



# Solving problems by searching

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# Steps in Problem Solving

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- Goal Formulation: based on current situation and the agent's performance measure is the first step in Problem Solving
- Problem Formulation-Deciding what actions and states to consider given a goal.
- Search for different possible actions and choose the best.
- A search algorithm takes problem as input and returns a solution in the form of an action sequence.



# Steps in Problem Solving

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- The action recommended by the solution can be carried out, known as Execution.



# Problem-solving agents

```
function SIMPLE-PROBLEM-SOLVING-AGENT(percept) returns an action
  static: seq, an action sequence, initially empty
           state, some description of the current world state
           goal, a goal, initially null
           problem, a problem formulation

  state  $\leftarrow$  UPDATE-STATE(state, percept)
  if seq is empty then do
    goal  $\leftarrow$  FORMULATE-GOAL(state)
    problem  $\leftarrow$  FORMULATE-PROBLEM(state, goal)
    seq  $\leftarrow$  SEARCH(problem)
  action  $\leftarrow$  FIRST(seq)
  seq  $\leftarrow$  REST(seq)
  return action
```

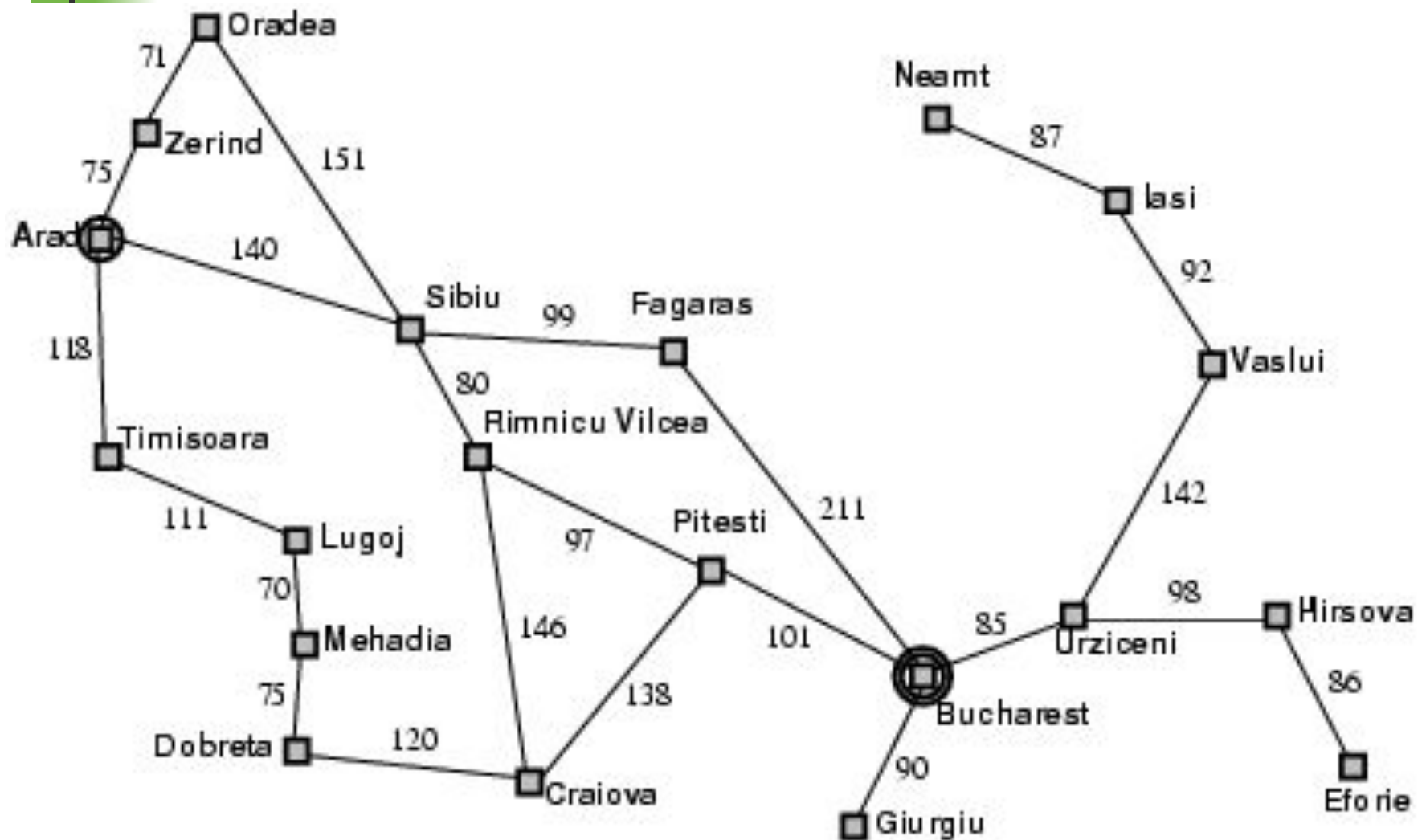


# Example: Romania

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- On holiday in Romania; currently in Arad.
- Flight leaves tomorrow from Bucharest
- Formulate goal:
  - be in Bucharest
- Formulate problem:
  - **states**: various cities
  - **actions**: drive between cities
- Find solution:
  - sequence of cities, e.g., Arad, Sibiu, Fagaras, Bucharest

# Example: Romania





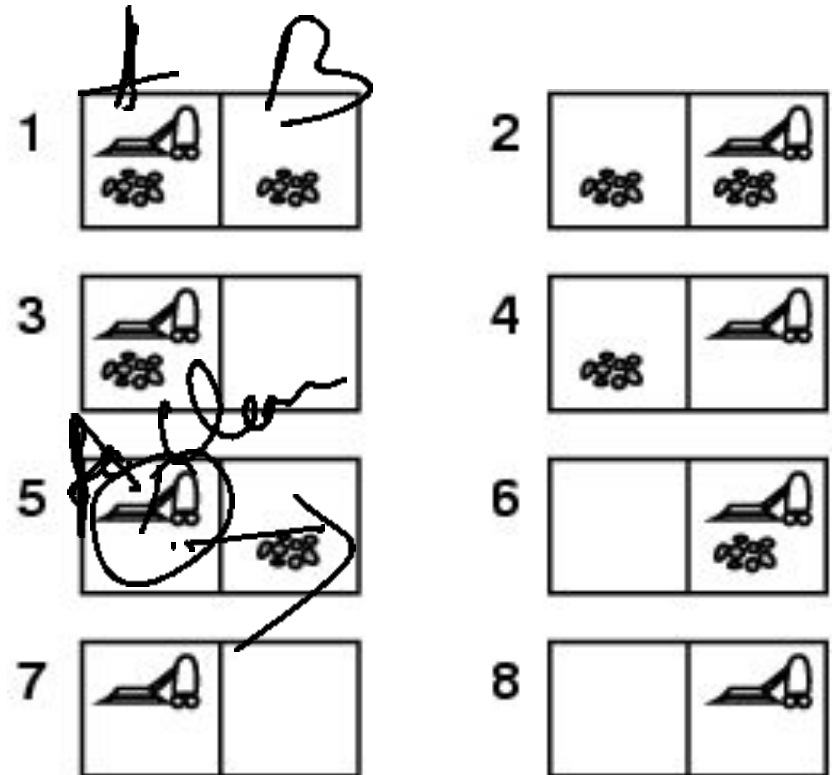
# Problem types

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- Deterministic, fully observable □ single-state problem
  - Agent knows exactly which state it will be in; solution is a sequence
- Non-observable □ sensorless problem (conformant problem)
  - Agent may have no idea where it is; solution is a sequence
- Nondeterministic and/or partially observable □ contingency problem
  - percepts provide new information about current state
  - often interleave} search, execution
- Unknown state space □ exploration problem

# Example: vacuum world

- Single-state, start in #5.  
Solution?

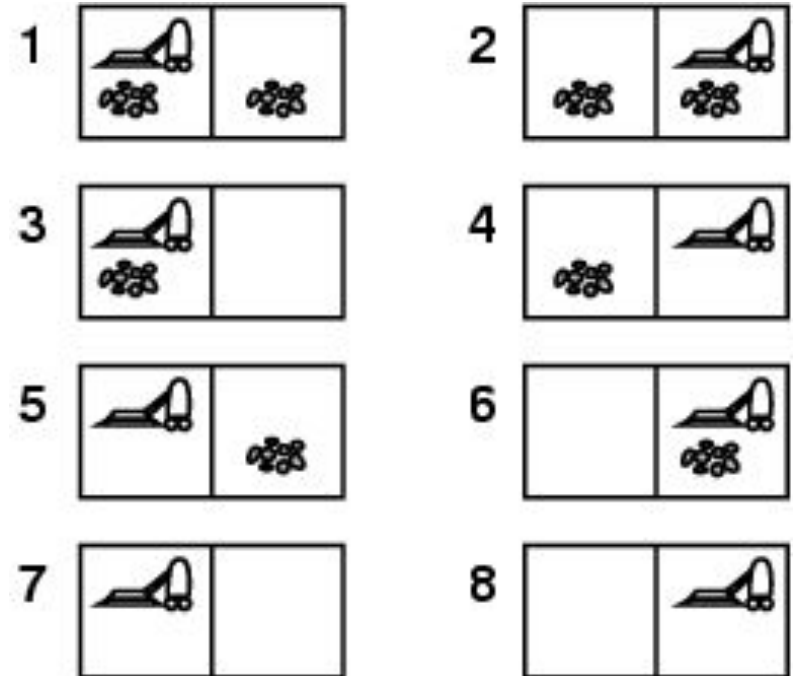




# Example: vacuum world

- Single-state, start in #5.  
Solution? [*Right, Suck*]

- Sensorless, start in  $\{1,2,3,4,5,6,7,8\}$  e.g., *Right* goes to  $\{2,4,6,8\}$   
Solution?



# Example: vacuum world

- Sensorless, start in  $\{1,2,3,4,5,6,7,8\}$  e.g.,  
*Right* goes to  $\{2,4,6,8\}$

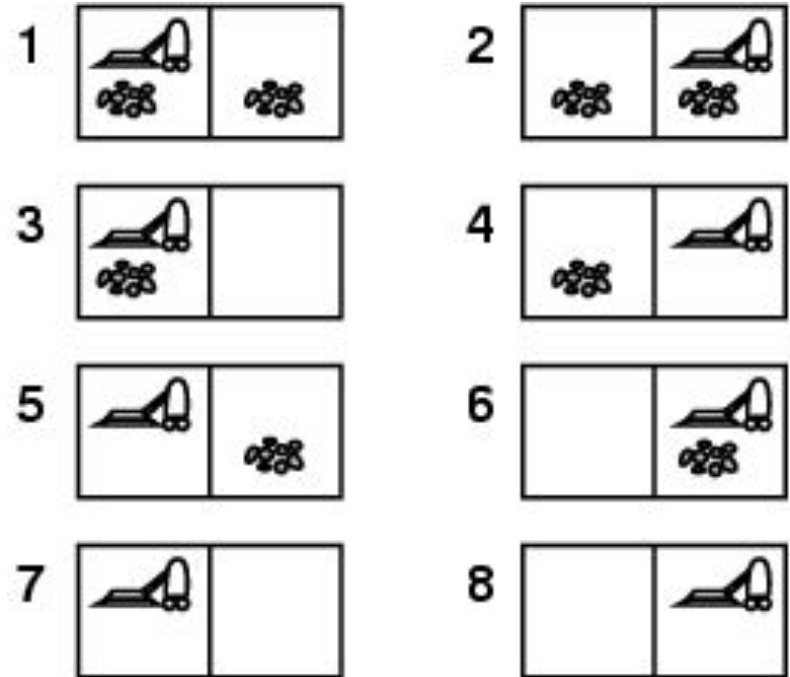
Solution?

*[Right, Suck, Left, Suck]*

- Contingency

- Nondeterministic: *Suck* may dirty a clean carpet
- Partially observable: location, dirt at current location.
- Percept:  $[L, \text{Clean}]$ , i.e., start in #5 or #7

Solution?

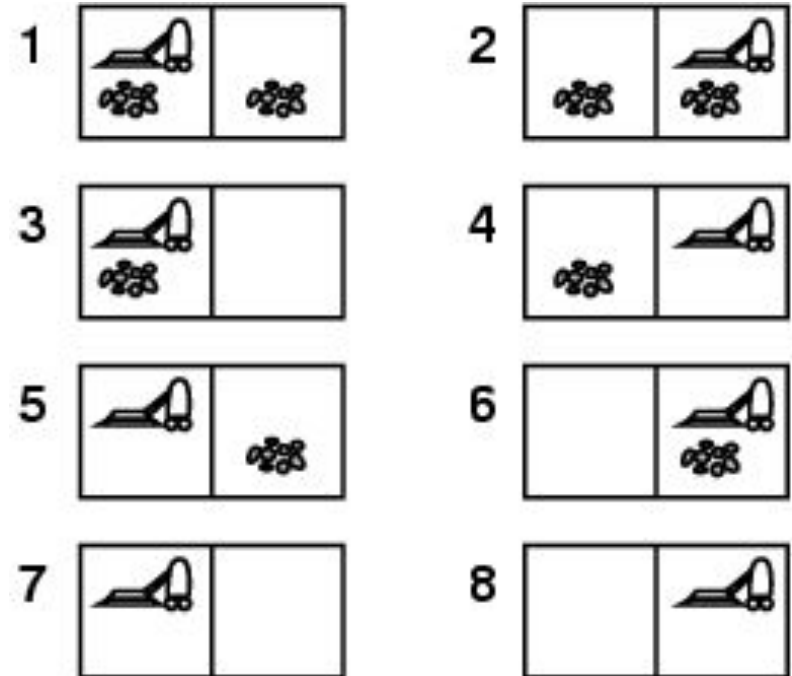


# Example: vacuum world

- Sensorless, start in  $\{1,2,3,4,5,6,7,8\}$  e.g.,  
*Right* goes to  $\{2,4,6,8\}$

Solution?

*[Right, Suck, Left, Suck]*



- Contingency

- Nondeterministic: *Suck* may dirty a clean carpet
- Partially observable: location, dirt at current location.
- Percept:  $[L, \text{Clean}]$ , i.e., start in #5 or #7

Solution? *[Right, **if** dirt **then** Suck]*



# Assumptions for the Environment

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- Observable : Knows the current state .
- Discrete: Only finitely many actions
- Known : Result of action is known.
- Deterministic : Each action has only one action.



# Single-state problem formulation

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A **problem** is defined by four items:

1. **initial state** e.g., "at Arad"
  2. **actions** or **successor function**  $S(x)$  = set of action–state pairs
    - e.g.,  $S(\text{Arad}) = \{ \langle \text{Arad} \rightarrow \text{Zerind}, \text{Zerind} \rangle, \dots \}$
  3. **goal test**, can be
    - **explicit**, e.g.,  $x = \text{"at Bucharest"}$
    - **implicit**, e.g.,  $\text{Checkmate}(x)$
  4. **path cost** (additive)
    - e.g., sum of distances, number of actions executed, etc.
    - $c(x, a, y)$  is the **step cost**, assumed to be  $\geq 0$
- 
- A **solution** is a sequence of actions leading from the initial state to a goal state

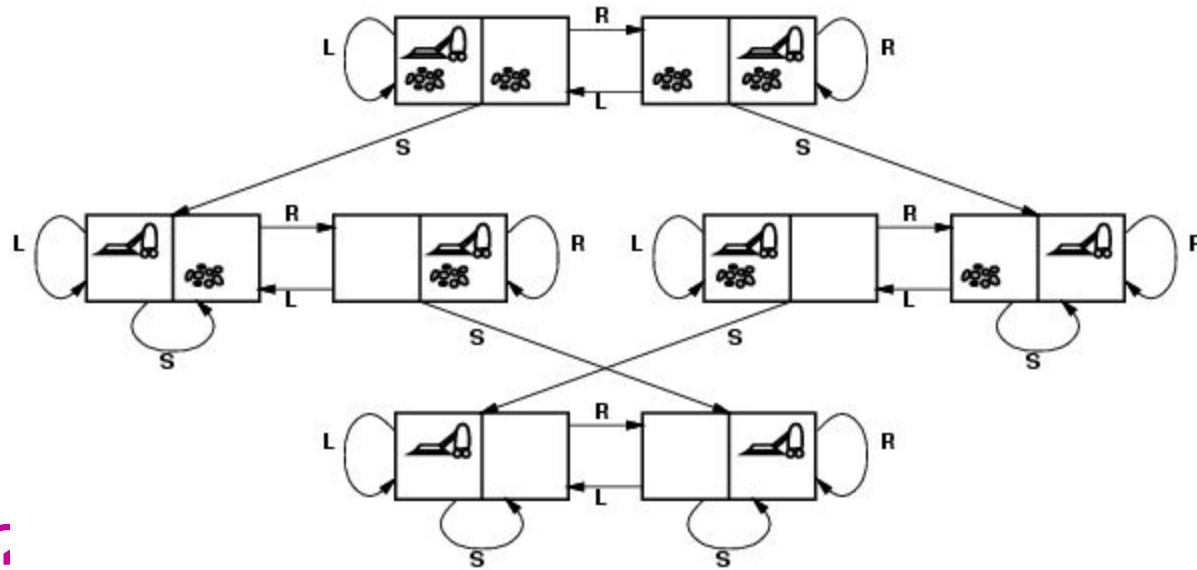


# Selecting a state space

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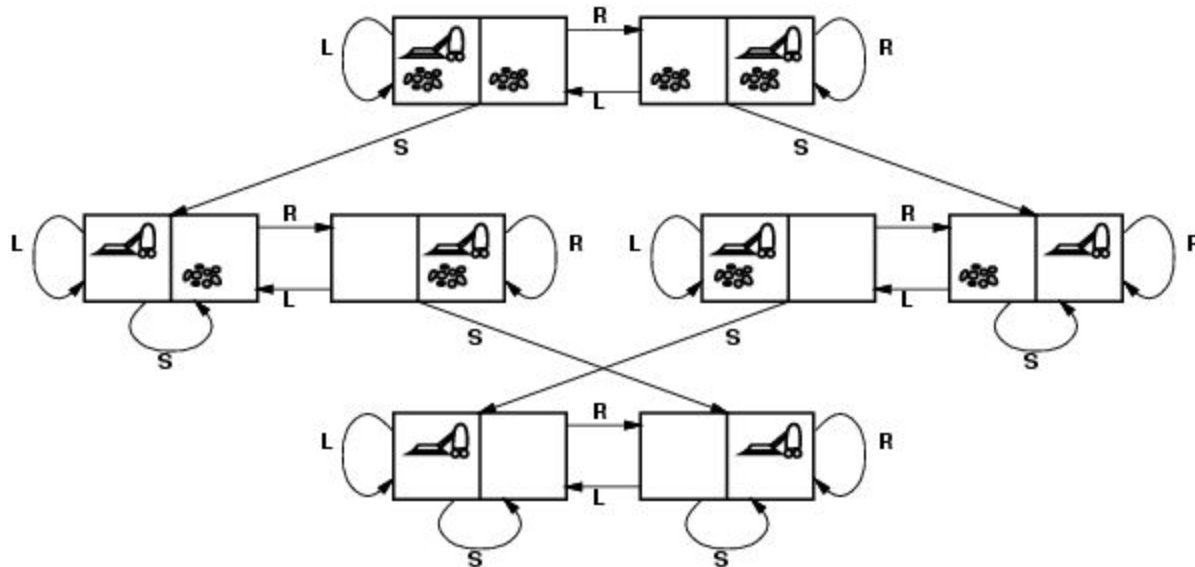
- Real world is absurdly complex
  - state space must be **abstracted** for problem solving
- (Abstract) state = set of real states
- (Abstract) action = complex combination of real actions
  - e.g., "Arad □ Zerind" represents a complex set of possible routes, detours, rest stops, etc.
- For guaranteed realizability, **any** real state "in Arad" must get to **some** real state "in Zerind"
- (Abstract) solution =
  - set of real paths that are solutions in the real world

# Vacuum world state space graph



- states?
- actions?
- goal test?
- path cost?

# Vacuum world state space graph



- states? integer dirt and robot location
- actions? *Left, Right, Suck*
- goal test? no dirt at all locations
- path cost? 1 per action



# Example: The 8-puzzle

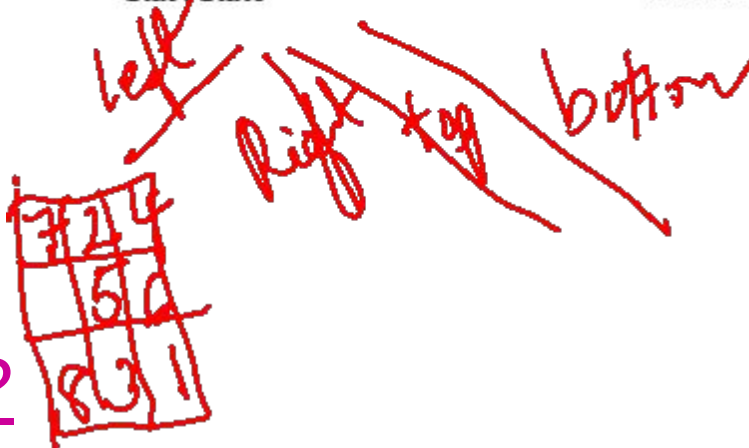
7	2	4
5		6
8	3	1

Start State

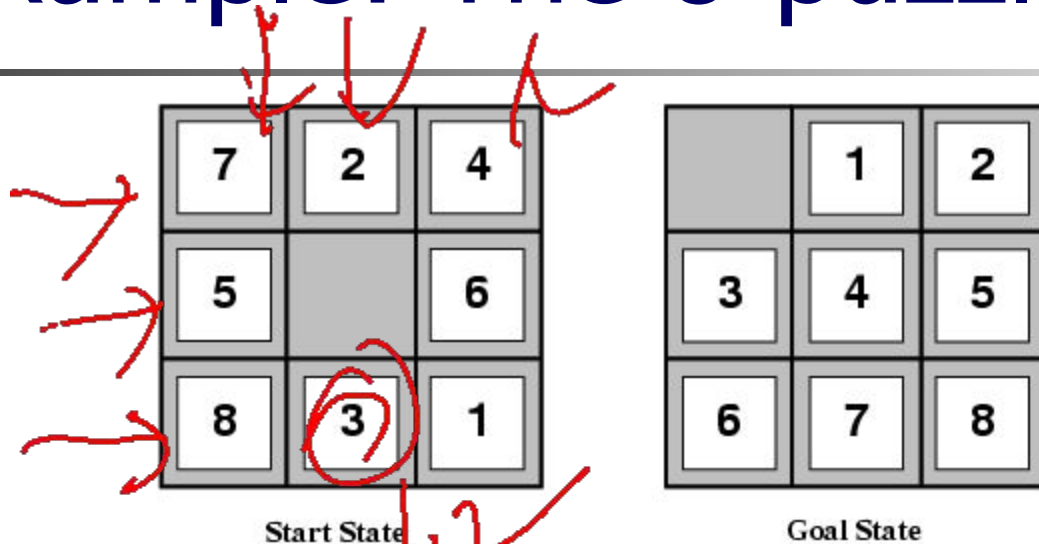
	1	2
3	4	5
6	7	8

Goal State

- states?
- actions?
- goal test?
- path cost?



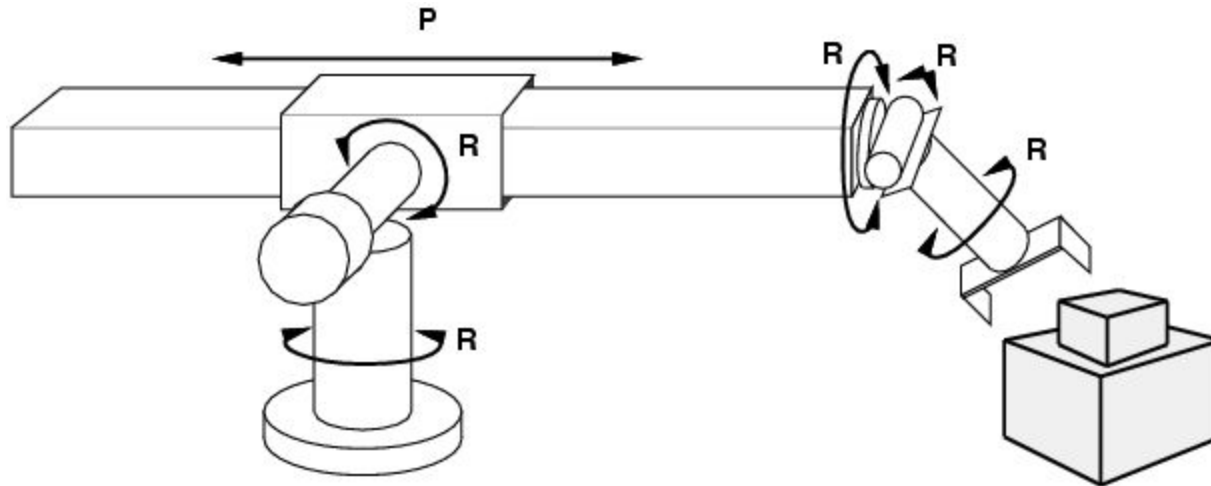
# Example: The 8-puzzle



- states? locations of tiles
- actions? move blank left, right, up, down
- goal test? = goal state (given)
- path cost? 1 per move

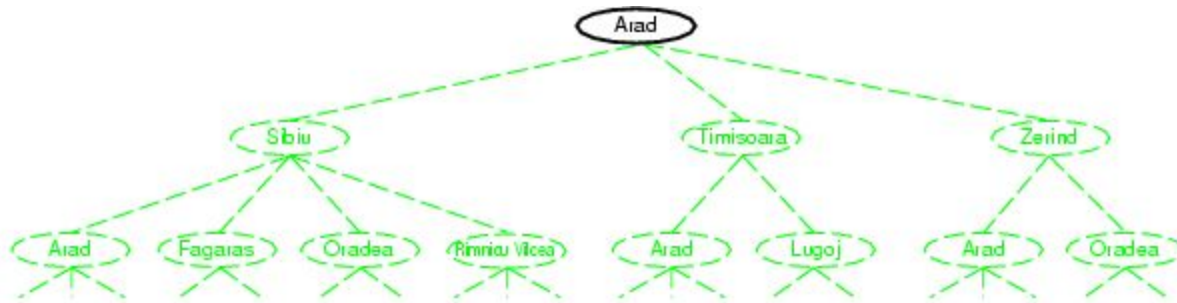
[Note: optimal solution of  $n$ -Puzzle family is NP-hard]

# Example: robotic assembly

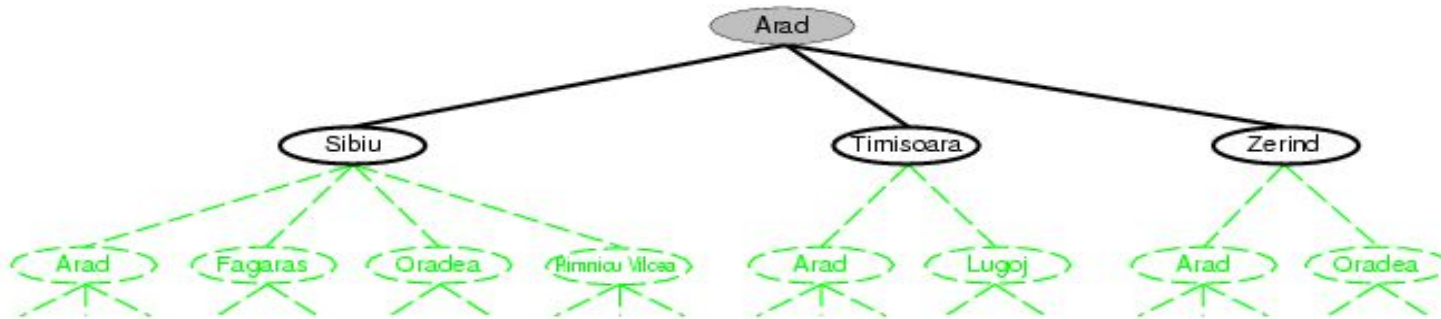


- states?: real-valued coordinates of robot joint angles parts of the object to be assembled
- actions?: continuous motions of robot joints
- goal test?: complete assembly
- path cost?: time to execute

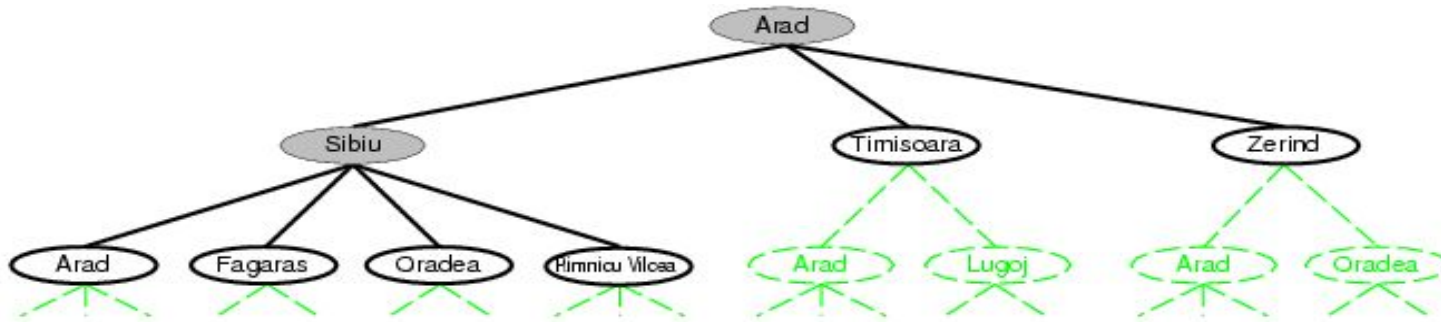
# Tree search example

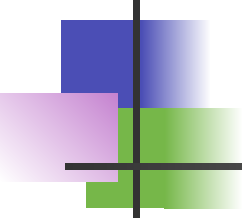


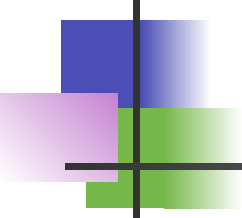
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# Tree search example



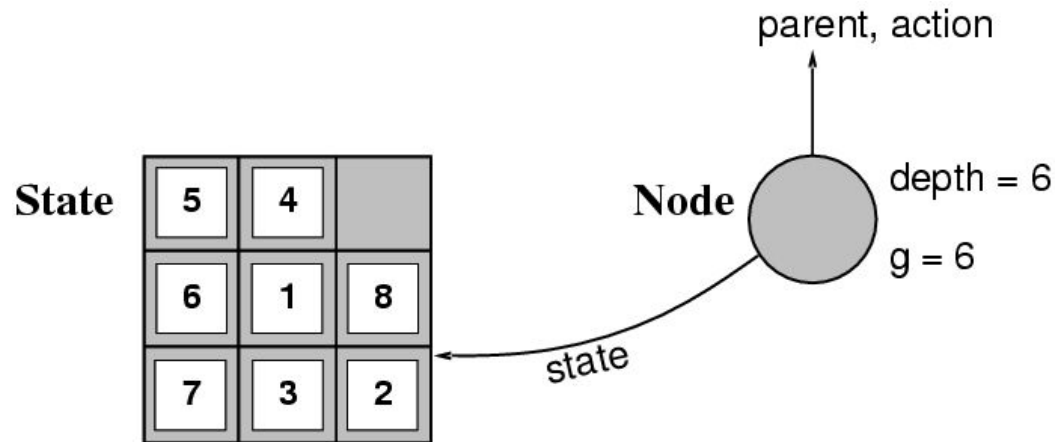
- 
- 
- function TREE-SEARCH(problem) returns a solution, or failure
  - initialize the frontier using the initial state of problem
  - loop do
  - if the frontier is empty then return failure
  - choose a leaf node and remove it from the frontier
  - if the node contains a goal state then return the corresponding solution
  - expand the chosen node, adding the resulting nodes to the frontier
- 
- \* frontier –The set of all leaf nodes available for expansion at any given point

- 
- 
- function GRAPH-SEARCH(problem) returns a solution, or failure
  - initialize the frontier using the initial state of problem
  - initialize the explored set to be empty
  - loop do
  - if the frontier is empty then return failure
  - choose a leaf node and remove it from the frontier
  - if the node contains a goal state then return the corresponding solution
  - add the node to the explored set
  - expand the chosen node, adding the resulting nodes to the frontier
  - only if not in the frontier or explored set

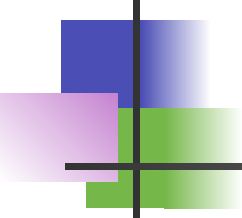


# Implementation: states vs. nodes

- A **state** is a (representation of) a physical configuration
- A **node** is a data structure constituting part of a search tree includes **state**, **parent node**, **action**, **path cost**  $g(x)$ , **depth**



- The `Expand` function creates new nodes, filling in the various fields and using the `SuccessorFn` of the problem to create the corresponding states.

- 
- function CHILD-NODE(problem, parent, action) returns a node
  - return a node with
  - $STATE = \text{problem.RESULT}(\text{parent.STATE}, \text{action}),$
  - $PARENT = \text{parent}, ACTION = \text{action},$
  - $PATH-COST = \text{parent.PATH-COST} + \text{problem.STEP-COST}(\text{parent.STATE}, \text{action})$



# Search strategies

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- A search strategy is defined by picking the **order of node expansion**
- Strategies are evaluated along the following dimensions:
  - **completeness**: does it always find a solution if one exists?
  - **time complexity**: number of nodes generated
  - **space complexity**: maximum number of nodes in memory
  - **optimality**: does it always find a least-cost solution?
- Time and space complexity are measured in terms of
  - $b$ : maximum branching factor of the search tree// **no.of children at each node**
  - $d$ : depth of the least-cost solution
  - $m$ : maximum depth of the state space (may be  $\infty$ )



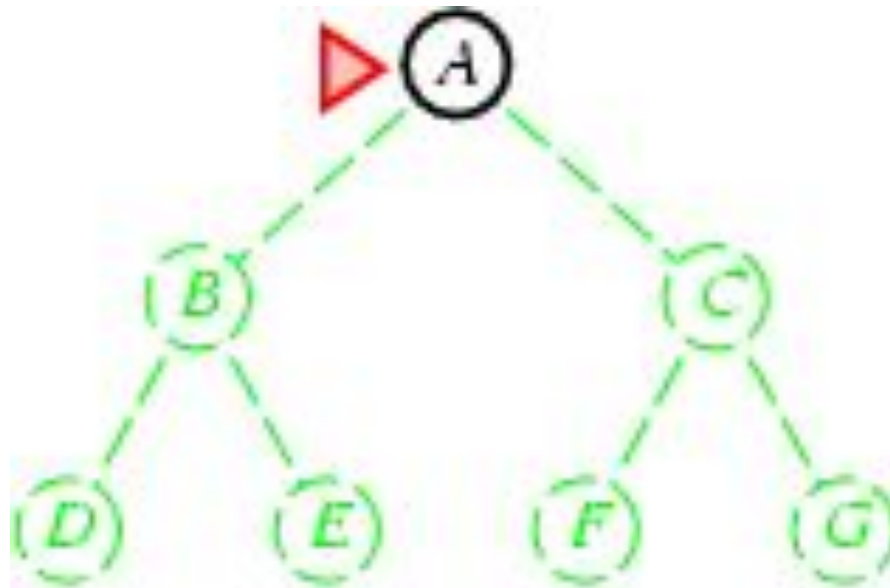
# Uninformed search strategies

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- **Uninformed** search strategies use only the information available in the problem definition
- Breadth-first search
- Uniform-cost search
- Depth-first search
- Depth-limited search
- Iterative deepening search

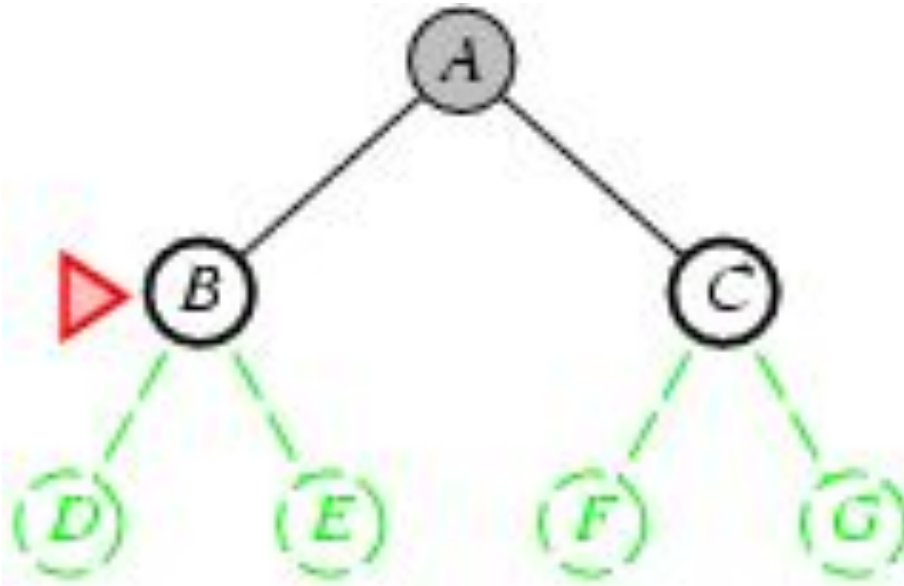
# Breadth-first search

- Expand shallowest unexpanded node
- Implementation:
  - *fringe* is a FIFO queue, i.e., new successors go at end



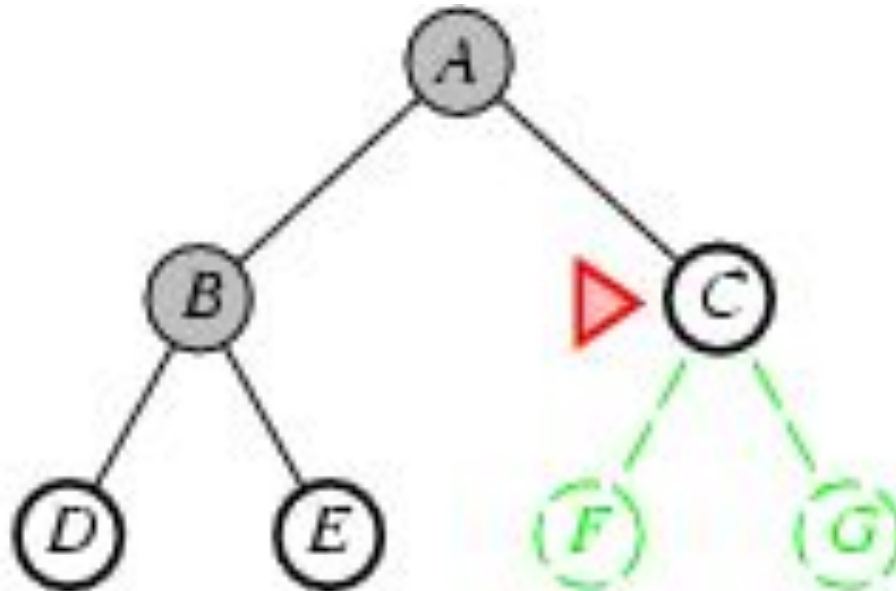
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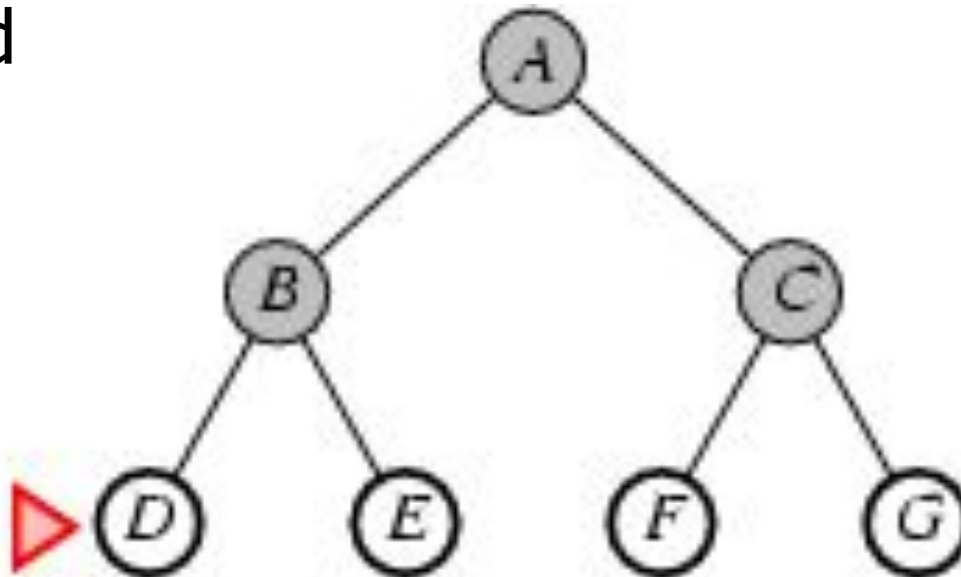
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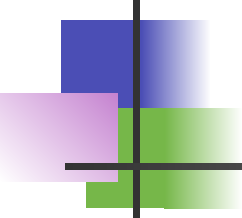




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function BREADTH-FIRST-SEARCH(problem) returns a solution, or failure

- node  $\leftarrow$  a node with STATE = problem.INITIAL-STATE, PATH-COST = 0
- if problem.GOAL-TEST(node.STATE) then return SOLUTION(node)
- frontier  $\leftarrow$  a FIFO queue with node as the only element
- explored  $\leftarrow$  an empty set
- loop do if EMPTY?(frontier ) then return failure
- node  $\leftarrow$  POP(frontier ) /\* chooses the shallowest node in frontier \*/

- 
- 
- add node.STATE to explored
  - for each action in problem.ACTIONS(node.STATE) do
  - child  $\leftarrow$  CHILD-NODE(problem, node, action)
  - if child.STATE is not in explored or frontier then
  - if problem.GOAL-TEST(child.STATE) then return SOLUTION(child)
  - frontier  $\leftarrow$  INSERT(child, frontier )



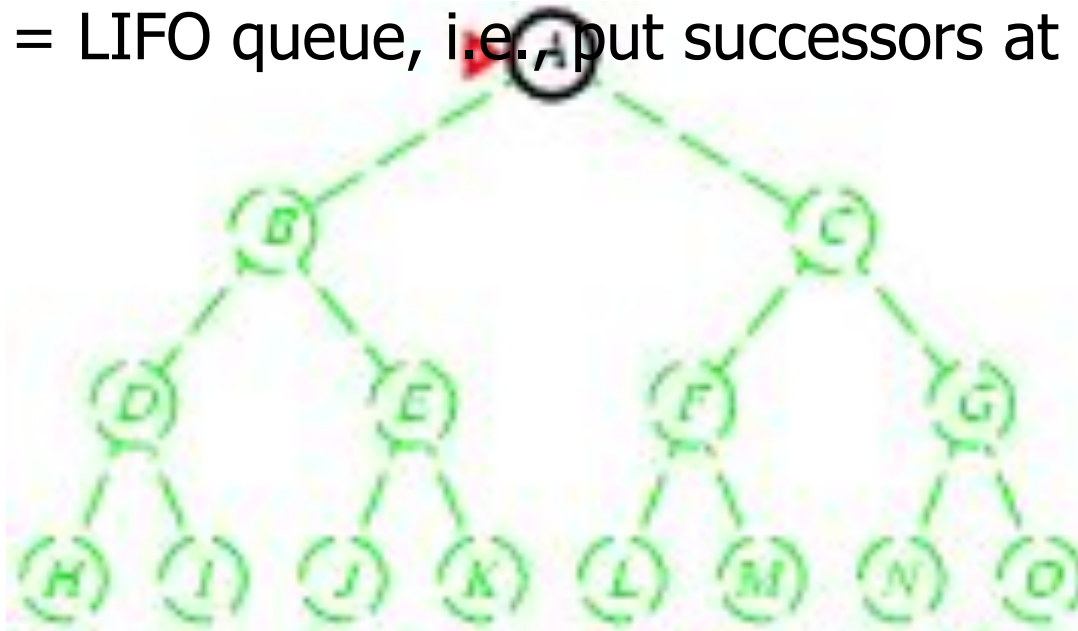
# Properties of breadth-first search

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- Complete? Yes (if  $b$  is finite)
- Time?  $1+b+b^2+b^3+\dots +b^d + b(b^d-1) = O(b^{d+1})$
- Space?  $O(b^{d+1})$  (keeps every node in memory)
- Optimal? Yes (if cost = 1 per step)
- **Space** is the bigger problem (more than time)

# Depth-first search

- Expand deepest unexpanded node
- Implementation:
  - *fringe* = LIFO queue, i.e., put successors at front

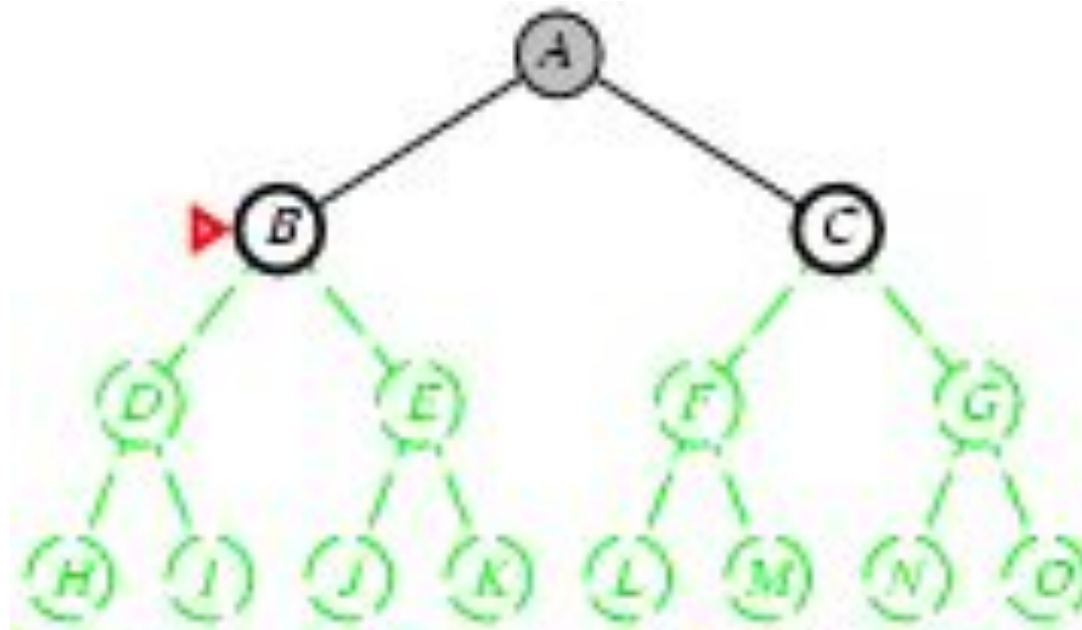


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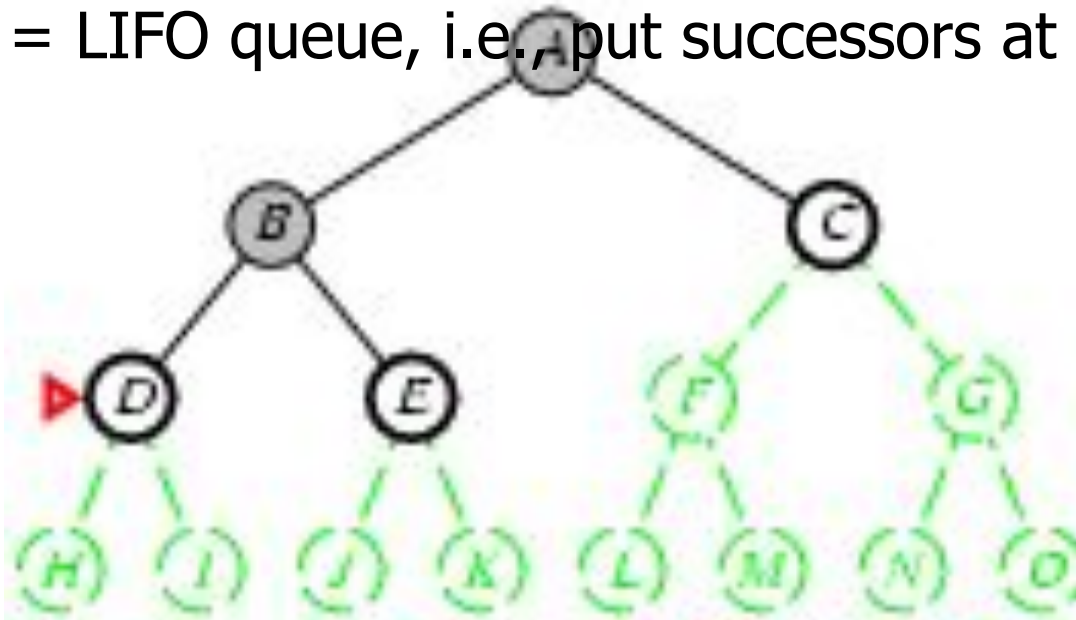
- Implementation:

- *fringe*



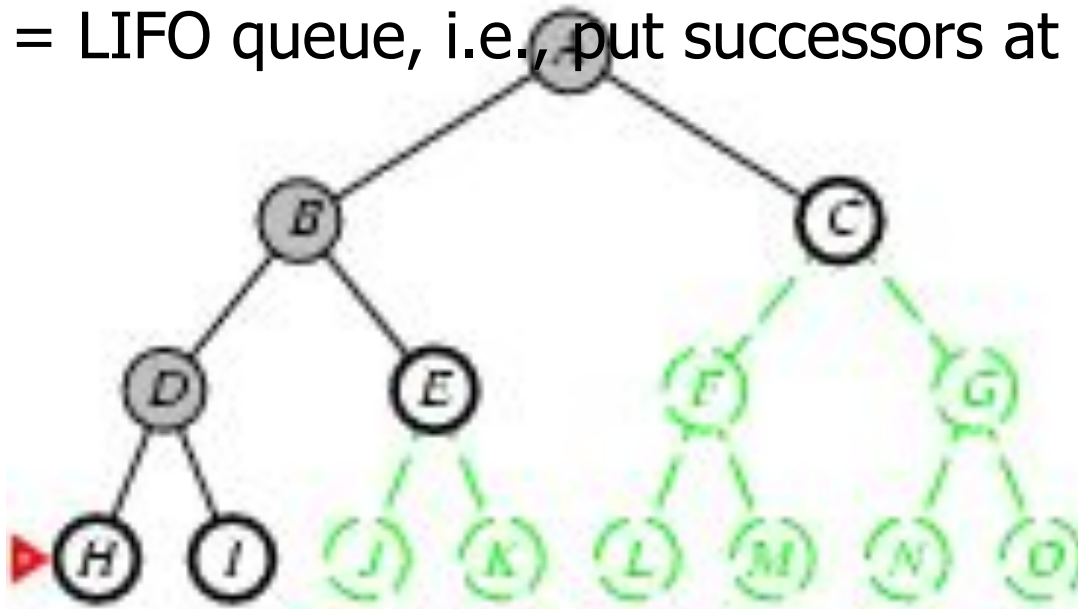
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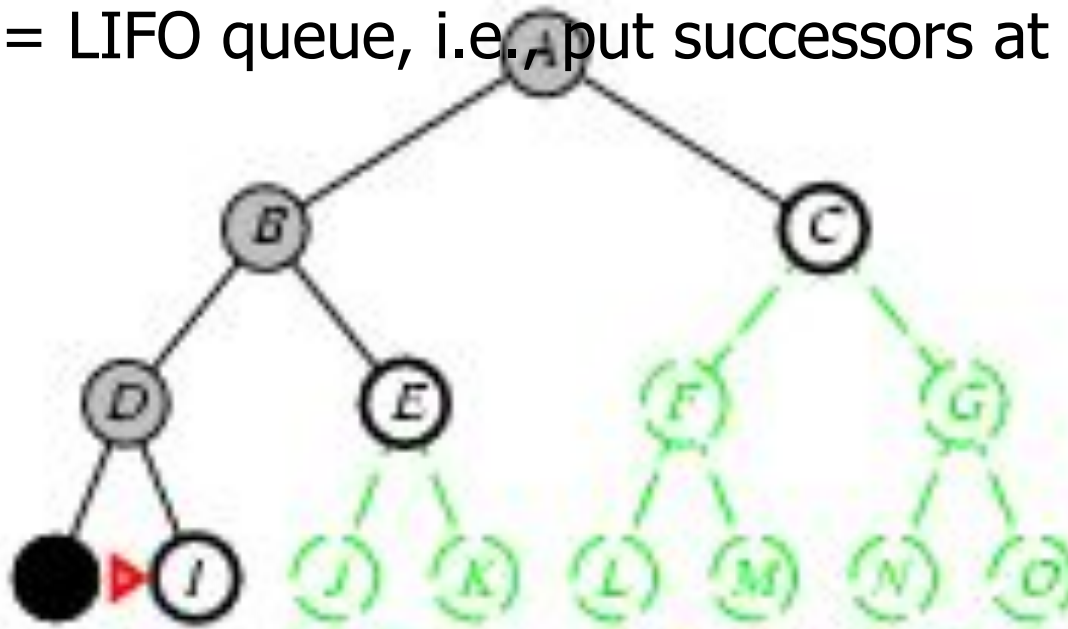
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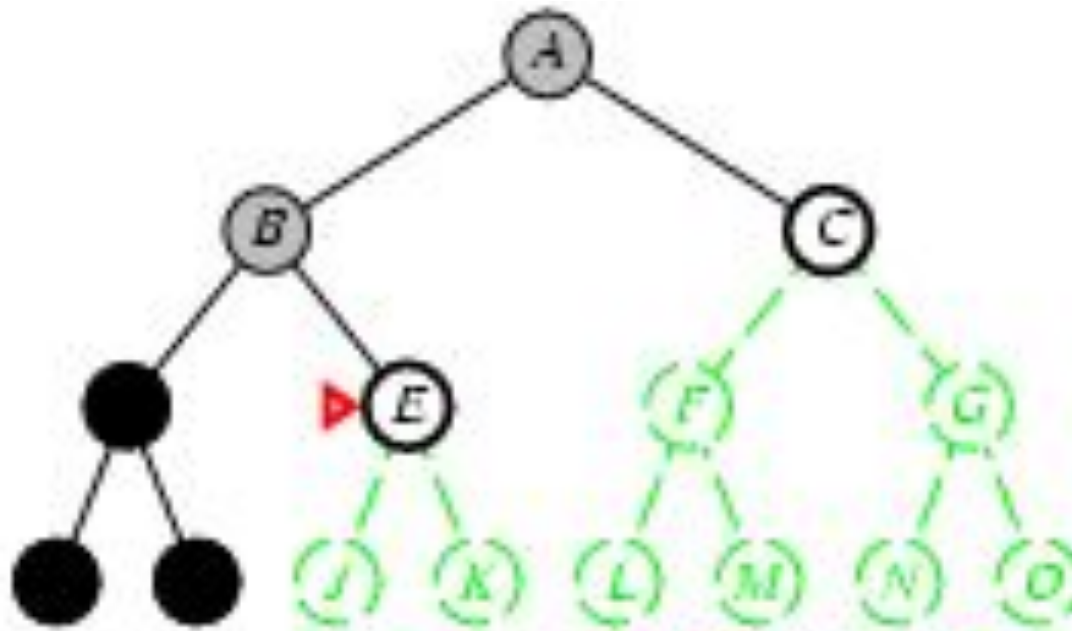


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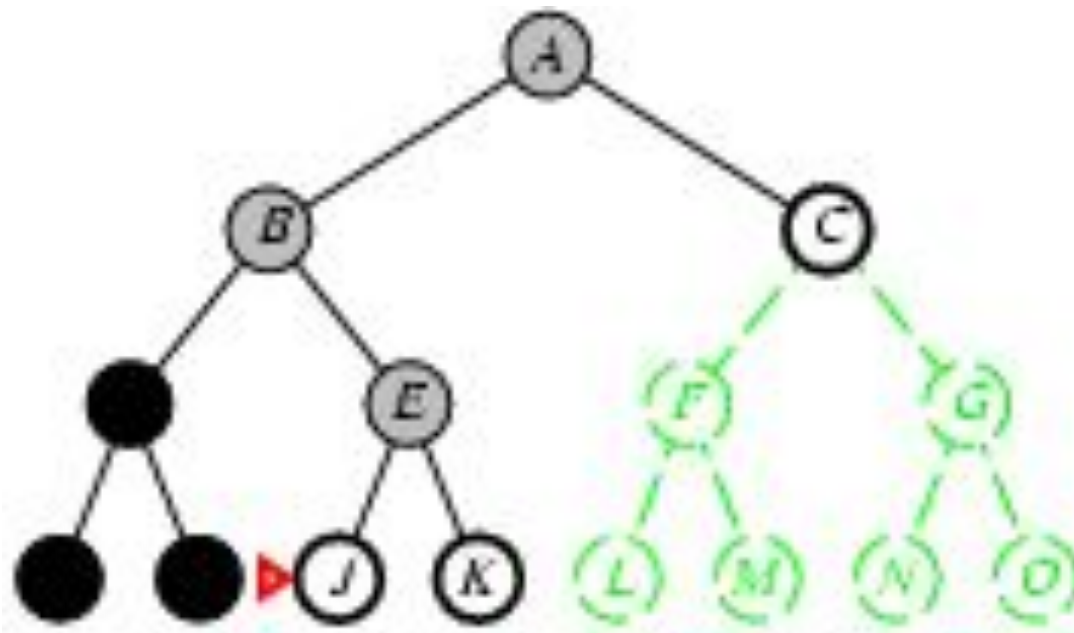


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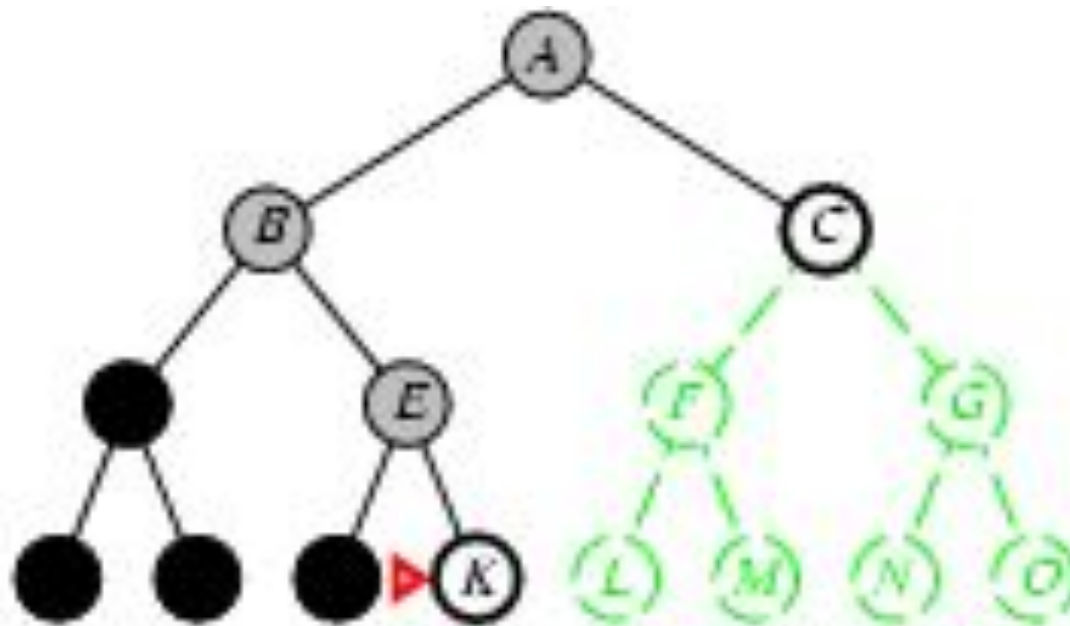
front

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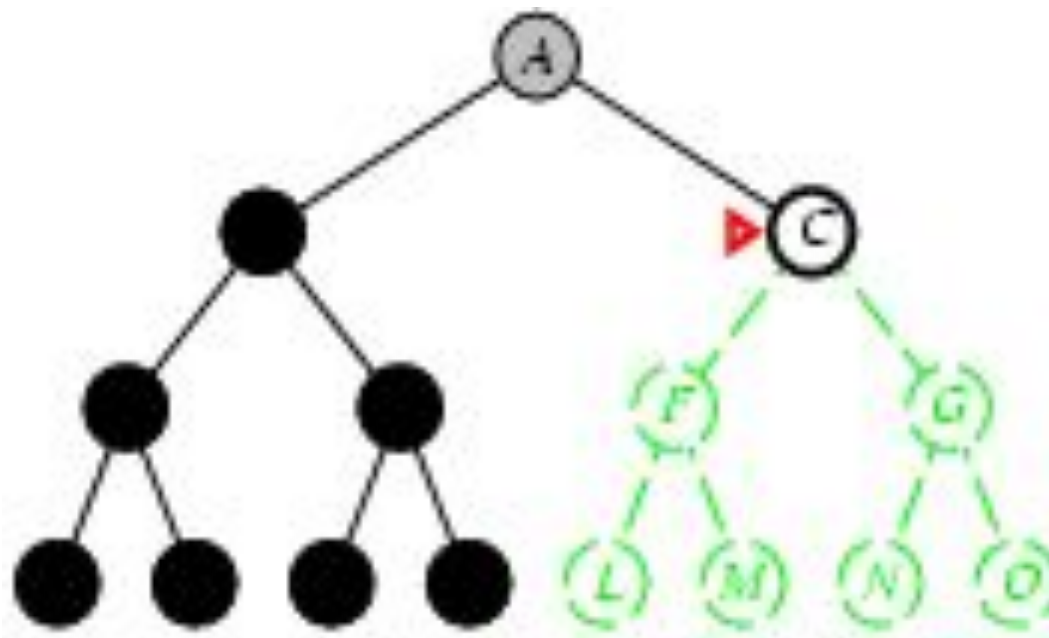


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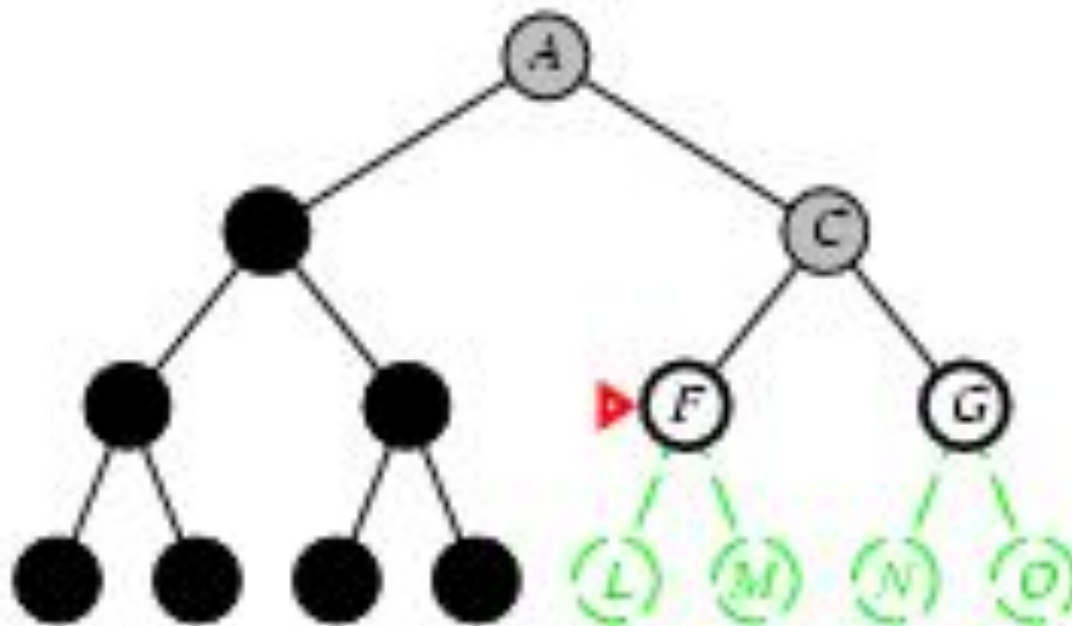


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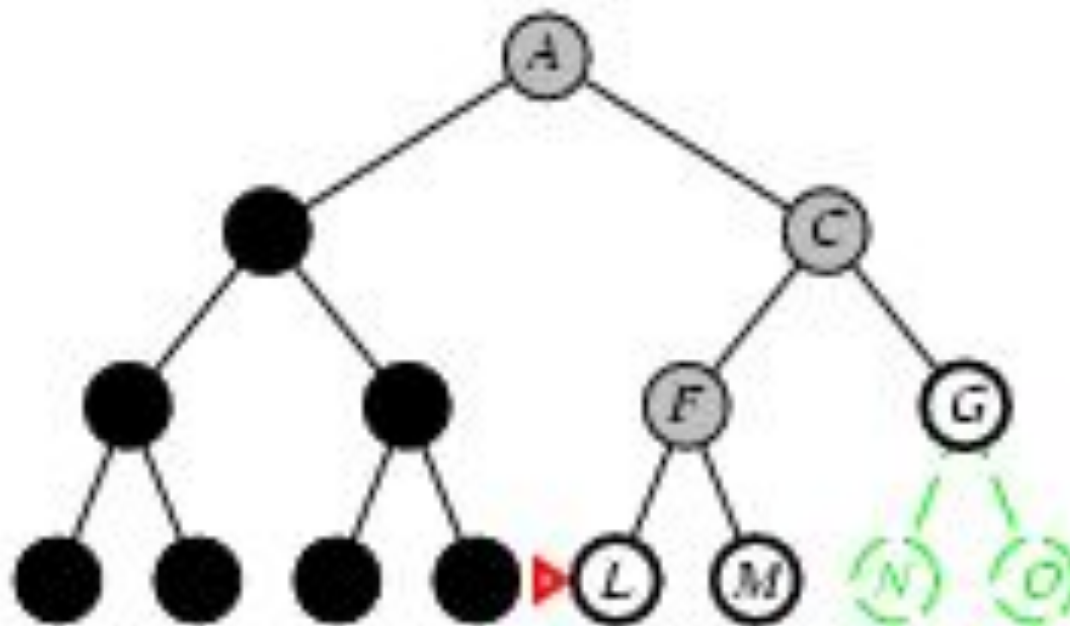


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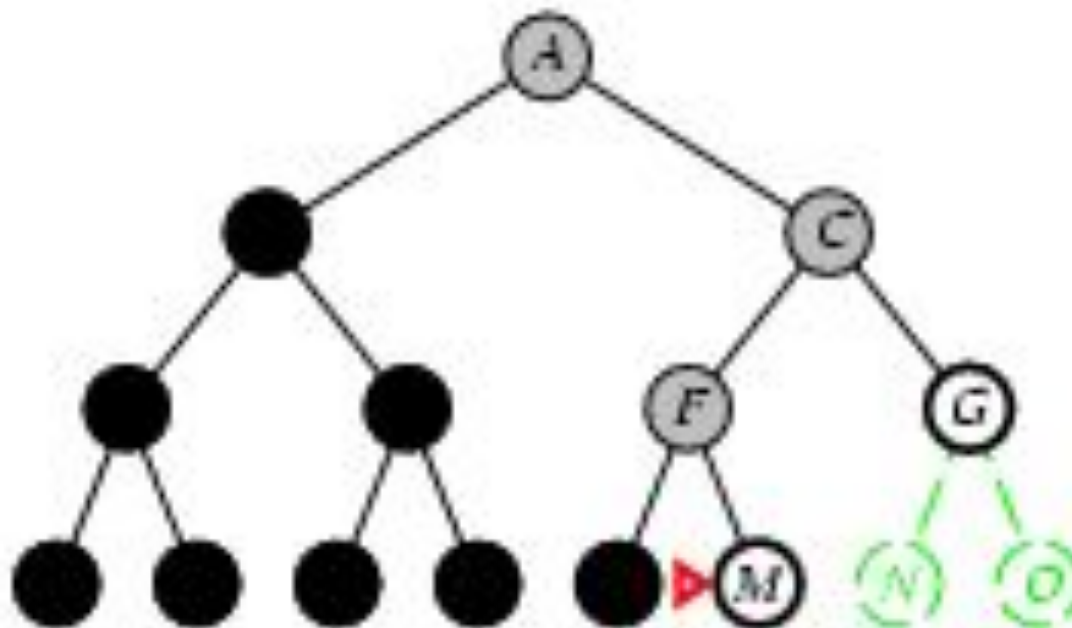


# Depth-first search

- Expand deepest unexpanded node

- Implementation:

- *fringe*





# Properties of depth-first search

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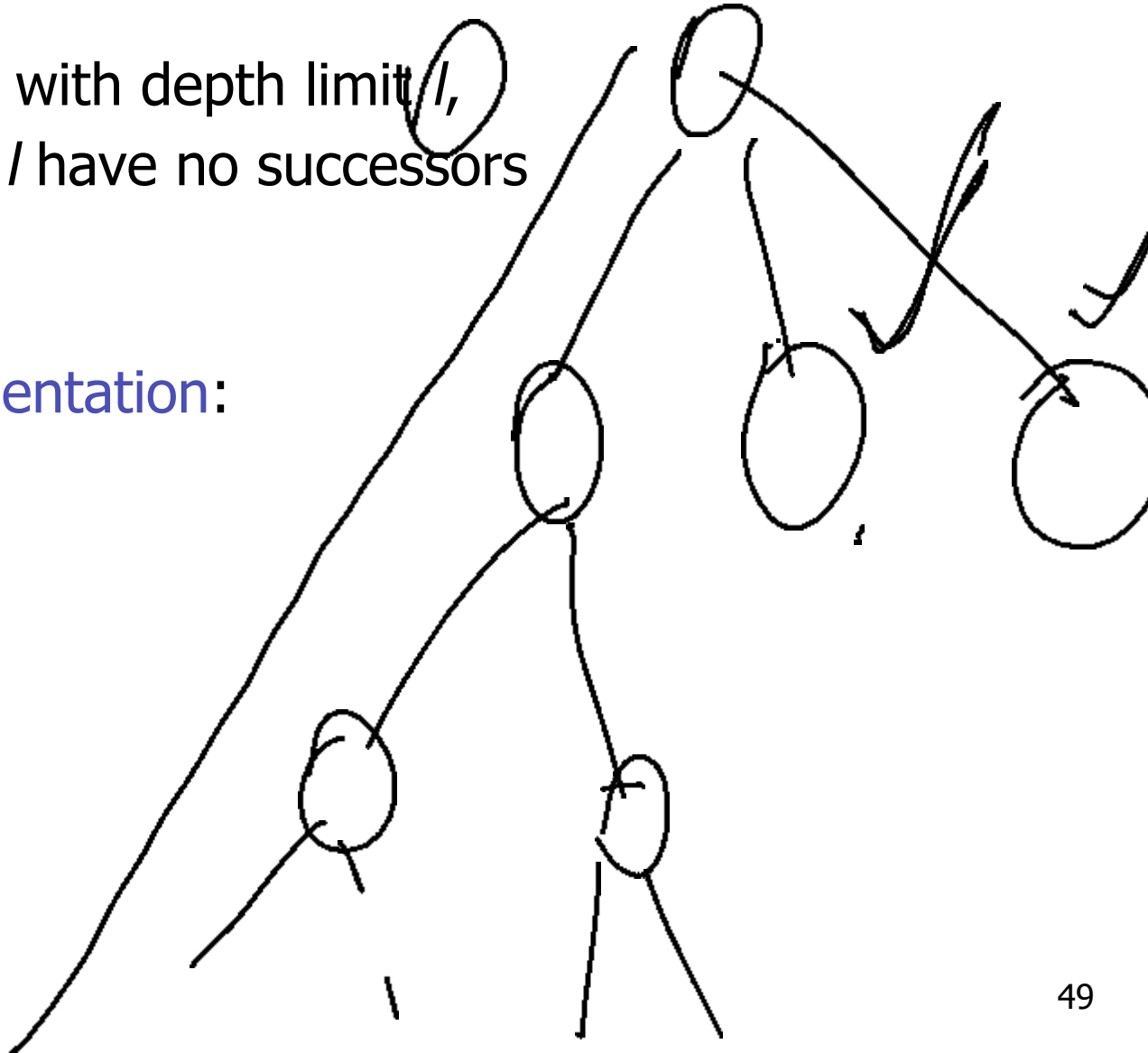
- Complete? No: fails in infinite-depth spaces, spaces with loops
  - Modify to avoid repeated states along path
    - ☐ complete in finite spaces
- Time?  $O(b^m)$ : terrible if  $m$  is much larger than  $d$ 
  - but if solutions are dense, may be much faster than breadth-first
- Space?  $O(bm)$ , i.e., linear space!

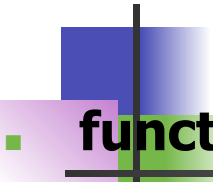


# Depth-limited search

= depth-first search with depth limit  $l$ ,  
i.e., nodes at depth  $l$  have no successors

- Recursive implementation:



- 
- **function** DEPTH-LIMITED-SEARCH(problem,limit) **returns** a solution, or failure/cutoff
  - **return** RECURSIVE-DLS(MAKE-NODE(problem.INITIAL-STATE),problem,limit)
  - **function** RECURSIVE-DLS(node,problem,limit) **returns** a solution, or failure/cutoff
  - **if** problem.GOAL-TEST(node.STATE) **then return** SOLUTION(node)
  - **else if** limit = 0 **then return** cutoff
  - **else**
  - cutoff occurred?  $\leftarrow$  false
  - **for** each action in problem.ACTIONS(node.STATE) **do**
  - child  $\leftarrow$  CHILD-NODE(problem,node,action)
  - result  $\leftarrow$  RECURSIVE-DLS(child,problem,limit - 1)
  - **if** result = cutoff **then**
  - cutoffoccurred?  $\leftarrow$  true
  - **else if** result  $\neq$  failure **then** return result
  - **if** cutoff occurred? **then** return cutoff **else** return failure



# Iterative deepening search

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```
function ITERATIVE-DEEPENING-SEARCH( problem) returns a solution, or fail-  
ure  
  inputs: problem, a problem  
  for depth  $\leftarrow$  0 to  $\infty$  do  
    result  $\leftarrow$  DEPTH-LIMITED-SEARCH( problem, depth)  
    if result  $\neq$  cutoff then return result
```

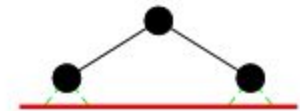
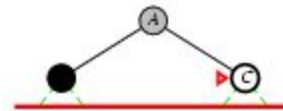
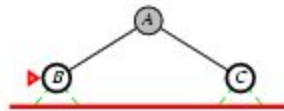
# Iterative deepening search / =0

Limit = 0



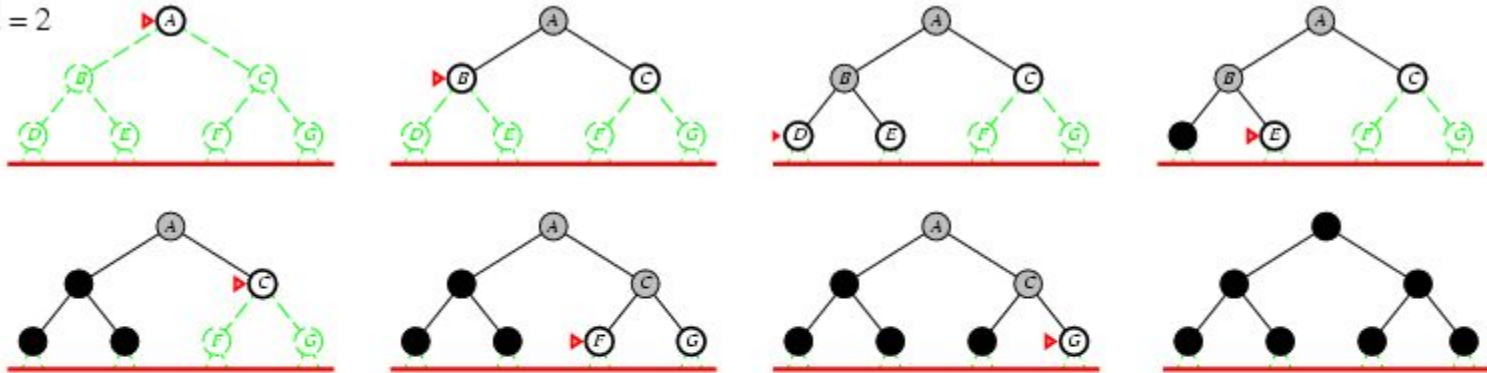
# Iterative deepening search / =1

Limit = 1



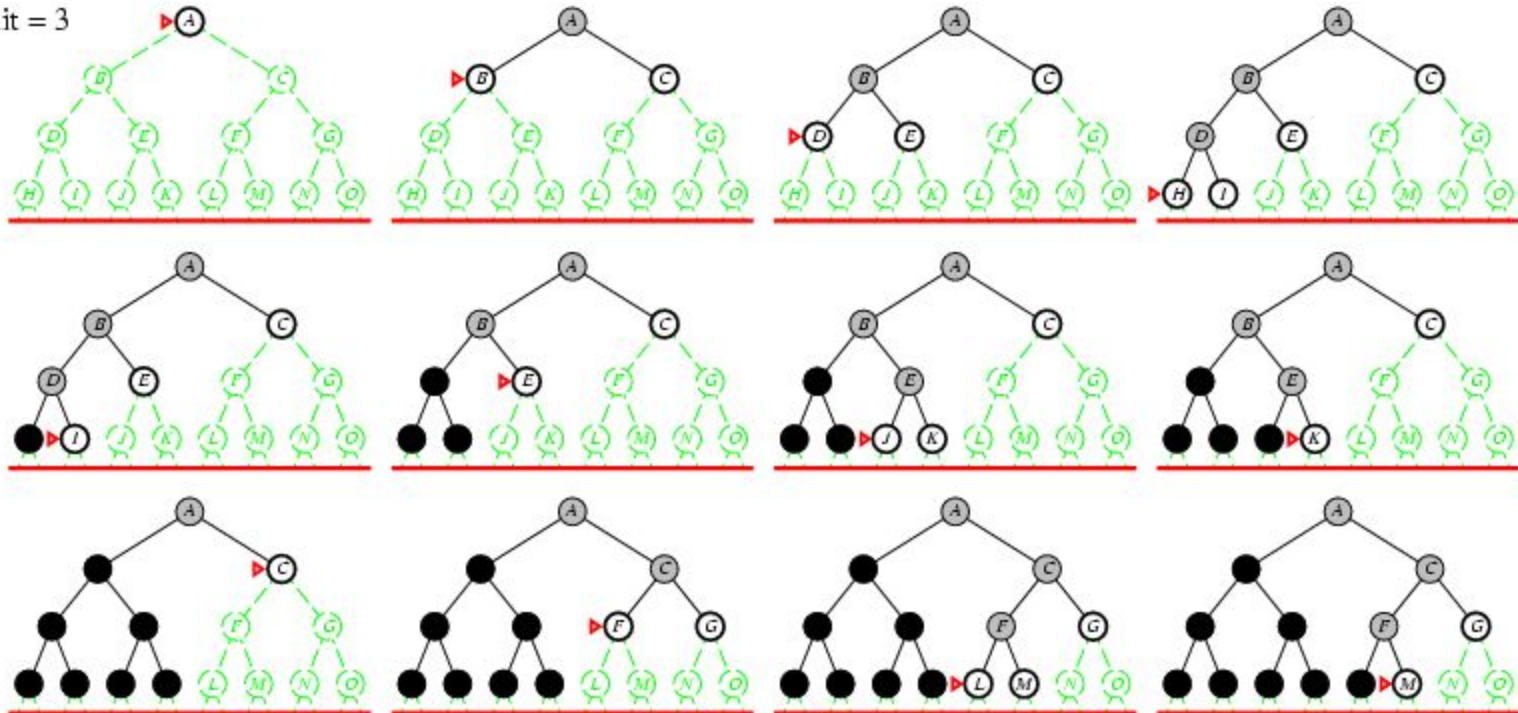
# Iterative deepening search / =2

Limit = 2



# Iterative deepening search / =3

Limit = 3



# Iterative deepening search

- Number of nodes generated in a depth-limited search to depth  $d$  with branching factor  $b$ :

$$N_{DLS} = b^0 + b^1 + b^2 + \dots + b^{d-2} + b^{d-1} + b^d$$

- Number of nodes generated in an iterative deepening search to depth  $d$  with branching factor  $b$ :

$$N_{IDS} = (d+1)b^0 + d b^1 + (d-1)b^2 + \dots + 3b^{d-2} + 2b^{d-1} + 1b^d$$

//last term is bottom level

- For  $b = 10$ ,  $d = 5$ ,

- $N_{DLS} = 1 + 10 + 100 + 1,000 + 10,000 + 100,000 = 111,111$

- $N_{IDS} = 6 + 50 + 400 + 3,000 + 20,000 + 100,000 = 123,456$



# Properties of iterative deepening search

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- Complete? Yes
- Time?  $(d+1)b^0 + d b^1 + (d-1)b^2 + \dots + b^d = O(b^d)$
- Space?  $O(bd)$
- Optimal? Yes, if step cost = 1