

# Labwork 4 - Conditional Statement

These labworks are automatically assessed based on an archive you have to deliver on time on the Moodle webpage. To make the archive, you have to type the command:

#### > make archive

This produces a file named archive.tgz that you have to deposit. The compilation labwork are roughly held each 2 weeks and the delivery date is usually on sunday before the next labwork week.

The three first labworks have introduced the whole procedure to extend our language:

- 1. declare the new required tokens in parser.mly,
- 2. generate these tokens in lexer.mll,
- 3. write the rules required by the extension in parser.mly,
- 4. add, in the rule action, the checks for the Semantic Analysis and the generation of the ASTs (in parser.mly),
- 5. update the module comp.ml with the code generation for the extended ASTs.

In addition, these labworks have dermonstrated the success of an incremental strategy to develop a compiler. A real programming language is often big and it is mostly impossible to detect all issues coming with the language. Instead, we start developing a small working core and then we add extensions step by step.

In this labwork, we will perform new steps in order to add conditional construction to AutoCell and therefore we will still apply the same strategy and method.

## 1 Simple Condition

In AutoCell, a condition has the following form:

```
\begin{array}{l} \mbox{if sum} = 3 \ \mbox{then} \\ \mbox{[0, 0]} := 1 \\ \mbox{end} \\ \\ \mbox{if } [0, 0] = 1 \ \mbox{then} \\ \mbox{if sum} < 2 \ \mbox{then} \\ \mbox{[0, 0]} := 0 \\ \end{array}
```

```
\begin{array}{rl} \mathrm{sum} \; := \; 0 \\ \mathrm{end} \end{array} end
```

The condition is very to usual languages: an if keyword, a conditional expression, a keyword then, a sequence (possibly empty) of statements and the keyword end to complete the condition construction. Clearly, a condition is a statement, not an expression as it does not produce any value. Its role is mainly to manage the execution flow according to the condition.

The condition is a comparison between two expressions. This supported comparators are: "=", "! =", "<", "<=", "<" and ">=".

**To Do** modify the Lexical Analysis and te Syntactic Analysis to support the conditional statement. Test it with autos/simpleif.auto.

**To Do** The AST module ast.ml already contains ASTs for conditional statement, IF\_THEN, and ASTs for conditions, cond. Change the actions in parser.mly to build the AST.

**To Do** Add the support for the else keyword to our condition statement.

```
\begin{array}{ll} \text{if sum} < 2 & \text{then} \\ & [0 \, , \, 0] \; := \; 0 \\ \text{else} \\ & [0 \, , \, 0] \; := \; 1 \\ \text{end} \end{array}
```

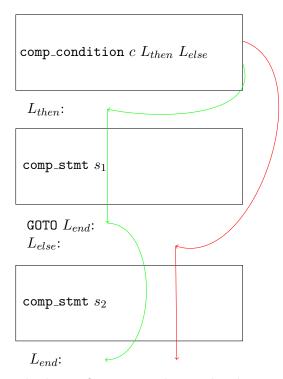
Test it with autos/ifelse.auto.

#### 2 Generation

In order to generate quadruplets for a conditional statement, we will to implement the condition and the only way is to use the conditional branches in the quadruplets:  ${\tt GOTO\_EQ}$ ,  ${\tt GOTO\_NE}$ , etc. Notice that a  ${\tt GOTO\_XXX}$  can be associated with each comparator  ${\tt COMP\_XXX}$ .

We need also to visualize several execution paths (when the condition true, when the condition is false) that will be rendered as a linear sequence of quadruplets. To help visualize the different execution flows, we advise to first represent them with a generation model. This model shows the quadruplets as a sequence (including sequences that are unknown corresponding to substatements) and with arrows showing the different execution paths.

The generation model below is proposed for the AST IF\_THEN $(c, s_1, s_2)$ :



The boxes figures out the quadruplets result of the compilation of c,  $s_1$  and  $s_2$ . The green edges represent the execution path taken when the condition is true, the red edges when the condition is false. Observe that:

- We need three labels  $L_{then}$ ,  $L_{else}$  and  $L_{end}$ ,
- The branch quadruplet is generated from comp\_cond and this is why  $L_{then}$  and  $L_{else}$  as passed as arguments.
- A GOTO quadruplet has been inserted between  $s_1$  and  $s_2$  quadruplets to avoid to execute the *else* part after the execution of the *then* part.

We are mostly done, we only need:

- a way to insert labels in the quadruplet: done with the pseudo quadruplet LABEL number;
- a way to generate unique label numbers: function new\_lab ().

**To Do** Implement the quadruplet generation fot the conditions in the function comp\_cond. Notice that only the condition COMP has to be implemented: the other types of condition might be implemented in a subsequent labwork. You have to keep in mind the goal of the quadruplets you have to generate: if the condition is true, a branch has to be done to the label  $L_{then}$ ; otherwise, a branch to the label  $L_{else}$  is performed. It is important to generate both branch quadruplets.

**Note** The quadruplets generated by comp\_cond are not very beautiful but this works. Yet, in a compiler, a subsequent very simple optimization (called *peephole optimization*) can nicefully fix this problem.

**To Do** Now, helped by the generation model proposed above, and the function comp\_cond, you can generate the quadruplets for a condition,  $\text{IF\_THEN}(c, s_1, s_2)$  in the function comp\_stmt. You can look and check the generated quadruplets with autos/simpleif.auto and autos/ifelse.auto.

### 3 Multiple Conditions

In this last exercise, we will extend the conditional statemenet with the elseif keyword, as exposed in the example below:

```
\begin{array}{ll} \mbox{if } [0 \;,\; 0] \; = \; 1 \; \mbox{then} \\ \mbox{if } \mbox{sum} \; < \; 2 \; \mbox{then} \\ \mbox{} [0 \;,\; 0] \; := \; 0 \\ \mbox{elsif sum} \; > \; 3 \; \mbox{then} \\ \mbox{} [0 \;,\; 0] \; := \; 0 \\ \mbox{end} \\ \mbox{else} \\ \mbox{if } \mbox{sum} \; = \; 3 \; \mbox{then} \\ \mbox{} [0 \;,\; 0] \; := \; 1 \\ \mbox{end} \\ \mbox{end} \\ \mbox{end} \\ \mbox{end} \end{array}
```

The *then* part of a conditional statement may be followed by else (achieved in the previous section) but also by a sequence of elsif. Each elsif has a condition, a then keyword and its own sequence of statements. Ultimately, this sequence of elseif can be completed by an else (but it is not mandatory).

**To Do** Modify the front-end of autocc to parse the elsif construction. Test if with autos/conway.auto.

**To Do** Modify parser.mly to generate the AST for supporting elsif: you do not need to modify the definition of ASTs for this!. Test if with autos/conway.auto.

**To Do** Now you can run the *Conway's Game Life* cellular automaton written in AutoCell:

```
> ./autocc autos/conway.auto
> ./autoas autos/conway.s
> ./autocell autos/conway.exe maps/MAP.map
```

You can test with maps: conway-1.map, conway-2.map, ..., conway-5.map and enjoy.