

Mid Term Exam Question Paper Subject: Design and Analysis of Algorithms (CSE 325) Max Marks = 50 Time = 1.5 hours

Note: Answer all parts of the question at one place only. Non-conformity will lead to answers being not corrected or deduction of marks. Also write your answers to the point in order to attempt all questions in time.

Q1. Fill in the blanks [5 marks] a. The upper bound for nu	mber of iterations in (Gale-Shapley algorithm for s	table matching
between n men and n w			
b. An algorithm is "efficie			
		class of problems	
			y
d. Bucket sort can be used	when the numbers to	be sorted are	in auxiliary
e. The 3 rd step (before re-	arranging) in counting	sort computes	III auxiliary
array.			
Q2. Answer the following que			
		times $T(n) = 5T\binom{n}{4} + n$	
$T(n) = aT(^{n}/_{8}) + n$	log n respectively. F	ind the maximum value of 's	a' such that B
algorithm is faster than	n A?	A THE TELESTICATION OF THE	
by To sort an array of int	egers with range from	1 to n ^c using radix sort in O	(n) time, find the
value of base in which	the numbers can be r	represented?	
c) For quick sort, let the	e be a procedure (with	h O(n) time) for finding a go	ood pivot element
which splits into two, each of which contains atleast one-third of elements. Write the			
recurrence for the wo			
0.1	C. 11in a gyantions	[2 6 19	
Q3. Attempt any three of the	f(n) = 2T(n / 2) + 1	$n \log_2 n$ using either substitution	ution or
recurrence tree metho	(n) = 21(n) = 3	titute $n=2^{m}$)	
b) Derive the asymptotic	cally tight bound for t	the BUILD MAX HEAP alg	gorithm (discussed
in class) which is used	to re-arrange any give	en array into a max heap arra	y. Note: you won't
get credit if your ansy	ver is not asymptotica	lly tight bound.	
c) Let f(n) and g(n) be	asymptotically nonne	gative functions. Using the	basic definition of

theta notation, prove that $\max(f(n), g(n)) = \Theta(f(n) + g(n))$

- d) How would you modify Strassen's algorithm to multiply $n \times n$ matrices in which n is not an exact power of 2 so that the algorithm still has a running time of Theta($n^{\log(7)}$)?(Hint: Think of extending the idea of block matrix multiplication)
- Q4. Although merge sort runs in $\Theta(\log n)$ worst-case time and insertion sort runs in $\Theta(n^2)$ worst-case time, the constant factors in insertion sort can make it faster in practice for small problem sizes on many machines. Thus, it makes sense to coarsen the leaves of the recursion by using insertion sort within merge sort when subproblems become sufficiently small. Consider a modification to merge sort in which n/k sublists of length k are sorted using insertion sort and then merged using the standard merging mechanism, where k is a value to be determined. [10 marks]
- a. What is the worst case time complexity for insertion sort to sort the n/k sublists, each of length k. [2 mark]
- b. Write the recurrence relation for the modified algorithm in worst case and solve it to determine the asymptotic running time. [5 mark]
- c. What is the largest value of k as a function of n for which the modified algorithm has the same running time as standard merge sort, in terms of Θ -notation? [2 mark]
- d. How should we choose k in practice? [1 mark]
- Q5. Given an array of N integers (which can be positive or negative, representing a person's emotional index on each day). Let the happiness in an interval be defined as the sum of the values in that interval. Design an O(N log N) algorithm to find the happiest interval of the person and its happiness value. For example, if the given array is $\{-2, -5, 6, -2, -3, 1, 5, -6\}$, then the happiest interval value is 7 (see highlighted elements). Write the pseudo code of the algorithm and justify its correctness and time complexity. [8 marks].