



FDS 2025 春夏

project2

Autograd for Algebraic Expressions

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1 Introduction

In this project, we need to write a program to do autograd for the given infix expression. And this report is mainly about the result of my program's performance and what is need to be done to improve it.

- (1). In Chapter 2, I will introduce the the main algorithm of process used in the program.
- (2). In Chapter 3, I present my test data , including input ,output ,and some analysis for the result.
- (3). In Chapter 4, I conduct an analysis of my program, including what function has been achieved and what is needed to be improved.
- (4). In Chapter 5, I include the complete code in the appendix.

2 Algorithm Specification

The following are the pseudocodes and explanations for main algorithm of process used in the program.

2.1 Infix Expression \rightarrow Postfix Expression

```
1      function infix_to_postfix(infix_expression):
2          create empty stack operator_stack
3          create empty list output
4
5          for each token in infix_expression:
6              if token is an operand (number or variable):
7                  append token to output
8              else if token is a left parenthesis '(':
9                  push token onto operator_stack
10             else if token is a right parenthesis ')':
11                 while operator_stack is not empty and top is not '(':
12                     pop from operator_stack and append to output
13                 pop the '(' from operator_stack (discard)
14             else if token is an operator:
15                 while operator_stack is not empty and
16                     precedence(top of operator_stack) >= precedence(token):
17                     pop from operator_stack and append to output
```

```
18         push token onto operator_stack
19
20     while operator_stack is not empty:
21         pop from operator_stack and append to output
22
23     return output (postfix expression)
```

To construct the infix expression as a tree ,we first need to transform it into a postfix expression .Here I deploy the traditional way of using a stack to save operators and do push & pop operations with the operators and brackets.

2.2 Postfix Expression \rightarrow Abstract Syntax Tree

```
1     div_con(arr, sum, left_p, right_p, count):
2         res = 0 //initialize result
3
4         //only begin search if the size is smaller than 20
5         if (right_p - left_p) > 20:
6             mid = (right_p + left_p) / 2 //calculate the mid num
7             res_left = div_con(arr, sum, left_p, mid, count) //recursiononly
8                 check the left part
9             if res_left == 1: //if find in the left half
10                 res = 1
11                 return res //target found
12
13             res_right = div_con(arr, sum, mid, right_p, count) //recursiononly
14                 check the right part
15             if res_right == 1: //if find in the left half
16                 res = 1
17                 return res //target found
18
19         //if the size if smaller than 20 or the smaller branch didn't find
20         the result
21         for i from left_p to right_p - 1:
22             for j from i + 1 to right_p - 1:
23                 if arr[i] + arr[j] == sum: //check the sum
24                     res = 1
25                     // printf("%d %d + %d = %d\n", count, arr[i], arr[j], sum)
```

```
23         return res //target found
24
25     return res //target not found
```

Then we need to transform the postfix expression into an ast .The basic algorithm is to use a stack to save the variable and number nodes .When it meet a operator ,it will pop the last two node of the stack as the left and right child of the operator node ,and push it into the stack.

2.3 Recursive Differentiation on AST

```
1     function derivative(node):
2         if node is a constant:
3             return a node representing constant 0
4         if node is a variable:
5             if node.name matches the differentiation variable:
6                 return a node representing constant 1
7             else:
8                 return a node representing constant 0
9
10        if node is an operator:
11            let u = node.left
12            let v = node.right (may be NULL for unary operators)
13
14            match node.operator:
15                case '+':
16                    return derivative(u) + derivative(v)
17                case '-':
18                    return derivative(u) - derivative(v)
19                case '*':
20                    return u * derivative(v) + v * derivative(u)
21                case '/':
22                    return (v * derivative(u) - u * derivative(v)) / (v^2)
23                case '^':
24                    if v is a constant:
25                        return v * u^(v - 1) * derivative(u)
26                    else:
27                        return u^v * [derivative(v) * ln(u) + v * derivative(u)]
```

```
                u]
28         case 'ln':
29             return derivative(u) / u
30
31     return the newly constructed node (expression subtree)
```

We use a recursive function, typically named `derivative(node)`, which takes the root of a subtree and returns a new subtree representing the derivative of that expression. The recursion works by decomposing a complex expression into smaller parts (its children), computing the derivative of each, and then combining the results according to differentiation rules.

(1). **Constant Rule**

$$\frac{d}{dx}(c) = 0$$

where c is a constant.

(2). **Variable Rule**

$$\frac{d}{dx}(x) = 1, \quad \frac{d}{dx}(y) = 0 \quad (\text{if } x \neq y)$$

(3). **Addition / Subtraction Rule**

$$\frac{d}{dx}(u \pm v) = \frac{du}{dx} \pm \frac{dv}{dx}$$

(4). **Multiplication Rule (Product Rule)**

$$\frac{d}{dx}(uv) = u \cdot \frac{dv}{dx} + v \cdot \frac{du}{dx}$$

(5). **Division Rule (Quotient Rule)**

$$\frac{d}{dx} \left(\frac{u}{v} \right) = \frac{v \cdot \frac{du}{dx} - u \cdot \frac{dv}{dx}}{v^2}$$

(6). **Power Rule**

- If exponent is constant n :

$$\frac{d}{dx}(u^n) = n \cdot u^{n-1} \cdot \frac{du}{dx}$$

- General case:

$$\frac{d}{dx}(u^v) = u^v \cdot \left(\frac{dv}{dx} \cdot \ln u + v \cdot \frac{1}{u} \cdot \frac{du}{dx} \right)$$

3 Testing Data

To analyze the performance of the autograd program , we try to input different infix expression and compare the output answer with expected answer.

3.1 Simply add and minus

Input:

```
1      a+b-c
```

Output:

```
1      a: 1
2      b: 1
3      c: -1
```

Expected output:

```
1      a: 1
2      b: 1
3      c: -1
```

3.2 Multiply and division

Input:

```
1      x*a-b/x
```

Output:

```
1      x: a-(-b)/(x*x)
2      a: x
3      b: -(x)/(x*x)
```

Expected output:

```
1      x:a+1/(x*x)
2      a:x
3      b:-1/x
```

Here we find that the true output is not completely the same as expected output. But if we do some simple calculation ,we can find that after some simplification the output is exactly the same as the expected output and we will analyze this in the next chapter.

3.3 Bracket and power

Input:

```
1      (x-1)^3+a^c
```

Output:

```
1      x: 3*(x-1)^2
2      a: a^c*(c*1/a)
3      c: a^c*(ln(a))
```

Expected output:

```
1      x: 3*(x-1)^2
2      a: a^(c-1)*c
3      c: a^c*ln(a)
```

3.4 Long variable

Input:

```
1      aa*bb-ab/aa
```

Output:

```
1      aa: bb-(-ab)/(aa*aa)
2      bb: aa
3      ab: -(aa)/(aa*aa)
```

Expected output:

```
1      aa: bb+ab/(aa*aa)
2      bb: aa
3      ab: -1/aa
```

3.5 Combination of situation

Input:

```
1      aa^(aa-b)-x/(d+b)*5
```

Output:


```
1      aa: aa^(aa-b)*(ln(aa)+(aa-b)*1/aa)
2      b: aa^(aa-b)*((-1)*ln(aa))-(-x)/((d+b)*(d+b))*5
3      x: -(d+b)/((d+b)*(d+b))*5
4      d: -(-x)/((d+b)*(d+b))*5
```

Expected output:

```
1      aa: aa^(aa-b)*(ln(aa)+(aa-b)/aa)
2      b: aa^(aa-b)*(-ln(aa))+x/((d+b)*(d+b))*5
3      x: -1/(d+b)*5
4      d: +x/((d+b)*(d+b))*5
```

4 Analysis and Comments

Firstly ,I'll briefly analyze the time and space complexity of the main algorithm used in the program. Then I will comment on the performance of the program.

4.1 Infix expression -> tree

(1). **Time complexity:**

- During this process, we have two step :1.infix->postfix,2.construct the tree.And for both the process ,we need to traverse all the characters in the expression. Therefore ,the time complexity of Infix expression -> tree process is $O(n)$.

(2). **Space complexity:**

- The main space cost is the stacks we use in the process and the node space for the tree.We only need $O(n)$ for the stack and tree ,therefore the total space complexity is $O(n)$.

4.2 Differentiation

(1). **Time complexity:**

- The differentiation process use the recursion to process the tree.When the worst case happens ,we need to process multiply and divide process ,which will cost 2-3 times of node .Therefore ,theoretically the worst case of time complexity is $O(2^n)$.

(2). Space complexity:

- The same as time complexity ,in worst case the space complexity will be $O(2^n)$.

4.3 Performance analysis

- The program can perform the operation of auto-differentiation for given infix expression.It can deal with add,minus,multiply,divide,bracket and power quite well.Even in some cases we find the output is not exactly the same as the expected output ,but if we artificially simplify the output expression ,we can find the output answer is right.
- The program also involves some simple simplification ,which can deal with operations that involves 0 and 1 ,which helps to simplify most cases like $a+0$, $a*0$, $a0$.etc

4.4 Things need to be improved

- Just as mentioned before ,the program can't deal with simplifying operations with same variable and combining same sub-expression.This will make some simple answer printed in a more complicated way.But yet I don't find a proper way to deal with it.

1 `xy: (-xx^2)/(xy*xy)*xy+xx^2/xy = xy: 0`

- This program also can't deal with any complicated math functions like $\ln(x)$, $\sin(x)$, $\text{pow}(x,y)$.etc This will involve a more complicated constructing rule for the tree.

5 Appendix

```
1  # include<stdio.h>
2  # include<stdlib.h>
3  # include<string.h>
4
5  typedef struct node{
6      char* variable;           //variable
7      char operator;           //operator
8      int num;                 //num
9      int type;                //1: variable 2:operator 3:
                                num
```

```
10     struct node* right;           //right child
11     struct node* left;           //left child
12 }node;
13
14 int precedence(char op);         //return the priority of
    operation
15 int is_right_associative(char op); //return if the operator is
    ^
16 int is_operator(char c);         //return if the char is an
    operator
17 node* create_node(char* token);   //create the node for each
    item
18 int get_root(char* exp ,char** var ,node** root); //construct the tree of
    the expression
19 int get_token(char* exp ,char** var); //divide the expression
    into operators and operands
20 node* autograd_for_var(char* var ,node* root); //calculate the grad of
    the tree
21 char* join(node* root ,char* expression); //simplify the tree for 0
    and 1 and()
22
23 node* new_operator_node(char op ,node* left ,node* right){ //create node
    of operator in grad
24     node* n = (node*)malloc(sizeof(node)); //prepare the space
    for new node
25     n->type = 2; //mark type as 2
26     n->operator = op;
27     n->left = left;
28     n->right = right;
29     return n;
30 }
31
32 node* new_variable_node(char* var){ //create node of
    variable in grad
33     node* n = (node*)malloc(sizeof(node)); //prepare the space
    for new node
34     n->type = 1; //mark type as 1
35     n->variable = var;
```

```

36     n->left = n->right = NULL;           //leaf node
37     return n;
38 }
39
40 node* new_number_node(int val){           //create node of
    number in grad
41     node* n = (node*)malloc(sizeof(node)); //prepare the space
    for new node
42     n->type = 3;                          //mark type as 3
43     n->num = val;
44     n->left = n->right = NULL;           //leaf node
45     return n;
46 }
47
48 int main(void){
49     char expression[100];                 //expression is the
    infix input
50     char** variables = (char**)malloc(sizeof(char*)*100); //to save the
    division first ,then save the variable list
51     node* ast_root = NULL;               //the root for the
    tree
52     scanf("%[^\n]" ,expression);
53     // printf("%s",expression);
54     // printf("\n");
55     //build the tree of the expression and find the variables
56     int len_of_var = get_root(expression ,variables ,&ast_root); //build
    the tree of expression and get the number of variable
57     printf("There's %d variables in the expression.\n",len_of_var); //
    print out the lenth of variable
58     for(int i = 0;i < len_of_var;i++){    //do grad
    for each variable
59         printf("%s",variables[i]);
60         printf(": ");
61         node* a_of_var;                  //to save
    the tree after grad
62         a_of_var = autograd_for_var(variables[i] ,ast_root); //
    calculate the grad
63         char* join_exp = (char*)calloc(1000 ,sizeof(char)); //simplify

```

```

        and turn it into a string
64     join_exp = join(a_of_var ,join_exp);           //print
        out the result
65     printf("%s" ,join_exp);
66
67     printf("\n");
68 }
69 return 0;
70 }
71
72 int precedence(char op){
73     switch(op){
74         case 'l':
75         case '^': return 3;           //^ is the highest
76         case '*':
77         case '/': return 2;          /* / is the 2nd
78         case '+':
79         case '-': return 1;          //+ - is the lowest
80         default: return 0;
81     }
82 }
83
84 int is_right_associative(char op){    //return if the operator is ^
85     return op == '^';
86 }
87
88 int is_operator(char c){              //return if the char is an
        operator
89     return c == '+' || c == '-' || c == '*' || c == '/' || c == '^';
90 }
91
92 node* create_node(char* token){
93     node* n = (node*)malloc(sizeof(node)); //create the new node
94     n->left = NULL;
95     n->right = NULL;
96     if(is_operator(token[0])){         //if the token is an operator
97         n->type = 2;                   //mark type as 2
98         n->operator = token[0];

```

```

99     }else if((token[0] >= 'a' && token[0] <= 'z') || (token[0] >= 'A' &&
    token[0] <= 'Z')){ //if the token is variable
100         n->type = 1;

        //mark type as 1
101         n->variable = token;
102     }else{

        //if the token is number
103         n->type = 3;

        //mark type as 3
104         int i = 0;
105         int result = 0;

        //

        turn the string number into a int type
106         while (token[i] != '\0'){
107             if(token[i] >= '0' && token[i] <= '9'){
108                 result = result * 10 + (token[i] - '0');
        //calculate the number
109             }else{
110                 printf("Invalid number: %s\n" ,token);
        //exit if it is not a
        number
111                 exit(1);
112             }
113             i++;
114         }
115         n->num = result;
116     }
117
118     return n;
119 }
120
121 int get_root(char* exp ,char**var ,node** root){
122     int var_num = 0;

    //

    count the number of variables

```

```

123 //transition of infix into postfix
124 char** stack = (char**)malloc(sizeof(char*) * 20);
                                //operator stack for in-post
                                transition
125 char** post = (char**)malloc(sizeof(char*) * 20);
                                //save postfix expression
126 int len = get_token(exp ,var);
                                //divide into
                                tokens
127 int post_p = 0;
128 int stack_p = 0;
129 int flag = 0;

    //count backects
130 for(int i = 0;i < len;i++){
131     if(is_operator(var[i][0])){
                                //if operator
                                push into stack or pop operator out
132     while (stack_p > 0 && stack[stack_p - 1][0] != '(' &&
133         (precedence(stack[stack_p - 1][0]) > precedence(var[i][0])
134         ||
135         (precedence(stack[stack_p - 1][0]) == precedence(var[i][0])
136             && !is_right_associative(var[i][0])))) {
137         post[post_p++] = stack[--stack_p];
                                //pop operator with
                                higher priority
138     }
139     stack[stack_p++] = var[i];
140 }else if(var[i][0] == '('){
                                //if ( push
                                into stack
141     stack[stack_p] = var[i];
142     stack_p++;
143     flag++;
144 }else if(var[i][0] == ')'){
                                //if ) pop
                                operator out until (
145     while(stack_p > 0 && stack[stack_p - 1][0] != '('){

```

```

144         post[post_p++] = stack[--stack_p];
                                           //pop operator out
145     }
146     if(stack_p > 0 && stack[stack_p - 1][0] == '('){
                                           //meet ( and throw it
147         --stack_p;
148     }else{
149         printf("Mismatched parentheses!\n");
                                           //if ( ) didn't match
150         exit(1);
151     }
152 }else if((var[i][0] >= 'A' && var[i][0] <= 'Z') || (var[i][0] >= '
a' && var[i][0] <= 'z') || (var[i][0] >= '0' && var[i][0] <= '9'
))){
153     post[post_p] = var[i];
                                           //if
                                           variable or num push into post
154     post_p++;
155 }else{
156     printf("Wrong input!");
                                           //if not
                                           allowed input
157     exit(1);
158 }
159 }
160 while(stack_p > 0){
161     // printf("%s", stack[stack_p]);
162     // printf("\n");
                                           //
                                           pop out the remaining operators in the stack
163     if(stack[stack_p - 1][0] == '('){
164         printf("Mismatched parentheses!\n");
                                           //if ( ) didn't match
165         exit(1);
166     }
167     post[post_p++] = stack[--stack_p];
                                           //push into post
168 }

```



```

169     len = len - flag * 2;
170     printf("Postfix expression:");
171     for(int i = 0;i < post_p;i++)printf("%s",post[i]);
172     printf("\n");
173     //now we start construct the ast tree
174     node** node_stack = (node**)malloc(sizeof(node*) * len);
175                                     //stack for elements
176     int n_stack_p = 0;
177     for(int i = 0;i < len;i++){
178         node* temp;
179
180         //create new tree node
181         temp = create_node(post[i]);
182         if(temp->type == 2){
183
184                                     //use
185
186         operator as root
187         temp->right = node_stack[--n_stack_p];
188                                     //pop the operand of the
189
190         operator from the stack
191         temp->left = node_stack[--n_stack_p];
192         node_stack[n_stack_p++] = temp;
193                                     //push the current
194
195         root into the stack
196     }else if(temp->type == 1){
197         node_stack[n_stack_p++] = temp;
198         int flag = 1;
199
200         //check if the variable has appeared
201         for(int i = 0;i < var_num;i++){
202             if(strcmp(temp->variable ,var[i]) == 0){
203                 flag = 0;
204             }
205         }
206         if(flag){
207
208             //if the first time count and push into varlist
209             var[var_num++] = temp->variable;
210         }

```

```

194     }else{
195         node_stack[n_stack_p++] = temp;
                                     //if number ,push
                                     into stack
196     }
197 }
198 *root = node_stack[0];
                                     //save
                                     the root for the tree
199 return var_num;
                                     //
                                     return the number of variables
200 }
201
202 int get_token(char* exp ,char** var){
203     int count = 0;
204
205     //count the elements
206     int i = 0;
207     var[count] = (char*)malloc(sizeof(char) * 20);
                                     //use var to save the divided
                                     element for the tree
208     int temp_p = 0;
209     while(exp[i] != '\0'){
                                     //go
                                     through every char
210         if((exp[i] >= 'A' && exp[i] <= 'Z') || (exp[i] >= 'a' && exp[i] <=
            'z') || (exp[i] >= '0' && exp[i] <= '9')){
211             var[count][temp_p] = exp[i];
                                     //if number or
                                     letter
212             temp_p++;
213         }else if(is_operator(exp[i]) || exp[i] == '(' || exp[i] == ')'){
            if(temp_p != 0){
                                     //
                                     save number or variable until meet an operator
214             var[count][temp_p] = '\0';
                                     //end it as a

```

```

                string
214         count++;
215         var[count] = (char*)malloc(sizeof(char) * 20);
216         temp_p = 0;
217     }
218     var[count][temp_p++] = exp[i];
                                                    //add the
                operator
219     var[count][temp_p++] = '\0';
220     count++;
221     var[count] = (char*)malloc(sizeof(char) * 20);
                                                    //start next element
222     temp_p = 0;
223 }
224 i++;
225 }
226 if(temp_p != 0){
                                                    //
        save tje last element
227     var[count][temp_p] = '\0';
228     count++;
229 }
230 return count;
231 }
232
233 node* autograd_for_var(char* var ,node* root){
234     if(!root){
235         return NULL;
                                                    //meet
        the leaf node
236     }
237     if(root->type == 3){
238         return new_number_node(0);
                                                    //for
        const ,grad is 0
239     }
240     if(root->type == 1){
                                                    //for
        var ,check if it is the target var for grad
241         if(strcmp(root->variable ,var) == 0){
242             return new_number_node(1);
                                                    //is

```

```

        return 1
243     }else{
244         return new_number_node(0);           //or
        return 0
245     }
246 }
247 if(root->type == 2){                          //for
    operator ,recursionly process the child
248     node *u = root->left ,*v = root->right;    //u,v
        as the operand of the operator
249     node *du = autograd_for_var(var ,u) ,*dv = autograd_for_var(var ,v
        );//du,dv as the grad of u,v
250     switch (root->operator){
251         case '+':                             //d(u+v
            )=du+dv
252             return new_operator_node('+', du ,dv);
253         case '-':                             //d(u-v
            )=du-dv
254             return new_operator_node('-', du ,dv);
255         case '*':                             //d(u*v
            )=du*v+u*dv
256             return new_operator_node('+',new_operator_node('*', du ,v)
                ,new_operator_node('*', u ,dv));
257         case '/':                             //d(u/v
            )=(du*v-u*dv)/(v*v)
258             return new_operator_node('/',new_operator_node('-',
                new_operator_node('*', du ,v),new_operator_node('*', u ,
                dv)) ,new_operator_node('*', v ,v));
259
260         case '^':{
261             int is_u_const = (u->type == 3);
262             int is_v_const = (v->type == 3);
263             if(is_v_const && !is_u_const){      //d(u^c
                )=c*u^(c-1)*du
264                 int c = v->num;
265                 node* new_exp = new_number_node(c - 1);
266                 node* u_pow_c_minus_1 = new_operator_node('^',u ,
                    new_exp);

```

```

267         return new_operator_node('*',new_operator_node('*',
268             new_number_node(c),u_pow_c_minus_1),du);
269     }else if(is_u_const && !is_v_const) { //d(c^v
270         )=c^v*ln(c)*dv
271         int c = u->num;
272         node* ln_c = new_operator_node('l' ,u ,NULL); //ln(u)
273         node* c_pow_v = new_operator_node('^' ,u ,v);
274         return new_operator_node('*',new_operator_node('*',
275             c_pow_v,ln_c),dv);
276     }else{ //d(u^v
277         )=u^v*(dv*ln(u)+v*du/u)
278         node* ln_u = new_operator_node('l' ,u ,NULL); //ln(u)
279         node* term1 = new_operator_node('*',dv ,ln_u);
280         node* u_div = new_operator_node('/',du ,u); //du/u
281         node* term2 = new_operator_node('*',v ,u_div);
282         node* sum = new_operator_node('+',term1 ,term2); //dv*
283             ln(u)+v*du/u
284         node* u_pow_v = new_operator_node('^' ,u ,v);
285         return new_operator_node('*',u_pow_v ,sum);
286     }
287 }
288 default:
289     printf("Unsupported operator: %c\n" ,root->operator); //
290         operator unsupported
291     exit(1);
292 }
293
294 return NULL;
295 }
296
297 char* join(node* root ,char* expression){
298     if(!root){ //meet the leaf node
299         return NULL;
300     }
301     if(root->type == 1){ //variable cat into
302         expression
303         strcat(expression ,root->variable);
304         return expression;

```

```

298     }
299     else if(root->type == 3){                                     //number cat into
        expression
300         char num_str[20];
301         sprintf(num_str ,"%d" ,root->num);                       //change int into
        str
302         strcat(expression ,num_str);
303         return expression;
304     }
305     else if(root->type == 2){
306         if(root->operator == 'l'){                                //opertor 'l' as ln
307             strcat(expression ,"ln(");
308             join(root->left ,expression);                         //put ln(x) into
        expression
309             strcat(expression ,")");
310             return expression;
311         }
312
313         char* left_str = (char*)calloc(100 ,sizeof(char)); //save the
        middle result for the two child
314         char* right_str = (char*)calloc(100 ,sizeof(char));
315         left_str = join(root->left ,left_str);                   //recursively
        calculate the result
316         right_str = join(root->right ,right_str);
317
318         if(root->operator == '+'){                                //simplify when 0+a,
        a+0, 0+0
319             if(strcmp(left_str ,"0") == 0){                     //0+a->a
320                 strcat(expression ,right_str);
321                 goto done;
322             }
323             if(strcmp(right_str ,"0") == 0){                     //a+0->a
324                 strcat(expression ,left_str);
325                 goto done;
326             }
327         }
328         else if(root->operator == '-'){                           //simplify when 0-a,
        a-0

```

```

329         if(strcmp(left_str , "0") == 0 && strcmp(right_str , "0") == 0){
330             strcat(expression , "0");           //0-0->0
331             goto done;
332         }
333         if(strcmp(left_str , "0") == 0){
334             strcat(expression , "-");           //0-a->-a
335             strcat(expression , right_str);
336             goto done;
337         }
338         if(strcmp(right_str , "0") == 0){
339             strcat(expression , left_str);       //a-0->a
340             goto done;
341         }
342     }
343     else if(root->operator == '*'){              //simplify when 0*a,
        a*0, 1*a, a*1
344         if(strcmp(left_str , "0") == 0 || strcmp(right_str , "0") == 0){
345             strcat(expression , "0");           //0*0->0
346             goto done;
347         }
348         if(strcmp(left_str , "1") == 0){
349             strcat(expression , right_str);       //1*a->a
350             goto done;
351         }
352         if(strcmp(right_str , "1") == 0){
353             strcat(expression , left_str);        //a*1->a
354             goto done;
355         }
356     }
357 }
358 else if(root->operator == '/'){                //simplify when 0/a,
        a/1
359     if(strcmp(left_str , "0") == 0){
360         strcat(expression , "0");           //0/a->0
361         goto done;
362     }
363     if(strcmp(right_str , "1") == 0){
364         strcat(expression , left_str);        //a/1->a

```

```

365         goto done;
366     }
367     }else if(root->operator == '^'){           //simplify when a^1,
        a^0
368         if(strcmp(right_str ,"1") == 0){
369             strcat(expression ,left_str);       //a^1->a
370             goto done;
371         }
372         if(strcmp(right_str ,"0") == 0){
373             strcat(expression ,"1");           //a^0->1
374             goto done;
375         }
376     }
377     if(root->left->type == 2 && precedence(root->operator) >
        precedence(root->left->operator)){
378         strcat(expression , "(");
379         strcat(expression ,left_str);           //add () if the left
            child operator is smaller than root
380         strcat(expression ,")");
381     }else{
382         strcat(expression ,left_str);           //else no ()
383     }
384     int len = strlen(expression);
385     expression[len] = root->operator;           //push the operator
386     expression[len + 1] = '\0';
387
388     if((root->right->type == 2 && precedence(root->operator) >
        precedence(root->right->operator))
389     || (root->right->type == 2 && root->operator == '/')){
390         strcat(expression , "(");
391         strcat(expression ,right_str);           //add () if the
            right child operator is smaller than root
392         strcat(expression ,")");
393     }else{
394         strcat(expression ,right_str);           //else no ()
395     }
396
397 done:

```



```
398     free(left_str);                                //free the temporal  
        string  
399     free(right_str);  
400     return expression;                            //return expression  
401 }  
402  
403     return NULL;                                    //return NULL when  
        something wrong  
404 }
```

6 Declaration

I hereby declare that all the work done in this project titled “Project 2 : Autograd for Algebraic Expressions ” is of my independent effort.