

## Agenda

1. Submatrix sum queries
2. Max submatrix sum
3. Sum of all submatrices

## Representing a submatrix

	0	1	2	3	4
0					
1					
2					
3					
4					
5					

TL  $\rightarrow 2, 2$

BR  $\rightarrow 4, 3$

Top-left and bottom right corners are required.

Q-1 Given a  $N \times M$  matrix and queries, for each query find submatrix sum.

queries

(TL) (BR)

$x_1$	$y_1$	$x_2$	$y_2$	ans
2	1	4	3	20
3	2	5	4	36

$A \Rightarrow$

	0	1	2	3	4
0	7	1	-6	3	12
1	10	5	-1	0	9
2	6	4	-3	8	11
3	13	-8	-5	12	4
4	3	2	1	9	8
5	4	3	-2	6	3

i) brute force : go on every query, and travel that submatrix to find sum.

$O(N \times M \times Q)$

ii) improvise : prefix sum

Prefix sum

$A \Rightarrow$

	0	1	2	3	4
0	7	1	-6	3	12
1	10	5	-1	0	9
2	6	4	-3	8	11
3	13	-8	-5	12	4
4	3	2	1	9	8
5	4	3	-2	6	3

	0	1	2	3	4
0	7			5	
1		23	16		
2					
3					
4					
5					

$ps[i][j] \rightarrow$  sum of matrix from 0,0 to i,j  
(TL) (BR)

A :

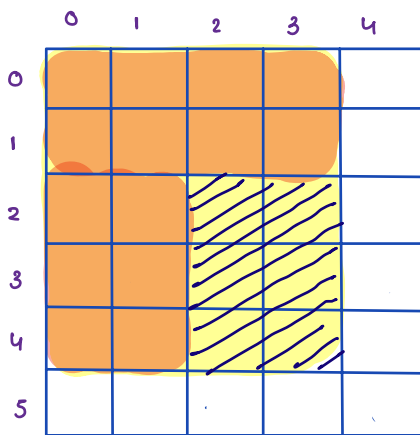
	0	1	2
0	a	b	c
1	d	e	f
2	g	h	i

apply row-wise  
Prefix sum →

	0	1	2
0	a	a+b	a+b+c
1	d	d+e	d+e+f
2	g	g+h	g+h+i

apply  
col-wise  
Prefix sum

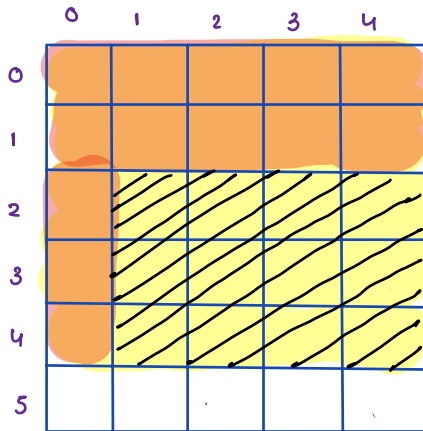
	0	1	2
0	a	a+b	a+b+c
1	a+d	a+b+d+e	a+b+c+d+e+f
2	a+d+g	a+b+d+e+g+h	a+b+c+d+e+f+g+h+i



TL  
(2,2)

BR  
(4,3)

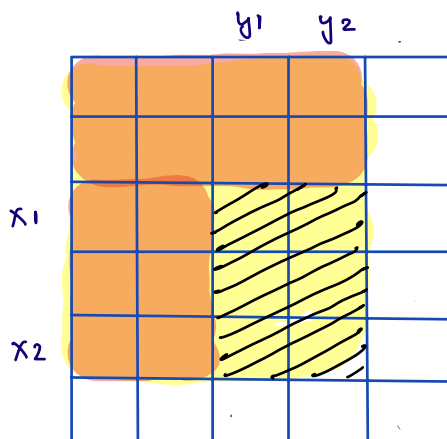
$$ps[4][3] - ps[1][3] - (ps[4][1] - ps[1][1])$$



TL  
(2,1)

BR  
(4,4)

$$ps[4][4] - ps[1][4] - (ps[4][0] - ps[1][0])$$



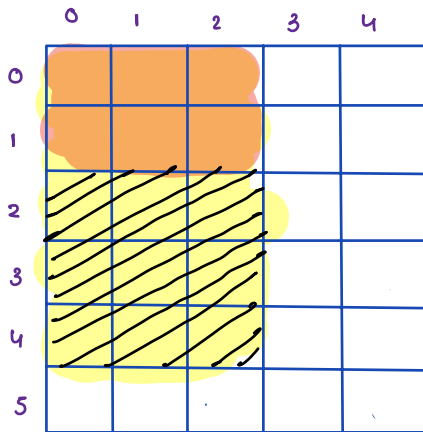
TL  
 $x_1, y_1$

BR  
 $x_2, y_2$

$$ps[x_2][y_2] - ps[x_1-1][y_2] - (ps[x_2][y_1-1] - ps[x_1-1][y_1-1])$$

$$ps[x_2][y_2] - ps[x_1-1][y_2] - (ps[x_2][y_1-1] - ps[x_1-1][y_1-1])$$

$$\Rightarrow ps[x_2][y_2] - ps[x_1-1][y_2] - ps[x_2][y_1-1] + ps[x_1-1][y_1-1]$$



TL

BR

(2,0)

(4,2)

$x_1 \ y_1$

$x_2 \ y_2$

$$ps[4][2] - ps[1][2] - \cancel{ps[4][1]} + \cancel{ps[1][1]}$$

```
void submatrixSumQuery( int [][ ] A, int [][ ] Q ) {
```

```
    int [][ ] ps = prefixSum2D (A);
```

```
    for(int i=0; i<Q.length; i++) {
```

```
        int x1 = Q[i][0];
```

```
        int y1 = Q[i][1];
```

```
        int x2 = Q[i][2];
```

```
        int y2 = Q[i][3];
```

iter:  $2(N \cdot M) + Q$

```
        int sum = 0;
```

Tc:  $O(N \cdot M + Q)$

```
        sum += ps[x2][y2];
```

Sc:  $O(N \cdot M)$

```
        if (x1 > 0)
```

```
            sum -= ps[x1-1][y2];
```

```
        if (y1 > 0)
```

```
            sum -= ps[x2][y1-1];
```

```
        if (x1 > 0 && y1 > 0)
```

```
            sum += ps[x1-1][y1-1];
```

```
        println(sum);
```

}

}

A :

	0	1	2
0	a	b	c
1	d	e	f
2	g	h	i

apply row-wise  
prefix sum →

	0	1	2
0	a	a+b	a+b+c
1	d	d+e	d+e+f
2	g	g+h	g+h+i

apply  
col-wise  
prefix sum

	0	1	2
0	a	a+b	a+b+c
1	a+d	a+b+d+e	a+b+c+d+e+f
2	a+d+g	a+b+d+e+g+h	a+b+c+d+e+f+g+h+i

```
int [][] prefixSum2D (int [][] A) {
```

```
    int n = A.length;
```

```
    int m = A[0].length;
```

```
    int [][] ps = new int [n] [m];
```

```
    // apply row by row ps
```

```
    for (int i=0; i<n; i++) {
```

```
        for (int j=0; j<m; j++) {
```

```
            if (j==0) {
```

```
                ps[i][j] = A[i][j];
```

```
            }  
            else {
```

```
                ps[i][j] = ps[i][j-1] + A[i][j];
```

```
            }
```

```
        }
```

```
    }
```

```
    // apply col by col ps
```

```
    for (int j=0; j<m; j++) {
```

```
        for (i=1; i<n; i++) {
```

```
            ps[i][j] = ps[i-1][j] + A[i][j];
```

```
        }
```

```
    }
```

```
    return ps;
```

```
}
```



Q.2 Given row wise and column wise sorted matrix, find maximum submatrix sum.

A =

	0	1	2	3
0	-20	-16	-4	8
1	-10	-8	2	14
2	-1	6	21	30
3	5	7	28	42

max element is always a part of our ans and max element is present at  $(n-1, m-1)$

A =

	0	1	2	3
0	-20	-16	-4	-1
1	-10	-8	-2	5
2	-4	2	4	8

In ans submatrix BR is fixed to

$(n-1, m-1)$

⇓

Go on all TL possible and find the best ans.

A =

	0	1	2
0	-50	-40	-30
1	-25	-20	-15
2	-14	-14	-3

TL

0,0      2,0

0,1      2,1

0,2      2,2

0,3      2,3

1,0

1,1

1,2

1,3

A =

	0	1	2	3
0	-20	-16	-4	-1
1	-10	-8	-2	5
2	-4	2	4	8

```
int maxSumSubmatrix (int [][] A) {
```

```
int [][] ps = prefixSum2D(A);
```

```
int n = A.length;
```

```
int m = A[0].length;
```

```
int x2 = n-1; } → BR
```

```
int y2 = m-1;
```

```
int ans = Integer.MIN_VALUE;
```

```
for (int i=0; i<n; i++) {
```

```
    for (int j=0; j<m; j++) {
```

```
        int x1 = i, y1 = j;
```

```
        // sum of submatrix TL: x1, y1 & BR: x2, y2
```

```
        int sum = 0;
```

```
        sum += ps[x2][y2];
```

```
        if (x1 > 0)
```

```
            sum -= ps[x1-1][y2];
```

```
        if (y1 > 0)
```

```
            sum -= ps[x2][y1-1];
```

```
        if (x1 > 0 && y1 > 0)
```

```
            sum += ps[x1-1][y1-1];
```

```
        ans = Math.max(ans, sum);
```

```
    }
```

```
return ans;
```

```
}
```

Q-3 Given a  $N \times M$  Submatrix, find sum of all submatrices sum.

	0	1
0	2	1
1	3	4

TL	BR	Sum
0,0	0,0	2
0,0	0,1	3
0,0	1,0	5
0,0	1,1	10
0,1	0,1	1
0,1	1,1	5
1,0	1,0	3
1,0	1,1	7
1,1	1,1	4
		+
		40

Improved idea { if  $A[i][j]$  is coming in  
 $\times$  no. of matrices then  
its contribution is final  
ans is:  $x * A[i][j]$

	0	1	2	3	4
0	TL	TL	TL		
1	TL	TL	TL		
2	TL	TL	TL		
3	TL	TL	TL BR	BR	BR
4			BR	BR	BR
5			BR	BR	BR

At  $i][j]$  is coming in how many submatrices.

valid TL  $\rightarrow 12$

valid BR  $\rightarrow 9$

$$12 * 9 = 108$$

				j		m-1
	0	1	2	3	4	
0	TL	TL	TL			
1	TL	TL	TL			
2	TL	TL	TL			
i	3	TL	TL	TL BR	BR	BR
4			BR	BR	BR	
n-1	5		BR	BR	BR	

valid TL:  $(i+1) * (j+1)$

valid BR:  $(n-i) * (m-j)$

$$\text{freq} \Rightarrow (i+1) * (j+1) * (n-i) * (m-j)$$

```
int solve (int r[] [] A) {
```

```
    int n = A.length;
```

```
    int m = A[0].length;
```

```
    for (int i = 0; i < n; i++) {
```

```
        |   for (int j = 0; j < m; j++) {
```

```
            |   int freq = (i+1)*(j+1) * (n-i)*(m-j);
```

```
            |   ans += freq * A[i][j];
```

```
            |   }
```

```
        }
```

```
    return ans;
```

```
}
```

Doubts

$$(i+1)^* (j+1) = (n-i)^* (m-j);$$

$$n=2, \quad m=2$$

	0	1
0	2	1
1	3	4

$(i, j)$	freq	ans
$0, 0$	$1 * 1 * 2 * 2 = 4$	8
$0, 1$	$1 * 2 * 2 * 1 = 4$	4
$1, 0$	$2 * 1 * 1 * 2 = 4$	12
$1, 1$	$2 * 2 * 1 * 1 = 4$	16

Q. No. of minimum swaps to make all ele  $\leq B$  together.

A: [ 1 10 15 2 4 13 4 ]      B = 5  
          0 1 2 3 4 5 6  
                                  ans = 2

A:  $[25 \quad 30 \quad 2 \quad 18 \quad 7 \quad 6 \quad 9 \quad 50 \quad 33]$   $B = 10$   
 $\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \end{matrix}$   $ans = 1$

A: [19 11 3 9 7 25 6 20 4] B=10  
0 1 2 3 4 5 6 7 8 ans = 2

A: [ 19 11 3 9 7 25 6 20 4 ]  
 0 1 2 3 4 5 6 7 8

B = 10

K = 5 ( $\leq B$ )

s e swaps

0 4 2 (no.  $\text{ele} > B$ , ans. of 1<sup>st</sup> window)

1 5 2

2 6 1

3 7 2

4 8 2

} if (A[s-1] > B) {  
 swaps--;  
 3  
 if (A[e] > B) {  
 swaps++;  
 3