



FDS 2025 春夏  
project2  
Autograd for Algebraic Expressions

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# 1 Introduction

In this project, we need to write a program to do autograd for the given infix expression. And this report is mainly about the result of my program's performance and what is need to be done to improve it.

- (1). In Chapter 2, I will introduce the the main algorithm of process used in the program.
- (2). In Chapter 3, I present my test data , including input ,output ,and some analysis for the result.
- (3). In Chapter 4, I conduct an analysis of my program, including what function has been achieved and what is needed to be improved.
- (4). In Chapter 5, I include the complete code in the appendix.

# 2 Algorithm Specification

The following are the pseudocodes and explanations for main algorithm of process used in the program.

## 2.1 Infix Expression → Postfix Expression

```
1     function infix_to_postfix(infix_expression):  
2         create empty stack operator_stack  
3         create empty list output  
4  
5         for each token in infix_expression:  
6             if token is an operand (number or variable):  
7                 append token to output  
8             else if token is a left parenthesis '(':  
9                 push token onto operator_stack  
10            else if token is a right parenthesis ')':  
11                while operator_stack is not empty and top is not '(':  
12                    pop from operator_stack and append to output  
13                    pop the '(' from operator_stack (discard)  
14            else if token is an operator:  
15                while operator_stack is not empty and  
16                    precedence(top of operator_stack) >= precedence(token):  
17                    pop from operator_stack and append to output
```

```

18     push token onto operator_stack
19
20     while operator_stack is not empty:
21         pop from operator_stack and append to output
22
23     return output (postfix expression)

```

To construct the infix expression as a tree ,we first need to transform it into a postfix expression .Here I deploy the traditional way of using a stack to save operators and do push & pop operations with the operators and brackets.

## 2.2 Postfix Expression → Abstract Syntax Tree

```

1     div_con(arr, sum, left_p, right_p, count):
2         res = 0 //initialize result
3
4         //only begin search if the size is smaller than 20
5         if (right_p - left_p) > 20:
6             mid = (right_p + left_p) / 2 //calculate the mid num
7             res_left = div_con(arr, sum, left_p, mid, count) //recursionly
8                 check the left part
9                 if res_left == 1: //if find in the left half
10                     res = 1
11                     return res //target found
12
13             res_right = div_con(arr, sum, mid, right_p, count) //recursionly
14                 check the right part
15                 if res_right == 1: //if find in the left half
16                     res = 1
17                     return res //target found
18
19             //if the size if smaller than 20 or the smaller branch didn't find
20             //the result
21             for i from left_p to right_p - 1:
22                 for j from i + 1 to right_p - 1:
23                     if arr[i] + arr[j] == sum: //check the sum
24                         res = 1
25                         // printf("%d %d + %d = %d\n", count, arr[i], arr[j], sum)

```

```

23         return res //target found
24
25     return res //target not found

```

Then we need to transform the postfix expression into an ast .The basic algorithm is to use a stack to save the variable and number nodes .When it meet a operator ,it will pop the last two node of the stack as the left and right child of the operator node ,and push it into the stack.

## 2.3 Recursive Differentiation on AST

```

1   function derivative(node):
2       if node is a constant:
3           return a node representing constant 0
4       if node is a variable:
5           if node.name matches the differentiation variable:
6               return a node representing constant 1
7           else:
8               return a node representing constant 0
9
10      if node is an operator:
11          let u = node.left
12          let v = node.right (may be NULL for unary operators)
13
14          match node.operator:
15              case '+':
16                  return derivative(u) + derivative(v)
17              case '-':
18                  return derivative(u) - derivative(v)
19              case '*':
20                  return u * derivative(v) + v * derivative(u)
21              case '/':
22                  return (v * derivative(u) - u * derivative(v)) / (v^2)
23              case '^':
24                  if v is a constant:
25                      return v * u^(v - 1) * derivative(u)
26                  else:
27                      return u^v * [derivative(v) * ln(u) + v * derivative(u)/

```

```

28         u]
29     case 'ln':
30         return derivative(u) / u
31
31     return the newly constructed node (expression subtree)

```

We use a recursive function, typically named `derivative(node)`, which takes the root of a subtree and returns a new subtree representing the derivative of that expression. The recursion works by decomposing a complex expression into smaller parts (its children), computing the derivative of each, and then combining the results according to differentiation rules.

#### (1). Constant Rule

$$\frac{d}{dx}(c) = 0$$

where  $c$  is a constant.

#### (2). Variable Rule

$$\frac{d}{dx}(x) = 1, \quad \frac{d}{dx}(y) = 0 \quad (\text{if } x \neq y)$$

#### (3). Addition / Subtraction Rule

$$\frac{d}{dx}(u \pm v) = \frac{du}{dx} \pm \frac{dv}{dx}$$

#### (4). Multiplication Rule (Product Rule)

$$\frac{d}{dx}(uv) = u \cdot \frac{dv}{dx} + v \cdot \frac{du}{dx}$$

#### (5). Division Rule (Quotient Rule)

$$\frac{d}{dx}\left(\frac{u}{v}\right) = \frac{v \cdot \frac{du}{dx} - u \cdot \frac{dv}{dx}}{v^2}$$

#### (6). Power Rule

- If exponent is constant  $n$ :

$$\frac{d}{dx}(u^n) = n \cdot u^{n-1} \cdot \frac{du}{dx}$$

- General case:

$$\frac{d}{dx}(u^v) = u^v \cdot \left( \frac{dv}{dx} \cdot \ln u + v \cdot \frac{1}{u} \cdot \frac{du}{dx} \right)$$

### 3 Testing Data

To analyze the performance of the autograd program , we try to input different infix expression and compare the output answer with expected answer.

#### 3.1 Simply add and minus

Input:

```
1      a+b-c
```

Output:

```
1      a: 1
2      b: 1
3      c: -1
```

Expected output:

```
1      a: 1
2      b: 1
3      c: -1
```

#### 3.2 Multiply and division

Input:

```
1      x*a-b/x
```

Output:

```
1      x: a-(-b)/(x*x)
2      a: x
3      b: -(x)/(x*x)
```

Expected output:

```
1      x:a+1/(x*x)
2      a:x
3      b:-1/x
```

Here we find that the true output is not completely the same as expected output. But if we do some simple calculation ,we can find that after some simplification the output is exactly the same as the expected output and we will analyze this in the next chapter.

### 3.3 Bracket and power

Input:

```
1      (x-1)^3+a^c
```

Output:

```
1      x: 3*(x-1)^2
2      a: a^c*(c*1/a)
3      c: a^c*(ln(a))
```

Expected output:

```
1      x: 3*(x-1)^2
2      a: a^(c-1)*c
3      c: a^c*ln(a)
```

### 3.4 Long variable

Input:

```
1      aa*bb-ab/aa
```

Output:

```
1      aa: bb-(-ab)/(aa*aa)
2      bb: aa
3      ab: -(aa)/(aa*aa)
```

Expected output:

```
1      aa: bb+ab/(aa*aa)
2      bb: aa
3      ab: -1/aa
```

### 3.5 Combination of situation

Input:

```
1      aa^(aa-b)-x/(d+b)*5
```

Output:

```

1      aa: aa^(aa-b)*(ln(aa)+(aa-b)*1/aa)
2      b: aa^(aa-b)*((-1)*ln(aa))-(-x)/((d+b)*(d+b))*5
3      x: -(d+b)/((d+b)*(d+b))*5
4      d: -(-x)/((d+b)*(d+b))*5

```

Expected output:

```

1      aa: aa^(aa-b)*(ln(aa)+(aa-b)/aa)
2      b: aa^(aa-b)*(-ln(aa))+x/((d+b)*(d+b))*5
3      x: -1/(d+b)*5
4      d: +x/((d+b)*(d+b))*5

```

## 4 Analysis and Comments

Firstly ,I'll briefly analyze the time and space complexity of the main algorithm used in the program. Then I will comment on the performance of the program.

### 4.1 Infix expression -> tree

#### (1). Time complexity:

- During this process, we have two step :1.infix->postfix,2.construct the tree.And for both the process ,we need to traverse all the characters in the expression. Therefore ,the time complexity of Infix expression -> tree process is  $O(n)$ .

#### (2). Space complexity:

- The main space cost is the stacks we use in the process and the node space for the tree.We only need  $O(n)$  for the stack and tree ,therefore the total space complexity is  $O(n)$ .

### 4.2 Differentiation

#### (1). Time complexity:

- The differentiation process use the recursion to process the tree.When the worst case happens ,we need to process multiply and divide process ,which will cost 2-3 times of node .Therefore ,theoretically the worst case of time complexity is  $O(2^n)$ .

## (2). Space complexity:

- The same as time complexity ,in worst case the space complexity will be  $O(2^n)$ .

## 4.3 Performance analysis

- The program can perform the operation of auto-differentiation for given infix expression.It can deal with add,minus,multiply,divide,bracket and power quite well.Even in some cases we find the output is not exactly the same as the expected output ,but if we artificially simplify the output expression ,we can find the output answer is right.
- The program also involves some simple simplification ,which can deal with operations that involves 0 and 1 ,which helps to simplify most cases like  $a+0,a*0,a0$ .etc

## 4.4 Things need to be improved

- Just as mentioned before ,the program can't deal with simplifying operations with same variable and combining same sub-expression.This will make some simple answer printed in a more complicated way.But yet I don't find a proper way to deal with it.

```
1      xy: (-xx^2)/(xy*xy)*xy+xx^2/xy = xy: 0
```

- This program also can't deal with any complicated math functions like  $\ln(x),\sin(x),\text{pow}(x,y)$  .etc This will involve a more complicated constructing rule for the tree.

## 5 Appendix

```
1 # include<stdio.h>
2 # include<stdlib.h>
3 # include<string.h>
4
5 typedef struct node{
6     char* variable;           //variable
7     char operator;           //operator
8     int num;                 //num
9     int type;                //1: variable 2:operator 3:
                           num
```

```

10     struct node* right;           //right child
11     struct node* left;           //left child
12 }node;
13
14 int precedence(char op);        //return the priority of
15                                operation
15 int is_right_associative(char op); //return if the operator is
16                                ^
16 int is_operator(char c);        //return if the char is an
17                                operator
17 node* create_node(char* token); //create the node for each
18                                item
18 int get_root(char* exp ,char** var ,node** root); //construct the tree of
19                                the expression
19 int get_token(char* exp ,char** var);      //divide the expression
20                                into operators and operands
20 node* autograd_for_var(char* var ,node* root); //calculate the grad of
21                                the tree
21 char* join(node* root ,char* expression);    //simplify the tree for 0
22                                and 1 and()
22
23 node* new_operator_node(char op ,node* left ,node* right){ //create node
24                                of operator in grad
24     node* n = (node*)malloc(sizeof(node));           //prepare the space
25                                for new node
25     n->type = 2;                                     //mark type as 2
26     n->operator = op;
27     n->left = left;
28     n->right = right;
29     return n;
30 }
31
32 node* new_variable_node(char* var){                //create node of
33                                variable in grad
33     node* n = (node*)malloc(sizeof(node));           //prepare the space
34                                for new node
34     n->type = 1;                                     //mark type as 1
35     n->variable = var;

```

```

36     n->left = n->right = NULL;           //leaf node
37     return n;
38 }
39
40 node* new_number_node(int val){          //create node of
41     number in grad
42     node* n = (node*)malloc(sizeof(node));    //prepare the space
43     for new node
44     n->type = 3;                         //mark type as 3
45     n->num = val;
46     n->left = n->right = NULL;           //leaf node
47     return n;
48 }
49
50
51 int main(void){
52     char expression[100];                  //expression is the
53     infix input
54     char** variables = (char**)malloc(sizeof(char*)*100); //to save the
55     division first ,then save the variable list
56     node* ast_root = NULL;                //the root for the
57     tree
58     scanf("%[^\\n]", expression);
59     // printf("%s", expression);
60     // printf("\n");
61     //build the tree of the expression and find the variables
62     int len_of_var = get_root(expression, variables, &ast_root); //build
63     the tree of expression and get the number of variable
64     printf("There's %d variables in the expression.\n", len_of_var); //
65     print out the length of variable
66     for(int i = 0; i < len_of_var; i++){      //do grad
67         for each variable
68         printf("%s", variables[i]);
69         printf(": ");
70         node* a_of_var;                      //to save
71         the tree after grad
72         a_of_var = autograd_for_var(variables[i], ast_root); ////
73         calculate the grad
74         char* join_exp = (char*)calloc(1000, sizeof(char)); //simplify

```

```

    and turn it into a string
64     join_exp = join(a_of_var ,join_exp);                                //print
    out the result
65     printf("%s" ,join_exp);
66
67     printf("\n");
68 }
69     return 0;
70 }

71
72 int precedence(char op){
73     switch(op){
74         case 'l':
75             case '^': return 3;           //^ is the highest
76             case '*':
77                 case '/': return 2;      /* / is the 2nd
78                 case '+':
79                     case '-': return 1;      //+ - is the lowest
80                     default: return 0;
81     }
82 }
83
84 int is_right_associative(char op){ //return if the operator is ^
85     return op == '^';
86 }
87
88 int is_operator(char c){          //return if the char is an
89     operator
90     return c == '+' || c == '-' || c == '*' || c == '/' || c == '^';
91 }

92 node* create_node(char* token){
93     node* n = (node*)malloc(sizeof(node)); //create the new node
94     n->left = NULL;
95     n->right = NULL;
96     if(is_operator(token[0])){           //if the token is an operator
97         n->type = 2;                   //mark type as 2
98         n->operator = token[0];

```

```

99     }else if((token[0] >= 'a' && token[0] <= 'z') || (token[0] >= 'A' &&
100        token[0] <= 'Z')){ //if the token is variable
101        n->type = 1;
102
103        //mark type as 1
104        n->variable = token;
105    }else{
106
107        //if the token is number
108        n->type = 3;
109
110        //mark type as 3
111        int i = 0;
112        int result = 0;
113
114        //turn the string number into a int type
115        while (token[i] != '\0'){
116            if(token[i] >= '0' && token[i] <= '9'){
117                result = result * 10 + (token[i] - '0');
118
119                //calculate the number
120            }
121
122            else{
123                printf("Invalid number: %s\n" ,token);
124
125                //exit if it is not a
126                //number
127                exit(1);
128            }
129
130            i++;
131        }
132
133        n->num = result;
134    }
135
136
137    return n;
138}
139
140
141 int get_root(char* exp ,char**var ,node** root){
142    int var_num = 0;
143
144    //count the number of variables

```

```

123 //transition of infix into postfix
124 char** stack = (char**)malloc(sizeof(char*) * 20);
125 //operator stack for in-post
126 transition
127 char** post = (char**)malloc(sizeof(char*) * 20);
128 //save postfix expression
129 int len = get_token(exp ,var);
130 //divide into
131 tokens
132 int post_p = 0;
133 int stack_p = 0;
134 int flag = 0;

135 //count backets
136 for(int i = 0;i < len;i++){
137     if(is_operator(var[i][0])){
138         //if operator
139         push into stack or pop operator out
140         while (stack_p > 0 && stack[stack_p - 1][0] != '(' &&
141             (precedence(stack[stack_p - 1][0]) > precedence(var[i][0])
142             ||
143             (precedence(stack[stack_p - 1][0]) == precedence(var[i][0])
144             && !is_right_associative(var[i][0])))) {
145             post[post_p++] = stack[--stack_p];
146             //pop operator with
147             higher priority
148         }
149         stack[stack_p++] = var[i];
150     }else if(var[i][0] == '('){
151         //if ( push
152         into stack
153         stack[stack_p] = var[i];
154         stack_p++;
155         flag++;
156     }else if(var[i][0] == ')'){
157         //if ) pop
158         operator out until (
159         while(stack_p > 0 && stack[stack_p - 1][0] != '('){
160

```

```

144         post[post_p++] = stack[--stack_p];
145                                         //pop operator out
146     }
147     if(stack_p > 0 && stack[stack_p - 1][0] == '('){
148                                         //meet ( and throw it
149         --stack_p;
150     }else{
151         printf("Mismatched parentheses!\n");
152                                         //if () didn't match
153         exit(1);
154     }
155     }else if((var[i][0] >= 'A' && var[i][0] <= 'Z') || (var[i][0] >= 'a' && var[i][0] <= 'z')|| (var[i][0] >= '0' && var[i][0] <= '9')){
156         post[post_p] = var[i];
157                                         //if
158                                         variable or num push into post
159         post_p++;
160     }else{
161         printf("Wrong input!");
162                                         //if not
163                                         allowed input
164         exit(1);
165     }
166 }
167 while(stack_p > 0){
168     // printf("%s",stack[stack_p]);
169     // printf("\n");
170                                         //
171                                         pop out the remaining operators in the stack
172     if(stack[stack_p - 1][0] == '('){
173         printf("Mismatched parentheses!\n");
174                                         //if () didn't match
175         exit(1);
176     }
177     post[post_p++] = stack[--stack_p];
178                                         //push into post
179 }
```

```

169     len = len - flag * 2;
170     printf("Postfix expression:");
171     for(int i = 0;i < post_p;i++)printf("%s",post[i]);
172     printf("\n");
173     //now we start construct the ast tree
174     node** node_stack = (node**)malloc(sizeof(node*) * len);
175                               //stack for elements
176
177     int n_stack_p = 0;
178     for(int i = 0;i < len;i++){
179         node* temp;
180
181         //create new tree node
182         temp = create_node(post[i]);
183         if(temp->type == 2){                                //use
184             operator as root
185             temp->right = node_stack[--n_stack_p];          //pop the operand of the
186                                         operator from the stack
187             temp->left = node_stack[--n_stack_p];
188             node_stack[n_stack_p++] = temp;                  //push the current
189                                         root into the stack
190         }else if(temp->type == 1){
191             node_stack[n_stack_p++] = temp;
192             int flag = 1;
193
194             //check if the variable has appeared
195             for(int i = 0;i < var_num;i++){
196                 if(strcmp(temp->variable ,var[i]) == 0){
197                     flag = 0;
198                 }
199             }
200             if(flag){
201
202                 //if the first time count and push into varlist
203                 var[var_num++] = temp->variable;
204             }

```

```

194     }else{
195         node_stack[n_stack_p++] = temp;
196         //if number ,push
197         into stack
198     }
199     *root = node_stack[0];
200     //save
201     the root for the tree
202     return var_num;
203     //
204     return the number of variables
205 }
206
207 int get_token(char* exp ,char** var){
208     int count = 0;
209
210     //count the elements
211     int i = 0;
212     var[count] = (char*)malloc(sizeof(char) * 20);
213     //use var to save the divided
214     element for the tree
215     int temp_p = 0;
216     while(exp[i] != '\0'){
217         //go
218         through every char
219         if((exp[i] >= 'A' && exp[i] <= 'Z') || (exp[i] >= 'a' && exp[i] <=
220             'z') || (exp[i] >= '0' && exp[i] <= '9')){
221             var[count][temp_p] = exp[i];
222             //if number or
223             letter
224             temp_p++;
225         }else if(is_operator(exp[i]) || exp[i] == '(' || exp[i] == ')'){
226             if(temp_p != 0){
227                 //
228                 save number or variable until meet an operator
229                 var[count][temp_p] = '\0';
230                 //end it as a

```

```

214         string
215         count++;
216         var[count] = (char*)malloc(sizeof(char) * 20);
217         temp_p = 0;
218     }
219     var[count][temp_p++] = exp[i];
220                                         //add the
221                                         operator
222     var[count][temp_p++] = '\0';
223     count++;
224     var[count] = (char*)malloc(sizeof(char) * 20);
225                                         //start next element
226     temp_p = 0;
227 }
228 i++;
229 }
230 if(temp_p != 0){
231                                         //
232                                         save the last element
233     var[count][temp_p] = '\0';
234     count++;
235 }
236 return count;
237 }

238 node* autograd_for_var(char* var ,node* root){
239     if(!root){
240         return NULL;
241                                         //meet
242                                         the leaf node
243     }
244     if(root->type == 3){
245         return new_number_node(0);
246                                         //for
247                                         const ,grad is 0
248     }
249     if(root->type == 1){
249                                         //for
250                                         var ,check if it is the target var for grad
251                                         if(strcmp(root->variable ,var) == 0){
252                                             return new_number_node(1);
253                                         //is

```

```

243         return 1
244     }else{
245         return new_number_node(0); //or
246     }
247 }
248 if(root->type == 2){ //for
249     operator ,recursionly process the child
250     node *u = root->left ,*v = root->right; //u,v
251     as the operand of the operator
252     node *du = autograd_for_var(var ,u) ,*dv = autograd_for_var(var ,v)
253     );//du,dv as the grad of u,v
254     switch (root->operator){
255         case '+': //d(u+v)
256             )=du+dv
257             return new_operator_node('+',du ,dv);
258         case '-': //d(u-v)
259             )=du-dv
260             return new_operator_node('-',du ,dv);
261         case '*': //d(u*v)
262             )=du*u+u*dv
263             return new_operator_node('+',new_operator_node('*',du ,v)
264                         ,new_operator_node('*',u ,dv));
265         case '/': //d(u/v)
266             )=(du*v-u*dv)/(v*v)
267             return new_operator_node('/',new_operator_node('-',new_operator_node('*',
268                     new_operator_node('*',du ,v),new_operator_node('*',u ,
269                     dv)) ,new_operator_node('*',v ,v)));
270
271         case '^':{
272             int is_u_const = (u->type == 3);
273             int is_v_const = (v->type == 3);
274             if(is_v_const && !is_u_const){ //d(u^c)
275                 )=c*u^(c-1)*du
276                 int c = v->num;
277                 node* new_exp = new_number_node(c - 1);
278                 node* u_pow_c_minus_1 = new_operator_node('^',u ,
279                     new_exp);

```

```

267         return new_operator_node('*' ,new_operator_node('*',
268                                     new_number_node(c),u_pow_c_minus_1),du);
269     }else if(is_u_const && !is_v_const) {                                //d(c^v
270         )=c^v*ln(c)*dv
271         int c = u->num;
272         node* ln_c = new_operator_node('l' ,u ,NULL);      //ln(u)
273         node* c_pow_v = new_operator_node('^' ,u ,v);
274         return new_operator_node('*' ,new_operator_node('*',
275                                     c_pow_v,ln_c),dv);
276     }else{                                         //d(u^v
277         )=u^v*(dv*ln(u)+v*du/u)
278         node* ln_u = new_operator_node('l' ,u ,NULL);      //ln(u)
279         node* term1 = new_operator_node('*' ,dv ,ln_u);
280         node* u_div = new_operator_node('/' ,du ,u);      //du/u
281         node* term2 = new_operator_node('*' ,v ,u_div);
282         node* sum = new_operator_node('+',term1 ,term2); //dv*
283         ln(u)+v*du/u
284         node* u_pow_v = new_operator_node('^' ,u ,v);
285         return new_operator_node('*' ,u_pow_v ,sum);
286     }
287 }
288 default:
289     printf("Unsupported operator: %c\n" ,root->operator); //operator unsupported
290     exit(1);
291 }
292 char* join(node* root ,char* expression){
293     if(!root){                                         //meet the leaf node
294         return NULL;
295     }
296     if(root->type == 1){                           //variable cat into
297         expression
298         strcat(expression ,root->variable);
299         return expression;

```

```

298     }
299     else if(root->type == 3){           //number cat into
300         expression
301         char num_str[20];
302         sprintf(num_str ,"%d" ,root->num);      //change int into
303         strcat(expression ,num_str);
304         return expression;
305     }
306     else if(root->type == 2){
307         if(root->operator == 'l'){           //opertor 'l' as ln
308             strcat(expression ,"ln(");
309             join(root->left ,expression);      //put ln(x) into
310             strcat(expression ,")");
311             return expression;
312         }
313         char* left_str = (char*)calloc(100 ,sizeof(char)); //save the
314             middle result for the two child
315         char* right_str = (char*)calloc(100 ,sizeof(char));
316         left_str = join(root->left ,left_str);      //recursively
317             calculate the result
318         right_str = join(root->right ,right_str);
319
320         if(root->operator == '+'){           //simplify when 0+a,
321             a+0,0+0
322             if(strcmp(left_str ,"0") == 0){      //0+a->a
323                 strcat(expression ,right_str);
324                 goto done;
325             }
326             if(strcmp(right_str ,"0") == 0){      //a+0->a
327                 strcat(expression ,left_str);
328                 goto done;
329             }
330         }
331         else if(root->operator == '-'){       //simplify when 0-a,
332             a-0

```

```

329     if(strcmp(left_str , "0") == 0 && strcmp(right_str , "0") == 0){
330         strcat(expression , "0");                                //0-0->0
331         goto done;
332     }
333     if(strcmp(left_str , "0") == 0){
334         strcat(expression , "-");                            //0-a->-a
335         strcat(expression , right_str);
336         goto done;
337     }
338     if(strcmp(right_str , "0") == 0){
339         strcat(expression , left_str);                      //a-0->a
340         goto done;
341     }
342 }
343 else if(root->operator == '*'){                         //simplify when 0*a,
344     a*0, 1*a, a*1
345     if(strcmp(left_str , "0") == 0 || strcmp(right_str , "0") == 0){
346         strcat(expression , "0");                                //0*0->0
347         goto done;
348     }
349     if(strcmp(left_str , "1") == 0){
350         strcat(expression , right_str);                      //1*a->a
351         goto done;
352     }
353     if(strcmp(right_str , "1") == 0){
354         strcat(expression , left_str);                      //a*1->a
355         goto done;
356     }
357 }
358 else if(root->operator == '/'){                         //simplify when 0/a,
359     a/1
360     if(strcmp(left_str , "0") == 0){
361         strcat(expression , "0");                                //0/a->0
362         goto done;
363     }
364     if(strcmp(right_str , "1") == 0){
365         strcat(expression , left_str);                      //a/1->a

```

```

365         goto done;
366     }
367 }else if(root->operator == '^'){
368     //simplify when a^1,
369     //a^0
370     if(strcmp(right_str , "1") == 0){
371         strcat(expression ,left_str);           //a^1->a
372         goto done;
373     }
374     if(strcmp(right_str , "0") == 0){
375         strcat(expression , "1");             //a^0->1
376         goto done;
377     }
378     if(root->left->type == 2 && precedence(root->operator) >
379         precedence(root->left->operator)){
380         strcat(expression , "(");
381         strcat(expression ,left_str);          //add () if the left
382         child operator is smaller than root
383         strcat(expression , ")");
384     }else{
385         strcat(expression ,left_str);          //else no ()
386     }
387     int len = strlen(expression);
388     expression[len] = root->operator;        //push the operator
389     expression[len + 1] = '\0';
390
391     if((root->right->type == 2 && precedence(root->operator) >
392         precedence(root->right->operator))
393     || (root->right->type == 2 && root->operator == '/')){
394         strcat(expression , "(");
395         strcat(expression ,right_str);          //add () if the
396         right child operator is smaller than root
397         strcat(expression , ")");
398     }else{
399         strcat(expression ,right_str);          //else no ()
400     }
401
402     done:

```

```
398     free(left_str);                      //free the temporal  
      string  
399     free(right_str);  
400     return expression;                  //return expression  
401 }  
402  
403 return NULL;                         //return NULL when  
      something wrong  
404 }
```

## 6 Declaration

I hereby declare that all the work done in this project titled “Project 2 : Autograd for Algebraic Expressions ” is of my independent effort.