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# Working Draft, C++ Extensions for Concepts

Note: this is an early draft. It's known to be incomplet and incorrekt, and it has lots of bad formatting.

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1 General [intro]

1.1 Scope [general.scope]

This technical specification describes extensions to the C++ Programming language (1.2) that enable the specification and checking of constraints on template arguments, and the ability to overload functions and specialize templates based on those constraints. These extensions include new syntactic forms and modifications to existing language semantics.

International Standard, ISO/IEC 14882, provides important context and specification for this Technical Specification. This document as written as a set of changes against that specification. Instructions to modify or add paragraphs are written as explicit instructions. Modifications made directly to existing text from the International Standard use <u>underlining</u> to represent added text and <u>strikethrough</u> to represent deleted text.

# 1.2 Normative references

[intro.refs]

- The following referenced document is indispensable for the application of this document. For dated references, only the edition cited applies. For undated references, the latest edition of the referenced document (including any amendments) applies.
  - ISO/IEC 1482:2014, Programming Languages - C++

**Editor's note:** The TS will formally refer to the ISO/IEC document defining the C++14 programming language. Until that document is published, the paper targets the current working draft NXXX

<sup>2</sup> ISO/IEC 1482:2014 is herein after called the C++ Standard. References to clauses within the C++ Standard are written as "C++ §3.2".

#### 1.3 Terms and definitions

[intro.defns]

<sup>1</sup> For the purposes of this document, the terms and definitions given in the C++ Standard and the following apply.

# 1.4 Implementation compliance

[intro.compliance]

Conformance requirements for this specification are the same as those defined in C++ §1.4. [ *Note:* Conformance is defined in terms of the behavior of programs. — *end note* ]

# 1.5 Acknowledgments

[intro.ack]

- <sup>1</sup> The design of this specification is based, in part, on a concept specification of the algorithms part of the C++ standard library, known as `The Palo Alto" TR (WG21 N3351), which was developed by a large group of experts as a test of the expressive power of the idea of concepts. Despite syntactic differences between the notation of the Palo Alto TR and this TS, the TR can be seen as a large-scale test of the expressiveness of this TS.
- <sup>2</sup> This work was funded by NSF grant ACI-1148461.

§ 1.5

# 2 Lexical conventions

[lex]

2.1 Keywords [lex.key]

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 $<sup>^1\,</sup>$  In C++ §2.12, Table 4, add the keywords concept and requires.

# 3 Expressions

[expr]

# 3.1 Primary expressions

[expr.prim]

<sup>1</sup> In C++ §5.1.1, add requires-expression to the rule, primary-expression.

primary-expression: requires-expression

### 3.1.1 Lambda expressions

[expr.prim.lambda]

- <sup>1</sup> Modify C++ §5.1.2/5.
- The closure type for a non-generic lambda-expression has a public inline function call operator (C++ §13.4) whose parameters and return type are described by the lambda-expression's parameter-declaration-clause and trailing-return-type respectively. For a generic lambda, the closure type has a public inline function call operator member template (C++ §14.5.2) whose template-parameter-list consists of one invented type template-parameter for each occurrence of auto or each unique occurrence of a constrained-type-name in the lambda's parameter-declaration-clause, in order of appearance. The invented type template-parameter is a parameter pack if the corresponding parameter-declaration declares a function parameter pack (C++ §8.3.5). If the decl-specifier-seq of the corresponding parameter-declaration includes a constrained-type-specifier, the invented type parameter is a constrained-parameter, whose constrained-type-specifier matches that of the parameter-declaration. (7.1). The return type and function parameters of the function call operator template are derived from the lambda-expression's trailing-return-type and parameter-declaration-clause by replacing each occurrence of auto in the decl-specifiers of the parameter-declaration-clause with the name of the corresponding invented template-parameter. [Example: Consider the following:

```
template<typename T> concept bool C = ...;
auto fun = [](const C& arg) { }
```

The parameter arg has a *constrained-type-specifier* C, signifying that it is a generic lambda whose closure type is like the following function object:

```
struct Fun {
  template<C __T>
   auto operator()(const __T& x) const { }
};
```

- end example ]
- <sup>3</sup> Also insert the following paragraph after C++ §5.1.2/5.
- <sup>4</sup> All placeholder types introduced using the same *concept-name* have the same invented template parameter. [ *Example:*

```
auto f = [](C a, C b) \{ \};

f(0, 0); // 0k

f(0, 'a'); // Error: template argument deduction failure
```

The second call to f results in a compiler error because the types of the deduced arguments cannot be unified. — end example ]

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### 3.1.2 Requires expressions

[expr.req]

A *requires-expression* provides a concise way to express syntactic requirements on template arguments.

```
requires-expression:
   requires requirement-parameter-list requirement-body
requirement-parameter-list:
   ( parameter-declaration-clause<sub>opt</sub> )
reauirement-body:
   { requirement-list }
requirement-list:
   requirement
   requirement-list requirement
requirement:
   simple-requirement
   compound-requirement
   type-requirement
   nested-requirement
simple-requirement:
   expression:
compound-requirement:
   constexpr<sub>opt</sub> { expression } noexcept<sub>opt</sub> trailing-return-type<sub>opt</sub> ;
type-requirement:
   typename-specifier;
nested-requirement:
   requires-clause;
```

- <sup>2</sup> A requires-expression has type bool.
- <sup>3</sup> A requires-expression shall not appear outside of a concept definition (4.1.4) or a requires-
- <sup>4</sup> [ *Example*: The most common use of *requires-expressions* is to define syntactic requirements in concepts (4.1.4) such as the one below:

```
template<typename T>
  concept bool R() {
    return requires (T i) {
      typename A<T>;
      {*i} -> const A<T>&;
    };
}
```

The concept is defined in terms of the syntactic and type requirements within the *requires-expression*. A *requires-expression* can also be used in a *requires-clause* templates as a way of writing ad hoc constraints on template arguments such as the one below:

```
template<typename T>
  requires requires (T x) { x + x; }
    T add(T a, T b) { return a + b; }
— end example ]
```

The requires-expression may introduce local arguments via a parameter-declaration-clause. These parameters have no linkage, storage, or lifetime. They are used only to write constraints within the requirement-body and are not visible outside the closing } of the requirement-body. The requirement-parameter-list shall not include an ellipsis.

§ 3.1.2

The *requirement-body* is a sequence of *requirements* separated by semicolons. These *requirements* may refer to local arguments, template parameters, and any other declarations visible from the enclosing context. Each *requirement* introduces a conjunction of one or more atomic constraints (7.6). The kinds of atomic constraints introduced by a *requirement* are:

- A valid expression constraint is a predicate on an expression. The constraint is satisfied
  if and only if the substitution of template arguments into that expression does not result
  in substitution failure. The result of successfully substituting template arguments into
  the dependent expression produces a valid expression.
- A valid type constraint is a predicate on a type. The constraint is satisfied if and only if
  the substitution of template arguments into that type does not result in substitution
  failure. The result of successfully substituting template arguments into the dependent
  type produces an associated type.
- A result type constraint is a predicate on the result type of a valid expression. Let E be a valid expression and X be a trailing-return-type. The constraint is satisfied if and only if E can be used as an argument to an invented function f, which has a single function parameter of type X and returning void. That is, the function call f(E) must be a valid expression. [ Note: Each template parameter referred to by X is a template parameter of the invented function f. If X contains a constrained-type-specifier or auto specifier, then f is a generic function (5.1.1). end note ]
- A constant expression constraint is satisfied if and only if a valid expression E is a constant expression (C++ §5.19).
- An *exception constraint* is satisfied if and only if, for a valid expression E, the expression noexcept(E) evaluates to true (C++ §5.3.7).
- <sup>7</sup> A requires-expression evaluates to true if and only the atomic constraints introduced by each requirement in the requirement-list are satisfied and false otherwise. The semantics of each kind of requirement are described in the following sections.

# 3.1.2.1 Simple requirements

[expr.req.simple]

A *simple-requirement* introduces a valid expression constraint for its *expression*. The expression is an unevaluated operand (C++ §3.2). [ *Example:* The following is requirement evaluates to true for all arithmetic types (C++ §3.9.1), and false for pointer types (C++ §3.9.2).

```
requires (T a, T b) {
  a + b; // A simple requirement
}
```

— end example 1

<sup>2</sup> If the expression would always result in a substitution failure, the program is ill-formed. [ *Example:* 

```
requires () {
  new T[-1]; // error: the valid expression well never be well-formed.
}
— end example ]
```

# 3.1.2.2 Type requirements

[expr.req.type]

A *type-requirement* introduces valid type constraint for its *typename-specifier*. [ *Note:* A type requirement requests the validity of an associated type, either as a nested type name, a class

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template specialization, or an alias template. It is not used to specify requirements for arbitrary type-specifiers. — end note ] [ Example:

<sup>2</sup> If the required type will always results in a substitution failure, then the program is ill-formed. [*Example:* 

```
requires () {
  typename int::X; // error: int does not have class type
  typename T[-1]; // error: array types cannot have negative extent
}
— end example ]
```

# 3.1.2.3 Nested requirements

[expr.req.nested]

A nested-requirement introduces an additional constraint expression 7.6 to be evaluated as part of the satisfaction of the requires-expression. The requirement is satisfied if and only if the constraint evaluates to value true. [Example: Nested requirements are generally used to provide additional constraints on associated types within a requires-expression.

```
requires () {
  typename X;
  requires C<X<T>>();
}
```

These requirements are satisfied only when substitution into X<T> is successful and when C<X<T>>() evaluates to true. — end example ]

#### 3.1.2.4 Compound requirements

[expr.reg.compound]

- <sup>1</sup> A *compound-requirement* introduces a conjunction of one or more constraints pertaining to its *expression*, depending on the syntax used. This set includes:
  - a valid expression constraint,
  - an optional associated type constraint
  - an optional result type constraint,
  - an optional constant expression constraint, and
  - an optional an exception constraint.

A *compound-requirement* is satisfied if and only if every constraint in the set is satisfied. The required valid expression is an unevaluated operand  $(C++\S3.2)$  except in the case when the constexpr specifier is present. These other requirements are described in the following paragraphs.

- <sup>2</sup> The brace-enclosed *expression* in a *compound-requirement* introduces a valid expression constraint. Let E be the valid expression resulting from successful substitution.
- <sup>3</sup> The presence of a trailing-return-type introduces a result type constraint on E.
- $^4$  If the constexpr specifier is present then a constant expression constraint is introduced for the valid expression E.

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<sup>5</sup> If the noexcept specifier is present, then an exception constraint is introduced for the valid expression E.

<sup>6</sup> [ Example:

```
template<typename I>
  concept bool Inscrutible() { ... }

requires(T x) {
  {x++}; #1
  {*x} -> typename T::r; #2
  {f(x)} -> const Inscrutible& #3
  {g(x)} noexcept -> auto& #4
  constexpr {T::value}; #5
  constexpr {T() + T()} -> T #6;
}
```

Each of these requirements introduces a valid expression constraint on the expression in its enclosing braces. Requirement #1 introduces no additional constraints. It is equivalent to a *simple-requirement* containing the same expression. Requirement #2 \*x introduces a result type constraint though its *trailing-return-type*, typename T::r. The required valid expression \*x must be usable as an argument to the invented function:

```
template<class T>
  void z1(typename T::r);
```

Requirement #3 also introduces a result type constraint on its required valid expression f(x). This expression must be usable as an argument to the invented generic function:

```
void z2(const Instrutible&)
```

Requirement #4 introduces a result type constraint and an exception constraint. The required valid expression g(x) must be usable as as an argument to the invented generic function:

```
void z3(auto&);
```

Additionally, g(x) must not propagate exceptions. Requirement #5 introduces a constant expression constraint: T::value must be a constant expression. The requirement in #6 introduces a result type constraint and a constant expression constraint. The required valid expression T() + T() must be usable as an argument to the invented function:

```
template<class T>
  void z4(T):
```

The valid expression must also be a constant expression. —  $end\ example\ ]$ 

§ 3.1.2.4

# 4 Declarations

[dcl.dcl]

4.1 Specifiers [dcl.spec]

1 Extend the decl-specifier production to include the concept specifier.

decl-specifier:

concept

# 4.1.1 Simple type specifiers

[dcl.type.simple]

#### 4.1.2 auto specifier

[dcl.spec.auto]

1 Modify C++ §7.1.6.4/1 so that the last sentence reads:

concept-name < template-argument-list >

The auto *type-specifier* is also used to signify that a lambda is a generic lambda or that a function is a generic function.

#### 4.1.3 Constrained type specifiers

[dcl.spec.constr]

- <sup>1</sup> A constrained-type-specifier designates a placeholder type that will be replaced later by deduction from a required valid expression in a compound-requirement. A constrained-type-specifier is also used to signify that a lambda is a generic lambda or that a function is a generic function.
- A constrained-type-specifier can appear in the trailing-return-type of a compound-requirement or in any context in which the auto type-specifier appears, except:
  - in the decl-specifier-seq of a variable declaration,
  - in the return type of a function declaration,
  - in the decltype(auto) type-specifier, or
  - a conversion-function-id.
- <sup>3</sup> If the *constrained-type-name* appears as one of the *decl-specifiers* of a *parameter-declaration* in a *template-parameter-list*, then the declared parameter is a *constrained-parameter*, and its meaning is defined in section 7.1. Otherwise, the meaning of *constrained-type-specifiers* is defined in this section. [ *Note:* A constrained template parameter can introduce type parameters as well as designate the type of a non-type template parameter. The meaning of those declarations are specified separately. *end note* ]

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<sup>4</sup> If the *constrained-type-specifier* appears as one of the *decl-specifiers* of a *parameter-declaration* in either a *lambda-expression* or function declaration then the lambda is a generic lambda 3.1.1 and the function is a generic function 5.1.1.

- <sup>5</sup> A *constrained-type-specifier* designates a placeholder type that will be replaced later, and it introduces an associated constraint on deduced type, called the *constrained type* within the enclosing declaration or *requires-expression*.
- <sup>6</sup> If the *constrained-type-specifier* appears in the *trailing-return-type* of a *compound-requirement*, then the constrained type is deduced from the required valid expression. Otherwise, the constrained type is deduced using the rules for deducing auto (4.1.2).
- <sup>7</sup> The *introduced constraint* is a constraint expression (7.6) synthesized from the *concept-name* or *partial-concept-id* in the *constrained-type-name*.
- When an identifier is a *concept-name*, it refers to one or more function concepts or a single variable concept. At least one concept referred to by the *constrained-type-name* shall be a type concept (4.1.4). [ *Example:* Function concepts can be overloaded to accept different numbers and kinds of template arguments. This is sometimes done to generalize a single concept for different kinds of arguments.

```
template<typename T>
  concept bool C() { ... }
template<typename T, typename U>
  concept bool C() { ... }
```

- end example ] The concept-name c refers to both concept definitions.
- <sup>9</sup> A partial-concept-id is a concept-name followed by a sequence of template arguments. A partial-concept-id does not refer to template specialization; the template argument list must be adjusted by adding a template argument before the first of the initial template arguments before the name refers to a template specialization. [ Example:

```
template<typename T, typename U>
  concept bool C = ...;

C<int> // A partial-concept-id
C<char, int> // A template-id
```

The first name is a *partial-concept-id* and can be used as part of constrained type name as part the type specifier of a parameter declaration or a template parameter. The second name is a *template-id* and determines whether the concept is satisfied for the given arguments. — *end example* ]

- <sup>10</sup> A partial-concept-id shall not have an empty list of template arguments.
- 11 An introduced constraint is formed by applying the following rules to each concept referenced by the *concept-name* in the *constrained-type-name*. Let c be a concept referred to by the *concept-name*. T be the constrained type, and Args be a sequence of template arguments. If the *constrained-type-name* is a *partial-concept-id*, then Args is its *template-argument-list*, otherwise Args is an empty sequence. The *candidate constraint* is a *template-id* having the form C<T, Args>. [ Note: If Args is empty, the resulting *template-id* is of the form C<T>. *end note* ] If C<T, Args> does not refer to a template specialization, the candidate constraint is rejected. [ Note: The expression C<T, Args> may not refer to a valid template specialization if Args contains too many or to few template arguments for C, or if Args do not match C's template parameters. *end note* ]
- <sup>12</sup> If, after constructing candidate constraints for each concept named by the *concept-name*, there are no candidates or more than one candidate, the program is ill-formed.

§ 4.1.3

The introduced constraint is constructed from the remaining candidate. If C is a function concept, then the resulting constraint is a function call of the form C<T, Args>(). Otherwise, the introduced constraint is the same as the remaining candidate.

<sup>14</sup> [ *Example:* The following unary and binary concepts are defined as variables and functions.

```
template<typename T>
concept bool V1 = ...;

template<typename T, typename U>
concept bool V2 = ...;

template<typename T>
concept bool F1() { return ...; }

template<typename T, typename T2>
concept bool F2() { return ...; }
```

Suppose x is a template parameter being declared, either explicitly or as an invented template parameter of a *parameter-declaration* in a generic function or generic lambda. The synthesized constraints corresponding to each declaration are:

```
V1 X // becomes V1<T>
V2<Y> X // becomes V2<X, Y>
F1 X // becomes F1<X>()
F2<Y> X // becomes F2<X, Y>()
```

— end example ]

The meaning of the introduced constraint depends on the context in which the *constrained-type-specifier* appears. If it appears in the *decl-specifiers* of a *parameter-declaration* of a generic lambda (3.1.1) or generic function (5.1.1), the the introduced constraint is associated with the corresponding template declaration (7). If it appears in *trailing-return-type* of a *compound-requirement*, the introduced constraint is evaluated as part of the enclosing *requires-expression* (3.1.2).

# 4.1.4 concept specifier

[dcl.concept]

- <sup>1</sup> The concept specifier shall be applied to only the definition of a function template or variable template. A function template definition having the concept specifier is called a *function concept*. A variable template definition having the concept specifier is called a *variable concept*. A *concept definition* refers to either a function concept and its definition or a variable concept and its initializer.
- <sup>2</sup> A *type concept* is a concept whose first template parameter is a *type-parameter*, but not a template template parameter. Otherwise, the concept is a *non-type concept*. A *variadic concept* is a concept whose first template parameter is a template parameter pack.
- <sup>3</sup> Every concept definition is also a constexpr declaration (C++ §7.1.5).
- <sup>4</sup> A function concept has the following restrictions:
  - The template must be unconstrained.
  - The result type must be bool.
  - The declaration shall have a *parameter-declaration-clause* equivalent to ().
  - The declaration shall be a definition.
  - The function shall not be recursive.

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— The function body shall consist of a single return statement whose expression shall be a *constraint-expression*.

# [ Example:

```
template<typename T>
    concept bool C1() { return true; } // OK

template<typename T>
    concept int C2() { return 0; } // error: must return bool

template<typename T>
    concept bool C3(T) { return true; } // error: must have no parameters

concept bool p = 0; // error: not a template

— end example ]
```

- <sup>5</sup> A variable template has the following restrictions:
  - The template must be unconstrained.
  - The declared type must be bool.
  - The declaration must have an initializer.
  - The initializer shall be a *constraint-expression*.

### [ Example:

```
template<typename T>
    concept bool D1 = has_x<T>::value; // OK

template<typename T>
    concept bool D2 = 3 + 4; // Error: initializer is not a constraint

template<Integral T>
    concept bool D3 = has_x<T>::value; // Error: constrained concept definition

— end example ]
```

<sup>6</sup> A program that declares an explicit or partial specialization of a concept definition is ill-formed. [ *Example:* 

```
template<typename T>
    concept bool C = is_iterator<T>::value;

template<typename T>
    concept bool C<T*> = true; // Error: partial specialization of a concept

— end example ]
```

<sup>7</sup> [ *Note:* The prohibitions against overloading and specialization prevent users from subverting the constraint system by providing a meaning for a concept that differs from the one computed by evaluating its constraints. — *end note* ]

 $\S~4.1.4~$ 

# 5 Declarators [dcl.decl]

- <sup>1</sup> Modify C++ §8/1 as follows:
- <sup>2</sup> A declarator declares a single variable, function, or type, within a declaration. The *init-declarator-list* appearing in a declaration is a comma-separated sequence of declarators, each of which can have an initializer have constraints, an initializer, or both.

init-declarator:

declarator requires-clause<sub>opt</sub> initializer<sub>opt</sub>

- 3 Insert the following paragraph after C++ §8/1
- <sup>4</sup> A *requires-clause* (7) shall only be present if the *declarator* declares a generic function (5.1.1). [ *Example*:

```
template concept bool C = \ldots;

void f1(auto x) requires C < \text{decltype}(x) > ; // 0k

void f2(int x) requires C < \text{int} > ; // Error: f2 is not a generic function auto n requires C < \text{decltype}(n) > = g(); // Error: cannot constrain variable declaration struct S \  requires C < S > ; // Error: cannot constrain a class declaration
```

— end example ]

# 5.1 Meaning of declarators

[dcl.meaning]

5.1.1 Functions [dcl.fct]

- <sup>1</sup> Add the following paragraphs after C++ §8.3.5/14.
- <sup>2</sup> A generic function is a function template whose *template-parameter-list* has a *parameter-declaration* whose *type-specifier* is either auto or a *constrained-type-name*. [ *Example:*

```
auto f(auto x); // Ok
void sort(C& c); // Ok (assuming C names a concept)
— end example ]
```

The declaration of a generic function has a *template-parameter-list* that consists of one invented type *template-parameter* for each occurrence of auto or each unique occurrence of a *constrained-type-name* in the function's *parameter-declaration-clause*, in order of appearance. The invented type of *template-parameter* is a parameter pack if the corresponding *parameter-declaration* declares a function parameter pack (C++ §8.3.5). If the *decl-specifier-seq* of the corresponding *parameter-declaration* includes a *constrained-type-specifier*, the invented type parameter is a *constrained-parameter*, whose *constrained-type-specifier* matches that of the *parameter-declaration*. (7.1). [ *Example:* The following generic function declarations are equivalent:

```
template<typenaem T>
  conxtexpr bool C() { ... }
auto f(auto x, const C& y);
```

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```
template<typename T1, C T2>
  auto f(T1 x, const T2& y);
```

The type of y is a type parameter constrained by  $c. - end \ example$ 

<sup>4</sup> All placeholder types introduced using the same *constrained-type-name* have the same invented template parameter. [ *Example:* The following generic function declarations are equivalent:

```
auto g(C a, C* b);
template<C T>
    auto g(T a, T* b);
— end example ]
```

<sup>5</sup> If an entity is declared by an abbreviated template declaration, then all its declarations must have the same form.

## 5.2 Function definitions

[dcl.fct.def]

# 5.2.1 In general

[dcl.fct.def.general]

1 Modify the function-definition syntax in C++ §8.4.1 to include a requires-clause.

function-definition:

attribute-specifier-seq $_{opt}$  decl-specifier-seq $_{opt}$  declarator virt-specifier-seq $_{opt}$  requires-clause $_{opt}$  function-body

- <sup>2</sup> Add the following paragraph at the end of C++ §8.4.1.
- A requires-clause (7) shall only be present if the declarator declares a generic function (5.1.1) or a member function definition (6.1). [ Note: Constraints for a function template or member function template are written after the template-parameter-list. end note ] [ Example:

```
template<typename T> concept bool C1 = ...;
template<typename T> concept bool C2 = ...;
template<typename T, typename U> concept bool D = ...;

void f(C1 a, C2 b) requires D<decltype(a), decltype(b)> { } // Ok

template<typename T>
    void f(const T& x) requires C<T>; // Error: f is declared as a template

template<typename T>
    struct S1 {
        S1(T&) requires C1<T> { } // Defines a constrained constructor
        void f() requires C2<T> { } // Defines a constrained member function
    };

struct S2 {
    void g(auto x) requires D<decltype(x)> { } // Ok
    };

— end example ]
```

<sup>4</sup> A function-definition shall not declare a destructor (C++ §12.4) with a requires-clause.

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6 Classes [class]

# 6.1 Class members [class.mem]

<sup>1</sup> In C++ §9.2, modify the *member-declarator* syntax.

member-declarator:

declarator virt-specifier-seqopt pure-specifier-seqopt requires-clauseopt

- <sup>2</sup> Insert the following paragraph after C++ §9.2
- <sup>3</sup> A *requires-clause* (7) shall only be present if the *declarator* declares a constrained member function of a class template (7.4.1.1) or a generic function (5.1.1). [ *Example:*

- end example ] [ Note: A constrained generic function declared at class scope is a member function template. - end note ]

<sup>4</sup> A destructor (C++ §12.4) shall not be declared with *requires-clauses*.

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7 Templates [temp]

Modify the template-declaration grammar in C++ §14/1.

template-declaration:

```
template < template-parameter-list > requires-clause<sub>opt</sub> declaration
nested-name-specifier<sub>opt</sub> concept-name { introduction-list } declaration
requires-clause:
    requires constraint-expression
introduction-list:
    introduced-parameter
    introduction-list, introduced-parameter
introduced-parameter:
...opt identifier
```

Add the following paragraphs after C++ §14/6.

- A template-declaration is written in terms of its template parameters. These parameters are declared explicitly in a template-parameter-list (7.1), or they are introduced by a concept introduction, a concept-name and following introduction-list.
- The concept designated by the *concept-name* is determined by the *introduction-list*. Let C be a *concept-name* and I1, I2, ..., In be a sequence of *identifiers* in the *introduced-parameters* of an *introduction-list*. If the *template-id*, C<I1, I2, ..., In>, refers to a single concept declaration, then that concept is the one designated by C. Otherwise, the program is ill-formed. [ *Example:*

```
template<typename T> concept bool Eq() { return true; } // #1 template<typename T, typename U> concept bool Eq() { return true; } // #2 Eq{T} void f1(T, T); // OK: Eq{T} designates #1 Eq{A, B} void f2(A, B); // void f2(A, B); // void f2(A, B); // void f2(A, B) void f2(A, B) void f2(A, B); // void f2(A, B) void f2(A, B); // void f3(A, B) void f2(A, B) void f2(A, B) void f3(A, B) void f4(A, B) void f
```

It is possible to overload function concepts in such a way that a *concept-name* can designate multiple concepts.

```
template<typename T> concept bool C() { return true; }
template<int N> concept bool C() { return true; }

C{X} void f(); // error: resolution of the concept C is ambiguous

— end example ]
```

- <sup>3</sup> Each *identifier*, I, in the *introduced-parameters* of the *introduction-list* is declared to be a template parameter that matches the corresponding template parameter, P, in the *template-parameter-list* of the concept designated by the *concept-name*.
  - If P is a template type-parameter declared with either the class or typename keyword, I is declared as a template type-parameter using the same keyword;
  - if P is a template type-parameter that declares a class template, I is declared as a class template with the template parameters of P;
  - if P is a non-type template-parameter, I is declared as a non-type template-parameter having the same type as P;
  - if P is a template parameter pack, the *identifier*, I, shall be preceded by an ellipsis, and
    is declared as a template parameter pack.

An *introduced-parameter* shall not contain an ellipsis if its corresponding template parameter does not declare a template parameter pack. [ *Example*:

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[ Note: A concept referred to be a concept-name may have template parameters with default template arguments. An introduction-list may omit identifiers for a corresponding template parameter if it has a default argument. However, only the introduced-parameters are declared as template parameters. [ Example:

There is no introduced-parameter that corresponds to the template parameter, B, in the Ineffable concept, so f(T) is declared with only one template parameter. — end example ] — end note ]

- <sup>6</sup> The *introduction-list* shall not be empty.
- <sup>7</sup> An introduced template parameter does not have a default template argument, even if its corresponding template parameter does. [ *Example*:

```
template<typename T, int N = -1> concept bool P() { return true; }

P{T, N} struct Array { };

Array<double, 0> s1; // OK
 Array<double> s2; // error: Array takes two template arguments

— end example ]
```

<sup>8</sup> [ *Note:* A constrained member function template of a constrained class template can be defined outside of its class definition by nested introductions. [ *Example:* 

```
template<typename T> concept bool C = true;
template<typename T> concept bool D = true;
C{T} struct X {
   D{U} void f();
};
C{T} D{U} void X<T>::f() { } // OK: definition of f()
```

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```
- end example ] - end note ]
```

<sup>9</sup> A template-declaration declared by a concept introduction can also be an abbreviated function (5.1.1). The invented template parameters introduced by the presence of auto type-specifiers or constrained-type-specifiers in the parameter-declaration-clause are added to the list of template parameters introduced by the the introduction-list. [Example:

```
template<typename T> concept bool C = true;
 template<typename T> concept bool D = true;
 C{T} void f(T, D);
 template<C T, D D> void f(T, D); // OK: redeclaration of f(T, D)
— end example ] [ Example:
 template<typename T> concept bool C = true;
 template<typename T> concept bool D = true;
 C{T} struct X {
   void f(D);
   D{U} void q(U, C);
 };
 C{T} void X{T}::f(D) { } // OK: definition of <math>X{T}::f(D);
                          // f is a function template with one invented
                          // template type-parameter
 C\{T\} D{U} void X<T>::q(U, C) { } // OK: definition of X<T>::q(U, C);
                                   // g is a function template with two template
                                   // type parameters: one introduced (U) and
                                   // one invented
```

— end example ]

The introduction of a sequence of template parameters, T1, T2, ..., Tn, by a concept-name, C, associates a constraint with the template-declaration. That constraint is C<T1, T2, ..., Tn> when C designates a variable concept and C<T1, T2, ..., Tn>() when C designates a function concept. If an introduced-parameter declares a template parameter pack, its corresponding template argument in the associated constraint is a pack expansion (C++ §14.5.3). [Example:

```
template<typename A, typename B, int C> concept bool C = true;
template<typename A, typename... Args> concept bool D = true;
C{X, Y, Z} struct S; // associates C<X, Y, Z> with S
D{P, ...Qs} struct T; // associates D<P, Qs...> with T
— end example ]
```

- <sup>11</sup> A template-declaration's associated constraints are a conjunction of all constraints introduced by
  - a concept introduction,
  - a requires-clause following a template-parameter-list,
  - any constrained template parameters (7.1) in the declaration's template-parameter-list,
  - any constrained-type-specifiers in the decl-specifier-seq of a parameter-declaration in a function declaration or definition (4.1.3),

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— a requires-clause appearing after the declarator of an init-declarator (5), function-definition (5.2.1), or member-declarator (6.1), or

some combination these.

A template-declaration,  $\tau$ , whose constraints are introduced using any combination of these mechanisms is equivalent to another template-declaration,  $\tau$ , whose template parameters are declared explicitly and as unconstrained template parameters, and  $\tau$  has a single requires-clause whose constraint-expression is equivalent to the associated constraints of  $\tau$  (7.4.5.2). [Note: This section describes how constrained template declarations can be equivalently written using alternative syntax in order to generate a canonical spelling of a template's associated constraints. [Example:

```
template<typename T> concept bool C = true;

// all of the following declarations are equivalent:
void g(C);
template<C T> void g(T);
C{T} void g(T);
template<typename T> requires C<T> void g(T);
```

The last declaration includes the canonical spelling of the associated constraints for all declarations of g(T) as the constraint-expression of its requires-clause. — end example ] The paragraphs below define the rules that make these declarations equivalent. — end note ]

When template-declaration is declared by a concept introduction, it is equivalent to a template-declaration whose template-parameter-list is defined according to the rules for introducing template parameters above, and the equivalent declaration has a requires-clause whose constraint-expression is equivalent to constraint associated by the concept introduction. [Example:

```
template<typename T, typename U> concept bool C1 = true;
template<typename T, typename U> concept bool C2() { return true; }
template<typename T, typename U = char> concept bool C3 = true;
template<typename... Ts> concept bool C4 = true;
C1{A, B} struct X;
C2{A, B} struct Y;
C3{P} void f(P);
C4\{...Qs\} void g(Qs\&\&...);
template<typename A, typename B>
  requires C1<A, B> // constraint associated by C1{A, B}
    struct X;
                  // OK: redeclaration of X
template<typename A, typename B>
  requires C2<A, B>() // constraint associated by C2{A, B}
    struct Y:
                      // OK: redeclaration of Y
template<class P>
  requires C3<P> // constraint associated by C3{P}
   void f(P); // OK: redeclaration of f(P)
template<typename... Qs>
                         // constraint associated by C4{...Qs}
  requires C4<0s...>
    void void g(Qs&&...); // OK: redeclaration of g(Qs&&...)
```

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- end example ]
- When a *template-declaration*, T, is explicitly declared with *template-parameter-list* that has constrained template parameters (7.1), it is equivalent to a *template-declaration*, E, with the same template parameters, except that all constrained parameters are replaced by unconstrained parameters matching the corresponding prototype parameter designated by the *constrained-type-specifier* (7.1). The declaration, E, has a *requires-clause* whose *constraint-expression* is the conjunction of the constraints associated by the constrained template parameters in T (7.1). The order in which the introduced constraints are evaluated is the same as the order in which the constrained template parameters are declared. If the original declaration, T, includes a *requires-clause*, its *constraint-expression* is evaluated after the constraints associated by the constrained template parameters in E. [ *Example:*

```
template<typename> concept bool C1 = true;
template<int> concept bool C2 = true;

template<C1 A, C2 B> struct S;
template<C1 A, C2 B> struct R;
template<C1 T> requires C2<sizeof(T)> void f(T);

template<typename X, int Y>
    requires C1<X> && C2<Y>
        struct S; // OK: redeclaration of S

template<typename X, int Y>
    requires C2<Y> && C1<X>
        struct R; // error: redeclaration of R with different constraints

template<typename T>
    requires C1<T> && C2<sizeof(T)>
        void f(T); // OK: redeclaration of f(T)

— end example ]
```

When the declaration is an abbreviated function, it is equivalent to a *template-declaration* whose template parameters are declared according to the rules in 5.1.1. [ *Example:* 

```
template<typename T> concept bool C = true;
template<typename T> concept bool D() { return true; }

void f(C, C, D);

template<C T, D U>
   void f(T, T, U); // OK: redeclaration of f(C, C, D)

template<typename T, typename U>
   requires C<T> && D<U>()
     void f(T, T, U): // OK: also a redeclaration of f(C, C, D)

— end example ]
```

An abbreviated function can also be declared as a *template-declaration*. The constraints associated by *constrained-type-specifiers* in the *parameter-declaration-clause* of the function declaration are evaluated after those introduced by *constrained-type-specifiers* in the *template-parameter-list* and the following *requires-clause*, if present. This is also the case for an abbreviated function that is declared is declared with a concept introduction. [ *Example:* 

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```
template<typename T> concept bool C = true;
template<typename T> concept bool D() { return true; }
template<typename T> concept bool P = true;

template<C T> requires P<T> void g1(T, D);
template<C T> void g2(T, D);

template<typename T, typename U> requires C<T> && P<T> && D<U>() void g1(T, U); // OK: redeclaration of g1(T, D)

template<C T, D U> requires P<T> // associated constraints are C<T> && D<U>() && P<T> void g1(T, U); // error: ill-formed, no diagnostic required;
C{T} void g2(T, D); // OK: redeclaration of g2(T, D)
```

The second declaration of g1(T, U) is ill-formed (no diagnostic required) because it is functionally equivalent to the first declaration, but not equivalent. — end example ]

<sup>17</sup> A trailing requires-clause is a requires-clause that appears after the declarator in an init-declarator (5), function-definition (5.2.1), or member-declarator (6.1). When a constrained function template or member function template declared with a trailing requires-clause is equivalent to a declaration in which the constraint-expression of the trailing requires-clause is evaluated after all other associated constraints. [Example:

```
template<C T> struct S {
  template<D U> void f(U) requires D<T>;
};

template<C T> template<typename U>
    requires D<U> && D<T>
    void S<T>::f(U) { } // OK: definition of S<T>::f(U)

template<C T> template<typename U, typename P>
    void S<T>::f(U) requires D<U> && D<T> { } // error: redefinition of S<T>::f(U)
```

The second definition if S<T>::f(U) is an error because its declaration is equivalent to the first. — end example ]

# 7.1 Template parameters

[temp.param]

Modify the *template-parameter* grammar in C++ §14.1/1 as follow.

```
template-parameter:
    parameter-declaration
    non-type-parameter
non-type-parameter:
    parameter-declarator
    parameter-declarator = initializer-clause
    parameter-declarator = type-id
    parameter-declarator = id-expression
```

Add the following paragraphs after C++ §14.1/15.

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<sup>2</sup> A constrained template parameter is a constrained-parameter whose decl-specifier-seq contains a constrained-type-specifier. A constrained-parameter defines its identifier to be a template parameter that matches in kind the first template parameter, called the *prototype parameter*, of the concept designated by the *constrained-type-specifier*. [ *Example*:

<sup>3</sup> If the prototype parameter is a type parameter (including template template parameters), then the *decl-specifier-seq* of the constrained parameter shall consist of only the *constrained-type-specifier*. [ *Example:* 

```
template<const C1> // Error: declares a const-qualified type parameter
    struct S;
— end example ]
```

<sup>4</sup> The declared *template-parameter* is a template parameter pack if the prototype parameter declares a template parameter pack. In such cases, the *declarator-id* or *abstract-declarator* of the *constrained-parameter* shall also include an ellipsis. [ *Example:* 

```
template<typename... Ts>
  concept bool X = ...;

template<X... Xs> void f(); // Xs is a parameter pack
template<X Xs> void g(); // Error: must X must include ...

— end example ]
```

<sup>5</sup> If the *constrained-parameter* declares a type parameter, then the *constrained-initializer* is parsed as a *type-id*. Otherwise, it is parsed as a *initializer-clause*. [ *Example*:

```
template<C1 T = int> void p(); // 0k
template<P N = 0> void q(); // 0k
template<P M = int> void r(); // Error: int is not an expression
— end example ]
```

<sup>6</sup> The declaration of a *constrained-parameter* introduces a new constraint on the template declaration. The constraint is formed by substituting the declared *template-parameter* as the first template argument of the concept declaration designated by the *constrained-type-specifier* in the *constrained-parameter* declaration. If the *constrained-type-specifier* is a *partial-concept-id*, its template arguments are substituted after the declared *template-parameter*. If the designated concept is a function concept, then the introduced constraint is a function call. [*Example:* 

```
template<C1 T> void f1(); // requires C1<T>
template<C2 U> void f2(); // requires C2<U>
template<P N> void f3(); // requires P<N>
```

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```
— end example ]
```

<sup>7</sup> If the *constrained-parameter* declares a template parameter pack, the formation of the constraint depends on whether the designated concept designated by the parameter's *constrained-type-specifier* is variadic. Let T be the declared parameter, C be the designated concept, and Args... be a sequence of template arguments from a *partial-concept-id*, possibly empty. If C is a variadic concept, then the associated constraint is a *template-id* of the form C<T..., Args...>. Otherwise, if C is not a variadic concept, the associated constraint is a conjunction of sub-constraints C<Ti, Args...> for each Ti in the parameter pack T. If C is a function concept, each introduced constraint or sub-constraint is adjusted to be a call expression of the form C<X, Args...>() where X is either the template parameter pack T or an element Ti. [ *Example*:

```
template<typename... Ts> concept bool P = ...;
template<typename T> concept bool U = ...;

template<P... Xs> void f4(); // requires P<Xs...>
template<U... Args> void f5(); // requires U<Args0> && U<Args1> && ... && U<Argsn>
```

Here, Args0, Args1, etc. denote elements of the template argument pack Args used as part of the introduced constraint. — end example ]

# 7.2 Template names

[temp.names]

- <sup>1</sup> Insert the following paragraphs after C++ §14.2/7.
- <sup>2</sup> If a *template-id* refers to a specialization of a constrained template declaration, the template's associated constraints are checked by substituting the *template-arguments* into the constraints and evaluating the resulting expression. If the substitution results in an invalid type or expression, or if the associated constraints evaluate to false, then the program is ill-formed. [ *Example:*

```
template<typename T> concept bool True = true;
template<typename T> concept bool False = false;
template<False T> struct S;
template<True T> using Ptr = T*;
S<int>* x;  // Error: int does not satisfy the constraints of False.
Ptr<int> z;  // Ok: z has type int*
```

- end example ] [ Note: Checking the constraints of a constrained class template does not require its instantiation. This guarantees that a partial specialization cannot be less specialized than a primary template. This requirement is enforced during name lookup, not when the partial specialization is declared. - end note ]

# 7.3 Template arguments

[temp.arg]

#### 7.3.1 Template template arguments

[temp.arg.template]

<sup>1</sup> Modify C++ §14.3.3.

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<sup>2</sup> A template-argument matches a template template-parameter (call it P) when each of the template parameters in the template-parameter-list of the template-argument's corresponding class template or alias template (call it A) matches the corresponding template parameter in the template-parameter-list of P, and the associated constraints of P shall subsume the associated constraints of A (7.6). [Example:

```
template<typename T>
  concept bool X = has_x<T>::value;
template<typename T>
  concept bool Y = X<T> && has_y<T>::value;
template<typename T>
  concept bool Z = Y<T> && has_z<T>::value;
template<template<Y> class C>
  class temp { ... };
template<X T> class x;
template<Z T> class z;
temp<x> s1; // OK: X is subsumed by Y
temp<z> s2; // Error: Z subsumes Y
```

The template x is a valid argument for temp because any template arguments satisfying Y will also satisfy X. That is, all uses of x by temp should result in well-formed programs. The template y is not valid because some template arguments satisfying Y may not satisfy Z. — end example ]

# 7.4 Template declarations

[temp.decls]

#### 7.4.1 Class templates

[temp.class]

- <sup>1</sup> Insert the following paragraph after C++ §14.5.1/3.
- <sup>2</sup> When a member of a constrained class template is defined outside of its class template definition, it shall be specified with the *template-parameters* and associated constraints of the class template.

[ Example:

```
template<typename T> concept bool Con = ...;

template<typename T> requires Con<T>
    struct S {
      void f();
      void g();
    }

template<typename T>
    requires Con<T>
      void S<T>::f() { } // Ok: parameters and constraints match

template<typename T>
    void S<T>::g() { } // Error: no declaration of g() in S<T>
      end example ]
```

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# 7.4.1.1 Member functions of class templates

[temp.mem.func]

- <sup>1</sup> Add the following paragraphs after C++ §14.5.1.1.
- A member function of a class template whose declarator contains a *requires-clause* is a *constrained member function*. [ *Example:*

```
template<typename T>
  class S {
    void f() requires C<T>();
};
```

- end example ]
- <sup>3</sup> Constraints on member functions are instantiated as needed during overload resolution, not when the class template is instantiated (C++ §14.7.1). [ *Note:* Constraints on member functions do not affect the declared interface of a class. That is, a constrained copy constructor is still a copy constructor, even if it will not be viable for a specialization of the class template.

   end note ]
- <sup>4</sup> A constrained member function of a class template may be defined outside of its class template definition. Its definition shall be specified with the constraints of its declaration. [Example: Consider possible definitions of the constrained member function S<T>f from above.

# 7.4.2 Member templates

[temp.mem]

- <sup>1</sup> Insert the following paragraph after C++ §14.5.2/1.
- <sup>2</sup> A constrained member template defined outside of its class template definition shall be specified with the *template-parameters* and constraints of the class template followed by the template parameters and constraints of the member template. [ *Example*:

```
template<typename T> concept bool Foo = ...;
template<typename T> concept bool Bar = ...; // Different than Foo

template<Foo T>
    struct S {
        template<Bar U> void f(U);
        template<Bar U> void g(U);
    };

template<Foo T> template<Bar U> void S<T>::f(U); // Ok
template<Foo T> template<Foo U> void S<T>::g(U); // Error: no g() declared in S
```

The template constraints in the definition of g do not match those in its declaration. — end example

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7.4.3 Friends [temp.friend]

- <sup>1</sup> Add the following paragraphs after C++ §14.5.4/9.
- A constrained friend of a class or class template is a constrained class template, constrained function template, a constrained ordinary or generic (non-member) function definition. [Example: When C is a type concept, all of the following are valid constrained friend declarations.

```
template<typename T>
  struct X {
   template<C U>
     friend void f(X x, U u) { } // Constrained function template

  template<C W>
     friend struct Z { }; // Constrained class template

  friend bool operator==(X a, X b) // Constrained ordinary function requires C<T>() { return true; }

  friend void g(X a, C b) { } // Constrained generic function };
```

Note that g is a generic function because the parameter b has a constrained-type-specifier. — end example ]

<sup>3</sup> A non-template friend function shall not be constrained unless the function's parameter or result type depends on a template parameter. [ *Example*:

```
template<typename T>
    struct S {
      friend void f(int n) requires C<T>(); // Error: cannot be constrained
    };
— end example ]
```

<sup>4</sup> A constrained non-template friend function shall not declare a specialization. [ Example:

```
template<typename T>
    struct S {
    friend void f<>(T x) requires C<T>(); // Error: declares a specialization

    friend void g(T x) requires C<T>() { } // OK: does not declare a specialization
    };

— end example ]
```

<sup>5</sup> As with constrained member functions, constraints on non-template friend functions are not instantiated during class template instantiation.

# 7.4.4 Class template partial specialization

[temp.class.spec]

# 7.4.4.1 Matching of class template partial specializations

[temp.class.spec.match]

- <sup>1</sup> Modify C++ §14.5.5.1/2.
- <sup>2</sup> A partial specialization matches a given actual template argument list if the template arguments of the partial specialization can be deduced from the actual template argument list

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(C++ §14.8.2), and the deduced template arguments satisfy the constraints of the partial specialization, if any (7.6).

### 7.4.4.2 Partial ordering of class template specializations

[temp.class.order]

- <sup>1</sup> Modify C++ §14.5.5.2/1.
- <sup>2</sup> For two class template partial specializations, the first is at least as specialized as the second if, given the following rewrite to two function templates, the first function template is at least as specialized as the second according to the ordering rules for function templates (C++ §14.5.6.2):
  - the first function template has the same template parameters <u>and constraints</u> as the first partial specialization and has a single function parameter whose type is a class template specialization with the template arguments of the first partial specialization, and
  - the second function template has the same template parameters <u>and constraints</u> as the second partial specialization and has a single function parameter whose type is a class template specialization with the template arguments of the second partial specialization.

## [ Example:

```
template<typename T>
  concept bool Integer = is_integral<T>::value;
template<typename T>
  concept bool Unsigned_integer = Integer<T> && is_unsigned<T>::value;
template<typename T> class S { };
template<Integer T> class S<T> { };  // #1
template<Unsigned_integer T> class S<T> { };  // #2

template<Integer T> void f(S<T>);  // A
template<Unsigned_integer T> void f(S<T>);  // B
```

The partial specialization #2 will be more specialized than #1 for template arguments that satisfy both constraints because B will be more specialized than A. —  $end\ example$ 

#### 7.4.5 Function templates

[temp.fct]

# 7.4.5.1 Template argument deduction

[temp.deduct]

- <sup>1</sup> Immediately after C++ §14.8.2/5, add the following paragraph:
- <sup>2</sup> If the template has associated constraints, the template arguments are substituted into those associated constraints and evaluated. If the substitution results in an invalid type or expression, or if the associated constraints evaluate to false, type deduction fails.

## 7.4.5.2 Function template overloading

[temp.over.link]

- <sup>1</sup> Modify C++ §14.5.6.1/6.
- <sup>2</sup> A function template can be overloaded either by (non-template) functions of its name or by (other) function templates of the same name. When a call to that name is written (explicitly, or implicitly using the operator notation), template argument deduction 7.4.5.1, and checking of

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any explicit template arguments C++ § , and checking of associated constraints 7.6 are performed for each function template to find the template argument values (if any) that can be used with that function template to instantiate a function template specialization that can be invoked with the call arguments. For each function template, if the argument deduction and checking succeeds, the template-arguments (deduced and/or explicit) are used to synthesize the declaration of a single function template specialization which is added to the candidate functions set to be used in overload resolution. If, for a given function template, argument deduction fails, no such function is added to the set of candidate functions for that template. The complete set of candidate functions includes all the synthesized declarations and all of the non-template overloaded functions of the same name. The synthesized declarations are treated like any other functions in the remainder of overload resolution, except as explicitly noted in C++ §.

- <sup>3</sup> Modify C++ §14.5.6.1
- <sup>4</sup> Two function templates are *equivalent* if they are declared in the same scope, have the same name, have identical template parameter lists, and have return types, and parameter lists, and constraints 7.6 that are equivalent using the rules described above to compare expressions involving template parameters.

# 7.4.5.3 Partial ordering of function templates

[temp.func.order]

- <sup>1</sup> Modify C++ §14.5.6.2/2.
- Partial ordering selects which of two function templates is more specialized than the other by transforming each template in turn (see next paragraph) and performing template argument deduction using the function type. The deduction process determines whether one of the templates is more specialized than the other. If so, the more specialized template is the one chosen by the partial ordering process. If the two templates have identical template parameter lists and equivalent return types and parameter lists, then partial ordering selects the template whose associated constraints subsume but are not equivalent to the associated constraints of the other 7.6. A constrained template is always selected over an unconstrained template.

# 7.5 Template instantiation and specialization

[temp.spec]

#### 7.5.1 Implicit instantiation

[temp.inst]

- <sup>1</sup> Insert the following paragraph after C++ §14.7.1/1.
- <sup>2</sup> The implicit instantiation of a class template does not cause the instantiation of the associated constraints of constrained member functions.

#### 7.5.2 Explicit instantiation

[temp.explicit]

- <sup>1</sup> Insert the following paragraph under C++ §14.7.2.
- An explicit instantiation of constrained template declaration (7) or constrained member function declaration (7.4.1.1) shall satisfy the associated constraints of that declaration (7.6). [Example:

```
template<typname T>
  concept bool C = requires(T t) { t.c(); };
template<typename T>
```

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```
requires C<T>
    struct X { }

template struct X<int>; // Error: int does not satisfy C.

— end example ]
```

# 7.5.3 Explicit specialization

[temp.expl.spec]

- <sup>1</sup> Insert the following paragraphs under C++ §14.7.3.
- <sup>2</sup> A constrained template declaration or constrained member function of a class template can be declared by a declaration introduced by template<>.
- <sup>3</sup> The *template arguments* of a *simple-template-id* that names an explicit specialization of a constrained template declaration must satisfy that template's associated constraints (7). [ *Example:* C is the type concept defined in the previous section.

```
template<C T>
    struct S1 { };

struct X { void c(); }

template<> S1<X> { };  // OK: X satisfies C
  template<> S1<int> { };  // Error: int does not satisfy C

- end example ]
```

<sup>4</sup> An explicit specialization of a constrained member function (7.4.1.1) shall not include a a *requires-clause*. [ *Example*:

```
template<typename T>
    struct S2 {
      void f(T) requires C<T>;
    };

template<> void S2<X>::f(T a) { } // OK
    template<> void S2<X>::f(T a) requires C<X> { } // Error: extra requires-clause
      — end example ]
```

# 7.6 Template constraints

[temp.constr]

- <sup>1</sup> Add this as a new section under C++ §14.
- <sup>2</sup> Certain contexts require expressions that satisfy additional requirements as detailed in this sub-clause. Expressions that satisfy these requirements are called *constraint expressions* or simply *constraints*.

constraint-expression: logical-or-expression

- <sup>3</sup> A *logical-or-expression* is a *constraint-expression* if, after substituting template arguments, the resulting expression
  - is a constant expression,
  - has type bool, and
  - both operands P and Q in every subexpression of a constraint of the form P || Q or P && Q have type bool.

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[ Note: A constraint-expression defines a subset of constant expressions over which certain logical implications can be proven during translation. The requirement that operands to logical operators have type bool prevents constraint expressions from finding user-defined overloads of those operators and possibly subverting the logical processing required by constraints.

— end note 1

- <sup>4</sup> A program that includes an expression not satisfying these requirements in a context where a *constraint-expression* is required is ill-formed.
- <sup>5</sup> [ Example: Let T be a dependent type, C be a unary function concept, P, Q, and R be value-dependent expressions whose type is bool, and M and N be integral expressions. All of the following expressions can be used as constraints:

An expression of the form M+N is not a valid constraint when the arguments have type int since the expression's type is not bool. Using this expression as a constraint would make the program ill-formed. — end example ]

- <sup>6</sup> A subexpression of a *constraint-expression* that calls a function concept or refers to a variable concept 4.1.4 is a *concept check*. A concept check is not evaluated; it is simplified according to the rules described in this section.
- Certain subexpressions of a constraint-expression are considered atomic constraints. A constraint is atomic if it is not:
  - a logical-or-expression of the form P || Q,
  - a logical-and-expression of the form P && 0.
  - a concept check,
  - a requires-expression, or
  - a subexpression of an atomic constraint.

The valid expression constraints, valid type constraints, result type constraints, and exception constraints introduced by a *requires-clause* are also atomic constraints.

[ Example:

```
has_trait<T>::value
M < N
M + N >= 0
true
false
```

— end example ]

[ *Note:* A concept check is not an atomic expression. — *end note* ]

<sup>8</sup> Constraints are *simplified* by reducing them to expressions containing only logical operators and atomic constraints. Concept checks and *requires-expressions* are replaced by simplified expressions. [ *Note:* An implementation is not required to normalize the constraint by rewriting in e.g., disjunctive normal form. — *end note* ]

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<sup>9</sup> A concept check that calls a function concept is simplified by substituting the explicit template arguments into the named function body's return expression. A concept check that refers to a variable concepts is simplified by substituting the template arguments into the variable's initializer.

- <sup>10</sup> A *requires-expression* is simplified by replacing it with the conjunction of constraints introduced by the *requirements* its *requirement-list*. [ *Note:* Certain atomic constraints introduced by a *requirement* have no explicit syntactic representation in the C++. *end note* ]
- 11 [ Example: Let P and Q be variable templates that are atomic constraints.

The associated constraints of f are simplified to the expression P<X> && Q<X>, and the associated constraints of g are simplified to P<X> || Q<X>. The associated constraints of h are:

```
P<X> && Q<X>
    && /* requires x.p() for all x of type X* /
    && /* requires that x.p() convert to int */
— end example ]
```

- <sup>12</sup> A constraint is *satisfied* if, after substituting template arguments, it evaluates to true. Otherwise, the constraint is *unsatisfied*.
- For a mapping M from a set X of atomic constraints to boolean values, let G(M) be the mapping from constraints to boolean values such that G(M)(C) is the result of substituting each atomic constraint A within C for M(A). For two constraints P and Q, let X be the set of all atomic constraints that appear in P and Q. P is said to S subsume Q if, for every mapping M from members of X to boolean values for which M(A) = M(B) whenever A and B are equivalent, either G(M)(P) is false or G(M)(Q) is true (or both).
- 14 Two constraint-expressions P and Q are logically equivalent if and only if P subsumes Q and Q subsumes P.

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