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## E - LCM Sequence (/contests/abc412/tasks/abc412\_e) Editorial by

en\_translator (/users/en\_translator)

Since  $A_n$  is monotonically increasing, the answer to the problem is

• 1+ (the number of integers i among  $L+1 \le i \le R$  and  $A_{i-1} < A_i$ )

(The first 1 corresponds to  $A_L$ .)

Which  $A_i$  satisfies  $A_{i-1} < A_i$ ?

We have  $A_n = \mathrm{LCM}(1,2,\ldots,n)$ . Consider how many times p divides  $A_n$ . Among  $1,2,\ldots,n$ , the integer that can be divided by p that most times is

• the maximum  $p^k$  with  $p^k < n$  (k is an integer),

and the number of times p divides the integer is k, so  $A_n$  is also divided by p, k times (using the integer k above).

Therefore, if and only if n can be represented as a power of a prime p (which is called a **prime power**),  $A_n$  can be divided by p one more time than  $A_{n-1}$ . For the other primes than p, the number of times it divides  $A_n$  is same as that for  $A_{n-1}$ , so we can say that  $A_n$  is p times  $A_{n-1}$ .

Conversely, if n is not a prime power, then for any prime p, the number of times p divides  $A_n$  is the same as that for  $A_{n-1}$ , so  $A_n = A_{n-1}$ .

By the discussion above, the answer to this problem turns out to be

• 1+ (the number of prime powers i such that  $L+1 \leq i \leq R$ )

2025-06-30 (Mon) 00:06:39 -04:00 It is known that one can find the prime factorizations of all integers  $L+1,\ldots,R$  using an algorithm called **segmented sieve** in  $O(M\log\log M)$  time, where  $M=\max(R-L,\sqrt{R})$ . This algorithm is already featured in past problems, so please refer to the editorial for the past problem (https://atcoder.jp/contests/abc227/editorial/2909)).

To roll up, the problem can be solved in a total of  $\mathrm{O}(M\log\log M)$  time, where  $M=\max(R-L,\sqrt{R})$ , which is fast enough.

• Sample code (C++)

```
Copy
1. #include <cmath>
2. #include <iostream>
3. #include <vector>
4. using namespace std;
5.
6. vector<int> prime_enumerate(int N) {
7.
     vector<bool> is_prime(N + 1, true);
    vector<int> primes;
9. if (N < 2) return primes;
     is_prime[0] = is_prime[1] = false;
10.
11.
     for (int i = 2; i * i <= N; ++i) {
      if (is_prime[i]) {
12.
         for (int j = i * i; j <= N; j += i) is_prime[j] = false;</pre>
13.
14.
      }
15.
     }
16.
     for (int i = 2; i <= N; ++i) {
17.
      if (is_prime[i]) primes.push_back(i);
18.
19.
     return primes;
20. }
21.
22. int main() {
     long long L, R;
23.
24.
     cin >> L >> R;
25.
     vector<int> vis(R - L);
26.
     int ans = 1;
     for (int p : prime_enumerate(sqrt(R) + 100)) {
27.
28.
      for (long long x = (L / p + 1) * p; x <= R; x += p) {
29.
         if (vis[x - (L + 1)]) continue;
30.
         vis[x - (L + 1)] = 1;
        long long y = x;
31.
32.
        while (y % p == 0) y /= p;
         if (y == 1) ans++;
33.
34.
       }
35.
36.
     for (int v : vis) ans += v == 0;
37.
     cout << ans << "\n";</pre>
                                                                         2025-06-30 (Moh)
38. }
                                                                          00:06:39 -04:00
```

posted: a day ago last update: a day ago

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