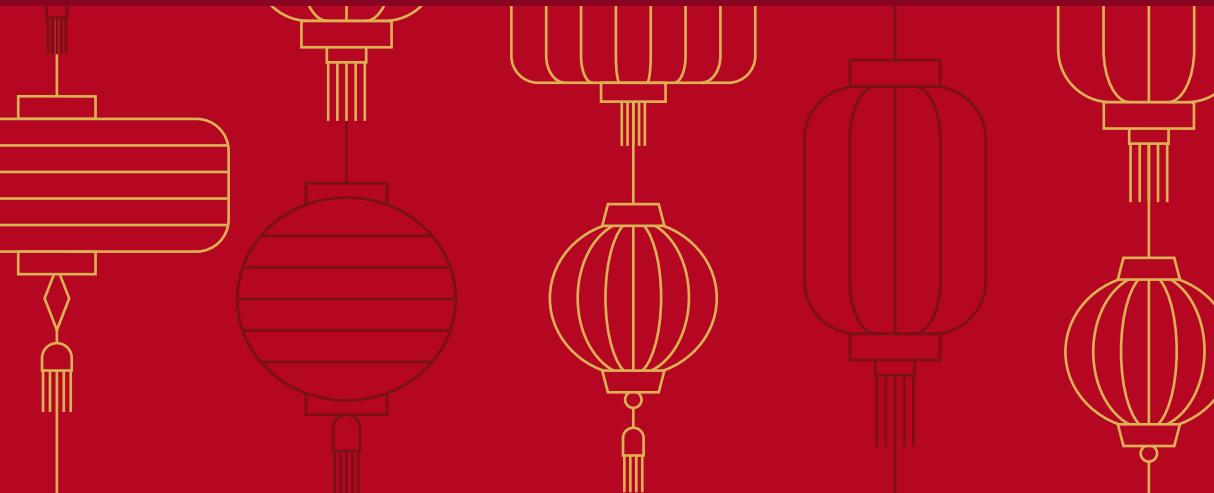


1 Heavis

~ Arbind

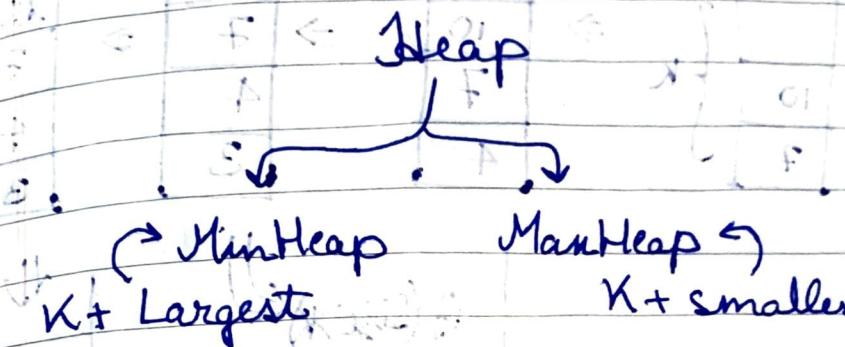


# Heap

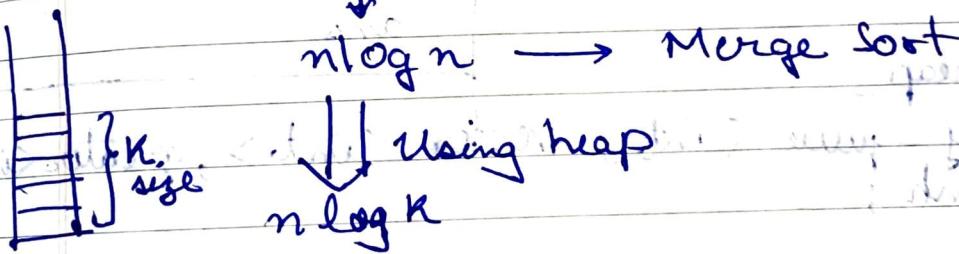
## Identification

- 1) K
- 2) smallest/largest

→ Heap



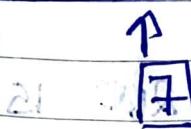
## Sorting



arr[]: [7 | 10 | 4 | 3 | 20 | 15]       $n = 3$

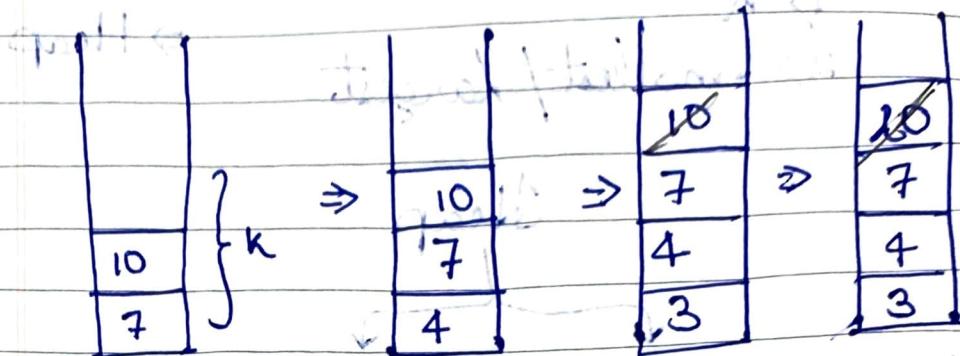
if we sort,

3	4	7	10	15	20
---	---	---	----	----	----



extra work done to sort this element which has no use

7	10	4	3	20	15
---	----	---	---	----	----



Code + Time Complexity

Maxheap

priority\_queue<int> maxh;

if (size > k)

pop()

→ top() is our ans

ans

Minheap

priority\_queue<int, vector<int>, greater<int> minh;

Using type def

type def pair<int, pair<int, int>> ppi;

K<sup>th</sup> Smallest Element = arr[0]

arr[] = {7, 10, 4, 3, 20, 15}

K = 3

## Code

```

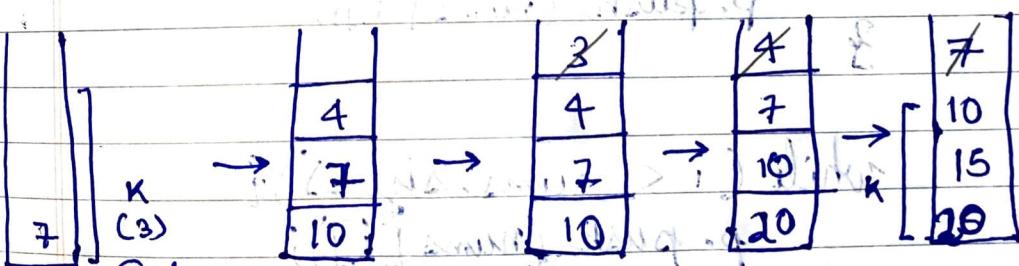
priority_queue<int> maxh;
for(int i=0; i< size; i++){
    maxh.push(arr[i]);
    if(maxh.size() > k)
        maxh.pop();
}
return maxh.top();

```

## K largest Element

arr[]: [ 7 | 10 | 4 | 3 | 20 | 15 ]  
 $k = 3$

Output: 20, 15, 10



Code:

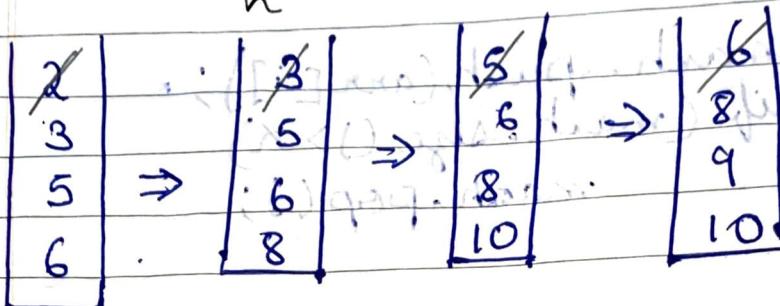
```

while(!minh.empty()){
    cout << minh.top() << " ";
    minh.pop();
}

```

# Sort a $k$ sorted array

array:  $\begin{matrix} 6 & 5 & 3 & 2 & 8 & 10 & 9 \end{matrix}$   $k = 3$



$2 \rightarrow 3 \rightarrow 5 \rightarrow 6 \rightarrow 8 \rightarrow 9 \rightarrow 10$

## Code:

```
priority_queue<int, vector<int>, greater<int>> p;
```

```
int i;
for(i=0; i<k; i++){
    p.push(nums[i]);
```

```
while(i < nums.size()){
    p.push(nums[i]);
```

```
cout << p.top() << " ";
p.pop();
```

```
p.pop();
```

while (!pq.empty()) {

cout << pq.top() << " ";  
pq.pop();

}

Kth Closet Number

arr[] : 5 6 7 8 9  
 $x = 7$

$K = 3$

$x = 7$

5 6 7 8 9  
7 7 7 7 7  
x = 7

| 0, 0 |

↓

we will use  
pair one will be  
abs(arr[i] - x) other  
is arr[i]

5 6 7 8 9  
| 2, 5 |  
| 1, 6 |  
| 0, 7 |

$x = 7$

$K = 3$

elements

$\epsilon \leftarrow 1$

$l \leftarrow \epsilon$

$h \leftarrow \epsilon$

priority-queue < pair<int, int> pq;

abs (arr[i] - x)

arr [i]

Code: ~~b((PriorityQ))~~

priority - queue <pair<int, int>> p;

for(int i=0; i < size; i++) {  
 p.push({abs(arr[i] - x), arr[i]});

if (p.size() > k)  
 p.pop();

}

while (!s.empty()) {

cout << p.top().second << " "  
p.pop();

}

### Top K Frequent Numbers

arr[]:

1	1	1	3	2	2	4
---	---	---	---	---	---	---

k = 2

1, 0

1 → 3

3 → 1

2 → 2

4 → 1

apply logic, we get 1 > 3 > 2 > 0

{1, 0}

(x + 1) mod 4

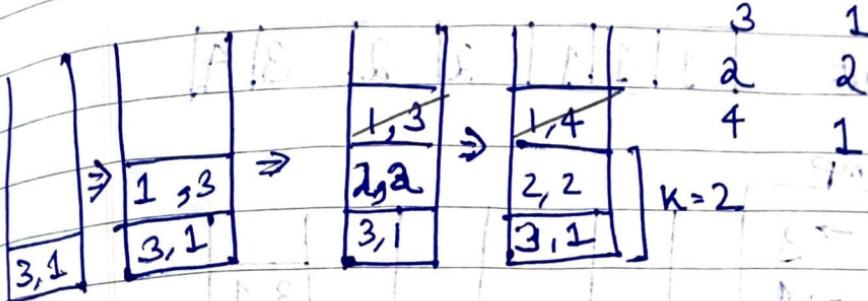
Largest  
highest  
Top

Smallest  
Lowest  
Closest

min-heap

max-heap

$k=2 \quad [1, 1, 1, 1, 3, 2, 2, 4] \Rightarrow [1, 3]$



unordered\_map<int, int> mp;  
priority\_queue<int, vector<int>, greater<int>> minh;  
for(int i=0 ; i<n ; i++)  
 mp[arr[i]]++

for( auto i = mp.begin(); i != mp.end(); i++)  
 minh.push({i->second, i->first});  
 if( minh.size() > k )  
 minh.pop();

while( ! minh.empty() ) {  
 cout << minh.top() << endl;  
 minh.pop();  
}

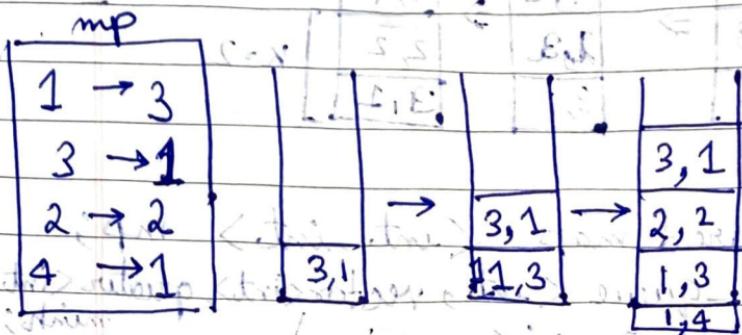
Notes:

In priority-queue int is replaced with pair

# Frequency Sort (Max, heap)

arr[]: [1 | 1 | 1 | 3 | 2 | 2 | 4]

Output [1 | 1 | 1 | 2 | 2 | 3 | 4]



## Code:

```

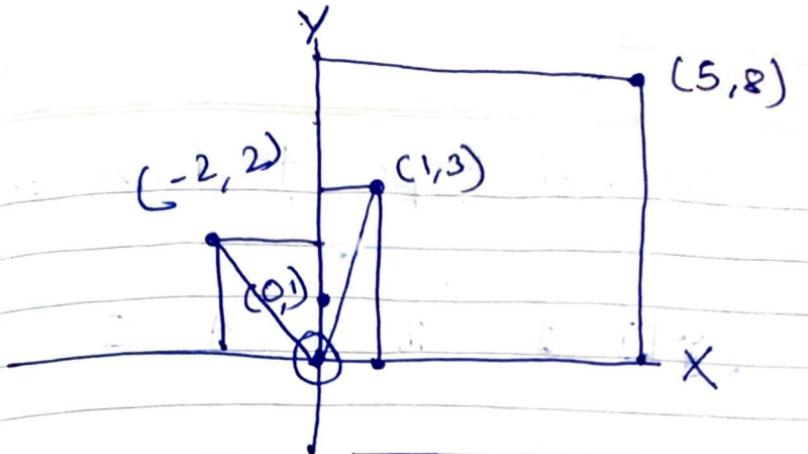
while ( !maxh.size() ) {
    int freq = maxh.top().first;
    int ele = maxh.top().second;
    for (int i=0; i< freq; i++) {
        cout << ele << " ";
        maxh.pop();
    }
}

```

## K - closest points to Origin

arr[]: =  $\begin{bmatrix} 1 & 3 \\ -2 & 2 \\ .5 & 8 \\ 0 & 1 \\ n & 4 \end{bmatrix}$

$$k=2$$



$$\text{distance} = \sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2}$$

$$= \sqrt{x^2 + y^2}$$

↓

if we sort on base of  
 $x^2 + y^2$  & remove root the  
 answer will be same

`pair<int, pair<int, int>> p;`

Code: `pair<int, pair<int, int>> p;`  
`priority_queue<pair<int, int>> m;`

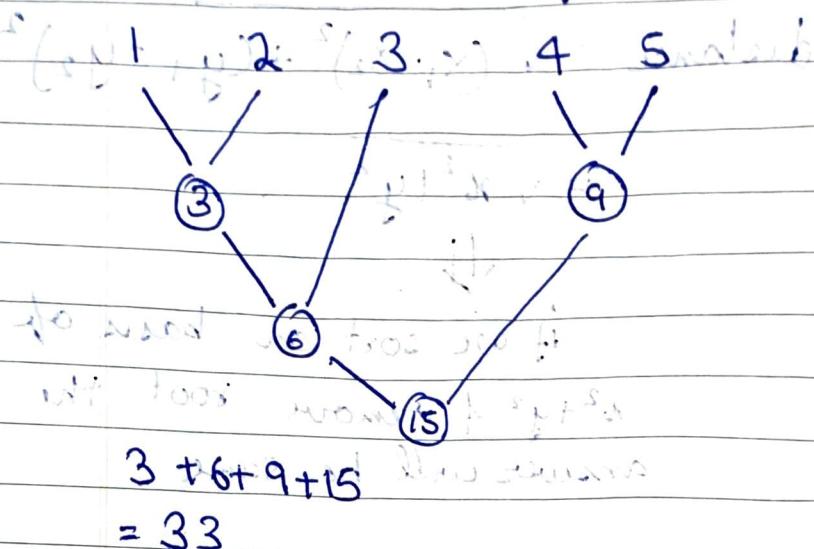
```
for(int i = 0 ; i < n ; i++){
    m.push({arr[i][0] * arr[i][0] +
            arr[i][1] * arr[i][1],
            {arr[i][0], arr[i][1]}});
}
```

```
if(m.size() > k){
    m.pop();
}
```

Print the whole queue

## Connect Ropes to Minimise the cost

arr[] : 1 2 3 4 5



Concept:

⇒ Minimal number will come when we take 2 ropes with min value.

1	$2+1=3$	$3+3=6$	$4+5=9$	8
2	3	4	5	9
3	3	5	6	
4	4	6		
5	5			

Total =  $3 + 6 + 9 + 15 = 33$

Code:

priority\_queue<int, vector<int>, greater<int>> q;

```
for(int i = 0; i < stones.size(); i++) {  
    q.push(stones[i]);
```

{ q < top > min. element

if(i > n - 1) i = 0 + (i - n);

int temp = 0, ans = 0;

```
while(q.size() != 1) {
```

```
    for(int i = 0; i < n; i++) {
```

temp += q.top();

q.pop();

{ q < top > min. element

q.push(temp);

ans += q.top();

temp = 0;

{

return ans;

Sum of Elements

arr[]: 

1	3	12	5	15	11
---	---	----	---	----	----

$$K_1 = 3$$

$$K_2 = 6$$

1	3	5	11	12	15
---	---	---	----	----	----

↑  
3rd

↑  
6th

## Code

```
int n = v.size();
if (n < k1 || n < k2)
    return -1;
```

```
if (k2 > k1) i += i;
```

```
    swap(k1, k2);
```

```
priority_queue<int> p;
for (int i = 0; i < n; i++)
```

```
    p.push(v[i]);
```

```
if (n < k1)
```

```
    p.pop();
```

```
int first = p.top();
```

```
for (int i = 0; i < k1 - k2; i++)
```

```
    p.pop();
```

```
return first + p.top();
```