#### **CS2109S: Introduction to AI and Machine Learning**

## Lecture 2: Solving Problems by Searching

21 January 2024

## DO NOT CLOSE YOUR POLLEVERYWHERE APP

There will be activities ahead

## New Midterm Timing

• Date/Time: Tuesday, 4 March (Week 7), 6.30pm - 8.00pm

• Venue: MPSH 2A & 2B

#### Consultation

#### • Walk-in:

- Office: COM2-02-06
  - If it's open, you can come in
  - If not and the light is on, you can knock
- Highest availability: Tuesday, 12-3pm

#### Appointment:

Send an email to rizki@nus.edu.sg



Muhammad Rizki Maulana rizki@nus.edu.sg

#### Consultation

#### • Walk-in:

- Office: COM2-04-29
- Monday, 2-4pm

#### Appointment:

• Send an email to cqtfpr@nus.edu.sg



Patrick Rebentrost cqtfpr@nus.edu.sg
Quantum Machine Learning

## For those who are just enrolled

- Please watch the recording or read lecture slides for Lecture 1
- Important: Course Policies (also available in Canvas)

#### Outline

- Designing an Agent
  - Problem solving agents
  - Problem formulation
- Search
  - Uninformed Search
    - Breadth-first, depth-first, and uniform-cost search
    - Depth-limited and iterative-deepening search
  - Informed Search
    - A\* search
    - Heuristics

#### Outline

#### Designing an Agent

- Problem solving agents
- Problem formulation

#### Search

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## Preliminary: Search Problems

- A search problem refers to a type of problem where the goal is to find a state (or a path to a state) from a set of possible states by exploring various possibilities.
- Examples:
  - Path finding
  - Puzzle solving



9	1	3			5		
6		7				2	4
	5		8			7	
	7	9					
		2	9			4	3
				4		9	
	4			1	9		
7		6		9			5
		1		6	4		7

## Preliminary: Search Algorithms

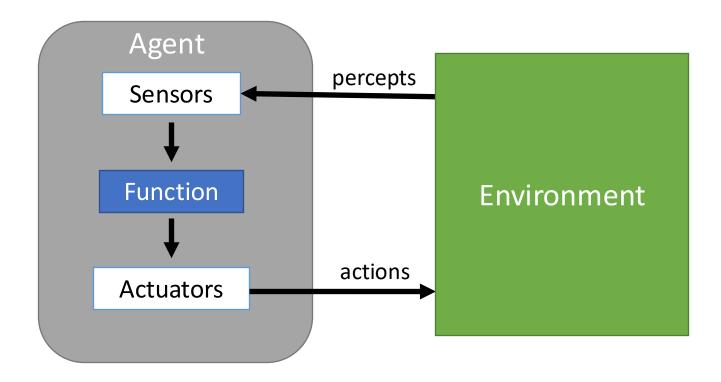
- A search algorithm takes in a search problem as input and returns a solution / failure.
- A solution of a search could be a sequence of actions or final state
  - An optimal solution is a solution with least cost, or least number of steps if there is no cost associated with actions
- Search algorithms
  - Linear search (bruteforce)
  - Djikstra (CS2040S)

• ...

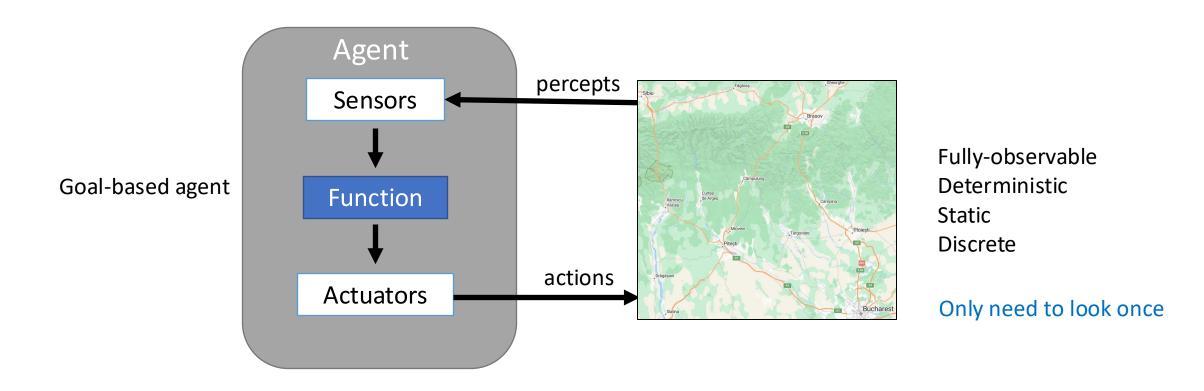


9	1	3			5		
6		7				2	4
	5		8			7	
	7	9					
		2	9			4	3
				4		9	
	4			1	9		
7		6		9			5
		1		6	4		7

## Designing an Agent



## Designing an Agent for problems that can be solved by searching. E.g., path finding

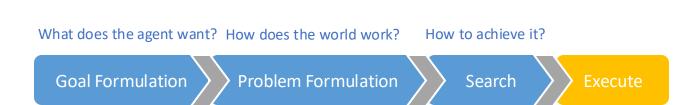


How is search used by the agent?

## Designing an Agent

To solve a problem using search, the agent needs to have:

- A goal, or a set of goals
- A model of the environment
- A search algorithm



**Problem Solving Steps** 

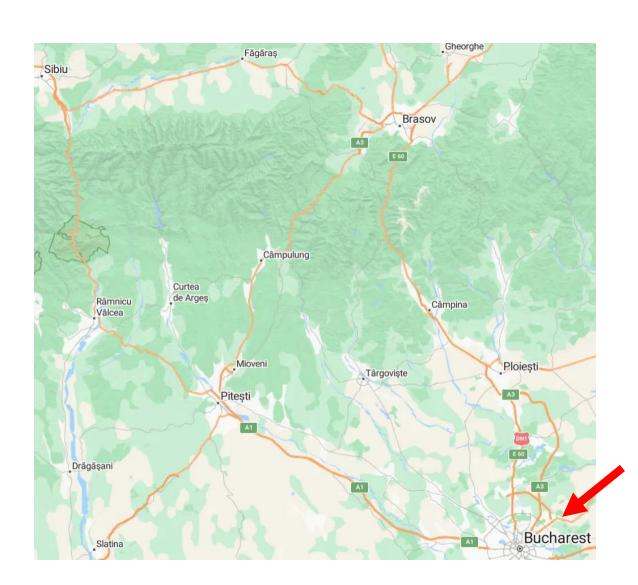
## Example: Romania

Goal Formulation

Problem Formulation

Search

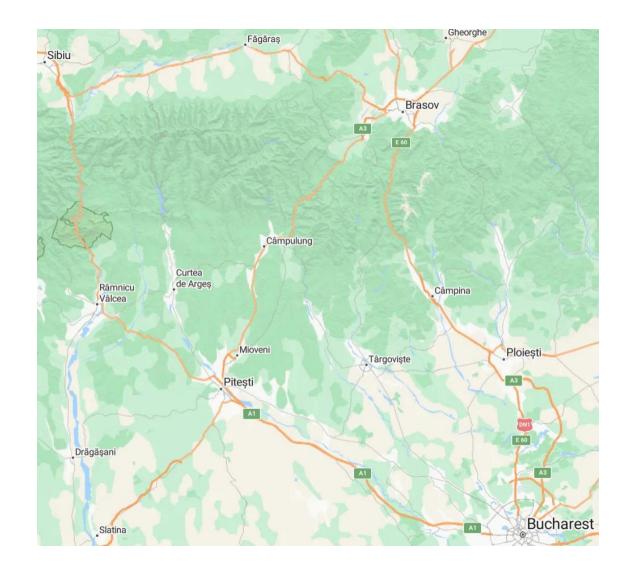
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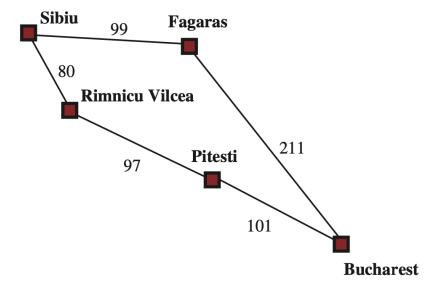


Go to Bucharest

## Example: Romania

**Problem Formulation** 





#### **Problem Formulation**

States

Initial state

Goal state(s)/test

Actions

Transition model

Action cost function See edges

Cities: {at Sibiu, at Pitesti, ...}

at Sibiu

at Bucharest

Go to a neighboring city

Current state = selected city

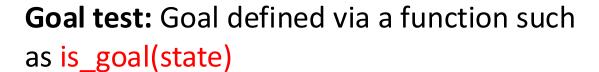
## Important facts

Goal Formulation Problem Formulation

Search

Execute

**Representation invariant:** ensure that the *abstract states* have corresponding *concrete states*.



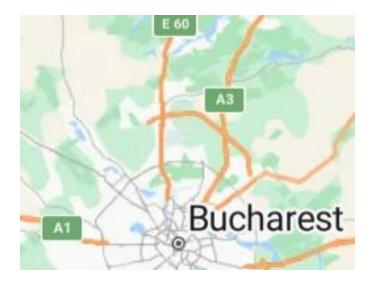
**Actions:** a set actions(state) with size |actions(state)| ≤ b (branching factor)

**Transition model:** new\_state = transition(state, action)





Credit: Passport & Plates



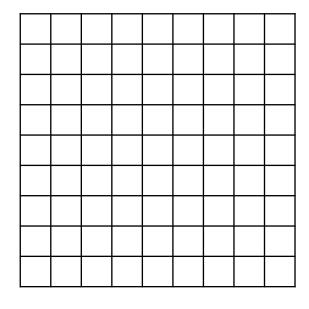
## Example: Sudoku

Goal Formulation Problem Formulation

Search

Execute

9	1	3			5		
6		7				2	4
	5		8			7	
	7	9					
		2	9			4	3
				4		9	
	4			1	9		
7		6		9			5
		1		6	4		7



9x9 Integer Matrix

#### **Problem Formulation**

States All possible assignments of

numbers to a 9x9 matrix

• Initial state Partially filled matrix

• Goal state(s)/test Check if all constraints are

satisfied

• Actions Fill matrix

Transition model Current state = filled matrix

Action cost function 1

#### Outline

- Designing an Agent
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  - Problem formulation

#### Search

- Uninformed Search
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  - Heuristics

#### Review: Data structures

# ATM

Credit: freepik.com

Queue
First-in-first-out (FIFO)



Credit: freepik.com

Priority Queue
High-priority-first-out

#### Required operations

- Add element ("add")
- Remove element ("pop")
- o Is empty



Credit: iStock

Stack
Last-in-first-out (LIFO)

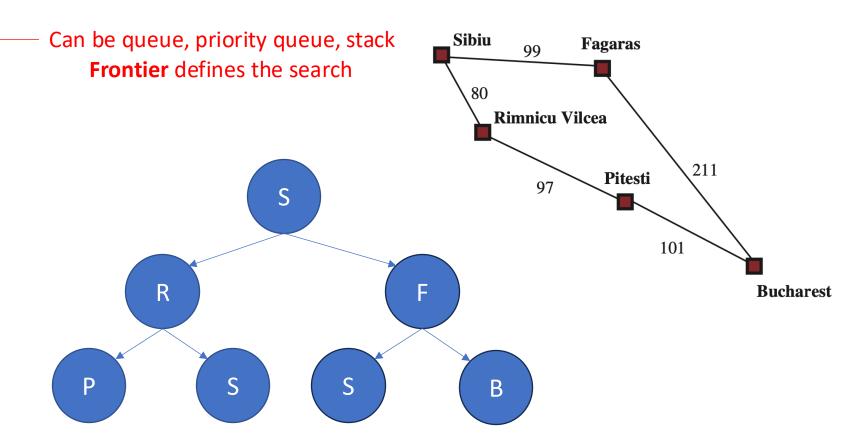
## Search AIAMA: tree search

Goal Formulation Problem Formula

Search

Execute

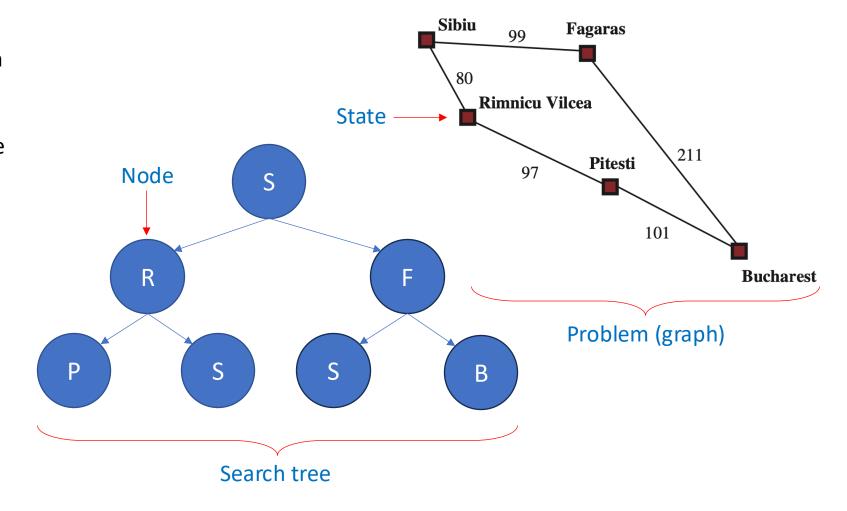
create frontier insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure



### Search Terms

Path cost is the cost of a path from any state to any state

Optimal path cost is the cost of the lowest-cost path from any state to any state



#### **Evaluation Criteria**

#### (Worst-case) Complexity:

- Time complexity: number of nodes generated or expanded
- Space complexity: maximum number of nodes in memory
- Measure: branching factor, optimal depth, maximum depth, ...

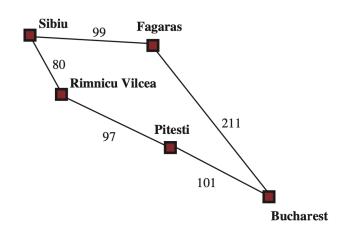
Disclaimer: CS2109S is not an algorithm-oriented course, so we will not discuss this in detail

#### Others:

- Completeness
- Optimality

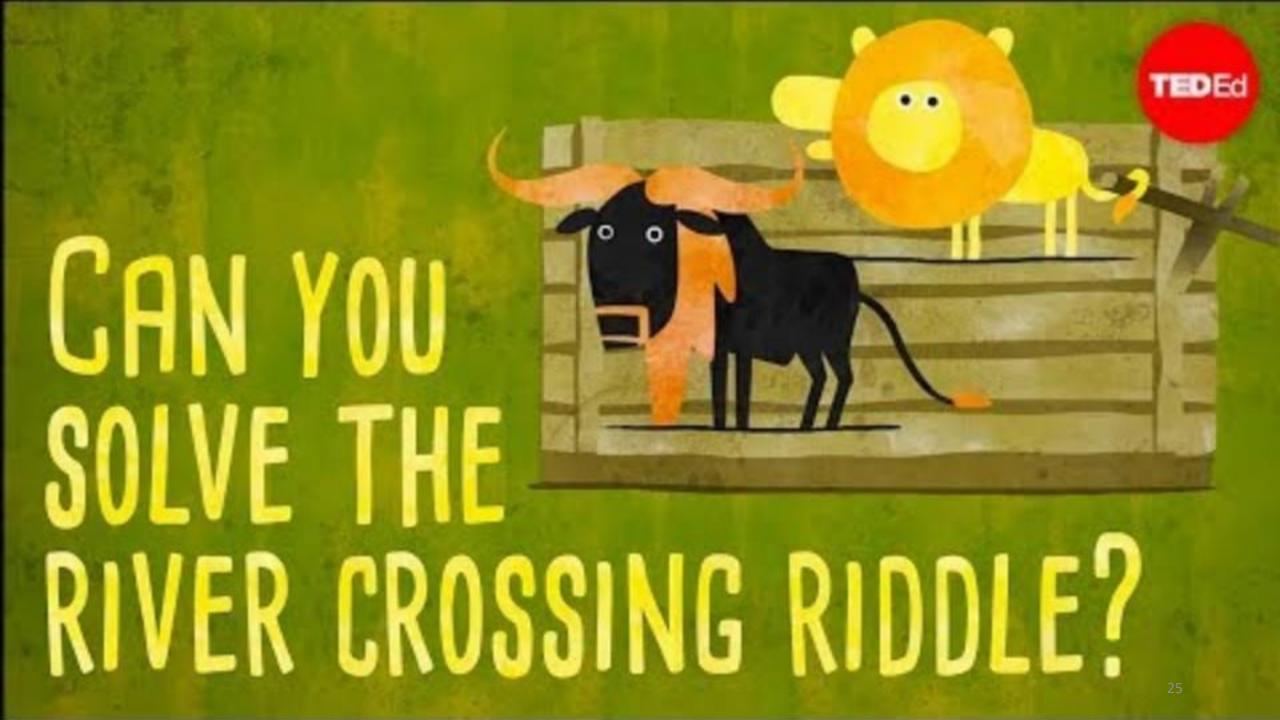
## Completeness and Optimality

- **Complete**: An algorithm is complete if, for every problem instance, it will find a solution if one exists. Conversely, an algorithm is incomplete if there exists at least one instance with a solution that the algorithm fails to find.
- **Optimal**: An algorithm is optimal if, for every instance where it produces a solution, that solution is the best possible. In other words, if the algorithm outputs a solution, it is guaranteed to be optimal.
- **Optimal and Incomplete**: An algorithm can be both optimal and incomplete. This means that there exists at least one instance where the algorithm fails to find a solution, even though, when it does find a solution, that solution is guaranteed to be optimal.



```
def search(state):
   if state == "Fagaras":
     return "go to Bucharest"
   return None
```

## Break

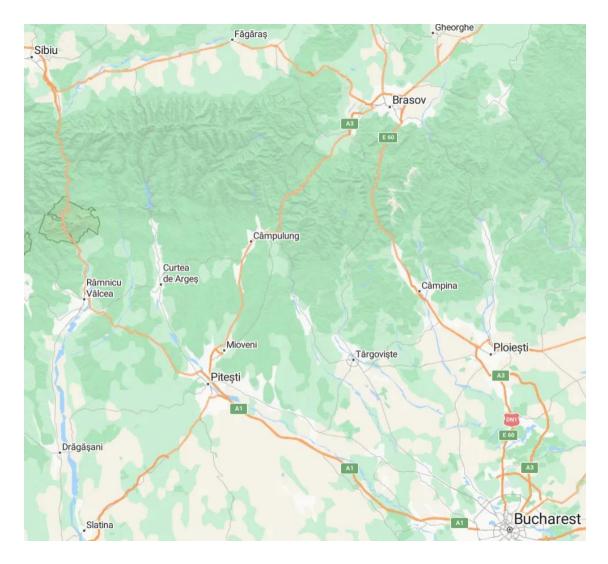


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## Uninformed Search Algorithms

No information that could guide the search



#### **Blind search**

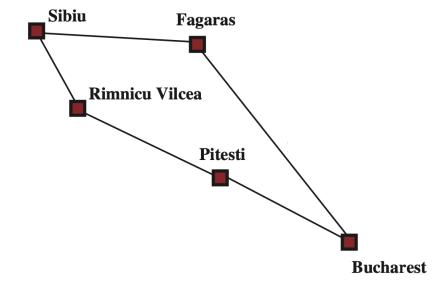
No clue how good a state is, e.g., how close it is to the goal

## Search

AIAMA: tree search

```
create frontier
insert Node(initial_state) to frontier
while frontier is not empty:
  node = frontier.pop()
  if node.state is goal: return solution
  for action in actions(node.state):
    next_state = transition(node.state, action)
    frontier.add(Node(next_state))
return failure
```

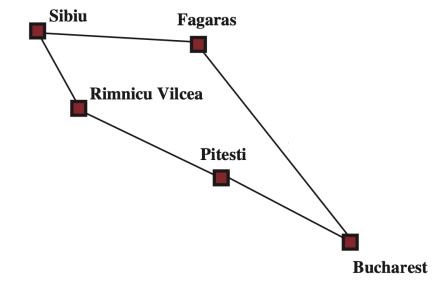
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create frontier : queue
insert Node(initial_state) to frontier
while frontier is not empty:
  node = frontier.pop()
  if node.state is goal: return solution
  for action in actions(node.state):
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    frontier.add(Node(next_state))
return failure
```



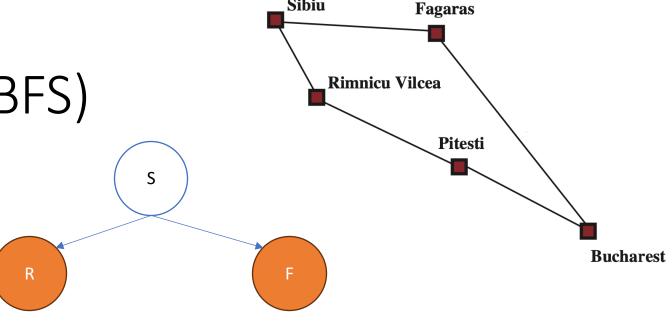
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Queue:

S

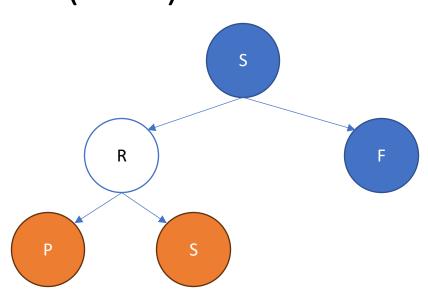


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Sibiu

create frontier : queue insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure



Sibiu

**Fagaras** 

**Pitesti** 

Rimnicu Vilcea

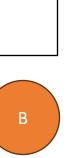


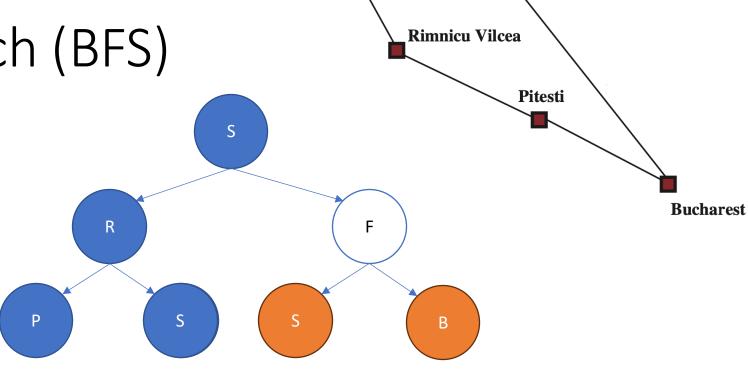
R F P S

**Bucharest** 

create frontier : queue insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure

Queue:

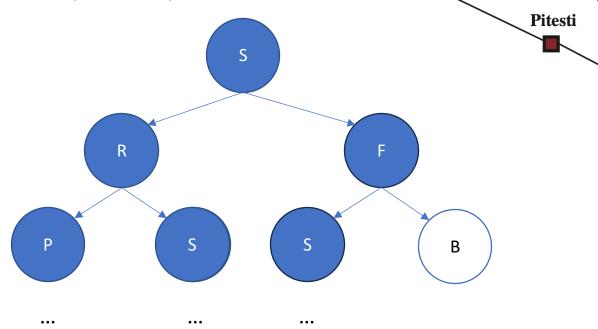




Sibiu

**Fagaras** 

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Sibiu

**Fagaras** 

Rimnicu Vilcea

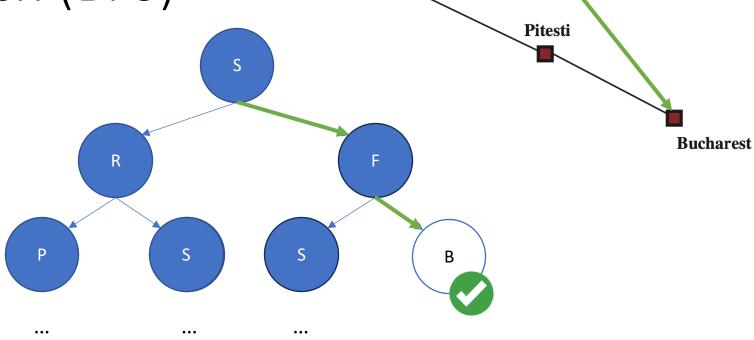
#### Queue:

•••

В

**Bucharest** 

create frontier : queue insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure



Sibiu

**Fagaras** 

Rimnicu Vilcea

#### Queue:

В ) …

## Breadth-first Search (BFS) - Analysis

```
create frontier : queue
insert Node(initial state) to frontier
while frontier is not empty:
  node = frontier.pop()
  if node.state is goal: return solution
  for action in actions(node.state):
    next state = transition(node.state, action)
    frontier.add(Node(next state))
return failure
```

Time complexity: (# nodes generated)

Exponential w.r.t. the depth of the optimal solution

**Space complexity: (size of frontier)** 

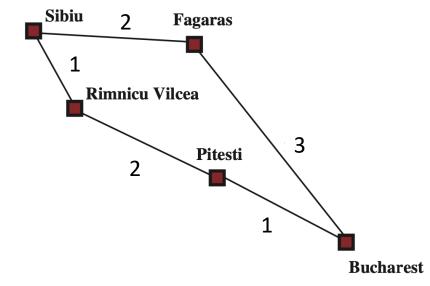
**Exponential** w.r.t. the depth of the optimal solution

**Complete:** Yes, if b is finite

**Optimal:** Yes, if step cost is the same everywhere

Cost from root to a state

```
create frontier: priority queue (path cost)
insert Node(initial_state) to frontier
while frontier is not empty:
  node = frontier.pop()
  if node.state is goal: return solution
  for action in actions(node.state):
    next_state = transition(node.state, action)
    frontier.add(Node(next_state))
return failure
```

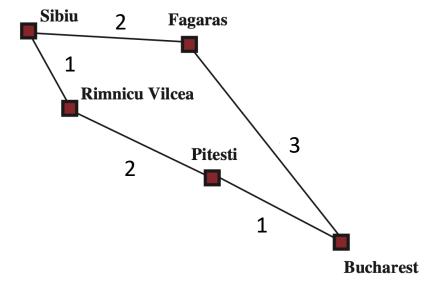


Cost from root to a state

create frontier: priority queue (path cost) insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure

#### **Priority Queue:**

S



Cost from root to a state

create frontier: priority queue (path cost)

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

if node.state is goal: return solution

for action in actions(node.state):

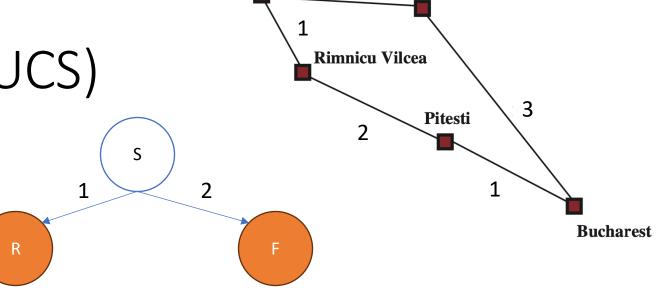
next\_state = transition(node.state, action)

frontier.add(Node(next\_state))

return failure

#### **Priority Queue:**

S R F



**Fagaras** 

Sibiu

Cost from root to a state

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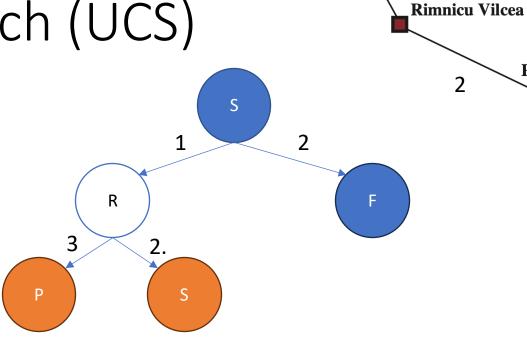
next\_state = transition(node.state, action)

frontier.add(Node(next\_state))

return failure

#### **Priority Queue:**

R F S P 2 3



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**Fagaras** 

**Pitesti** 

3

**Bucharest** 

Cost from root to a state

create frontier: priority queue (path cost)

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

if node.state is goal: return solution

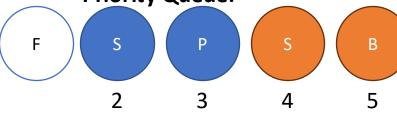
for action in actions(node.state):

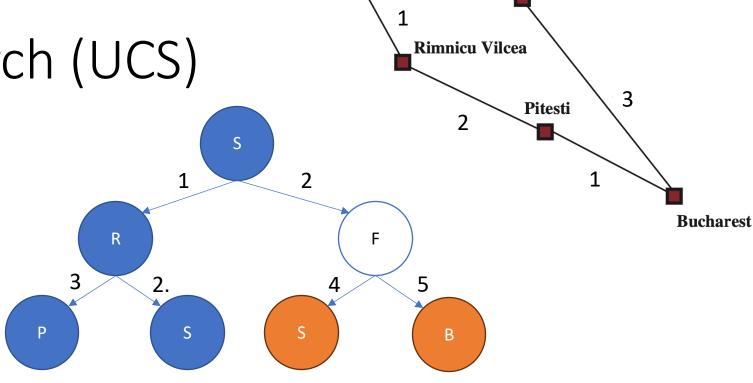
next\_state = transition(node.state, action)

frontier.add(Node(next\_state))

return failure







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**Fagaras** 

Cost from root to a state

create frontier: priority queue (path cost)

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

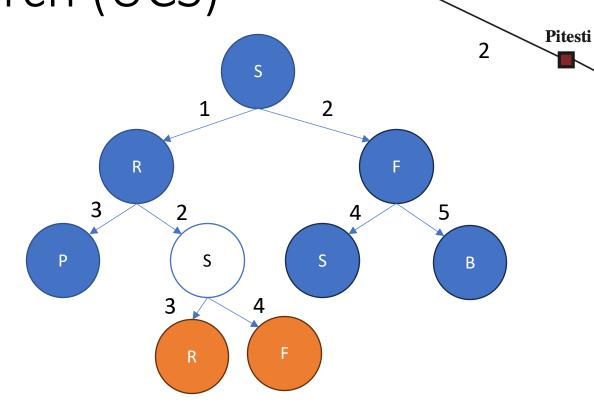
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frontier.add(Node(next\_state))

return failure



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**Fagaras** 

3

**Bucharest** 

Rimnicu Vilcea

#### **Priority Queue:**

S P R F S B 3 3 4 4 5

Cost from root to a state

create frontier: priority queue (path cost)

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

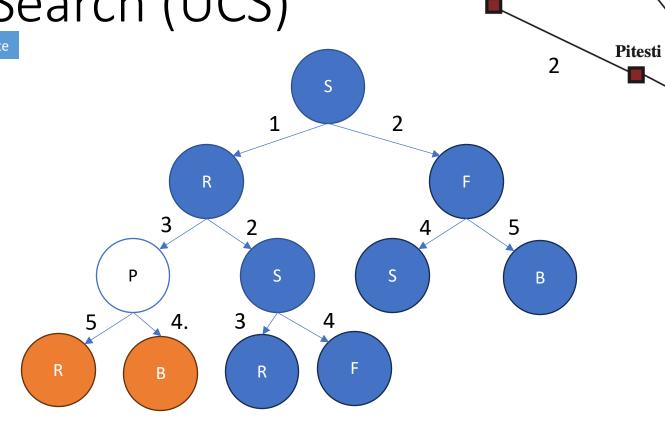
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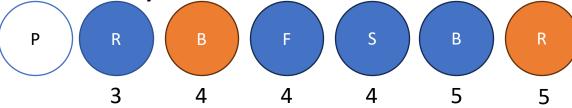
**Fagaras** 

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**Bucharest** 

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#### **Priority Queue:**



create frontier: priority queue (path cost)

insert Node(initial state) to frontier

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node = frontier.pop()

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for action in actions(node.state):

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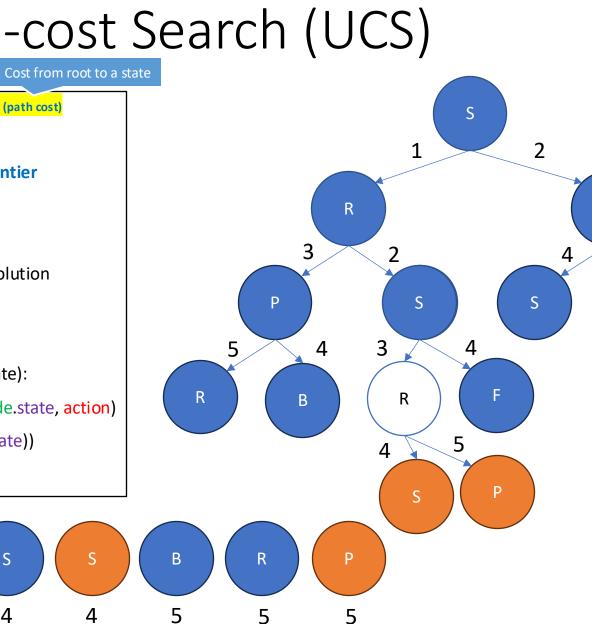
frontier.add(Node(next\_state))

**Priority Queue:** 

return failure

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**Fagaras** 

**Pitesti** 

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**Bucharest** 

Cost from root to a state

create frontier: priority queue (path cost)

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

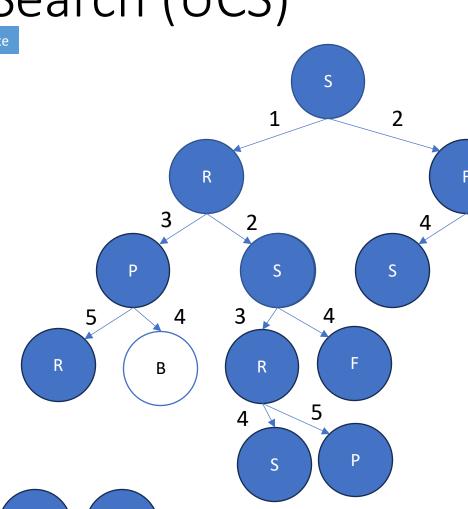
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frontier.add(Node(next\_state))

return failure



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**Fagaras** 

**Pitesti** 

3

**Bucharest** 

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**Priority Queue:** 

B F S S B R P 4 4 4 5 5 5 5

create frontier: priority queue (path cost)

insert Node(initial state) to frontier

while **frontier** is not empty:

node = frontier.pop()

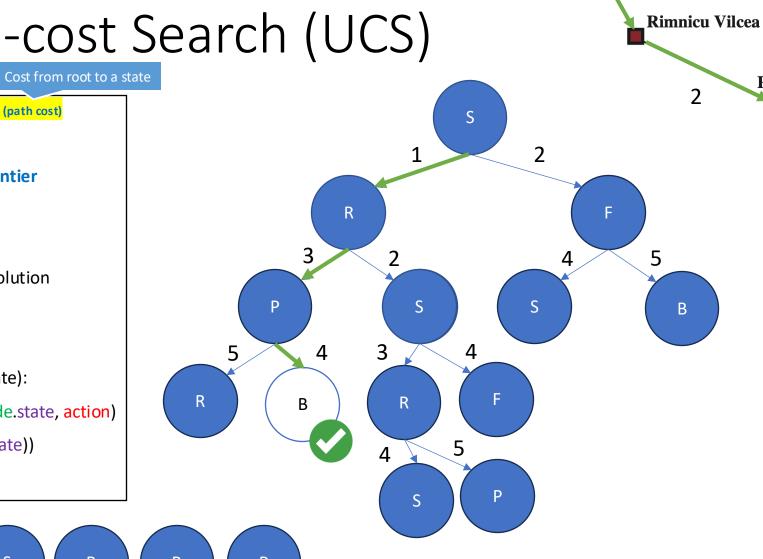
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Sibiu

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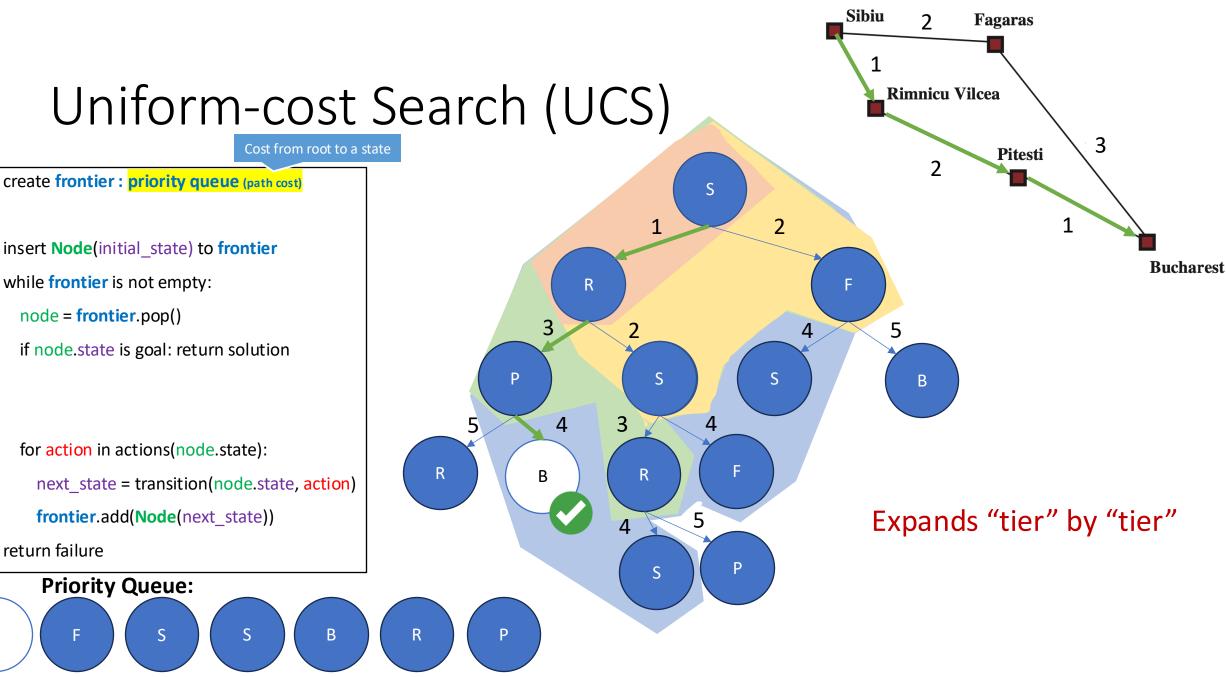
**Pitesti** 

3

**Bucharest** 

**Priority Queue:** 

В 5 5



В

## Uniform-cost Search (UCS) - Analysis

Cost from root to a state

```
create frontier: priority queue (path cost)
```

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

if node.state is goal: return solution

for action in actions(node.state):

next state = transition(node.state, action)

frontier.add(Node(next\_state))

return failure

Time complexity: (# nodes generated)

Exponential w.r.t. the "tier" of the optimal solution

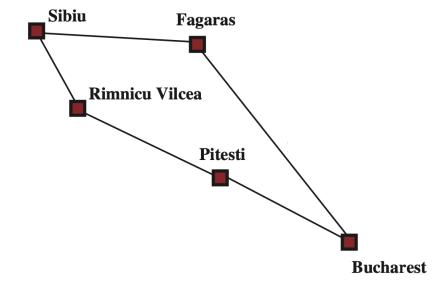
**Space complexity: (size of frontier)** 

**Exponential** w.r.t. the "tier" of the optimal solution

**Complete:** Yes, if positive step cost everywhere and finite total cost

**Optimal:** Yes, if positive step cost everywhere

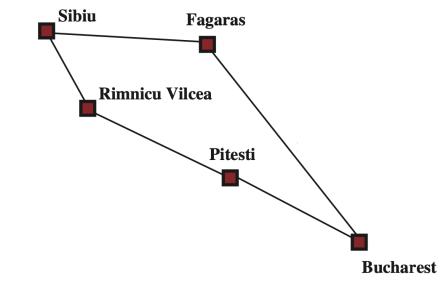
```
create frontier: stack
insert Node(initial_state) to frontier
while frontier is not empty:
  node = frontier.pop()
  if node.state is goal: return solution
  for action in actions(node.state):
    next_state = transition(node.state, action)
    frontier.add(Node(next_state))
return failure
```



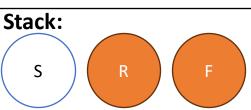
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```

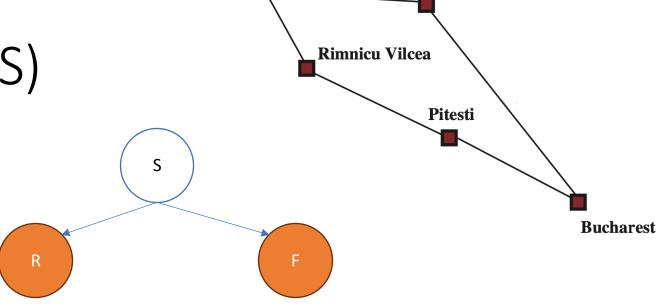
Stack:

S



create frontier: stack insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure

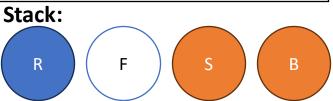


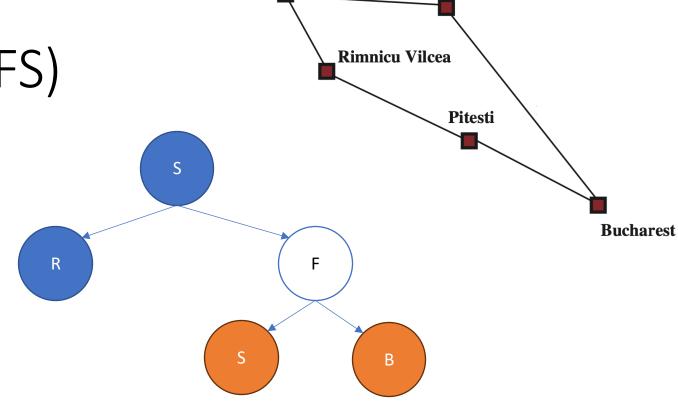


**Fagaras** 

Sibiu

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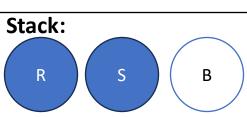


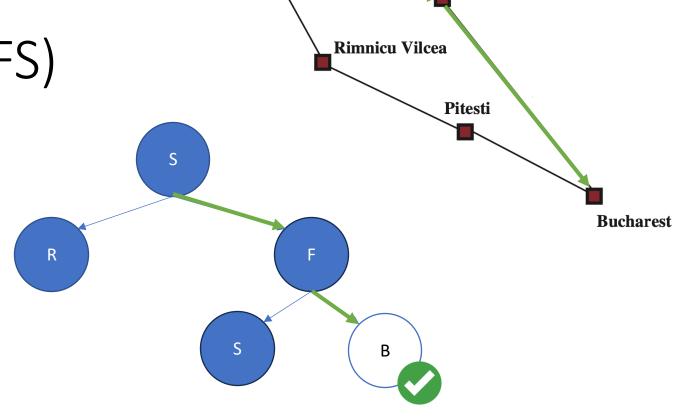


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**Fagaras** 

create frontier: stack insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure





Sibiu

**Fagaras** 

## Depth-first Search (DFS) - Analysis

```
create frontier: stack
insert Node(initial state) to frontier
while frontier is not empty:
  node = frontier.pop()
  if node.state is goal: return solution
  for action in actions(node.state):
    next state = transition(node.state, action)
    frontier.add(Node(next state))
return failure
```

Time complexity: (# nodes generated)

**Exponential** w.r.t. the maximum depth of the search tree

**Space complexity: (size of frontier)** 

Polynomial w.r.t. the maximum depth of the search tree

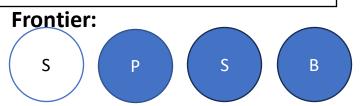
**Complete:** No, when the depth of the search tree is infinite (e.g., when the action is reversible, causing the state to go back and forth)

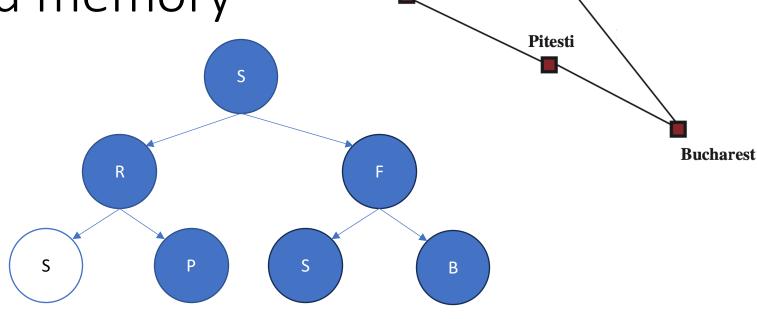
Optimal: No, the optimal solution may be in a shallower depth

## Search with visited memory

AIAMA: graph search

create frontier create visited insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution if node.state in visited: continue visited.add(state) for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure

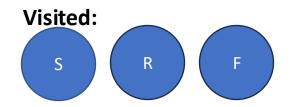




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**Fagaras** 

Rimnicu Vilcea



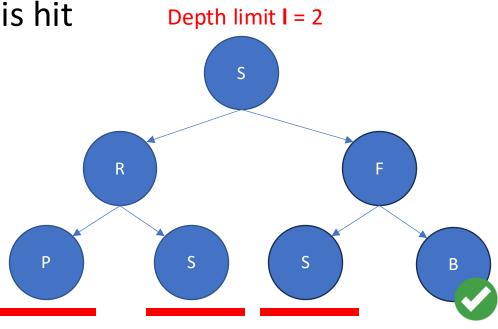
#### Outline

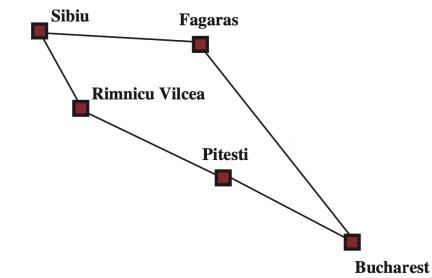
- Designing an Agent
  - Problem solving agents
  - Problem formulation
- Search
  - Uninformed Search
    - Breadth-first, depth-first, and uniform-cost search
    - Depth-limited and iterative-deepening search
  - Informed Search
    - A\* search
    - Heuristics

# Depth-limited Search (DLS)

Limit the search to depth l

Backtrack when limit is hit





# Depth-limited Search (DLS)

Limit the search to depth l Backtrack when limit is hit

Time complexity: (# nodes generated)

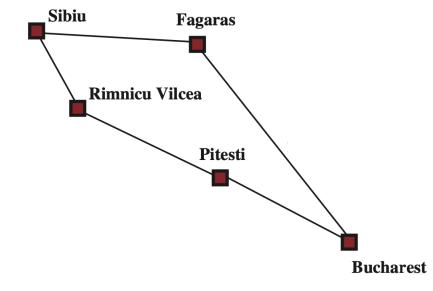
Exponential w.r.t. the depth limit

**Space complexity: (size of frontier)** 

*Polynomial* w.r.t. the depth limit if used with DFS

**Complete: No** 

**Optimal:** No, if used with DFS



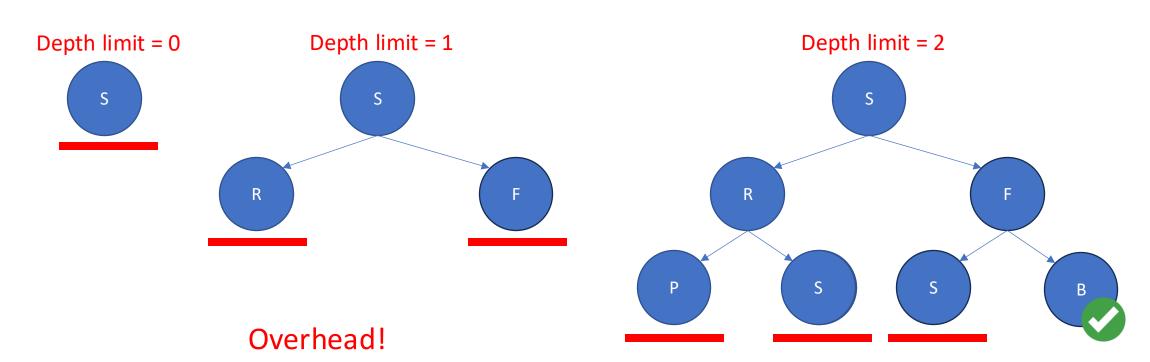
# Iterative Deepening Search (IDS)

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Search with depth limit  $0, ..., \infty$ Return solution if found



# Iterative Deepening Search (IDS)

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Pitesti

Bucharest

Search with depth limit  $0, ..., \infty$ Return solution if found

Time complexity: (# nodes generated)

Exponential w.r.t. the depth of optimal solution (with overhead)

**Space complexity: (size of frontier)** 

**Polynomial** if used with DFS

**Complete: Yes** 

**Optimal:** Yes, if step cost is the same everywhere

### Summary: Uninformed Search Algorithms

#### Search algorithms

- Breadth-first search queue, explore layer by layer
- Uniform-cost search priority queue (path cost)
- Depth-first search **stack**, go deep first then backtrack

#### • Variants:

- Depth limited search limit max depth of the search
- Iterative deepening search try DLS with depth limit  $0, ..., \infty$

## Summary: Uninformed Search Algorithms

#### Worst-case

Name	Time Complexity*	Space Complexity*	Complete?	Optimal?
Breadth-first Search	Exponential	Exponential	Yes	Yes
Uniform-cost Search	Exponential	Exponential	Yes	Yes
Depth-first Search	Exponential	Polynomial	No	No
Depth-limited Search	Exponential	Polynomial**	No**	No**
Iterative Deepening Search	Exponential	Polynomial**	Yes	Yes

<sup>\*)</sup> In terms of some notion of depth/tier

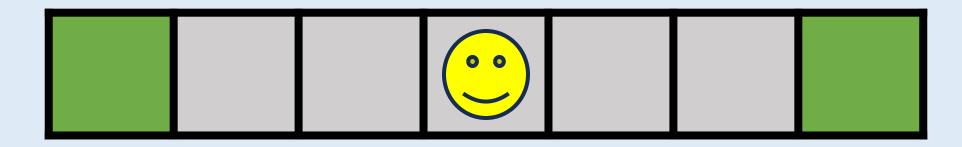
<sup>\*\*)</sup> If used with DFS

### Further Reading (Optional)

- Full analysis of BFS, UCS, DFS, DLS, and IDS (AIAMA 3.4)
- Backward Search (AIAMA 3.4.6)
- Bidirectional Search (AIAMA 3.4.6)

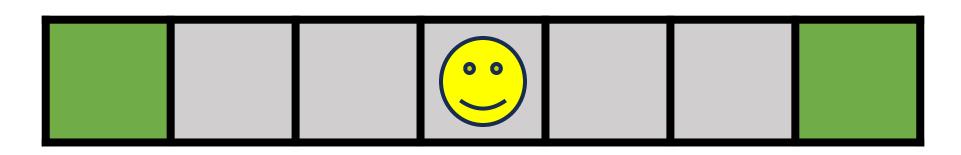
#### Activity:

Given a maze where the agent is in the middle and there are two goals at the ends, which search algorithm(s) with visited memory is/are optimal and the most space and time efficient (in worst-case)?



Note: we care about overhead in space and time complexities

Name	Time Complexity*	Space Complexity*	Complete?	Optimal?
Breadth-first Search	Exponential	Exponential	Yes	Yes
Uniform-cost Search	Exponential	Exponential	Yes	Yes
Depth-first Search	Exponential	Polynomial	No	No
Depth-limited Search	Exponential	Polynomial**	No**	No**
Iterative Deepening Search	Exponential	Polynomial**	Yes	Yes



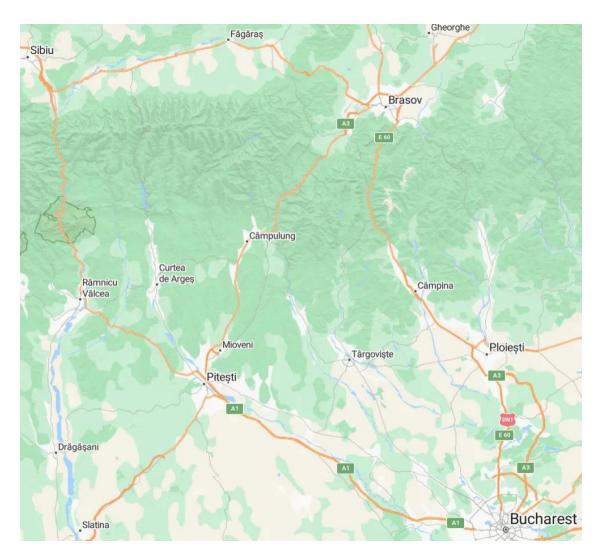
Note: we care about overhead in space and time complexities

#### Outline

- Designing an Agent
  - Problem solving agents
  - Problem formulation
- Search
  - Uninformed Search
    - Breadth-first, depth-first, and uniform-cost search
    - Depth-limited and iterative-deepening search
  - Informed Search
    - A\* search
    - Heuristics

### Uninformed Search Algorithms

No information that could guide the search

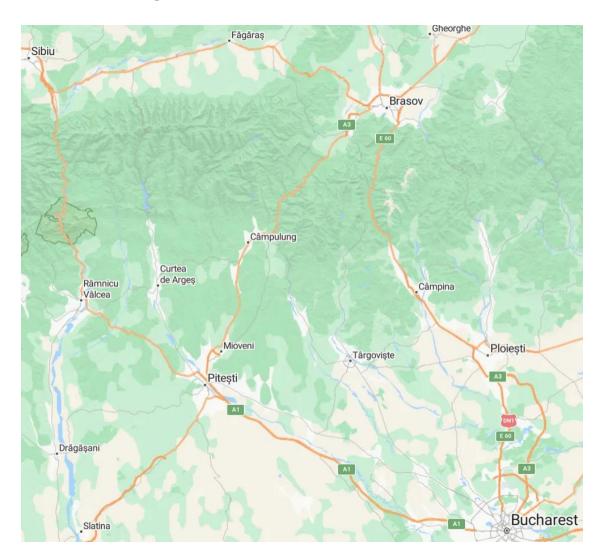


#### **Blind search**

No clue how good a state is, e.g., how close it is to the goal

#### Informed Search Algorithms

Use information to guide the search



#### Search with extra info

<u>Has</u> a clue how good a state is,e.g., how close it is to the goal, estimated using a heuristic function

#### Heuristic

- Recall that:
  - Path cost is the cost of a path from any state to any state
  - Optimal path cost is the cost of the lowest-cost path from any state to any state

A heuristic is an estimate of the optimal path cost from any state to the

goal state

- Examples (in path finding):
  - Straight-line distance  $h_{SLD}$
  - Manhattan distance  $h_{MD}$

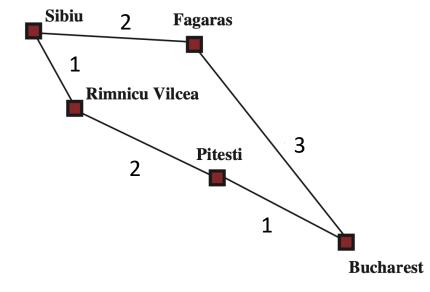


#### Best-first Search

**Evaluation function** 

f(n) estimates how good a state is

create frontier: priority queue (f(n)) insert Node(initial\_state) to frontier while **frontier** is not empty: node = frontier.pop() if node.state is goal: return solution for action in actions(node.state): next\_state = transition(node.state, action) frontier.add(Node(next\_state)) return failure



#### A\* Search

Evaluation functior

create frontier: priority queue (f(n))

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

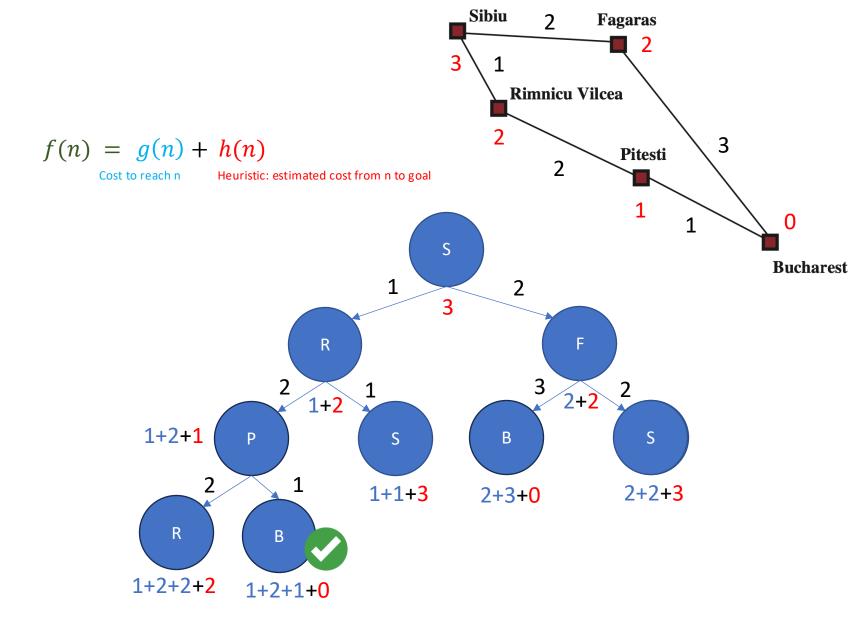
if node.state is goal: return solution

for action in actions(node.state):

next\_state = transition(node.state, action)

frontier.add(Node(next\_state))

return failure



## A\* Search - Analysis

**Evaluation function** 

create frontier: priority queue (f(n))

insert Node(initial\_state) to frontier

while **frontier** is not empty:

node = frontier.pop()

if node.state is goal: return solution

for action in actions(node.state):

next\_state = transition(node.state, action)

frontier.add(Node(next\_state))

return failure

$$f(n) = g(n) + h(n)$$

Cost to reach n

Heuristic: estimated cost from n to goal

Time complexity: (# nodes generated)

Exponential, good heuristic can improve

**Space complexity: (size of frontier)** 

Exponential

#### **Complete:**

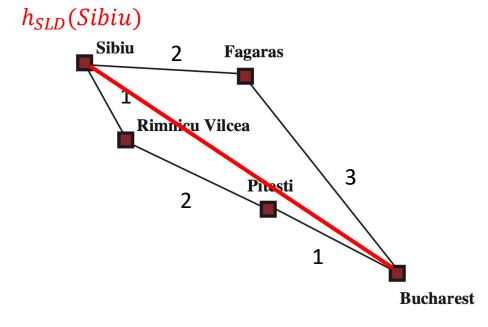
Yes, if edge costs are positive and branching factor is finite (Hart, Nilsson and Raphael, 1968)

**Optimal:** Depends on the heuristic

#### Admissible Heuristics

A heuristic h(n) is admissible if for every node n,  $h(n) \le h^*(n)$ , where  $h^*(n)$  is the **optimal path cost** to reach the goal state from n.

An admissible heuristic **never over-estimates** the cost to reach the goal, i.e., it is an **optimistic estimate**.

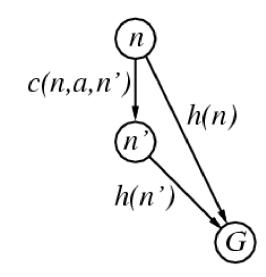


**Example:**  $h_{SLD}(n)$  never overestimates the actual road distance

**Theorem** (Hart, Nilsson and Raphael, 1968): if h(n) is admissible, A\* search (without visited memory) is optimal

#### Consistent Heuristics

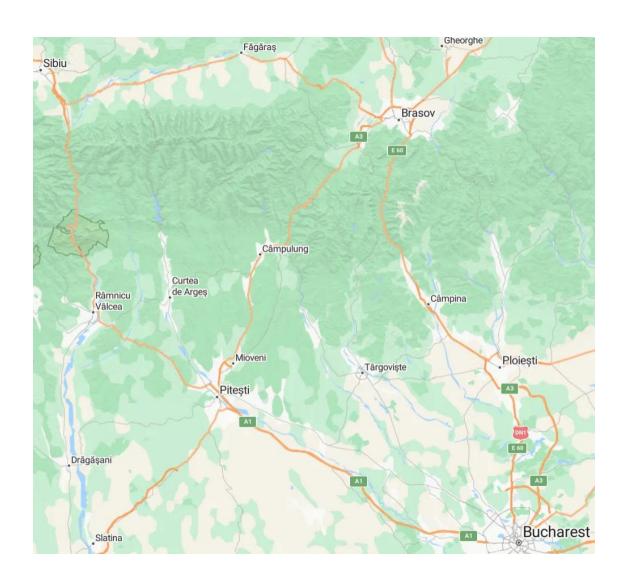
A heuristic h(n) is consistent if for every node n, every successor n' of n generated by any action a,  $h(n) \le c(n,a,n') + h(n')$ , and h(G) = 0.



Theorem (Hart, Nilsson and Raphael, 1968):

If h(n) is consistent, A\* search with visited memory is optimal

#### "Inventing" Admissible Heuristics



A problem with **fewer restrictions** on the actions is called a <u>relaxed problem</u>. The cost of an optimal solution to a relaxed problem is an **admissible heuristic** for the original problem.

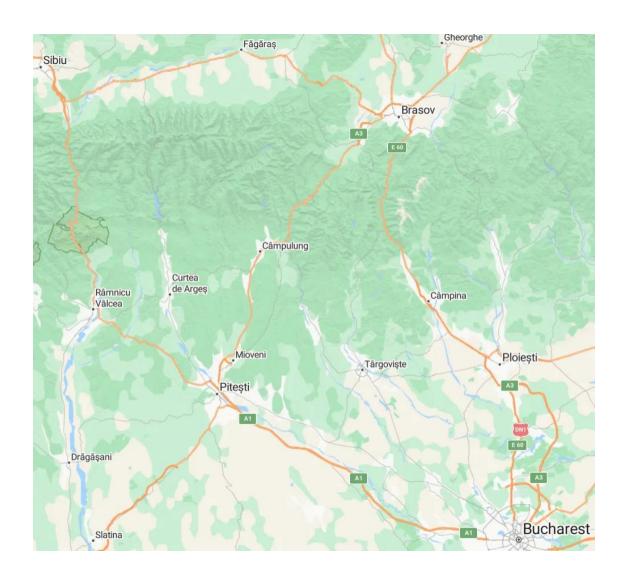
#### **Original:**

An agent can only move along the road

#### **Relaxations:**

- An agent can move off-road (fly)
  - $h_1$  straight line distance
- An agent can teleport
  - $h_2 = 1$

#### Dominance



If  $h_1(n) \ge h_2(n)$  for all n, then  $h_1$  dominates  $h_2$ .

•  $h_1$  is better for search if admissible.

#### **Heuristics**

- *h*<sub>1</sub> straight-line distance
- $h_2 = 1$

 $h_1$  dominates  $h_2$ 

#### Summary: Informed Search Algorithms

- Informed search: guide search with extra information
- Best-first search
  - A\* search
    - f(n) = g(n) + h(n), cost so far + heuristic
- Heuristics
  - Admissible:  $h(n) \le h^*(n)$
  - Consistent:  $h(n) \le c(n,a,n') + h(n')$
  - Dominant: if  $h1(n) \le h2(n)$ , h2 dominant
- Creating admissible heuristic: optimal path cost of the relaxed problem

## Further Reading (Optional)

(Hart, Nilsson and Raphael, 1968)

- Proof: admissible heuristic → optimal search without visited memory
- Proof: consistent heuristic → optimal search with visited memory (AIAMA 3.5.2)
- Greedy best-first search (AIAMA 3.5.1)
- Variants of A\* search (AIAMA 3.5.3)
  - Iterative Deepening A\* (IDA\*)
  - Simplified Memory-bounded A\* (SMA\*)

#### Summary

- Problem solving agents:
  - Goal formulation -> problem formulation -> search -> execute
- Uninformed search
  - Breadth-first search (BFS)
  - Uniform-cost Search (UCS)
  - Depth-first Search (DFS)
  - Depth-limited & Iterative deepening search (DLS & IDS)
- Informed search
  - A\* Search
  - Heuristics: admissible, consistent, dominant, relaxed problem

## Coming Up Next Week

- Local search
  - Hill climbing
  - ...
- Adversarial search
  - Games
  - Minimax
  - Alpha-beta pruning
  - •

#### To Do

- Lecture Training 2
  - +250 EXP
  - +100 Early bird bonus
- Problem Set 1

Will be released later today!