

GrAlgo

Group 5

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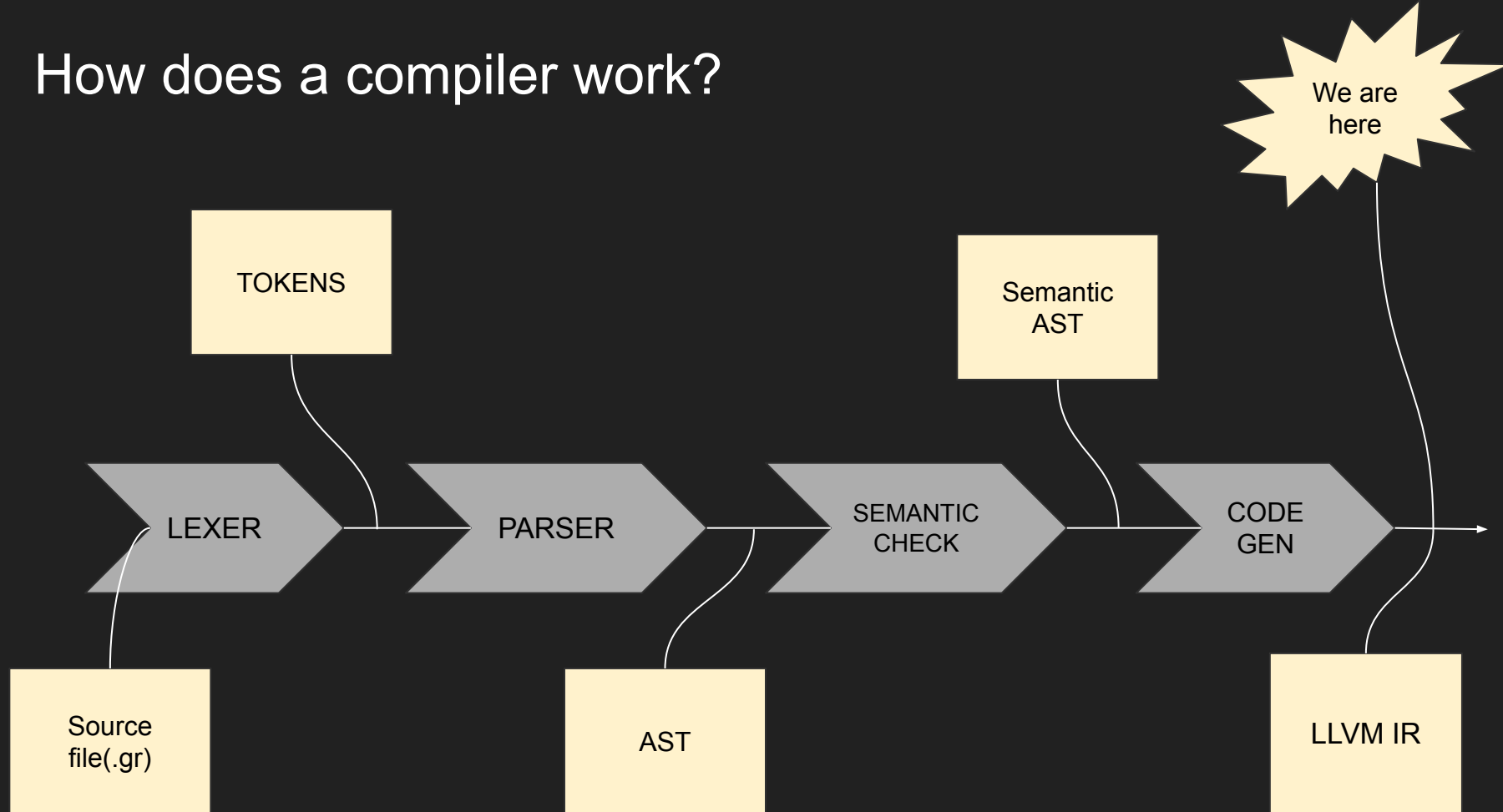
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How does a compiler work?



Code Generation

- Code generation is the final phase of the compiler in which we turn Intermediate representation of the source code into a form of machine code that can be executed by the target machine.
- There are multiple types of **Intermediate Representations** such as Three-Address Code, Postfix Notation, Syntax tree and LLVM IR.
- This phase takes in **AST** or **Symbol Table** as the input and gives out an **IR** which is then converted to target code.
- Generally we optimise the code while generating IR on whatever the user may desire such as power, time or memory.

LLVM IR

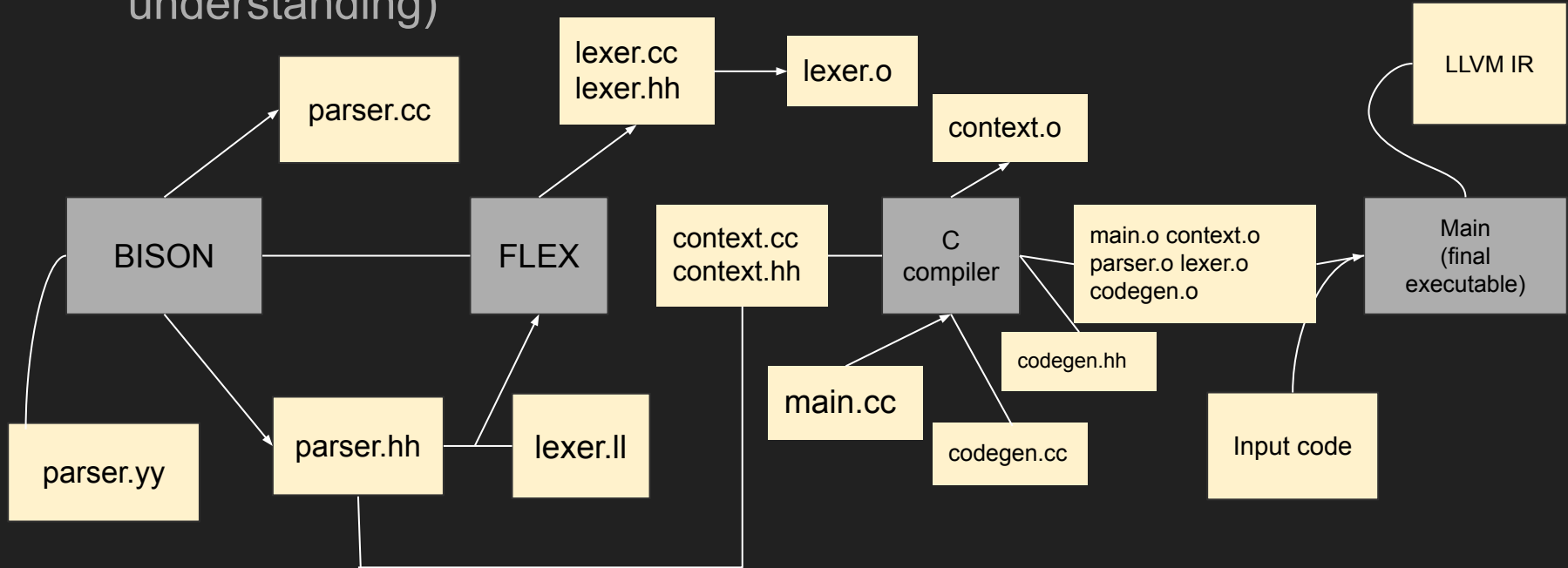
- LLVM is a Static Single Assignment based Language that provides type safety, low-level operations, flexibility, and the capability of representing 'all' high-level languages cleanly.
- If we are able to generate LLVM IR then we can use the LLVM compiler to convert this IR into target machine code easily using a single command.
- That command is of the following format: `llc [options] [filename]` for example if we have `new.ll` file and we want assembly code in say `sparc` we can run the following command: `llc -march=sparc new.ll -o new.s` and we get the assembly code for `sparc` in `new.s`

Code Generator

- Using AST and Symbol table generated in the Parsing stage we generate LLVM IR.
- As in the parsing phase of the compiler we have stored our AST and our symbol table in `main.cc` file we call `codegen` function which performs the actual `codegen` process.
- We stored AST and the symbol table in `main` and from that call the `codegen` function which generates an `.ll` file which is an extension for LLVM IR files and using LLVM we can convert this `.ll` file into machine code of our choice.

Code Generator

- Parser with semantic analysis uses **BISON** and **FLEX** for compiling. Context.cc file stores the Symbol Table. (See chart below for better understanding)



Using our Code Generator

- The folder “Semantics” contains the following files.
 - parser.yy lexer.ll GrFlexLexer.hh
 - context.cc context.hh types.hh
 - codegen.cc codegen.hh
 - main.cc
 - Makefile
 - input<number>.gr in inputs folder (Test Cases)
- “make” command is used to run these files and generate the output executable file which takes .gr file as input and generates the AST and also does the semantic analysis.
- To get the output run the command `./main < input.gr`

Example



```
1  int foo(int f, int c)
2  {
3      return 2+3;
4  }
5
6  int main()
7  {
8      int x = 2+3-3;
9
10     return x;
11 }
```

Code Snippet

LLVM IR

```
1 |; ModuleID = 'gralgo'
2 source_filename = "gralgo"
3
4 define i32 @foo(i32 %f, i32 %c) {
5 entry:
6     %c2 = alloca i32, align 4
7     %f1 = alloca i32, align 4
8     store i32 %f, i32* %f1, align 4
9     store i32 %c, i32* %c2, align 4
10    %addtmp = add i32 2, 3
11    ret i32 %addtmp
12 }
13
14 define i32 @main() {
15 entry:
16    %0 = alloca i32, align 4
17    %addtmp = add i32 2, 3
18    %addtmp1 = add i32 %addtmp, -3
19    store i32 %addtmp1, i32* %0, align 4
20    %x = load i32, i32* %0, align 4
21    ret i32 %x
22 }
```




Thank You!!