


- Welcome to PS + DSA Module 
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- Recoding (DSA Instructor + Content creator)
- 2+ years of Teaching Experience

- Programming constructs
- Problem solving (efficient)

Q.1 Count of factors

24 → 1, 2, 3, 4, 6, 8, 12, 24

10 → 1, 2, 5, 10

N, factors : 1 to N

N=10

```
int countFactors (int N) {
    int count = 0;
    for (int i=1; i<=N; i++) {
        if (N%i == 0) {
            // i is factor of N
            count++;
        }
    }
    return count;
}
```

i → 1 2 3 4 5 6 7 8 9 10

count = ~~1~~ ~~2~~ ~~4~~

it's: N

Execution time

→ value of N

→ system configuration

Assume

10^8 iteration per second

N	iterations (N)	time
10^8	10^8	1 Sec
10^9	10^9	10 Sec
10^{18}	10^{18}	10^{10} Sec = 317 years

$$10^8 \text{ itr} \rightarrow 1 \text{ Sec}$$

$$1 \text{ itr} \rightarrow \frac{1}{10^8} \text{ Sec}$$

$$10^9 \text{ itr} \rightarrow \frac{1}{10^8} \times 10^9$$
$$= 10 \text{ Sec}$$

$$1 \text{ itr} \rightarrow \frac{1}{10^8} \text{ Sec}$$

$$10^{18} \text{ itr} \rightarrow \frac{1}{10^8} \times 10^{18}$$
$$= 10^{10} \text{ Sec}$$

Improvisation

$i \times j = N$ (both i and j are factors of N)

$j = \frac{N}{i}$ (both i and N/i are factors of N)

$N = 24$

i	$\frac{N}{i}$
1	24
2	12
3	8
4	6
<hr/>	
6	4
8	3
12	2
24	1

$$i \leq \frac{N}{i}$$

$$i \times i \leq N$$

$$i \leq \sqrt{N}$$

$N = 100$

i	$\frac{N}{i}$
1	100
2	50
4	25
5	20
10	10
<hr/>	
20	5
25	4
50	2
100	1

Observations a) all the factor of n are present in the 1st part.
b) we are in 1st part till $i \leq \sqrt{N}$.

```
int countFactors (int N) {
```

```
    int count = 0;
```

```
    for (int i = 1; i <=  $\sqrt{N}$ ; i++) {
```

```
        if (N % i == 0) {
```

```
            if (i == N/i) {
```

```
                count++;
```

```
            }
```

```
        else {
```

```
            // both i and N/i are factors
```

```
            count += 2;
```

```
        }
```

```
    }
```

```
    return count;
```

```
}
```

iterations $\rightarrow \sqrt{N}$

$N = 15$

$\sqrt{15} \approx 3$

i	count	
1	2	(1, 15)
2		
3	4	(3, 5)

$N = 4$

$\sqrt{4} = 2$

i	count	
1	2	(1, 4)
2	3	(2, 2)

N	itr (\sqrt{N})	time
10^{18}	10^9	10 sec

$10^8 \text{ itr} \rightarrow 1 \text{ sec}$

$10^9 \text{ itr} \rightarrow \frac{1}{10^8} \times 10^9$
 $= 10 \text{ sec}$

Q.2 check if no. is prime or not.

prime no. \rightarrow two factors (1 and no. itself)

{10, 11, 23, 2, 25, 27, 31}

\rightarrow primes: 11, 23, 2, 31

```
boolean is_prime(int N) {  
    if (count_factors(N) == 2) {  
        return true;  
    }  
    else {  
        return false;  
    }  
}
```

}

\sqrt{N} iterations

Q.3 Given an Array, reverse it.

A = { 10, 20, 30, 40 }

0 1 2 3

40 30 20 10

{ ~~10~~, ~~20~~, ~~30~~, ~~40~~ }

0 1 2 3

s	e
0	3
1	2
2	1

A = { 10, 20, 30, 40, 50, 60, 70 }

0 1 2 3 4 5 6

70 60 50 40 30 20 10

{ ~~10~~, ~~20~~, ~~30~~, ~~40~~, ~~50~~, ~~60~~, ~~70~~ }

0 1 2 3 4 5 6

s
e

s	e
0	6
1	5
2	4
3	3

work till s < e

```
void reverse (int [ ] A) {
```

```
    int s = 0;
```

```
    int e = A.length - 1;
```

```
    while (s < e) {
```

```
        // swap A[s] and A[e]
```

```
        int temp = A[s];
```

```
        A[s] = A[e];
```

```
        A[e] = temp;
```

```
        s++; e--;
```

```
    }
```

0	1	2	3	4
10 50	20 40	30	40 20	50 10

0	1	2	3
10 40	20 30	30 20	40 10

Q.4 Reverse part of an array.

A = { 10, 19, 40, 50, 25, 64, 38, 37 }
0 1 2 3 4 5 6 7

S = 2

e = 5

{ 10, 19, 64, 25, 50, 40, 38, 37 }
0 1 2 3 4 5 6 7

void reversePart (int A[], int S, int e) {

while (S < e) {

// swap A[S] and A[e]

int temp = A[S];

A[S] = A[e];

A[e] = temp;

S++; e--;

}

}

A = { 10, 19, 25, 24, 38, 30, 18 }
0 1 2 3 4 5

S
e

S = 1

e = 3

Q-5 Given an array, rotate it from last to first k times.

hoogle, amazon

constraints : (i) don't use extra array

(ii) do it in efficient way

A = 10 20 30 40 50 60

$k=3$

↓

60 10 20 30 40 50

↓

50 60 10 20 30 40

↓

40 50 60 10 20 30

A =

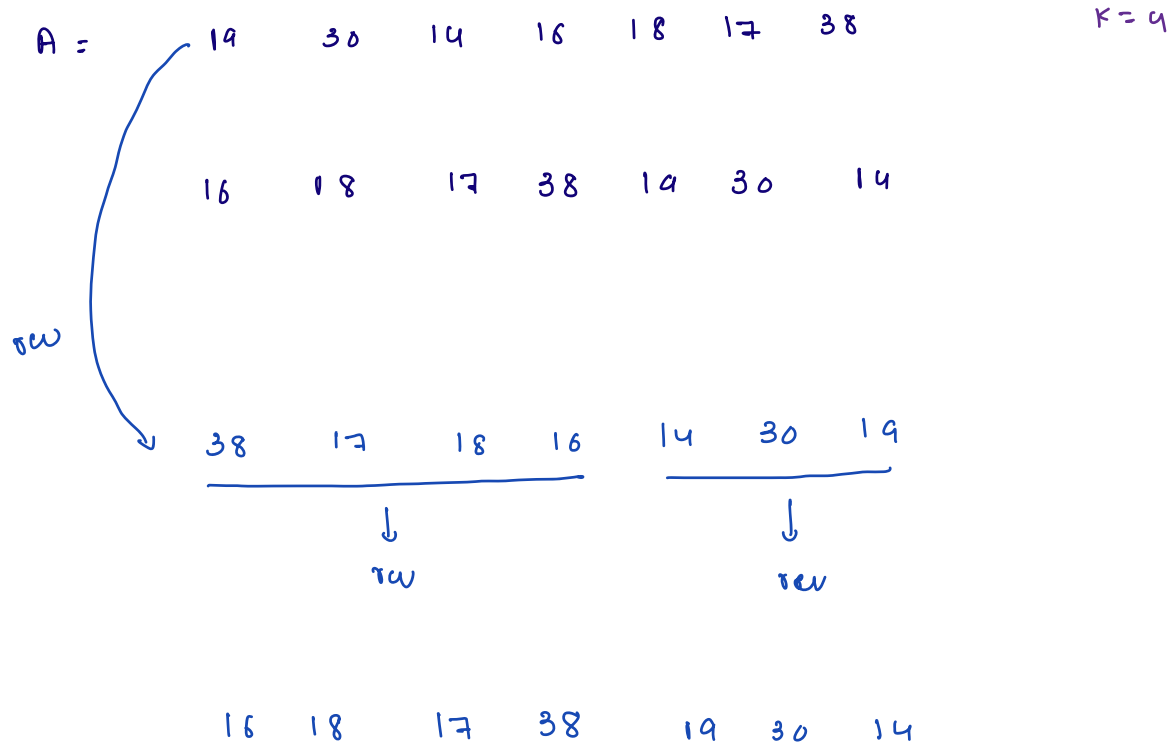
14 30 14

16 18 17 38

$k=4$

16 18 17 38

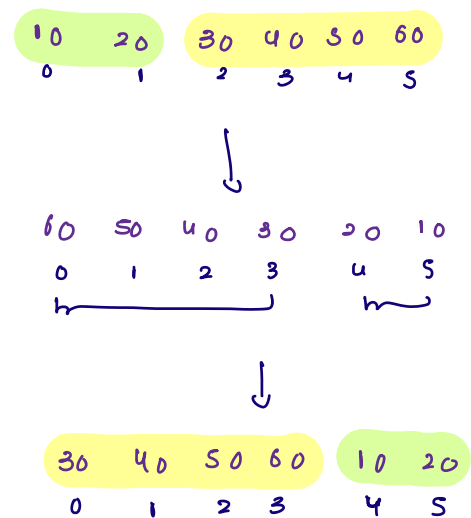
14 30 14



i) reverse entire array

(ii) reverse the first k elements

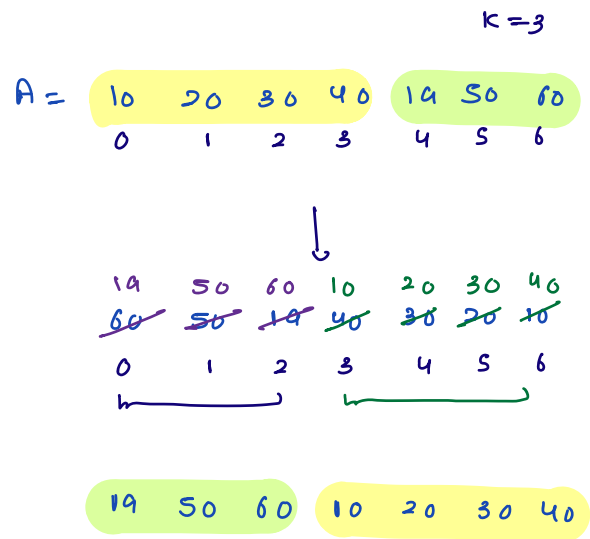
(iii) rev. the remaining elements



```
void rotate (int [] A, int k) {
    int n = A.length;
```

```
    reversePart (A, 0, n-1);
    reversePart (A, 0, k-1);
    reversePart (A, k, n-1);
```

5



```
void rotate (int [] A, int k) {
    int n = A.length;
```

```
    reversePart (A, 0, n-1); ✓
```

```
    reversePart (A, 0, k-1);
```

```
    reversePart (A, k, n-1);
```

5

Array
index
out of bound

what if $k > n$

A = 10 20 30 40 50

k = 12

10 20 30 40

$k = 4$

$n = 4$



40 10 20 30



30 40 10 20



20 30 40 10



10 20 30 4

$n = 5$

$nx = 2$

$k = 12$

$n = 6$

$nx = 3$

$k = 27$

void rotate (int [] A, int k) {

int n = A.length;

k = k % n;

reversePart (A, 0, n-1);

reversePart (A, 0, k-1);

reversePart (A, k, n-1);

log basics

$$\log_b a$$

b power what is equals to a

$$\log_b a = c$$

$$b^c = a$$

$$\log_2 64 = 6$$

$$\log_2 8 = 3$$

$$\log_{10} 10000 = 4$$

$$\log_3 81 = 4$$

$$i) N = 2^k$$

$$k = \log_2 N$$

$$ii) \log_b b^N = N$$

Problem Solving and DSA

- Time complexity
- Arrays: Prefix sum, subarrays, sliding window, 2D matrices.
- Bit manipulation
- Hashing
- Recursion
- Sorting
- Searching
- 2 pointer technique
- Strings
- Linked list
- Trees, BST, heaps
- Dynamic programming
- Graphs

52 classes

doubts
=

$$\log_b a = ?$$

$$b^? = a$$

$$\log_2 31 = \underline{\underline{4.954}}$$

↳ int : 4

$$\log_2 32 = 5$$

$$2^{\cancel{5}} = 32$$

$$\log_2 16 = 4$$

$$2^{\cancel{4}} = 16$$

$$\log_3 81 = 4$$

$$3^{\cancel{4}} = 81$$