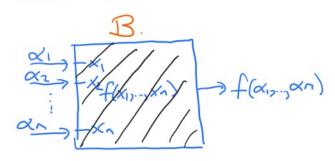
Assignment #6 due 11pm Monday Nov. 27th Assignment #7 due 11pm Monday Dec 4th.

The "Black box" representation for polynomials.

Let  $f \in R[x_1, ..., x_n]$ , R an integral domain eg. Z,  $\mathbb{O}(\alpha)$ ,  $\mathbb{F}_{2}$ .

Sparse representation:  $f = \underbrace{\xi}_{i=1} \underbrace{\alpha_i \cdot M_i(x_1, -\infty x_n)} \underbrace{\alpha_i \in \mathbb{R} \setminus \xi \delta \xi}_{monomial}$ 

Black-box representation:  $B:R^n \to R$  is a computer program that on input of  $\alpha \in R^n$  computes  $f(\alpha)$  i.e.  $B(\alpha) = f(\alpha)$ .



od; od;

We cannot see inside B.

All we can do is evaluate

B at a point  $\alpha \in \mathbb{R}^n$ .

We "probe" the black-box

at a point  $\alpha \in \mathbb{R}^n$ .

Let d=deg(f), t=#f, for R=Z let  $h=\|f\|_{\infty}=\max_{1\leq i\leq t}\|a_i\|$ . We may or may not know bounds  $D \geqslant d$ ,  $T \geqslant t$ ,  $H \geqslant h$ .

Example.  $f = \det(T_2) = \det(\left[\begin{matrix} u \lor w \\ \lor u \lor u \end{matrix}\right]) \in \mathbb{Z}[u_i v_i w].$ 

B:= proc(alpha: list(integer)) Notice  $deg(f) \le 3 = D$  local T3, i, j; uses Linear Algebra;  $lifloo \le 3! = T$   $lifloo \le 3! = H$   $lifloo \le 3! = H$ 

Determinant (T3);

How can we multiply two polynomials f, g ∈ R[sc, -, xn] given by black boxes Bf: R^->R and Bg: Ru->R? Let h=f.g.

Bh BBmultiply := proc(Bf, Bg)  $\alpha \rightarrow Bf(\alpha) \cdot Bg(\alpha) \rightarrow h(\alpha)$ end

This costs O(1).

For R=Z it is needful to use the Chinese remainder theorem. A modular black-box representation for  $f\in Z[x_0, x_1]$ is a black-box B:(Zp,p) that for  $x\in Zp$  computes f(x) mod p.

of f(x1,7-)Xn) -> f(ox) mod p.

B := proc(alpha:: list(integer), p::prime)
local T;
uses Linear Algebra;
T:=ToeplitzMatrisc(alpha, symmetric);
Det(T) mod P;
end;

Eiven a black-box  $B: R^n \rightarrow R$  for  $f \in R[x_1, \dots, x_n]$ Is f = 0? What is deg(f)?  $Aeg(f, x_i)$ ? What is t = f? What is t = f?

Interpolate f, i.e., find a cell and Michingxin).