

Coq cheat sheet

Notation

| Propositions | Coq |
|---|---------------------------|
| \top, \perp | True, False |
| $p \wedge q$ | p /\ q |
| $p \Rightarrow q$ | p -> q |
| $p \vee q$ | p \/ q |
| $\neg p$ | ~ p |
| $\forall x \in A. p(x)$ | forall x:A, p x |
| $\forall x, y \in A. \forall u, v \in B. q$ | forall (x y:A) (u v:B), q |
| $\exists x \in A. p(x)$ | exists x:A, p x |

| Sets | Coq |
|----------------------------|------------------------|
| 1 | unit |
| $A \times B$ | prod A B or A * B |
| $A + B$ | sum A B or A + B |
| B^A or $A \rightarrow B$ | A -> B |
| $\{x \in A \mid p(x)\}$ | {x:A p x} |
| $\sum_{x \in A} B(x)$ | {x:A & B x} or sig A B |
| $\prod_{x \in A} B(x)$ | forall x:A, B x |

| Elements | Coq |
|--|--------------------------------|
| $\star \in 1$ | tt : unit |
| $x \mapsto f(x)$ or $\lambda x \in A. f(x)$ | fun (x : A) => f x |
| $\lambda x, y \in A. \lambda u, v \in B. f(x)$ | fun (x y : A) (u v : B) => f x |
| $(a, b) \in A \times B$ | (a,b) : A * B |
| $\pi_1(t)$ where $t \in A \times B$ | fst t |
| $\pi_2(t)$ where $t \in A \times B$ | snd t |
| $\pi_1(t)$ where $t \in \sum_{x \in A} B(x)$ | projT1 t |
| $\pi_2(t)$ where $t \in \sum_{x \in A} B(x)$ | projT2 t |
| $\iota_1(t) \in A + B$ where $t \in A$ | inl t |
| $\iota_2(t) \in A + B$ where $t \in B$ | inr t |
| $t \in \{x \in A \mid p(x)\}$ because ρ | exist t ρ |
| $\iota(t)$ where $\iota : \{x \in A \mid p(x)\} \hookrightarrow A$ | projT1 t |

Basic tactics

| When the goal is ... | ... use tactic |
|-----------------------|---------------------------|
| very simple | auto, tauto or firstorder |
| $p \wedge q$ | split |
| $p \vee q$ | left or right |
| $p \rightarrow q$ | intro |
| $\sim p$ | intro |
| $p \leftrightarrow q$ | split |
| an assumption | assumption |
| forall x , p | intro |
| exists x , p | exists t |

| To use hypothesis H ... | ... use tactic |
|---------------------------|---------------------------------------|
| $p \vee q$ | destruct H as [H_1 H_2] |
| $p \wedge q$ | destruct H as [H_1 H_2] |
| $p \rightarrow q$ | apply H |
| $p \leftrightarrow q$ | apply H |
| $\sim p$ | apply H or elim H |
| False | contradiction |
| forall x , p | apply H |
| exists x , p | destruct H as [x G] |
| $a = b$ | rewrite H or rewrite $\leftarrow H$ |

| If you want to ... | ... then use |
|--|---------------------|
| prove by contradiction $p \wedge \neg p$ | absurd p |
| simplify expressions | simpl |
| prove via intermediate goal p | cut p |
| prove by induction on t | induction t |
| pretend you are done | admit |
| import package P | Require Import P |
| compute t | Eval compute in t |
| print definition of p | Print p |
| check the type of t | Check t |
| search theorems about p | SearchAbout p |

Inductive definitions

Inductive definition of X

```
Inductive X args :=  
  | constructor1 : args1 -> X  
  | constructor2 : args2 -> X  
  ...  
  | constructorN : argsN -> X.
```

Coq generates induction and recursion principles `X_ind`, `X_rec`, `X_rect`.

Construction of an object by cases

```
match t with  
  | case1: result1  
  | case2: result2  
  ...  
  | caseN: resultN  
end
```

Recursive definition of f

```
Fixpoint f args := ...
```