

An example of a thesis on the subject of your degree

by

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Abstract

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Keywords: thesis template; Simon Fraser University; L^AT_EX time travel paradoxes

Dedication

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Chapter 1

Introduction

1.1 Blackboxes and Evaluation homomorphism

1.1.1 Blackboxes

Definition 1.1.1. *A blackbox is a function that takes an input and produces an output. The internal workings of the function are not known to the user. The user can only interact with the blackbox by providing inputs and observing the outputs.*

1.1.2 Evaluation homomorphism

Definition 1.1.2. *Let K be a field and $K[x] \in K(x)$ be the ring of polynomials and field of rational functions respectively. Let $\alpha \in K$ be a point in the field. The evaluation homomorphism is a map*

$$\begin{aligned}\phi: K[x] &\longrightarrow K[x]/(x - \alpha) \cong K \\ f(x) &\longmapsto f(\alpha)\end{aligned}$$

Chapter 2

Maximal Quotient Rational Function Reconstruction

2.1 Univariate Rational function reconstruction

We are given a black box for a rational function $F(x) = \frac{f(x)}{g(x)} \in K(x)$ where $f(x), g(x) \in K[x]$ and $g(x) \neq 0$. We need to recover the polynomials $f(x), g(x)$ upto a constant factor. We can evaluate the black box at n distinct points $\alpha_1, \dots, \alpha_n$ to get the values y_1, \dots, y_n .

2.1.1 Interpolation

Given a set of points

2.1.2 Extended Euclidean Algorithm

Algorithm 1 Extended Euclidean Algorithm

```
1:  $r_0 \leftarrow f, \quad s_0 \leftarrow 1, \quad t_0 \leftarrow 0$ 
2:  $r_1 \leftarrow g, \quad s_1 \leftarrow 0, \quad t_1 \leftarrow 1$ 
3:  $i \leftarrow 1$ 
4: while  $r_i \neq 0$  do
5:    $q_i \leftarrow r_{i-1} / r_i$ 
6:    $r_{i+1} \leftarrow r_{i-1} - q_i r_i$ 
7:    $s_{i+1} \leftarrow s_{i-1} - q_i s_i$ 
8:    $t_{i+1} \leftarrow t_{i-1} - q_i t_i$ 
9:    $i \leftarrow i + 1$ 
10: end while
11:  $l \leftarrow i - 1$ 
12: return  $r_l, s_l, t_l \in K[x]$ 
```

$$\begin{array}{r} \frac{f}{g} = \frac{g^{-1}f}{g) \quad f} \\ \quad \quad \quad \frac{-f}{0} \end{array}$$

Theorem 2.1.1. *Bezout's identity*

Let $f, g \in K[x]$ such that $f, g \neq 0$

Let $d = \gcd(f, g)$ then $\exists s, t \in K[x]$ such that

$$s_i f + t_i g = d$$

when $\gcd(f, g) = 1 \implies s_l f + t_l g = 1$

$$\implies t_l g \equiv 1 \pmod{f} \implies t_l = \frac{1}{g} \pmod{f}$$

2.1.3 Chinese Remainder Theorem

Theorem 2.1.2. *[?] Let I_1, I_2, \dots, I_k be ideals in a ring R . The map*

$$\begin{aligned} \phi: R &\longrightarrow R/I_1 \times R/I_2 \times \cdots \times R/I_k \\ r &\longmapsto (r + I_1, r + I_2, \dots, r + I_k) \end{aligned}$$

is a ring homomorphism with kernel $I_1 \cap I_2 \cap \cdots \cap I_k$. If for each $i, j \in \{1, 2, \dots, k\}$ with $i \neq j$, the ideals I_i, I_j are called comaximal and the map ϕ is surjective. Moreover,

$$I_1 \cap I_2 \cap \cdots \cap I_k = I_1 I_2 \dots I_k$$

$$\implies R/(I_1 I_2 \dots I_k) = R/(I_1 \cap I_2 \cap \cdots \cap I_k) \cong R/I_1 \times R/I_2 \times \cdots \times R/I_k.$$

2.1.4 Chinese remainder theorem and evaluating black box for polynomials

We know the following result from algebra that,

Theorem 2.1.3. *Let $K[x]$ be a polynomial ring over a field K and let $I \subset K[x]$ be an ideal of the ring then $K[x]/I$ is a field if and only if I is a maximal ideal.*

Given a black box B for a polynomial $f(x) \in K[x]$ and $x = \alpha \in K$ we can evaluate $f(\alpha)$ by querying the black box B at α .

$$\begin{aligned} \phi: K[x] &\longrightarrow K[x]/(x - \alpha) \\ f(x) &\longmapsto f(\alpha) \end{aligned}$$

Thus, evaluating the polynomial at n points is querying the black box at n points $\alpha_1, \dots, \alpha_n$ is the following homomorphism.

$$\begin{aligned}\phi: K[x] &\longrightarrow K[x]/(x - \alpha_1) \times K[x]/(x - \alpha_2) \times \dots \times K[x]/(x - \alpha_n) \\ f(x) &\longmapsto (f(\alpha_1), f(\alpha_2), \dots, f(\alpha_n))\end{aligned}$$

From the Chinese remainder theorem, we know that

$$K[x]/(x - \alpha_1) \times K[x]/(x - \alpha_2) \times \dots \times K[x]/(x - \alpha_n) \cong K[x]/((x - \alpha_1)(x - \alpha_2) \dots (x - \alpha_n))$$

$$\text{Let } [\alpha_1, \dots, \alpha_n] \in K^n$$

such that

$$\alpha_i \neq \alpha_j \forall i \neq j \text{ and } n > 0.$$

$$\text{Let } F(x) = \frac{f(x)}{g(x)} \in K(x), \text{ where } f(x), g(x) \in K[x],$$

$$\text{Let } y_i = F(\alpha_i) = \frac{f(\alpha_i)}{g(\alpha_i)}, \forall 1 \leq i \leq n, \text{ such that } g(\alpha_i) \neq 0.$$

$$\exists! u(x) \in K[x] \text{ of degree } < n \text{ such that } u(\alpha_i) = y_i$$

$$\implies u(x) \equiv y_i \pmod{(x - \alpha_i)}$$

$$\implies F(x) = \frac{f(x)}{g(x)} \equiv u(x) \pmod{(x - \alpha_i)}$$

$$\text{Let } \overline{m}(x) = \prod_{i=1}^n (x - \alpha_i)$$

2.1.5 Main idea

1. We are given the black box of rational polynomial $F(x) \in K(x)$. We have the option to choose n distinct points $\alpha_1, \dots, \alpha_n$, such that $n > \text{degree}_{\text{numerator}} + \text{degree}_{\text{denominator}}$ and evaluate the black box at these points to get y_1, \dots, y_n .
2. Construct two polynomials m, u out of these two sets of points.
3. $m(x) = \prod_{i=1}^n (x - \alpha_i)$ is the Chinese remainder theorem polynomial.
4. $u(x) = \text{Interpolation}(\alpha_i, y_i)$. Note that the degree of $u(x) = n - 1$ and the degree of $m(x) = n$.
5. We take these two polynomials m, u and apply the extended euclidean algorithm to get $f(x), g(x) \implies f, g$ appear as remainder and coefficient in the division algorithm.

6. We can think of any rational function $\frac{f}{g}$ as members of an equivalence class. (localization?) with other elements.
7. What MQRFR says is that the $f = r_i$ and $g = t_i$ for the i^{th} iteration of the extended euclidean algorithm such that $\deg(q_i)$ is max.
8. There does seem to be some relationship between Univariate rational functions, the interpolated polynomial u the product polynomial m and the extended euclidean algorithm(m, u).
9. CRT from 265 Dummit and Foote.
- 10.

Chapter 3

Ben-Or and Tiwari's Multivariate Polynomial Interpolation

3.1 Introduction

Ben-Or Tiwari is an multivariate interpolation algorithm that interpolates all the variables of the polynomial simultaneously as opposed to Zippel's multivariate interpolation algorithm which interpolates the variables one by one.

Algorithm 2 Ben-Or & Tiwari Multivariate Interpolation

Require: Black-box function f (evaluates the target polynomial), number of variables n , degree bound D , finite field \mathbb{Z}_p

Ensure: Sparse polynomial $P(x_1, x_2, \dots, x_n)$ as a list of [coefficient, exponents] pairs.

```
1:  $T_{\max} \leftarrow (D + 1)^n$  ▷ Upper bound on number of terms
2:  $m \leftarrow 2 \cdot T_{\max}$  ▷ Number of evaluation points
3: Step 2: Generate evaluation points
4:  $evaluations \leftarrow$  empty list
5: for  $k \leftarrow 1$  to  $m$  do
6:    $x_k \leftarrow [\alpha^{k^i} \text{ for } i = 1 \dots n]$ 
7:    $v_k \leftarrow f(x_k)$ 
8:   Append  $v_k$  to  $evaluations$ 
9: end for
10: Step 3: Construct minimum characteristic polynomial using Berlekamp-Massey algorithm
11:  $\Lambda \leftarrow \text{Berlekamp\_Massey}(evaluations, p)$ 
12: Step 5: Factor  $\Lambda(z)$  to find roots
13:  $roots \leftarrow \text{FIND\_ROOTS}(\Lambda)$ 
14: Step 6: Recover exponents from roots
15:  $terms \leftarrow$  empty list
16: for each root in  $roots$  do
17:    $exponents \leftarrow \text{TRIAL\_DIVISION}(root, n)$ 
18:   Append  $exponents$  to  $terms$ 
19: end for
20: Step 7: Recover coefficients via linear system
21:  $A \leftarrow \text{CONSTRUCT\_VANDERMONDE\_MATRIX}(terms, evaluations)$ 
22:  $b \leftarrow evaluations[1 \dots T_{\max}]$ 
23:  $coefficients \leftarrow \text{SOLVE\_LINEAR\_SYSTEM}(A, b)$ 
24: Step 8: Validate with additional evaluations
25:  $validated \leftarrow \text{TRUE}$ 
26: for  $k \leftarrow m + 1$  to  $m + 10$  do
27:    $x_k \leftarrow [f(\alpha)^{k^i} \text{ for } i = 1 \dots n]$ 
28:   if  $f(x_k) \neq \text{EVALUATE\_POLYNOMIAL}(terms, coefficients, x_k)$  then
29:      $validated \leftarrow \text{FALSE}$ 
30:     break
31:   end if
32: end for
33: if not  $validated$  then
34:   Increase  $T_{\max}$  and restart
35: end if
36: return  $\text{SPARSE\_POLYNOMIAL}(terms, coefficients)$ 
```

Chapter 4

Multivariate Rational function interpolation

4.1 Introduction

Algorithm 3 Multivariate rational function interpolation

Require: Black box for rational function $\frac{ff(x_1, x_2, \dots, x_n)}{gg(x_1, x_2, \dots, x_n)}, p$, where $ff, gg \in K(x_1, x_2, \dots, x_n)$.

```
1: while true do
2:    $num \leftarrow []$ 
3:    $den \leftarrow []$ 
4:    $T \leftarrow 4$ 
5:   for  $i \leftarrow 0$  to  $T$  do
6:      $\Sigma \leftarrow [[2^i, 3^i, \dots, \Psi^i]]$ , where  $\sigma_i \leftarrow [2^i, 3^i, \dots, \Psi^i] \in \mathbb{Z}_p^n$ 
7:     while true do
8:        $t \leftarrow T$ 
9:       Pick random vector  $\alpha_i = [\alpha_{i1}, \dots, \alpha_{it}] \in \mathbb{Z}_p^t$ 
10:      Pick random vector  $\beta_i = [\beta_{i1}, \dots, \beta_{i(n-1)}] \in \mathbb{Z}_p^{n-1}$ 
11:       $m_i(x) = \prod_{k=1}^t (x - \alpha_{ik})$ , where  $m_i(x) \in \mathbb{Z}_p[x]$ .
12:       $[[\alpha_{ik}, \phi(\alpha_{ik})]]$  where  $[\alpha_{ik}, \phi(\alpha_{ik})] \in \mathbb{Z}_p^n$  and  $\phi(x) \leftarrow \beta_{ij}(x - \sigma_{i1}) + \sigma_{i(j+1)}, \forall 1 \leq j \leq n-1$ 
13:       $Y_i \leftarrow [y_i, \dots, y_{it}]$ , where  $y_{ik} \leftarrow B(\alpha_{ik}, \phi(\alpha_{ik}), p) \forall 1 \leq k \leq t$ 
14:       $u_i(x) \leftarrow \text{Interpolate}(\alpha_i, Y_i, x) \bmod p$ 
15:       $f_i(x), g_i(x), deg\_q_i \leftarrow \text{MQRR}(m_i, u_i) \bmod p$ 
16:      if  $deg\_q_i > 1$  then
17:         $num.insert(f_i(\sigma_{i1})) \bmod p$ 
18:         $den.insert(g_i(\sigma_{i1})) \bmod p$ 
19:        break
20:      else
21:         $t \leftarrow 2t$ 
22:      end if
23:    end while
24:  end for
25:  Construct minimum characteristic polynomial using Berlekamp-Massey algorithm
26:   $\Lambda_n \leftarrow \text{Berlekamp\_Massey}(num, p)$ 
27:   $\Lambda_d \leftarrow \text{Berlekamp\_Massey}(den, p)$ 
28:  Find number of terms in denominator and numerator
29:   $terms_n \leftarrow \text{degree}(\Lambda_n)$ 
30:   $terms_d \leftarrow \text{degree}(\Lambda_d)$ 
31:  Factor  $\Lambda_n(z)$  and  $\Lambda_d(z)$  to find roots
32:   $roots_n \leftarrow \text{ROOTS}(\Lambda_n)$ 
33:   $roots_d \leftarrow \text{ROOTS}(\Lambda_d)$ 
34:  Check if number of terms and roots are equal
35:  if  $terms_n \neq roots_n$  or  $terms_d \neq roots_d$  then
36:     $T \leftarrow 2T$ 
37:  else
38:    break
39:  end if
40: end while
41: Recover monomials from roots using trial division
42: Recover coefficients via Zippel Vandermonde solver
43:  $coeff_n \leftarrow \text{Zippel\_Vandermonde\_solver}(num, terms_n, roots_n, \Lambda_n, p)$ 
44:  $coeff_d \leftarrow \text{Zippel\_Vandermonde\_solver}(den, terms_d, roots_d, \Lambda_d, p)$ 
```

Bibliography

- [1] D.S. Dummit and R.M. Foote. *Abstract Algebra*. Wiley, 2003.
- [2] Frank Mittelbach and Michel Goossens. *The L^AT_EX Companion*. Addison-Wesley, 2004.

Appendix A

Code

Appendices should be used for supplemental information that does not form part of the main research. Remember that figures and tables in appendices should not be listed in the List of Figures or List of Tables.