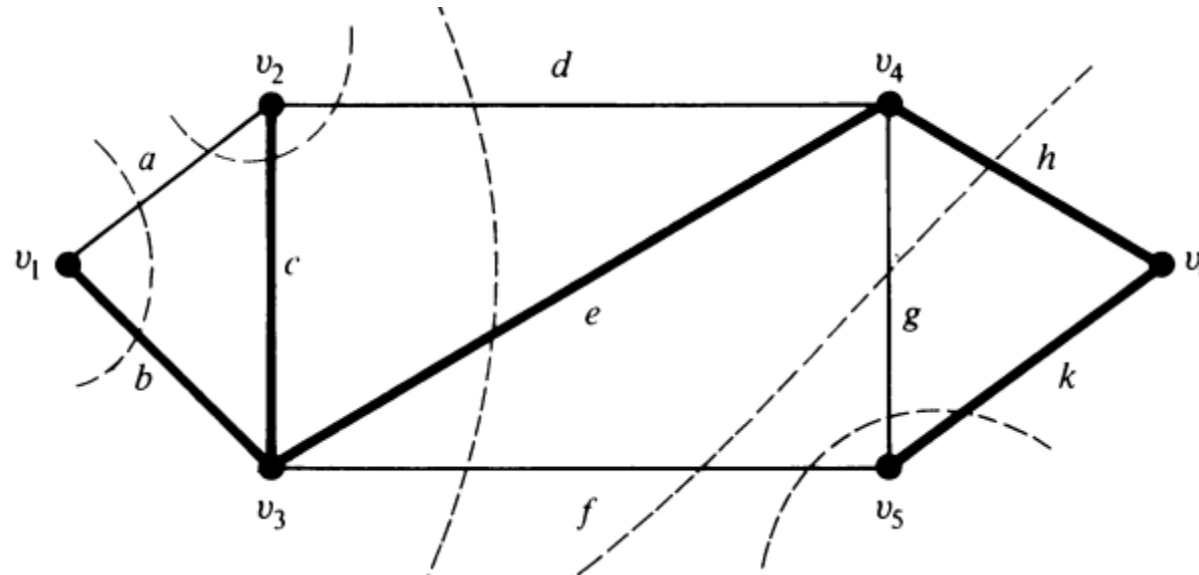
- Theorem 4.3 – Every circuit has an even number of edges in common with any cut-set.



Circuit  $\Gamma$  shown in heavy lines, and is traversed along the direction of the arrows

# All Cut-Sets in a graph

- Just as a spanning tree is essential for defining a set of fundamental circuits, so is a spanning tree essential for a set of *fundamental cut-sets*.
- A cut-set  $S$  containing exactly one branch of a tree  $T$  is called a *fundamental cut-set* with respect to  $T$ .

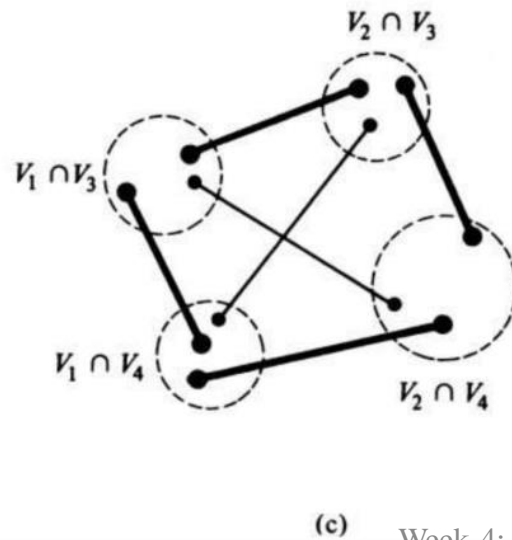
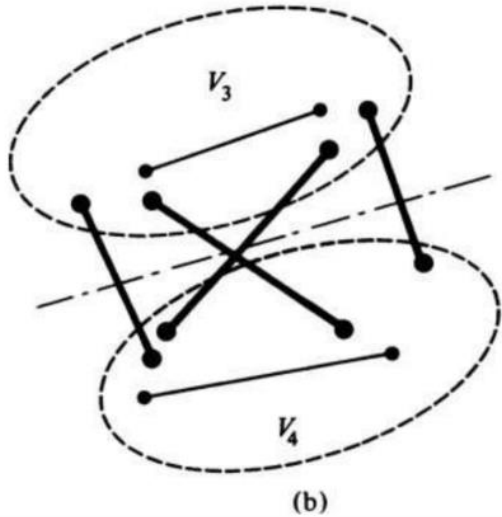
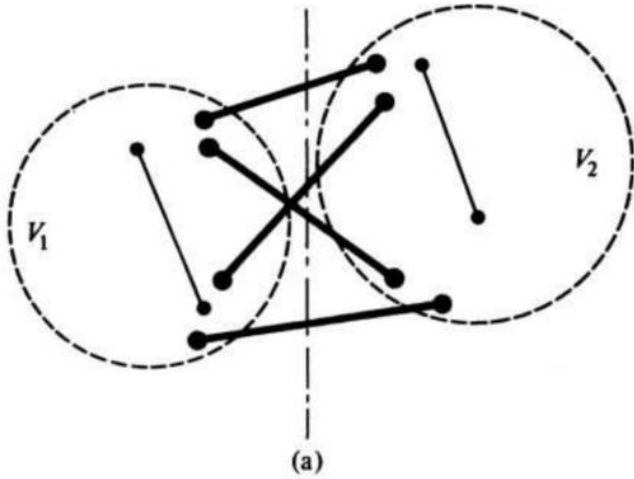


# All Cut-Sets in a graph

- Just as every chord of a spanning tree defines a *unique* fundamental circuit, every branch of a spanning tree defines a *unique* fundamental cut-set.
- Both fundamental cut-set and fundamental circuit has meaning only with respect to a *given* spanning tree.
- Theorem 4.4 – The ring sum of any two cut-sets in a graph is either a third cut-set or an edge-disjoint union of cut-sets.



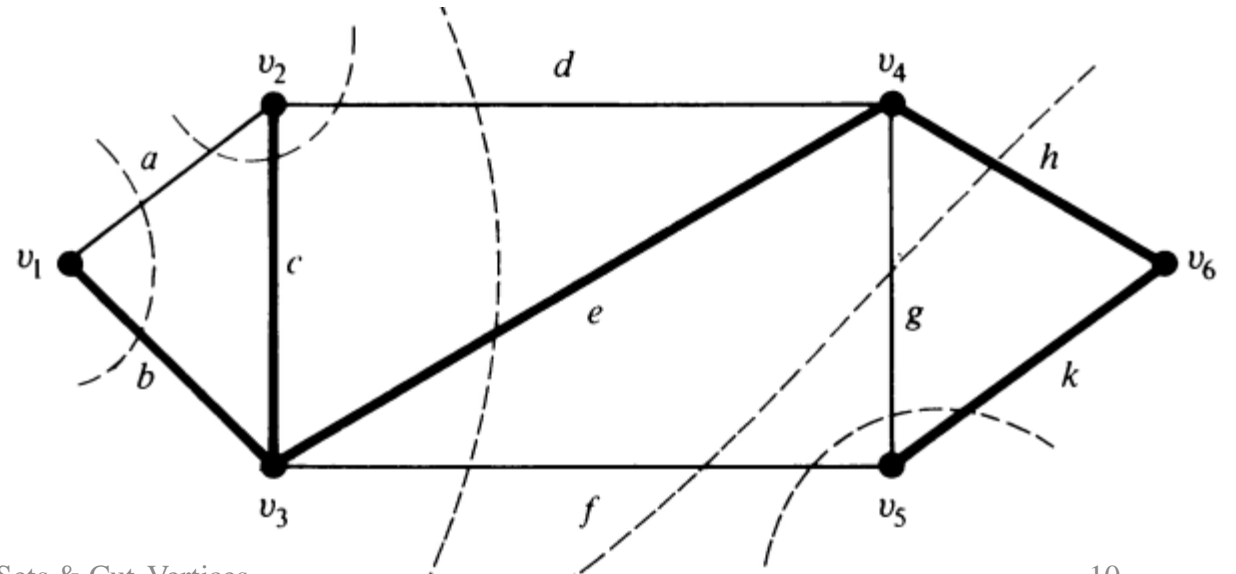
# All Cut-Sets in a graph



# All Cut-Sets in a graph

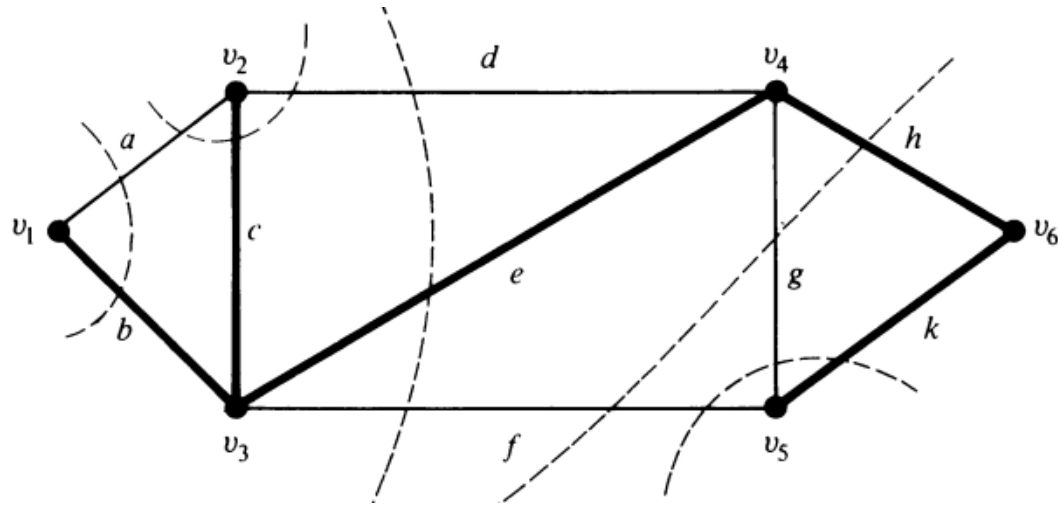
- $\{d,e,f\} \oplus \{f,g,h\} = \{d,e,g,h\}$ ; another cut-set
- $\{a,b\} \oplus \{b,c,e,f\} = \{a,c,e,f\}$ ; another cut-set
- $\{d,e,g,h\} \oplus \{f,g,k\} = \{d,e,f,h,k\}$   
 $= \{d,e,f\} \cup \{h,k\}$ ; an edge-disjoint union of cut-sets.

So, we found a way to generate more cut-sets



# Fundamental Circuits & Cut-Sets

- Theorem 4.5 – With respect to a given spanning tree  $T$ , a chord  $c_i$  that determines a fundamental circuit  $\Gamma$  occurs in every fundamental cut-set associated with the branches in  $\Gamma$  and in no other.
- $T = \{b, c, e, h, k\}$
- $\Gamma = \{f, e, h, k\}$ ; taking  $f$
- Cut-Sets produced
  - $\{d, e, f\}$
  - $\{f, g, h\}$
  - $\{f, g, k\}$



# Fundamental Circuits & Cut-Sets

- Theorem 4.6 – With respect to a given spanning tree  $T$ , a branch  $b_i$ , that determines a fundamental cut-set  $S$  is contained in every fundamental circuit associated with the chords in  $S$ , and in no others.

- $T = \{b, c, e, h, k\}$
- Fundamental Cut-set by  $e = \{d, e, f\}$
- F.C. for chord  $d = \{d, c, e\}$
- F.C. for chord  $f = \{e, h, k, f\}$

