

Q.1 Given a string, Find longest substring with distinct characters.

↓  
non-repeated

str: a b c a b c d d      ans: 4      (abcd)

str: s i p p e r      ans: 3      (sip | per)

str: a b c g h e g k l m h a b k      ans: 8

i) Idea 1

travel on all substrings with the help of s and e. check if substring from s to e is containing distinct chars or not with the help of Hashset.

T.C:  $O(n^3)$

0	1	2	3	4	5
a	b	c	e	b	f
s				e	

for (s → 0 to n-1) {

    for (e → s to n-1) {

        Hashset < c, g > hs = ....;

        for (s to e) {

            // to check whether content from s to e is  
            // is distinct or not

        }

        ans = Math.max(ans, hs.size());      based on a check

    }

}

b	c
e	

ii) Idea 2 : Expected TC  $\rightarrow O(n^2)$

S  
a b c g h e g k l m h a b k  
0 1 2 3 4 5 6 7 8 9 10 11 12 13  
e

S = 0

e  $\rightarrow$  1 to 13

b c g  
h e

start from every start point possible and extend till the moment you can (getting non-repeated char).

```
int solve(String str) {
    int n = str.length();
    int ans = 0;

    for (int s = 0; s < n; s++) {
        HashSet<Character> hs = new HashSet<>();
        for (int e = s; e < n; e++) {
            if (hs.contains(str.charAt(e)) == true) {
                break;
            }
            else {
                hs.add(str.charAt(e));
            }
        }
        ans = Math.max(ans, hs.size());
    }
    return ans;
}
```

<sup>S</sup>  
 a b c g h e g k d m h a b k  
 0 1 2 3 4 5 6 7 8 9 10 11 12 13  
e

```

for (int s=0; s<n; s++) {
    HashSet<character> hs= new HashSet<>();
    for (int e=s; e<n; e++) {
        if (hs.contains(str.charAt(e)) == true) {
            break;
        }
        else {
            hs.add(str.charAt(e));
        }
    }
    ans = Math.max(ans, hs.size());
}

```

e g k  
 d m h  
 a b

hs

ans = ~~6~~ 8

Idea 3: TC  $\rightarrow O(n)$

<sup>S</sup>  
 a b c g h e g k d m h a b k  
 0 1 2 3 4 5 6 7 8 9 10 11 12 13  
e

```

int solve (String str) {
    int n = str.length();
    HashSet < Character > hs = new HashSet<>();

    int s = 0, e = 0;
    int ans = 0;

    while (e < n) {
        if (hs.contains (str.charAt(e)) == false) {
            hs.add (str.charAt(e));
            e++;
        }
        else {
            // get rid of repeated char
            hs.remove (str.charAt(s));
            s++;
        }
        ans = Math.max(ans, hs.size());
    }

    return ans;
}

```

s

a	b	c	g	h	e	g	k	u	m	h	a	b	k
0	1	2	3	4	5	6	7	8	9	10	11	12	13

e

while (e < n) {

if (hs.contains(str.charAt(e)) == false) {

hs.add(str.charAt(e));

e++;

} acquire

}

else {

// get rid of repeated char

hs.remove(str.charAt(s));

s++;

} release

}

ans = Math.max(ans, hs.size());

}

<del>x</del>	<del>x</del>	<del>x</del>	<del>x</del>
<del>x</del>	<del>x</del>	<del>x</del>	<del>x</del>
u	m	h	a
b	k		

hs

ans = ~~0~~ ~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~ ~~7~~

8

max itr: 2n

Tc: O(n)

Sc: O(1)

Sc: O(26) ⇒ O(1)

HashSet < character > hs

char → lowercase  
chars

HashMap < character, Integer > map

→ Sc: O(26) ⇒ O(1)

Q.2 Given two strings, check if they are **anagrams** or not.

↓  
permutation

eg1  $\begin{cases} A = \text{doodle} \\ B = \text{ooddle} \end{cases}$

ans: true

(diff. rearrangement  
of same content)

eg2  $\begin{cases} A = \text{givei} \\ B = \text{vijee} \end{cases}$

ans: false

Don't use hashset (freq matters)

i) create freq map of str1  $\rightarrow$  map1

ii) create freq map of str2  $\rightarrow$  map2

iii) areMapSame(map1, map2)

str1 = a a b c a c

map1  $\Rightarrow$   $\begin{aligned} a &\rightarrow 3 \\ b &\rightarrow 1 \\ c &\rightarrow 2 \end{aligned}$

str2 = b a c c a a

map2  $\Rightarrow$   $\begin{aligned} b &\rightarrow 1 \\ a &\rightarrow 3 \\ c &\rightarrow 2 \end{aligned}$

check same: travel a to z and check freq  
of ch in both maps.

Q.3 Count no. of substrings of **B** which are permutations of A.

↓  
anagrams

eg1  $\left[ \begin{array}{l} A = a b c \\ B = \underline{a b c} \underline{b a c} \underline{a b c} \end{array} \right.$

ans = 3

eg2  $\left[ \begin{array}{l} A = a a b \\ B = \underline{a b a} \underline{a b e} \underline{b a b a} \underline{a a d} \end{array} \right.$

ans = 5

Sliding window

Window length:  $A.length()$

A = a a b  
0 1 2

B = a b a a b e b a b a a d  
0 1 2 3 4 5 6 7 8 9 10 11

K = 3

a → 2  
b → 1

Amap

{ never going to  
change }

	s	e	rem	add
	0	2		
	1	3	0	3
	2	4	1	4
	3	5	2	5
	4	6	3	6
	5	7	4	7
	6	8	5	8
→	7	9	6	9
	⋮			

a → 2  
b → 1

wmap

{ current  
window }

if (areMapSame(wmap, Amap)) {

ans++;

}

ans = 1 2 3 4

TC:  $O(n)$

doubt5

A = [ abcdef, abcg, abckt, abtmno ]

max ans  $\rightarrow$  length of smallest

$\rightarrow$  slen

ans = ab

idx = 0, ch = a

idx = 1, ch = b

idx = 2, ch = c

ans = "";

for (int idx = 0; idx < slen; idx++) {

    // check if the char at idx is same in all of

    // string or not

        char ch = A[0].charAt(idx);

        for (k = 1; k < A.length; k++) {

            if (A[k].charAt(idx) != ch) {

                return ans;

        }

    }

        ans += ch;

}

return ans;



a b c a d a c  
0 1 2 3 4 5 6

for (s → 0 to n-1) {

for (e → s to n-1) {

HashSet < c, g > hs = ....;

for (s to e) {

|| to check whether content from s to e is

|| is distinct or not

}

ans = Math.max(ans, hs.size());

}

}

↳ check based

s	e	
0	0	a
0	1	a b
0	2	a b c