What does $O(\langle expr \rangle)$ mean?	What does $\Theta()$ mean?
What does $\Omega()$ mean?	Say that the input represents a positive integer, x , what is the size of n ?
What does it mean by $O(1)$?	What would the pseudo code be for Euclid's algorithm?
What would the pseudo code be for Fast Modular Exponentiation?	What are some of the advantages of ElGamal encryption?

The complexity (i.e. space/running time) has the complexity proportional to $\langle expr \rangle$.

The complexity (i.e. running time/space) is bounded by the < expr >.

2

1

 $\lfloor \log_b x \rfloor + 1$ Where b is the number representation, usually binary (so 2).

The complexity (i.e. running time/space) is at least by the $\langle expr \rangle$.

4

3

```
// Assume a>=b
hcf(a,b)
if b = 0
return a
r = amodb
return hcf(b,r)
```

It takes a constant time, no matter the amount of data, to perform the operation.

 $Euclid's\ algorithm$

6

5

Sender Verification
Private key remains with owner
Public key is freely distributable
No secret channel needed at any point
No need for pre-shared keys

```
fme(a,b,k)
d = a
e = b
s = 1
While e > 0
if e is odd
s = (s.d)modk
d = d^2modk
e = \lfloor e/2 \rfloor
return s
```

 $Fast\ Modular\ Exponentiation$

What is the basic procedure for an encryption and decryption using publik key cryptography if Alice wants to send a message to Bob?	Describe public key generation in ElGamal encryption using p as the Prime Modulus and g as the Primitive root (as described in the COMP26120 lab)
Describe the encryption procedure used in the $ElGamal\ cryptosystem\ given\ that\ person\ B\ wants\ to$ $send\ message\ M\ to\ preson\ A$	Describe the decryption process used in the ElGamal cryptosystem given that person A has received cyphertext (γ, δ) from person B , encrypted encrypted using the public key (p, g, g^a)
Consider the equation $a^x = y \mod p$. If a is a primitive root of modulo p , then for every $y(1 \le y < p)$, such an $x(1 \le x < p)$ exists. What is x ?	The is the inverse of exponentiation.
Why can a private key in the ElGamal cryptosystem not, in practice, be recovered using the public key when p is large?	What is one way you can argue correctness of Euclid's algorithm?

Generate a large p and a g in $1 \le g < p$ Generate a random integer a in $1 \le a \le p-2$ Compute $g^a \mod p$. The public key is

 (p, g, g^a)

The private key is a

Alice generates a private random integer a and Bob generates a private random integer b

Alice generates her public value $g^a \mod p$ Bob generates his public value $g^b \mod p$ Alice computes $g^{ab} = (g^a)^b \mod p$ Bob computes $g^{ba} = (g^b)^a \mod p$ Now they have a shared secret k since $k = g^{ab} = g^{ba}$

10

9

Use private key a to compute $(\gamma^{p-1-a}) \mod p$ $NOTE\ THAT: (\gamma^{p-1-a}) = \gamma^{-a} = g^{-ak}$ Recover the message M by computing $(\gamma^{-a} \cdot \delta \mod p)$ Note that this evaluates to $(g^{-ak} \cdot g^{ak} \cdot M \mod p)$ or $1 \cdot M \mod p$ Obtain A's public key (p, g, g^a) Represent the message M as integers in the range 0,...,p-1Select a random integer k from $1 \le k \le p-2$ Compute $\gamma = g^k \mod p$ and $\delta = m \cdot (g^a)^k$ Send ciphertext $c = (\gamma, \delta)$ to A

12

11

The discrete logarithm is the inverse of exponentiation.

X is the **discrete logarithm** of y with base a, modulo p.

14

13

Let r = amodb. hcf(a, b) = hcf(b, r) because all factors of a and b are also factors of b and r and vice versa. If they have the same factors, they have the same highest common factor.

To calculate a public key, y, a private key, x is needed. The equation for modular exponentiation can be used to generate the public key: $y = g^x \mod p$ where g is a primitive root of the modulus p. It is considered a one-way, or trapdoor function - easy to compute, hard to invert. For a large p, one of the few ways to figure out the private key x would be to calculate $g^x \mod p$ for every x in $1 \le x < p$ and find when one of these results matches y

What would half the correctness proof be for Euclid's algorithm?	(a.b)modk =
Let p be a prime number. What is meant by a primitive root modulo p?	What are the best, average and worst case complexities of Bubble Sort?
What are the best, average and worst case complexities of Merge Sort?	Give pseudo code for merging 2 sorted lists, as part of merge sort.
$Give\ pseudo\ code\ for\ MergeSort(L).$	What are the best, average and worst case complexities of Quick Sort?

```
As r = a \mod b, \exists q \text{ such that } a = bq + r,
                                                                                           \therefore r = a - bq.
                                                                      Suppose x is a factor of a and b, then \exists yandz such
          (a.b)modk = (amodk.bmodk)modk
                                                                                       that a = xy, b = xz.
                                                                              Hence: r = xy - xzq, r = x(y - zq).
                                                                           \therefore x is a factor of r (and also of b and r).
                                                        18
                                                                                                                              17
                                                                        The numbers r_x between 1 and p-1 that, when
                      Best: O(n),
                                                                      raised by the numbers between 1 and p-1 compute
                    Average: O(n^2),
                                                                      all the numbers between 1 and p-1 in some order
                      Worst: O(n^2)
                                                                                        with no repetitions.
                                                        20
                                                                                                                              19
Merge(L_1, L_2)
   if L_1 = [] return L_2
   if L_2 = [] return L_1
   x_1 = L_1[0]
   x_2 = L_2[0]
                                                                                        Best: O(n \log_2 n),
   L'_1 = L_1[1:|L_1|-1]

L'_2 = L_2[1:|L_2|-1]
                                                                                       Average: O(n \log_2 n),
                                                                                        Worst: O(n \log_2 n)
   if x_1 \leq x_2
       return [x_1] + Merge(L'_1, L_2)
   return [x_2] + Merge(L_1, L'_2)
                 Merge two sorted lists
                                                        22
                                                                                                                              21
                                                                     MergeSort(L)
                                                                         if |L| \leq 1
                   Best: O(n \log_2 n),
                                                                            return\ L
```

Average: $O(n \log_2 n)$,

Worst: $O(n^2)$

24 23

Split L into roughly equal halves, L_l and L_r

 $return Merge(MergeSort(L_l), MergeSort(L_r))$

MergeSort(L)

What would the pseudo code be for Quick Sort?	What is the minimum time for any sorting algorithm that uses only number comparisons?
What does saying that algorithm A runs in time g mean?	What is a permutation of a set?
What do we mean by a composition of two permutations?	What is the number of possible permutations on an n-element set?
In the context of a permutation, what do we mean by a transposition?	Convert this pair of simultaneous equations into matrix form $a_{1,1}x_1+a_{1,2}x_2=b_1 \\ a_{2,1}x_2+a_{2,2}x_2=b_2$
a transposition:	

 $n\log_2 n$

 $\begin{aligned} &quicksort(L) \\ &if\ length\ of\ L \leq 1 \\ &return\ L \\ &remove\ the\ first\ element,\ x,\ from\ L \\ &L_{\leq} := \ elements\ of\ L\ less\ than\ or\ equal\ to\ x \\ &L_{>} := \ elements\ of\ L\ greater\ than\ x \\ &L_{l} := \ quicksort(L_{\leq}) \\ &L_{r} := \ quicksort(L_{>}) \\ &return\ L_{l}\ + \ [x]\ + \ L_{r} \end{aligned}$

Quick Sort

25

26

A 1-to-1 map of the set onto itself. In basic terms, it is a set mapped to another order of itself. i.e $[0,1,2,3,4] \mapsto [2,4,1,0,3]$

Given an input of size n, the number of operations executed by A is bounded above by g(n).

28 27

The composition is the product of two permutations, α and β , on a set n, given by $\alpha \cdot \beta(n)$ or $\beta(\alpha(n))$

30 29

A transposition is a special kind of permutation where only 2 elements in a set are affected (they are swapped). On a set X a transposition $\sigma = (i, j)$ is given by

$$\sigma(k) = \begin{cases} j & \text{if } k = i \\ i & \text{if } k = j \\ k & \text{ow.} \end{cases}$$

n!

 $\begin{pmatrix} a_{1,1} & a_{1,2} \\ a_{2,1} & a_{2,2} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$

32

What is the determinant of the matrix: $\begin{pmatrix} a_1 & a_2 \\ a_3 & a_4 \end{pmatrix}$	What is an upper triangular matrix and how do you calculate its determinant?
Which 4 operations have no effect on a matrix's determinant?	In a tree, if node u is the parent (ancestor) node of v, then v is a (ancestor) of u. Two children of the same parent are (ancestor).
What is a tree?	In a tree, an external node is known as It has no
In a tree, an internal node has one or more.	What do we do here? Left Left Case Root 5 A 40

It is a matrix where all of its entries below the diagonal are zero.

$$\begin{pmatrix} a_{1,1} & a_{1,2} & \cdots & a_{1,n} \\ 0 & a_{2,2} & \cdots & a_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & a_{n,n} \end{pmatrix}$$

Its determinant is calculated by taking the product of the entries on the diagonal. i.e $a_{1,1} \cdot a_{2,2} \cdot ... \cdot a_{n,n}$

34

36

 $a_1a_4 - a_2a_3$ Often denoted as:

$$\begin{vmatrix} a_1 & a_2 \\ a_3 & a_4 \end{vmatrix}$$

The original system of equations to which the matrix corresponds only has a unique solution if the determinant is non-zero.

33

In a tree, if node u is the **parent (ancestor)** node of v, then v is a child (descendent) of u. Two children of the same parent are siblings.

Transposing two rows
Transposing two columns
Adding a multiple of one row to another
Adding a multiple of one column to another
Also note that if all entries in any row or column
are 0 then the determinant is 0

35

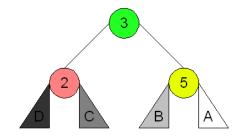
In a tree, an external node is known as a leaf node.

It has no children.

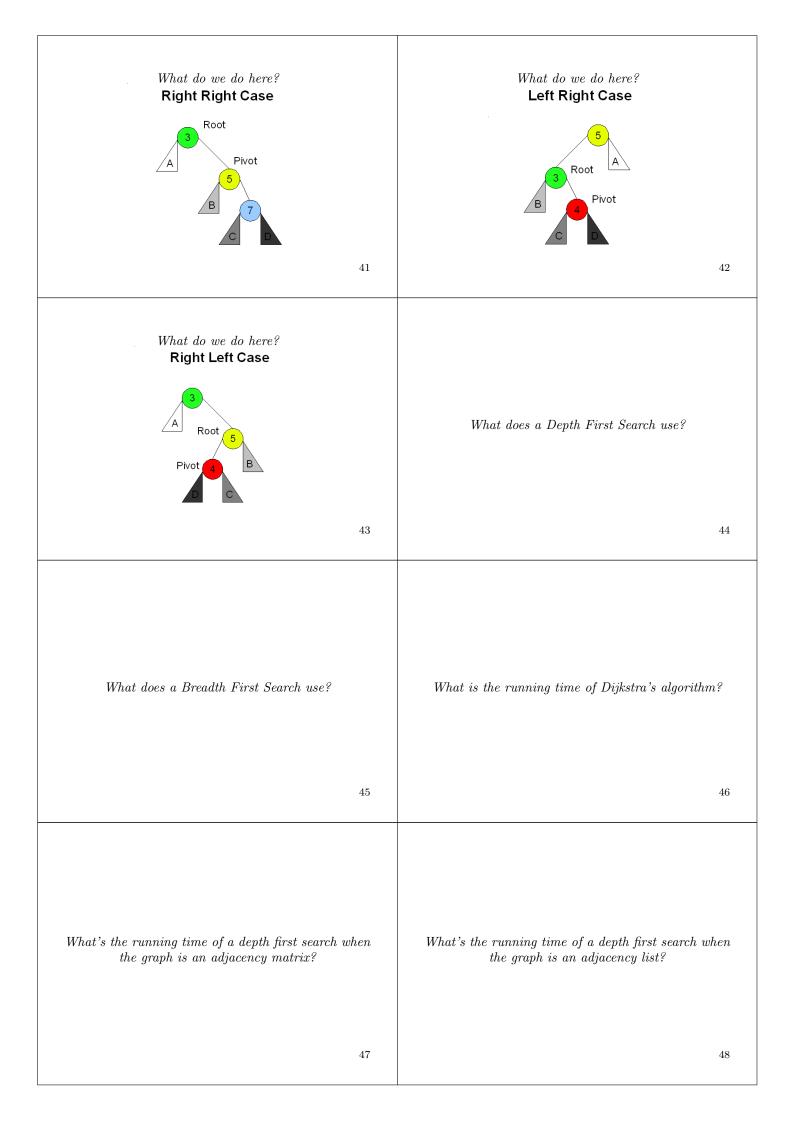
A tree T is a non-empty set of nodes storing useful information in a parent-child relationship with the following properties:

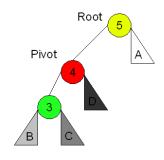
T has a special node r referred to as the root. Each node v of T different from r has a parent node u.

38

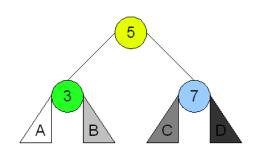


In a tree, an internal node has one or more children.



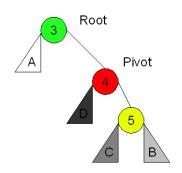


Then a left-left!



42 41

 $A \ stack!$



Then a right-right!

43

44

 $O(E + V \log(v))$ A queue!

46 45

O(V+E) $O(V^2)$ since finding neighbours takes O(V) time.

How do we insert an element into a heap?	How do we remove the smallest element from a $(min)\ heap?$
49	50
Dijkstras algorithm is a algorithm for graph structures.	Give the pseudocode for Dijkstra's algorithm w/ priority queue
51	52

We move the last element from the heap to the first element (we can override the first element since we've removed it). Now we 'down heap' by swapping the moved node with its smallest child until it is smaller than both its children or it has no children.

First, you insert it at the next space in the heap (last element of current row, or a new row), then you keep swapping it with its parent if the parent is larger than it.

50 49

```
function \ Dijkstra(Graph, \ Source): \\ dist[source] = 0, \ create \ Q \\ for \ each \ vertex \ v \ in \ Graph: \\ if \ v \ != \ source, \ dist[v] = INFINITY \\ prev[v] = UNDEFINED \\ Q.addWithPriority(v, \ dist[v]) \\ while \ Q \ != \ empty \\ u = Q.extractMin() \\ for \ each \ neighbour \ v \ of \ u: \\ alt = \ dist[v] + \ length(u, \ v) \\ if \ alt \ i \ dist[v] \\ dist[v] = \ alt, \ prev[v] = u \\ Q.decreasePriority(v, \ alt) \\ return \ dist, \ prev
```

Dijkstras algorithm is a shortest path algorithm for graph structures.

 $Code\ for\ dijkstras$