

# Virtual Machine 2.0 Specification

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Batburger

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## 1 Definitions

**Value** A *value* shall be a 32bit number.

**Stack** The *stack* is a data structure, from which one may only add a *value* or may read, and than always remove, the lastly added *value*.<sup>1</sup>

**Pushing** *pushing* refers to the act of adding a new value onto the stack.

**Poping** *poping* refers to the act of reading the lastly pushed *value* from the stack and removing it.

**Register** A *register* stores one *value* at a time. One can push the stored *value* onto the *stack* or pop a value from the *stack* and store the popped value in the register.

**Address** An *address* shall be a 32bit number which describes the location of stored data.

**Linear Memory** *linear memory* refers to a data structure where one can write arbitrary many *values* to an arbitrary *address*<sup>2 3</sup>. One can also push any *value* stored in the *linear memory* on to the stack.

**State** The *state* of the vm refers to the data which is stored in the *linear memory*, *registers*, *instruction pointer* and *stack*.

**Instruction** An *instruction* shall be callable and modify the state of the vm, when called.

**ByteCode** The *bytecode* is an array of 8bit numbers which refer to instructions or static *values*<sup>4</sup>.

**Opcode** An *opcode* is a 8bit number<sup>5</sup>, which uniquely defines an *instruction*.

**Argument** A *argument* shall be a *value* which is stored on the *stack* and gets popped by an instruction.<sup>6</sup>

**Instruction Pointer** The *instruction pointer* is a pointer which points into the *bytecode*.

**Float** A *float* refers to single-precision float as defined by the IEEE Std 754-2008 - IEEE STANDARD FOR FLOATING-POINT ARITHMETIC<sup>7</sup>

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<sup>1</sup>We may refer to the *value* which was added lastly as the *top value*.

<sup>2</sup>Address are numbered starting from 0.

<sup>3</sup>The *linear memory's address's* are in no way related to the *address's* of other data structure.

<sup>4</sup>4 8bit values are read and sticked together to build one 32bit number.

<sup>5</sup>Usually represented in hex.

<sup>6</sup>We say the popped *value* is the *argument* of the instruction popping it.

<sup>7</sup><http://standards.ieee.org/findstds/standard/754-2008.html>

## 2 High Level Function

The vm keeps track of the Instruction Pointer (IP). The IP points into the bytecode starting at the first byte. The vm functions in a loop, executing the instruction which is defined by the opcode, to which the IP is currently pointing. Every instruction may or may not modify the state of the vm, but must always modify the ip. The vm stops executing when the ip points to the instruction 0x11.

## 3 Technical notes

Every instruction is defined to be one byte in size. There are separate instructions for the different number types, but values are not strictly typed; one instruction could interpret a value to be unsigned, and the next instruction interprets the same value (at the same address) to be a float, for instance. The VM only exposes the values stored on the stack to the instructions, but internally the vm stores, besides the value, what instruction pushed that value onto the stack. (This is required so the *return* instruction knows what value was pushed onto the stack by the *call* instruction.)

## 4 Input Parameters

Any command line arguments added to the VM executable after the file to execute<sup>8</sup> will be pushed onto the stack in order of appearance.

## 5 Instruction Set

We use "..." to specify a range of numbers. The first and last values are both included in the range.

Opcode	Name	Description
0xa0...0xa9	readRegister0...readRegister9	Pushes the value stored in a register onto the stack.
0xb0...0xb9	setRegister0...setRegister9	Writes one argument to the given register.
0xc0	setSize	Takes one argument, sets by how many multiples of 32 the linear memory gets expanded when calling 0xc3 (alloc) <sup>9</sup> . YOU CAN CALL THIS INSTRUCTION EXACTLY ONCE.
0xc1	move	Writes the first argument into linear memory at the address given by the second argument.
0xc2	read	Pushes the value, specified by the address given as the first argument, from linear memory onto the stack.
0xc3	alloc	Expands the writeable address range in the linear memory by the size that was set with the instruction 0xc0 a.k.a. setSize. <sup>10</sup>
0xd0	push	Pushes a value, specified by the next 4 bytes in the bytecode, onto the stack.
0xd1	remove	Removes a value from the stack.

<sup>8</sup>e.g. the number 1024 in "vm ./path/to/my/file 1024".

<sup>9</sup>Addresses are 32bit in size, therefore if you were to set this value to 32 you would, every time you call alloc, increase the linear memory by 1024bit (1 Kibibit).

<sup>10</sup>If you want to write to linear memory, you have to call this instruction at least once; there would be no address to which you could write to otherwise.

0xe0	uadd	Adds two arguments and pushes the result onto the stack. (All used values are interpreted to be unsigned)
0xe1	sadd	Adds two arguments and pushes the result onto the stack. (All used values are interpreted to be signed)
0xe2	fadd	Adds two arguments and pushes the result onto the stack. (All used values are interpreted to be floats)
0xe3	usub	Subtracts two arguments and pushes the result onto the stack. (All used values are interpreted to be unsigned)
0xe4	ssub	Subtracts two arguments and pushes the result onto the stack. (All used values are interpreted to be signed)
0xe5	fsub	Subtracts two arguments and pushes the result onto the stack. (All used values are interpreted to be floats)
0xe6	umult	Multiplies two arguments and pushes the result onto the stack. (All used values are interpreted to be unsigned)
0xe7	smult	Multiplies two arguments and pushes the result onto the stack. (All used values are interpreted to be signed)
0xe8	fmult	Multiplies two arguments and pushes the result onto the stack. (All used values are interpreted to be floats)
0xe9	udiv	Divides two arguments and pushes the result onto the stack. (All used values are interpreted to be unsigned)
0xea	sdiv	Divides two arguments and pushes the result onto the stack. (All used values are interpreted to be signed)
0xeb	fdiv	Divides two arguments and pushes the result onto the stack. (All used values are interpreted to be floats)
0xec	utof	One argument is transformed to be stored in float representation. <sup>11</sup> (Value is interpreted to be unsigned)
0xed	stof	One argument is transformed to be stored in float representation. <sup>11</sup> (Value is interpreted to be signed)
0xee	abs	One argument is turned into a positive number.
0xef	ucmp	Compares two argument; then pushes a value onto the stack, which represents if $arg1 < arg2$ , or if $arg1 > arg2$ , or if $arg1 = arg2$ . <sup>12</sup> (All values are interpreted to be unsigned)
0xf0	scmp	Compares two argument; then pushes a value onto the stack, which represents if $arg1 < arg2$ , or if $arg1 > arg2$ , or if $arg1 = arg2$ . <sup>12</sup> (All values are interpreted to be signed)

<sup>11</sup> The instruction will store the number represented in its first argument with as much precision as a float allows. If precision is lost, no further action is taken.

<sup>12</sup> 0 represents  $arg1 = arg2$ , 1 represents  $arg1 < arg2$ , 2 represents  $arg1 > arg2$ .

0xf1	fcmp	Compares two argument; then pushes a value onto the stack, which represents if $\text{arg1} < \text{arg2}$ , or if $\text{arg1} > \text{arg2}$ , or if $\text{arg1} = \text{arg2}$ . <sup>12</sup> (All values are interpreted to be floats)
0x01	jmp	Sets the ip to the address given by the first argument of this instruction.
0x02	jless	Functions like the <i>jmp</i> instruction if, and only if, the last <i>cmp</i> that $\text{arg1} < \text{arg2}$ (Arguments of <i>cmp</i> ).
0x03	jgreater	Functions like the <i>jmp</i> instruction if, and only if, the last <i>cmp</i> detriment that $\text{arg1} > \text{arg2}$ (Arguments of <i>cmp</i> ).
0x04	jequal	Functions like the <i>jmp</i> instruction if, and only if, the last <i>cmp</i> detriment that $\text{arg1} = \text{arg2}$ (Arguments of <i>cmp</i> ).
0x05	jNequal	Functions like the <i>jmp</i> instruction if, and only if, the last <i>cmp</i> detriment that $\text{arg1} \neq \text{arg2}$ (Arguments of <i>cmp</i> ).
0x06	call	Sets the ip to the address specified in its first argument and pushes the address of this instruction onto the stack.
0x07	return	Looks through the stack, top to bottom, and jumps to the address which was put there by a previous <i>call</i> instruction.
0x10	int	Triggers the interrupt specified by its first argument. <sup>13 14</sup>
0x11	halt	Stops the vm.
0x12	noop	Iterates the IP by one with no other changes to the state of the vm.
0x13	unsignedOut	Prints one argument to standard out (The Value is written in unsigned decibel form.).
0x14	signedOut	Prints one argument to standard out (The Value is written in signed decibel form.).
0x15	floatOut	Prints one argument to standard out (The Value is written as float.).
0x16	charOut	Prints one argument to standard out (The Value is interpreted as a char.).

Table 1: Instruction Set

## 6 Interrupts

<sup>13</sup>See Section 6

<sup>14</sup>It is not sure which, if any, interrupts will exist.