

Algorithmics	Student information	Date	Number of session
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Activity 1. Basic recursive models

- Subtraction 1: there is one recursive call, so $a = 1$, therefore complexity is $O(n^{(k+1)})$, where $k = 0$, so $O(n)$.
- Subtraction 2: as with the previous case, $a = 1$, however $k = 1$ in this case, making the complexity $O(n^2)$.
- Subtraction 3: now, $a = 2$, bigger than 1, so it follows the $(a^{(n/b)})$ pattern, where $b = 1$, so the complexity is $O(2^{(n)})$.
- Subtraction 4: target complexity is $O(3^{(n/2)})$. For this, a has to be bigger than 1, and b has to be 2. For this I made 3 recursive call where I subtract 2.
- Division 1: $a < b^k$, being $a = 1$, $b = 3$, $k = 1$, therefore the complexity is $O(n)$.
- Division 2: $a = b^k$, where $a = 2$, $b = 2$, and $k = 1$, making the complexity $O(n \cdot \log n)$.
- Division 3: $a > b^k$, because $a = 2$ and $b = 2$, but $k = 1$, so $2 > (2^1 = 1)$, the complexity is then $O(n^{\log 2})$.
- Division 4: target complexity is $O(n^2)$ having a number of subproblems (a) of 4. The condition for the $O(n^k)$ pattern is $a < b^k$. To fulfill that $4 < b^k$, I used a complexity of $O(n^2)$ so that $k = 2$, and made the recursive call divide by 3.