

What does $O(< expr >)$ mean?

1

What does $\Theta(< expr >)$ mean?

2

What does $\Omega(< expr >)$ mean?

3

*What are the best, average and worst case complexities of **Bubble Sort**?*

4

*What are the best, average and worst case complexities of **Merge Sort**?*

5

Give pseudo code for merging 2 sorted lists, as part of merge sort.

6

Give pseudo code for MergeSort(L).

7

*What are the best, average and worst case complexities of **Quick Sort**?*

8

The complexity (i.e. space/running time) has the complexity proportional to $\langle \text{expr} \rangle$.

The complexity (i.e. running time/space) is bounded by the $\langle \text{expr} \rangle$.

2

1

Best: $O(n)$,
Average: $O(n^2)$,
Worst: $O(n^2)$

The complexity (i.e. running time/space) is at least by the $\langle \text{expr} \rangle$.

4

3

Merge(L_1, L_2)
 if $L_1 = []$ return L_2
 if $L_2 = []$ return L_1
 $x_1 = L_1[0]$
 $x_2 = L_2[0]$
 $L'_1 = L_1[1 : |L_1| - 1]$
 $L'_2 = L_2[1 : |L_2| - 1]$
 if $x_1 \leq x_2$
 return $[x_1] + \text{Merge}(L'_1, L_2)$
 return $[x_2] + \text{Merge}(L_1, L'_2)$
 Merge two sorted lists

Best: $O(n \log \log_2 n)$,
Average: $O(n \log \log_2 n)$,
Worst: $O(n \log \log_2 n)$

6

5

Best: $O(n \log \log_2 n)$,
Average: $O(n \log \log_2 n)$,
Worst: $O(n^2)$

MergeSort(L)
 if $|L| \leq 1$
 return L
 Split L into roughly equal halves, L_l and L_r
 return *Merge*(*MergeSort*(L_l), *MergeSort*(L_r))

MergeSort(L)

8

7

<p><i>What would the pseudo code be for Quick Sort?</i></p> <p>9</p>	<p><i>Say that the input represents a positive integer, x, what is the size of n?</i></p> <p>10</p>
<p><i>What does it mean by $O(1)$?</i></p> <p>11</p>	<p><i>What is the minimum time for any sorting algorithm that uses only number comparisons?</i></p> <p>12</p>
<p><i>What would the pseudo code be for Euclid's algorithm?</i></p> <p>13</p>	<p><i>What would the pseudo code be for Fast Modular Exponentiation?</i></p> <p>14</p>
<p><i>Consider the equation $a^x = y \bmod p$. If a is a primitive root modulo p, then for every $y(1 \leq y < p)$, such an $x(1 \leq x < p)$ exists. What is x?</i></p> <p>15</p>	<p><i>The is the inverse of exponentiation.</i></p> <p>16</p>

$\lfloor \log_b x \rfloor + 1$ Where b is the number representation,
usually binary (so 2).

```
quicksort(L)
  if length of L ≤ 1
    return L
  remove the first element, x, from L
  L≤ := elements of L less than or equal to x
  L> := elements of L greater than x
  Ll := quicksort(L≤)
  Lr := quicksort(L>)
  return Ll + [x] + Lr
```

Quick Sort

10

9

$n \log_2 n$

*It takes a constant time, no matter the amount of
data, to perform the operation.*

12

11

```
fme(a,b,k)
  d = a
  e = b
  s = 1
  While e > 0
    if e is odd
      s = (s.d)modk
      d = d2modk
      e = ⌊e/2⌋
  return s
```

Fast Modular Exponentiation

```
// Assume a ≥ b
hcf(a,b)
  if b = 0
    return a
  r = a mod b
  return hcf(b,r)
```

Euclid's algorithm

14

13

*The discrete logarithm is the inverse of
exponentiation.*

*X is the **discrete logarithm** of y with base a,
modulo p.*

16

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Why can the private key not, in practice, be recovered from the public key when p is large?

17

What is one way you can argue correctness of Euclid's algorithm?

18

What would half the correctness proof be for Euclid's algorithm?

19

$(a.b) \bmod k =$

20

Let p be a prime number. What is meant by a primitive root modulo p ?

21

What does saying that algorithm A runs in time g mean?

22

What is a permutation of a set?

23

What do we mean by a composition of two permutations?

24

Let $r = a \bmod b$. $\text{hcf}(a, b) = \text{hcf}(b, r)$ because all factors of a and b are also factors of b and r and vice versa. If they have the same factors, they have the same highest common factor.

18

To calculate a public key, y , a private key, x is needed. The equation for modular exponentiation can be used: $y = g^x \bmod p$

It is considered a one-way function - easy to compute, hard to invert. For a large p , the only way to figure out the private key would be to use brute force, which would take a large amount of time.

17

$$(a.b) \bmod k = (a \bmod k . b \bmod k) \bmod k$$

As $r = a \bmod b$, $\exists q$ such that $a = bq + r$, $\therefore r = a - bq$. Suppose x is a factor of a and b , then $\exists y$ and z such that $a = xy$, $b = xz$.

Hence: $r = xy - xzq$, $r = x(y - zq)$.
 $\therefore x$ is a factor of r (and also of b and r).

20

19

Given an input of size n , the number of operations executed by A is bounded above by $g(n)$.

The numbers r_x between 1 and $p - 1$ that, when raised by the numbers between 1 and $p - 1$ compute all the numbers between 1 and $p - 1$ in some order with no repetitions.

22

21

The composition is the product of two permutations, α and β , on a set n , given by $\alpha \cdot \beta(n)$ or $\beta(\alpha(n))$

A 1-to-1 map of the set onto itself. In basic terms, it is a set mapped to another order of itself. i.e
 $[0, 1, 2, 3, 4] \mapsto [2, 4, 1, 0, 3]$

24

23

What is the number of possible permutations on an n -element set?

25

In the context of a permutation, what do we mean by a transposition?

26

Convert this pair of simultaneous equations into matrix form

$$a_{1,1}x_1 + a_{1,2}x_2 = b_1$$

$$a_{2,1}x_1 + a_{2,2}x_2 = b_2$$

27

What is the determinant of the matrix:

$$\begin{pmatrix} a_1 & a_2 \\ a_3 & a_4 \end{pmatrix}$$

28

What is an upper triangular matrix and how do you calculate its determinant?

29

Which 4 operations have no effect on a matrix's determinant?

30

A transposition is a special kind of permutation where only 2 elements in a set are affected (they are swapped). On a set X a transposition $\sigma = (i, j)$ is given by

$$\sigma(k) = \begin{cases} j & \text{if } k = i \\ i & \text{if } k = j \\ k & \text{ow.} \end{cases}$$

$$n!$$

$$a_1a_4 - a_2a_3$$

Often denoted as:

$$\begin{vmatrix} a_1 & a_2 \\ a_3 & a_4 \end{vmatrix}$$

The original system of equations to which the matrix corresponds only has a unique solution if the determinant is non-zero.

$$\begin{pmatrix} a_{1,1} & a_{1,2} \\ a_{2,1} & a_{2,2} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$$

It is a matrix where all of its entries below the diagonal are zero.

Transposing two rows
Transposing two columns

Adding a multiple of one row to another

Adding a multiple of one column to another

Also note that if all entries in any row or column are 0 then the determinant is 0

$$\begin{pmatrix} a_{1,1} & a_{1,2} & \cdots & a_{1,n} \\ 0 & a_{2,2} & \cdots & a_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & a_{n,n} \end{pmatrix}$$

Its determinant is calculated by taking the product of the entries on the diagonal. i.e $a_{1,1} \cdot a_{2,2} \cdot \dots \cdot a_{n,n}$