

# Transactions.

[Talk - 5 min]

## Agenda

- ① Transactions
- ② Perspective of transactions
  - A
  - C
  - I
  - D
- ③ Read & Write ops
- ④ Commit & Rollbacks
- ⑤ Transaction Isolation level
  - lost updates
  - dirty reads
  - phantom reads
  - Non-repeatable reads
- ⑥ Deadlocks

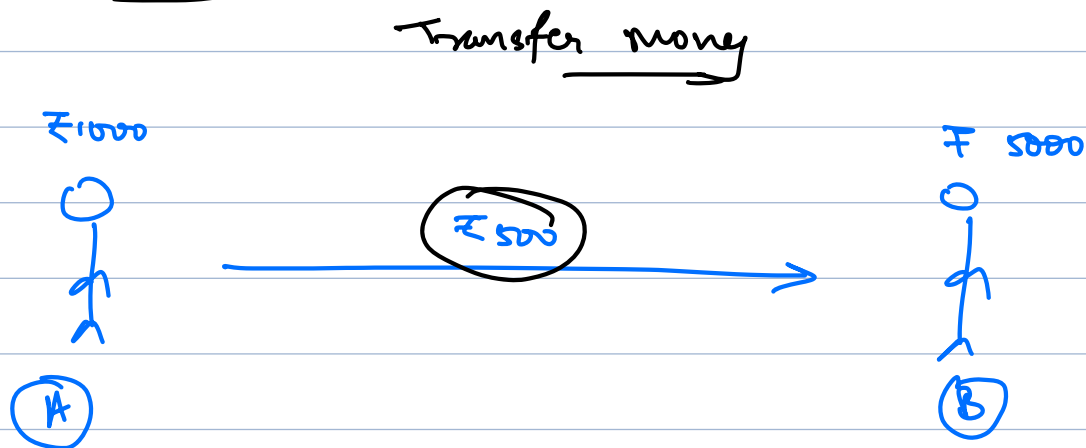
(5 Talk)

5 ✓  
4 ✓  
3 ✓  
2 ✓  
1 ✓

\_\_\_\_\_

What are transactions.

Bank Example.



Accounts

id	name	balance	created_At	branch_id
1	A	₹1000	-	-
1	B	₹5000	-	-

- ① Check the balance A ✓ DB call
- ② Check > 500 ✗ App work
- ③ Reduce balance of A by 500 ✓ DB call
- ④ Increase balance of B by 500 ✓ DB call

← transferMoney (from, to, amount)  
≡ if ( bal >= amount )  
≡

else  
f  
,

⇒ Transfer Money may be getting executed by multiple people at the same time.

$A \rightarrow B$  ( $\leq 500$ )

$C \rightarrow B$  ( $\geq 500$ )

- ① Check the balance A → Read ✓
- ② Check > 500 → Nothing
- ③ Reduce balance of A by 500 → write
- ④ Increase balance of B by 500 → read & write

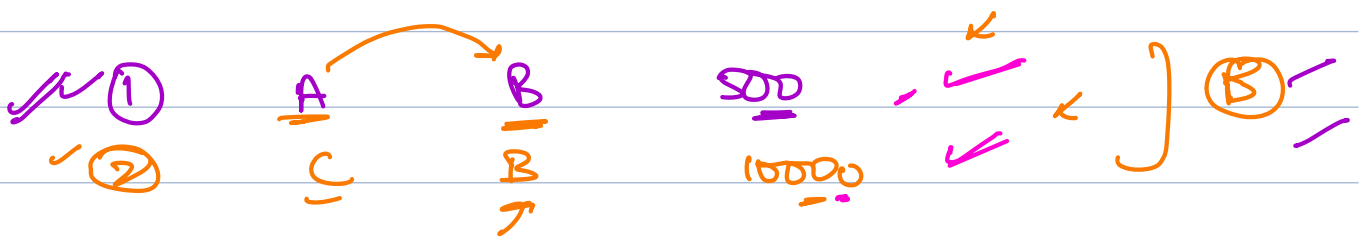
$B += 500$

↓

$B = B + 500$

This is NOT a single operation

Read B → temp  
temp → temp + 500  
Write ← temp



A → ~~10000~~ 500      transfer Money (a, b amount) {

B → ~~5000~~ ~~500~~ 15000

C → ~~15000~~ 5000  
21000      20500

Read A → x    ① → 10000    ② → 15000

if (x != amount) {

write A ← x - amount    ① → 500    ② → 5000

Read B → x    ① → 5000    ② → 15000

x = x + amount

write B ← x    ① → 5500    ② → 15000

① B → ~~5000~~ ~~5500~~ 15000

① Inconsistent / illogical state

② Entire operation may be incomplete

Transaction → A set of DB operations logically grouped together to perform a task

Transfer Money

A ← x
A ← x -    -
B ← x
B ← x +    -

# Expectations / Properties of a transaction

## ( ACID Properties )

A → Atomicity  
C → Consistency  
I → Isolation  
D → Durability

### Atomicity

↓  
from the word Atom → smallest indivisible unit.

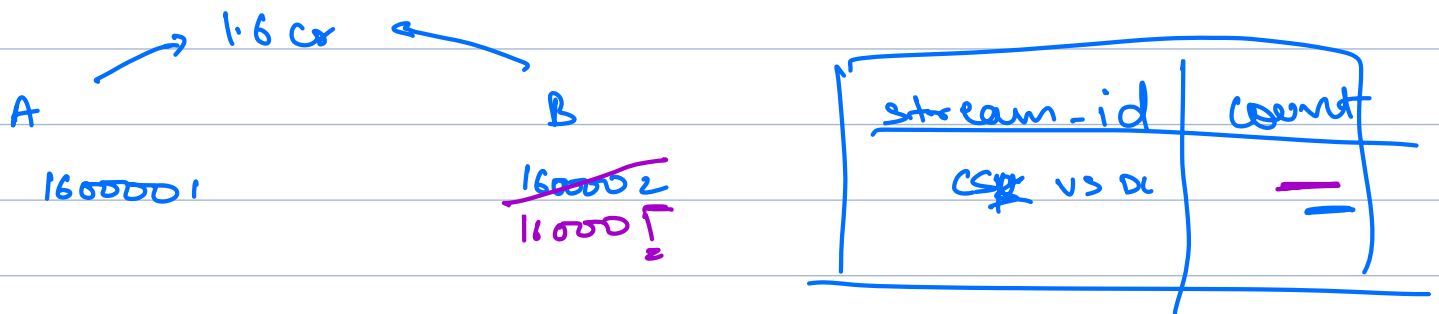
- Transaction should appear indivisible to an end user
- to an outside it should feel that either nothing has happened or everything is done.
- A transaction should never end in an intermediary state.

### Consistency

- correctness
- exactness
- accuracy
- logical correctness

Example where consistency is not needed.

Hotfix

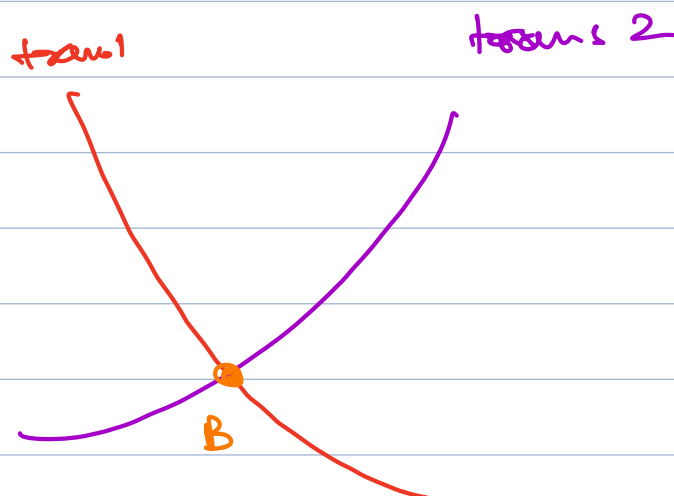


consistency hampered  
But performance improved.

The state of DB trans started and also after trans completed should be accurate.

## ISOLATION

→ one transaction shouldn't affect other transactions running at the same time on the same data.



## Durability .

once a transaction is completed we would want its work to stay persistent .

↳ Store it on disk .

Break → 5 min

8:14

↓

8:23 AM

## Commit and Roll back

Students			
id	name	prp	b-id
1	ujjwal	80	2

[ update students  
set prp = 100  
when id = 1 ]

↓

select \* from Students

↓

prp = 100;

## Auto commit

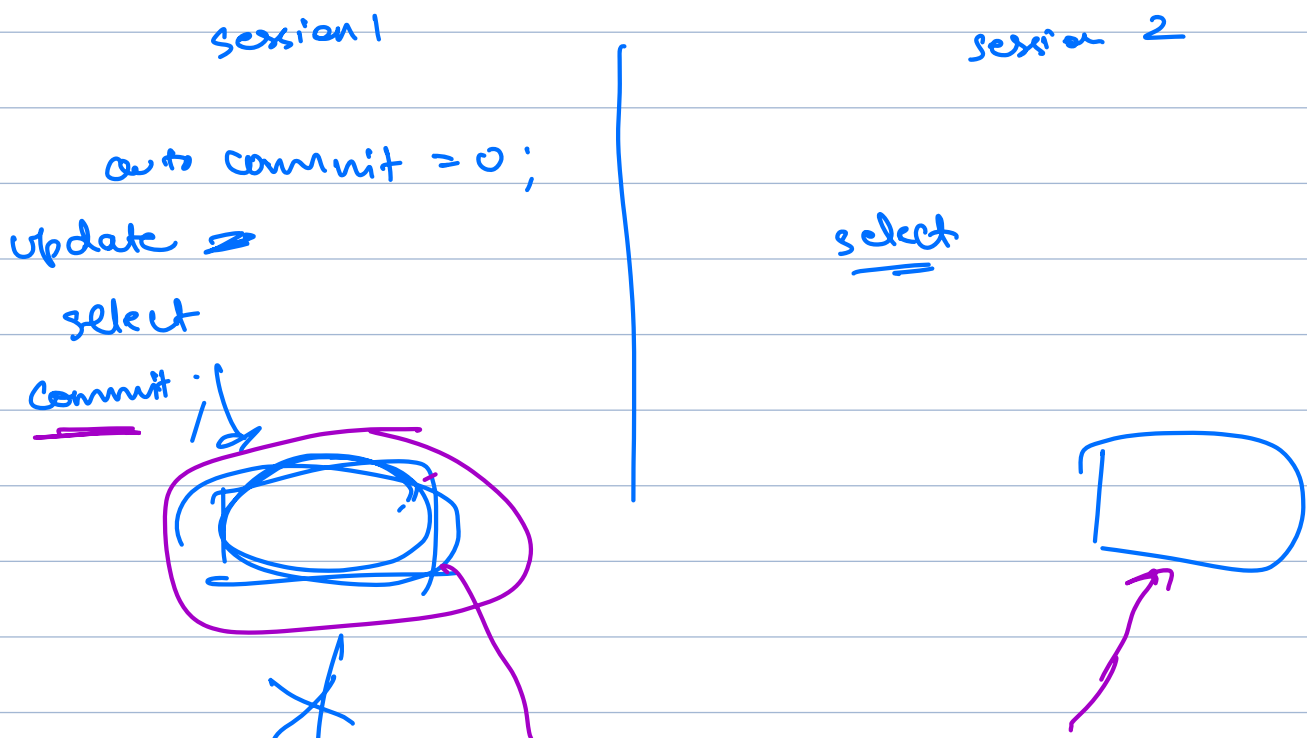
- every SQL query starts a transaction
- execution itself.
- saves the changes // commit

- ① Start transaction
- ② Execute operation
- ✓ ③ Commits

↳ statement that is used to persist the result of an SQL query.

★ Commit takes care of durability

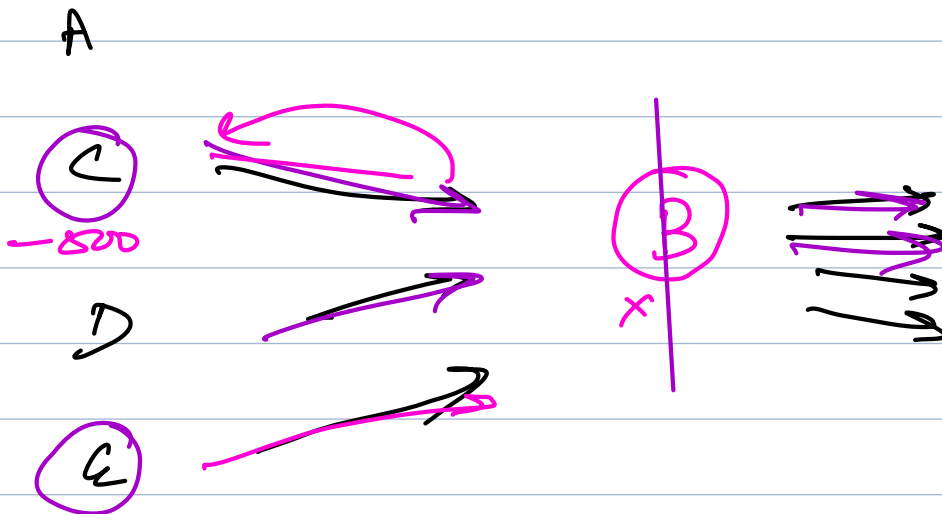
Auto commit = true By default.





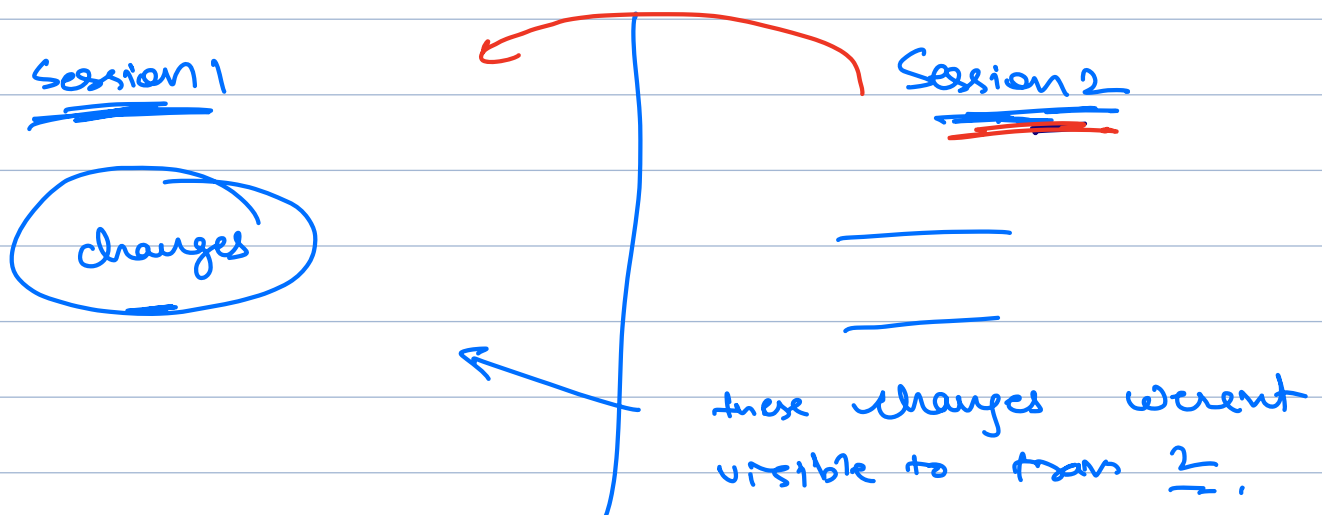


Roll back.



undo all changes done since last commit;

→ You cannot rollback after a commit is done.



ISOLATED

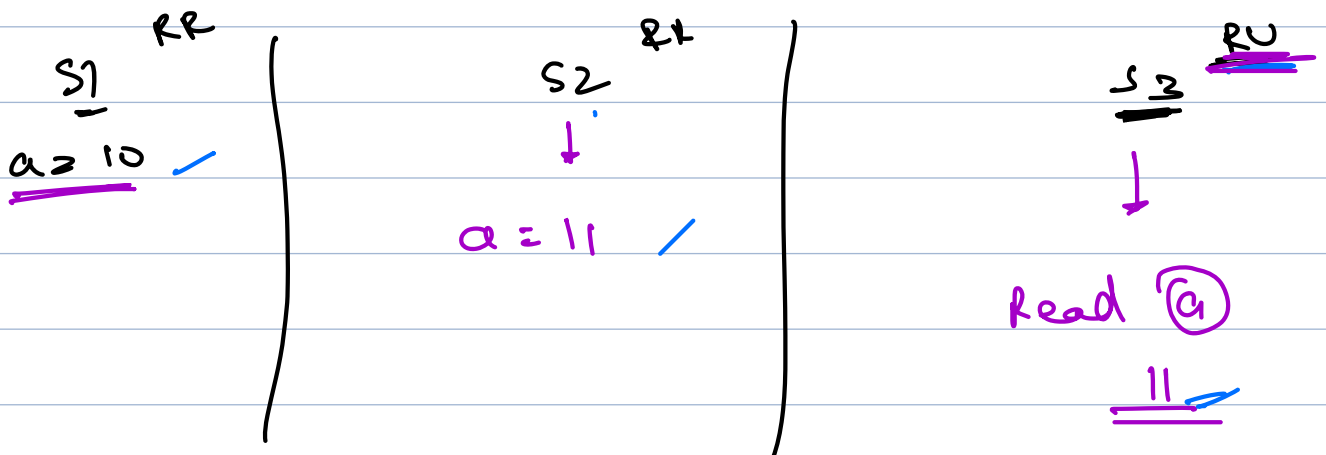
## 4 levels of isolation.

- 1) Read uncommitted
  - 2) Read committed
  - 3) Repeatable read
  - 4) serializable
- relaxed
- ↓  
inc order of severity.
- Default for MySQL → strict

\* Isolation levels are talking about what data is read, not what is updated.

### Read uncommitted

- Allows a transaction to read even uncommitted data from another transaction
- latest uncommitted data also



\* Isolation level only talks about how your session will act (read data)

S<sub>1</sub>  
a = 10 ✓

S<sub>1</sub>  
a = 11 ✓  
commit

S<sub>3</sub> RU  
a  
11

PROS.

Fast → performance

CONS.

A → ~~2000~~ 2090  
B → ~~2000~~ 1900

Trans 1 (RR)  
A → B + 10

Trans 2 (RU)  
B → A + 100

① Read A → x (2000)

②  $x \leftarrow x - 10$  (1990)

③ Write A ← x (1990)

⑦ Read B → x (RU → 2000, ~~RU → 1900~~)

$x \leftarrow x + 10$

Write B ← x

commit ;

④ Read B → x (2000)

⑤  $x \leftarrow x - 100$  (1900)

⑥ Write B ← x (1900)

⑧ Read A → x (1990)

⑨  $x \leftarrow x + 100$  (2090)

⑩ Write A ← x (2090)

⑪ commit ;

At line 8 → trans 2 read data that was not committed and rolled back.

## DIRTY READ

when a trans reads data that is not committed.  
That dat might change / rolled back / not yet committed.