

#1
(Lec-12)

Chapter 7 (TCP/IP)

IPv4

28-June

Monday

Important frame/packet format \rightarrow Wired LAN, IPv4, IPv6

Syllabus (7.1-7.5), 7.7, 7.8.

IP \rightarrow connectionless protocol

Trying to move from datagram to virtual circuit switching.

- IGMP \rightarrow Internet Group Management Protocol
- ICMP \rightarrow Internet Control Message Protocol
- ARP \rightarrow Address Resolution Protocol

\rightarrow All these are helper protocols and these packets are encapsulated by IP packet.

Ex- ~~the~~ the physical address generated by ARP is encapsulated by ~~IP~~ ^{IP}

\rightarrow If a router receives an unfragmentable packet exceeding MTU (maximum transfer unit) then the packet is dropped and an error message is sent.

ICMP packets are used for error messages and the ICMP packet is encapsulated by IP.

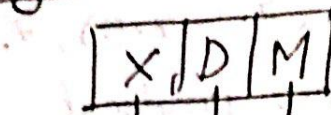
IP provides encapsulation to other helper protocols.

Datagram (IP packet)

IP header (20-60 bytes)

- * The protocol field in an IP packet determines receiving protocol.
- * Time to live \rightarrow maximum no. of hops (routers) visited by packet.
 - (The packet decrements by 1 in each hop).
 - Usually 2 times max no. of routers between 2 hosts
 - When $ttho = 0$, then packet is dropped.

→ 3-bit-flag



Reserved

Do not fragment

More fragment

[More fragments are coming]
 → 1 (not last fragment)
 → 0 (last fragment)

a) Version

→ Valid 4/6 [0100 or 0110]

b) Header Length

No. of bytes of packet [4 bits]

→ Multiplying by 4 we get header length

If the packet is

0100, 1110

Ip4

14

len

→ $(14 \times 3) = 42$ bytes

Header field can have (5 to 15) values.

c) ~~Frame offset~~

⊛ Identification number is counted in the same session.
 The fragmented packets have same identification no.

Chapter-7 (TCP/IP) IPv4 (continued)

30-June
Wednesday

Service Type [3 bits | 4 bit | X] 8 bit

[For priority] →

3 bits for precedence
4 bit TOS (type of service)

- Delay
- Throughput
- Reliability
- Cost

D T R C

0000 (Normal)
1000 Minimise Delay
0100 Maximise throughput
0010 Maximize Reliability
0001 Minimise cost

Newer version of service type

6 bits | X | X

$2^6 = 64$ code-points

Differential Service Interpretation

→ This kind of assignment is done for backwards compatibility

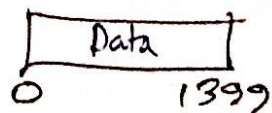
Category

1 [0] Internet
2 [11] Local
3 [01] Temporary

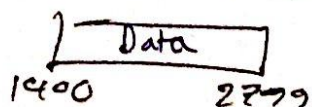
[*] Only data portion of the packet can be fragmented. The header is just copied.

0, 2, 4, 6, ...
3, 7, 11, ...
1, 5, 9, ...

Offset → the starting value of the whole data frame Divide by 8



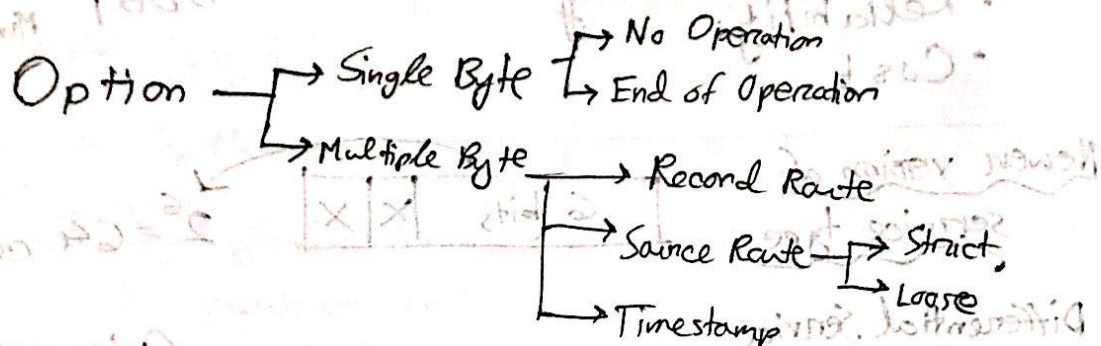
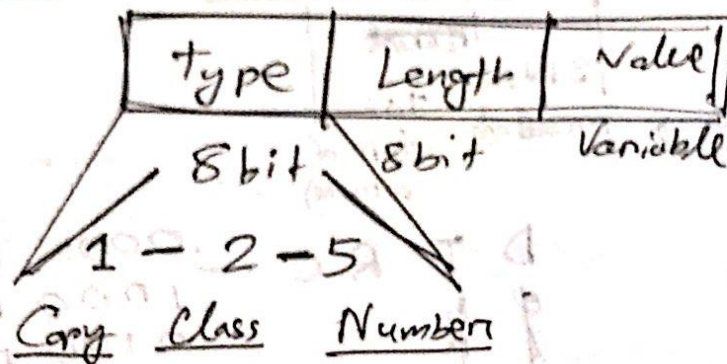
offset value → $0/8 = 0$



offset value → $1400/8 = 175$

By multiplying offset value by 8, we get actual value.

Option Format (TLV)



- Strict → All the addresses must be visited but no other addresses can be visited
- Loose → All addresses must be visited, more addresses can be visited if required.

In case of source routing,
destination IP address changes.

Record Route

<u>Type</u>	<u>Length</u>	<u>Value Pointer</u>	<u>Value</u>
7	(No. of bytes for recording) Single IP address is 4 bytes (32 bits)	At Pointing to the next empty space	

If length is 15 then it can store 15 IP addresses

- The IP addresses are stored in value for Record route when a new router is visited
- For source route, at first all IP addresses are stored and only those are visited (for strict)

For timestamp, the router puts their time-stamp in the value field. Overflow determines the no. of routers who couldn't put timestamp value ~~in~~ because value ~~is~~

- Timestamp can be done with strict routing.

IP Package → The modules responsible for packaging the IP address.

- Read the ~~new~~ pseudocode.

A diagram showing a rectangular frame containing two boxes. The first box is labeled 'IP' and the second box is labeled 'ARP'. An arrow points from the 'IP' box to the 'ARP' box, and another arrow points from the 'ARP' box downwards, indicating a process flow.

broadcast

Static Mapping Disadvantages

→ Changing Network Interface Card (NIC) changes P.A.

But using ARP, and physical address, we solve this problem.

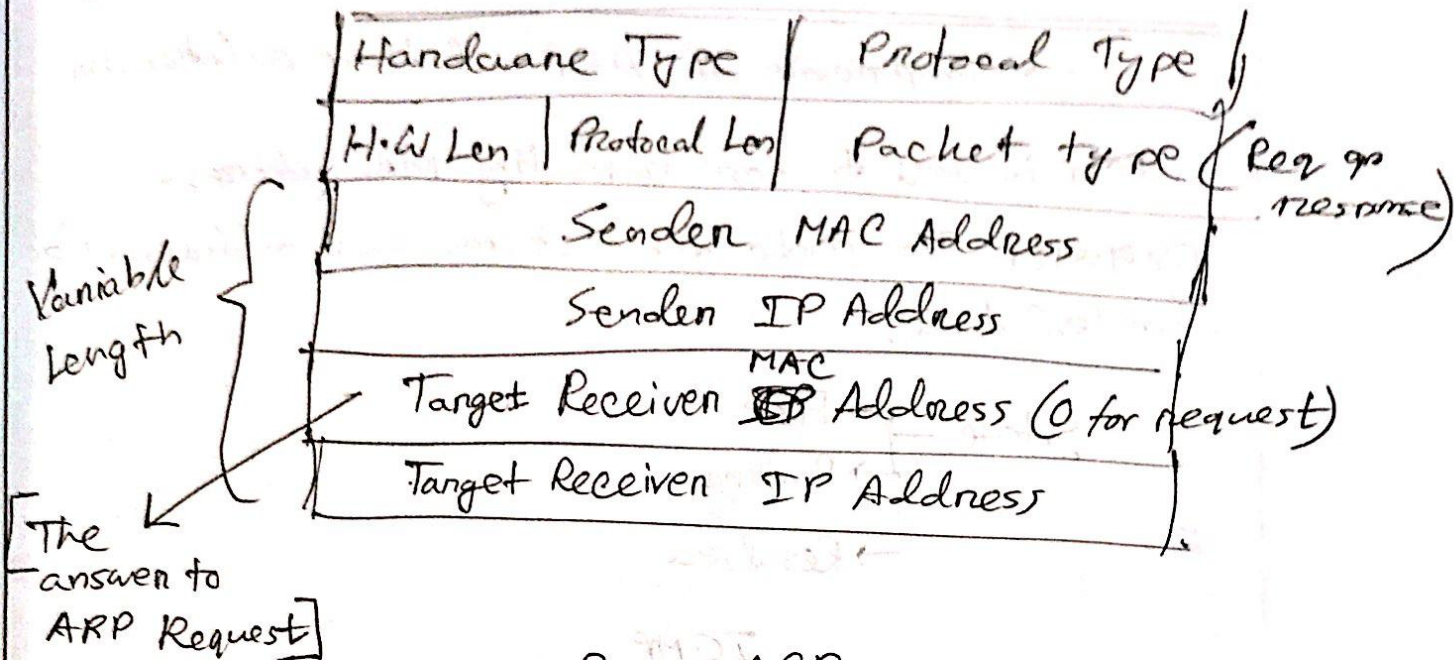
ARP Packet. — $\left\{ \begin{array}{l} \rightarrow \text{Query Packet [Broadcast]} \\ \rightarrow \text{Response Packet [Unicast]} \end{array} \right.$

⊛ ARP requests to find target

IP address's physical address.

* In the ethernet frame used for ARP query or request packet, the physical address is broadcast address.

ARP Packet



Proxy ARP

→ For subnetting (to restrict broadcast domain) effect

- * A proxy ARP router replies to any ARP request, for those particular destinations, with its own physical address.

(a set of hosts)

I.e. reply packet has ^{proxy router's} ARP physical address in sender MAC address portion.

When the router receives the actual ^{IP} packet, it forwards to the actual sender.

ARP Package

→ Components of ARP and their relationships

Cache is used to store repeating MAC address requests. The cache table stores such addresses for limited time.

State →

- Free
- Pending
- Resolved

ICMP

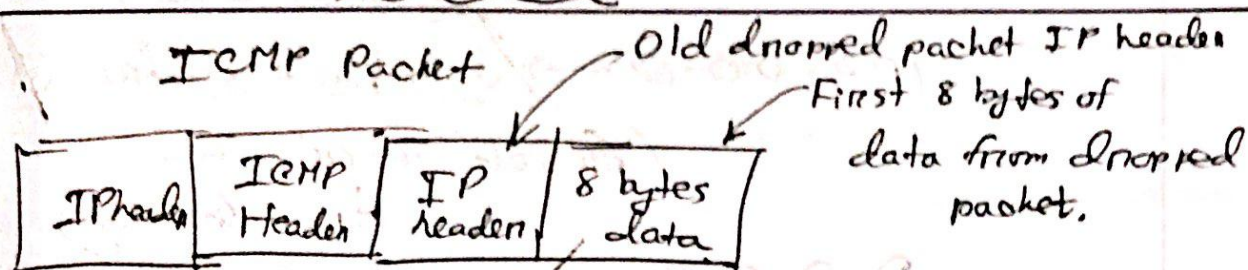
→ For error ^{reporting} messages

→ query messages (shortcoming of IP for unreliability)

[Encapsulated in an IP packet]
(Goes to upper layer)

ICMP

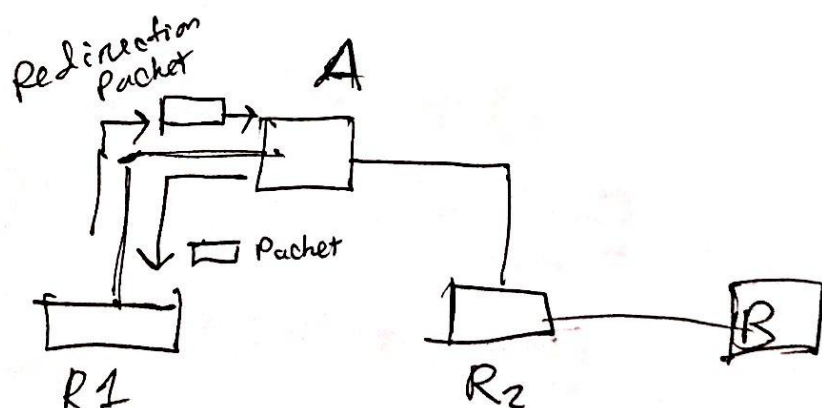
Chapter-9



ICMP packet encapsulated in an IP packet.

Redirection Concept

→ A host machine has no routing capability and sends it to wrong router, The wrong router redirects the message to host and tells it to send to appropriate router.



Cases where ICMP is not generated

- In response to datagram carrying ICMP error.
- A fragmented datagram not the first fragment.
- For special addresses 127.0.0.0 or 0.0.0.0

127.0.0.0 → Loop back address

(for self-referencing and troubleshooting)

0.0.0.0 → unspecified address

→ used during packet generation, assigned by a device after it boots up,

