Stratis Software Design: Version 0.6^*

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Introduction

Stratis is a local storage solution that lets multiple logical filesystems share a pool of storage that is allocated from one or more block devices. This is roughly equivalent to what volume-managing filesystems like ZFS and Btrfs do, but differs in implementation. Instead of an entirely in-kernel approach, Stratis uses a hybrid user/kernel approach that builds upon existing block capabilities like devicemapper, existing filesystem capabilities like XFS, and a user space daemon for monitoring and control.

The goal of Stratis is to provide the conceptual simplicity that volume-managing filesystems offer, many of their features, but also surpass them in areas like monitoring and notification, automatic reconfiguration, and integration with higher-level storage management frameworks.

Asking Questions and Making Changes to this Document

This document can be found in the stratis-docs repo, and is written using LyX 2.2.1. Please ask any questions by opening an issue, and propose changes as pull requests.

1 Conceptual Model

Stratis's conceptual model consists of blockdevs, pools, and filesystems. A pool is created from one or more blockdevs (block devices), after which filesystems may be created from the pool. Filesystems are mountable hierarchical collections of files that allocate backing storage from the pool on an as-needed basis. The key difference between a Stratis filesystem and a conventional Unix filesystem is that Stratis filesystem sizes are not managed by the user, but by Stratis.

Since a single system may have multiple pools, each pool has a name, as does each filesystem within a pool. Each pool has a data redundancy level, which may be none, raid1, raid5, raid6, or raid10. The default is none. The pool's cache also has a redundancy level that also defaults to none. This redundancy level applies to all storage allocated within the pool.

Stratis supports effectively unlimited blockdevs in a pool, and additional blockdevs may be added after the pool is created. Blockdevs may also be removed from a pool, if certain preparations are performed and conditions are met.

Stratis supports up to 2^{20} filesystems. New filesystems may be created from an existing filesystem (i.e. a read/write snapshot), or created from scratch. After creation, Stratis treats all filesystems identically; a filesystem and its snapshot have no dependencies on each other, but simply take up less space in the pool due to each initially referencing the same set of data blocks. However, to aid the user informationally, Stratis tracks and reports snapshot parents.

2 User Experience

Stratis has a command-line tool that enables the administrator to create a Stratis pool from one or more blockdevs, and then allocate filesystems from the pool.

See reference implementation at https://github.com/stratis-storage/stratis-cli for the most up-to-date status of the CLI design.

3 D-Bus Programmatic API

The Stratis service process exposes a D-Bus interface, for other programs to integrate support for Stratis. This is considered the primary Stratis interface. The command-line tool uses the D-Bus API.

TODO: Create a separate D-Bus API Reference Doc and link to it here.

3.1 D-Bus Access Control

- 3.1.1 Security Policy
- 3.1.2 Prevent Spoofing

4 Scalability and Performance Considerations

Stratis assumes three usage models:

- 1. Disk only
- 2. Disk with SSD caching
- 3. SSD only

The use cases that might do(1) are not terribly concerned with IOPs, but are concerned about resiliency, integrity, and ease of management. This might include large media files that are accessed mostly sequentially, or other "cold" data.

If random IO performance is needed, adding a caching tier based on SSDs(2) is going to be better than any optimization of a disk-only solution. Stratis therefore doesn't optimize case(1) for IOPs, but supports setting up faster devices in a caching tier, for either read caching or read/write caching (for which users likely want redundancy in the cache tier).

In case(3), the cache tier isn't needed; the pool should have high IOPs. Stratis adds value through resiliency, integrity, ease of use, and ease of management.

5 Software Components

Stratis consists of a command-line tool, stratis, and a service, stratisd.

stratis implements the command-line interface, and converts commands into D-Bus API calls to stratisd.

stratisd implements the D-Bus interface, and manages and monitors Stratis internal pool blockdevs, as described below. It is started by (udev? systemd?) and continues to run as long as Stratis pools or blockdevs are present in the system.

6 Internals

Stratis internals are opaque to the user. This allows its implementation maximum flexibility to do whatever it needs in Stratis version 1, as well as to be extended in later versions without violating any user-visible expectations.

6.1 Cache Tier Requirements

Caching may be configured for redundancy, or no redundancy.

Caching may be configured for write-back and write-through modes.

Stratis concatenates all cache blockdevs and uses the resulting device to cache the thinpool device (L4). This lets all filesystems benefit from the cache.

Cache blocksize should match thinpool datablock size.

The cache tier can be entirely reconfigured without data loss, once flushed and in write-through mode, because it's just a cache. To handle configuration changes, we can just drop the cache and recreate it with its new configuration.

(more details to fill in here.)

6.2 Cache Tier Metadata

6.2.1 Cache Tier Metadata Requirements

- 1. Identify all blockdevs part of the pool's cache tier, the configured redundancy level, and other cache-specific configuration parameters (e.g. WT/WB, block size, cache policy)
- 2. Cache tier supports up to 8 devices.

6.3 Data Tier Requirements

The data tier of Stratis must manage blockdevs on behalf of the user to provide the following:

- 1. Managed filesystems that consume only as much space as the files they contain
- 2. Fast snapshots of existing filesystems
- 3. The ability to add (and eventually remove) individual blockdevs to grow the available space available to filesystems
- 4. User-selectable redundancy level (per-pool granularity)⁴
- 5. Integrity checking⁵

6.4 Data Tier

The data tier achieves these requirements by layering Devicemapper (DM) devices on top of the pool's blockdevs, culminating in a thin devices allocated from a thinpool. Stratis initializes these with a filesystem, and manages the DM devices and filesystems to meet the above requirements.

6.4.1 Layer 0: blockdevs

This layer is responsible for discovering existing blockdevs in a pool, initializing and labeling new blockdevs unambiguously as part of the pool, setting up any disk-specific parameters (TLER?), and storing pool metadata on each blockdev. Blockdevs may be in the following states: good-unused, good-used, bad, spare.

6.4.2 Layer 1: Integrity⁶

Identity-mapped immediately above each blockdev in the pool, Stratis uses a to-be-developed DM target that allows the detection of incorrect data as it is read, by using extra space to record the results of checksum/hash functions on

⁴in a future version

 $^{^5 \}mathrm{in}$ a future version

 $^{^6}$ in a future version

the data blocks, and then compare the results with what the blockdev actually returned. This will enable Stratis to detect data corruption when the pool is non-redundant, and to repair the corruption when the pool is redundant. Such a DM target could use DIF information if present, but would also need to work with non-DIF blockdevs, perhaps requiring some space for its metadata.

6.4.3 Layer 2: Redundancy⁷

A Stratis pool may optionally be configured to spread data across multiple physical disks, so that the loss of any one disk does not cause data loss. Stratis uses conventional RAID technology (1, 5, 6, 10) as specified, and converts Layer 0 blockdevs into a smaller-sized amount of storage with the specified raid properties.

Since Stratis supports more blockdevs than are RAID-able (generally 8 or fewer is best for a raidset IIRC), and differently-sized blockdevs, a redundant Stratis pool will likely contain multiple raid sets (all of the same global type) Depending on layout, there may be some amount of space in a pool's blockdevs that cannot be used becauseit cannot be used in a RAID set. Stratis will intensively manage raidsets, extending them across newly added blockdevs or creating new raidsets; and handling the removal of blockdevs⁸. Stratis may use dm-raid's reshape capabilities when possible.

A pool may also be configured with redundancy level of 'none'. In this case, Stratis may opt to create a striped blockdev ('raid0') across all or some part of the blockdevs, if there are more than one.

Stratis cannot support redundancy with a single disk, but we may wish to reserve the small space for raid metadata and other uses even on one-disk Stratis pools. This will allow the pool to be made redundant (in the version when we support this) without encountering some ugly edge cases.

6.4.4 Layer 3: Flexibility

Whether blockdevs are part of raidsets or used directly, pools need to cope with the addition or removal⁹ of them.

Stratis needs to allow adding a blockdev to an existing pool, and use the new available to allow its the pool's allocated space to grow. (This may involve reshaping one or more raidsets to be wider, or just knowing to extend the thinpool onto the blockdev when previous blockdevs are full.)

In a future version, Stratis needs to be able assign states to each blockdev that reflect whether a) it can be removed from the pool with no effect, b) removed with loss of redundancy, or c) not removed without data loss. It should also be able to move active data off a perhaps-failing disk (or refuse, if space on other disks in the pool is not sufficient) so that the disk then can be removed from the pool. blockdevs may also be marked as 'spare', so that they do not

⁷in a future version

 $^{^8 \}mathrm{in}$ a future version

 $^{^9 \}mathrm{in}$ a future version

contribute to the total available space of the pool, but can immediately replace a failed disk.

The flexibility layer will support two linear DM devices made up of segments from lower-level devices. It will track what lower-level devices they are allocated from, allow the two devices to grow, and handle the online movement of segments in these devices when lower-level devices come and go.

6.4.5 Layer 4: Thin Provisioning

The two linear targets from L3 are used as metadata device and data device for a DM thin pool. Stratis manages the thinpool device by extending the two L3 subdevices when either runs low on available blocks. If the pool approaches a point where the pool no longer has empty lower-level space to extend onto, Stratis alerts the user and does what it must to avoid data corruption.

6.4.6 Layer 5: Thin Volumes

From the thin pool Stratis creates one or more thin volumes. It will automatically give a new volume a default size, and format it with a filesystem, and make it available to the user.

Stratis keeps snapshot origin info in its metadata for user informational purposes.

6.4.7 Layer 6: Encryption¹⁰

Stratis could enable per-filesystem encryption between the thin device and the filesystem.

6.4.8 Layer 7: Filesystem

Stratis monitors each filesystem's fullness vs its capacity and automatically extends (using e.g. xfs_growfs) them without user intervention. Stratis also periodically runs fstrim on each one as well.

6.5 Data Tier Metadata

Stratis must track the blockdevs that make up the data tier of the pool (L0), integrity parameters (L1), the raidsets that are created from the data blockdevs (L2), the two linear targets that span the L2 devices (L3), the thinpool device (L4) and the attributes of the thin devices (L5) and filesystems created from the thinpool (L7).

6.5.1 Requirements

1. Uniquely identify a blockdev as used by Stratis, which pool it is a member of, and parameters needed to recreate all layers

 $^{^{10}{}m in}$ a future version

- 2. Detect incomplete or corrupted metadata
- 3. Allow for blockdevs being expanded underneath Stratis
- 4. Redundant on each blockdev to tolerate e.g. accidental dd
- 5. Redundant across blockdevs to handle missing, damaged etc. members. Can provide number & characteristics of missing blockdevs
- 6. Handle thousand+ blockdevs in a pool
- 7. Handle million+ file systems in a pool and updates without writing to each block dev
- 8. Extensible/upgradeable metadata format

6.5.2 Design Overview

Stratis metadata is in three places:

- 1. Blockdev Data Area
 - (a) Signature Block
 - (b) Metadata Area (MDA)
- 2. Thin Device 0

6.5.3 BlockDev Data Area (BDA)

The BDA consists of a fixed-length Static Header of eight sectors, which contains the Signature Block; and the metadata area (MDA), whose length is specified in the Signature Block. These are written to the beginning and end of the blockdev as described below.

Sectors are 512 bytes in length¹¹.

Stratis reserves the first 8 sectors of each blockdev for the Static Header:

sector offset	length (sectors)	contents
0	1	unused
1	1	Signature Block
2	6	unused

The Signature Block consists of:

¹¹Historically this is the minimum storage unit of a hard drive. Many Linux kernel APIs assume this value is constant (as does this document), and use another term such as 'block size' for dealing with cases where the minimum storage unit is different.

byte offset	length (bytes)	description
0	4	CRC32 of signature block (bytes at offset 4 length 508)
4	16	Stratis signature: '!Stra0tis\x86\xff\x02^\x41rh'
20	8	Device size in 512-byte sectors (u64)
28	4	flags
32	32	ASCII hex UUID for the Stratis pool
64	32	ASCII hex UUID for the blockdev
96	8	MDAa UNIX timestamp (seconds since Jan 1 1970) (u64)
104	4	MDAa nanoseconds (u32)
108	4	MDAa used length in bytes (u32)
112	4	MDAa CRC32 over used length
116	12	unused
128	8	MDAb UNIX timestamp (seconds since Jan 1 1970) (u64)
136	4	MDAb nanoseconds (u32)
140	4	MDAb used length in bytes (u32)
144	4	MDAb CRC32 over used length
148	12	unused
160	4	sector length of blockdev metadata area (MDA) (u32)
164	4	sector length of reserved space (u32)
168	344	unused

- No flags are yet defined, so 'flags' field is zeroed.
- All 'unused' fields are zeroed, and are reserved for future use.
- Blockdev metadata area (offset 160) must be an even number of at least 2040 combined minimum length of overall BDA (static header and MDA) is 2048 sectors (1 MiB).
- Reserved space (offset 164) is used by Stratis for temporary storage or other undefined purposes.
- The BDA is written to the beginning and end of the blockdev. When writing to the beginning of the blockdev, the Static Header is placed first (LBA 0), then the MDA starts at LBA 8, then reserved space. When writing to the end of the blockdev, the Static Header occupies the last 8 sectors of the blockdev, and the MDA precedes it. (The end-of-blockdev BDA has no reserved space.)
- The BDA at the start of the blockdev is followed immediately by reserved space, whose size is specified in the signature block (offset 164).
- The MDA is evenly divided into two areas: MDAa and MDAb.
- Data within the metadata areas is stored in JSON format.
- Metadata updates write to the older of the MDAa and MDAb areas. This is determined by lowest timestamp, and then lowest nanoseconds if timestamps are equal.

- The procedure for updating metadata is:
 - 1. Determine which area in the MDA to use (MDAa or MDAb) as described above. Call this area 'MDAx'.
 - 2. The JSON metadata is written to MDAx.
 - 3. MDAx's timestamp, nanosecond, length, and CRC32 fields in the signature block are updated: timestamp gets the UNIX timestamp. Nanoseconds is the current nanoseconds within the second. Length is set to the length of the metadata, and the CRC32 is calculated on the data up to that length.
 - 4. The signature block's CRC32 is calculated.
 - 5. The signature block is written and flushed to disk. The implementation may choose to write the entire 8 sectors of the static header.
 - 6. The identical MDAx, and then signature block is written to the end-of-blockdev BDA, followed by a flush/FUA.
- Multiple blockdevs being updated with the same metadata must use identical values for MDAx values in the signature block, but which area (MDAa or MDAb) is used may vary, if they have received differing numbers of metadata updates over time. (Corollary: sig block CRC32 must be recalculated for each blockdev.)

6.5.4 Metadata Area (MDA)

The MDA contains a JSON object that represents the pool's overall configuration of blockdevs, from L0 to L4.

Top level objects:

name	json type	required	description
name	string	у	the name of the Stratis pool
block_devs	object	У	the block devices in the pool (see
			below)
integrity_devs	object	n	info on integrity devices & layout
			(TBD)
$raid_devs$	object	n	layout of the raid devices (TBD)
$flex_devs$	object	У	layout of the data and metadata
			linear devices (see below)
thinpool_dev	object	У	parameters of the thinpool device
			(see below)
thin_devs	object	у	parameters for thin devices (see
			below)

[&]quot;block_devs" is an object whose strings are 32-byte ASCII hex unids and whose values are objects with the following strings:

name	json type	required	description
dev	string	n	the entry in /dev where this blockdev
			was last available
size	number	У	the blockdev's size in sectors
location	string	n	user-provided information for tracking
			the device
disk_id	string	n	uniquely identifying information for
			the block device, such as SCSI VPD83
			or serial number

integrity_devs: TBD raid devs: TBD

"flex_devs" is an object that contains two strings: "meta_dev" and "data_dev". Each of these values are arrays of objects. These define the linear segments that make up each. Each segment object has the following strings:

	name	json type	required	description
Ī	parent	string	У	32 byte ASCII hex uuid of the
				lower-layered device the segment is in
Ī	start	number	У	the starting sector offset within the parent
				device
Ī	length	number	У	the length in sectors of the segment

If neither "integrity_devs" or "raid_devs" sections are present, "parent" is a valid blockdev uuid. If only "integrity_devs" is present, parent refers to an integrity dev uuid. If only "raid_devs" are present, "parent" refers to a raid dev uuid. If both "integrity_devs" and "raid_devs" are present, "parent" refers to a raid dev uuid.

"thinpool dev" is an object with the following values:

· —	v		e e e e e e e e e e e e e e e e e e e
name	json type	required	description
data_block_size	number	У	the size in sectors of the
			thinpool data block size
low_water_blocks	number	у	the number of free data blocks
			that dropping below will result
			in extending the data device

"thin_devs" is an object whose strings are ASCII strings (base 10) of 24-bit unsigned values, representing in-use thin device numbers. The values are objects with these fields:

name	json type	required	description
fs	string	У	the filesystem used to format the thin device. The only valid value is "xfs"
name	string	У	the name the user has given this filesystem
size	number	у	the size in sectors of the thin device

An entry for thin dev "0" is required, and is used by Stratis for its purposes (see below). Its "name" field is an empty string.

6.5.5 Thin Device 0

Thin device 0 is an XFS-formatted thin filesystem that is used by Stratis to store information on user-created filesystems (L5). These are stored in the filesystem in TBD format.

6.5.6 The MDA and Very Large Pools

Stratis pools with very large numbers of blockdevs will encounter two issues. First, updating the metadata on all blockdevs in the pool may become a performance bottleneck. Second, the default MDA size may become inadequate to contain the information required.

To solve the first issue, Stratis caps the number of blockdevs that receive updated metadata information. A reasonable value for this cap might be in the range of 6 to 10, and should try to spread metadata updates across pathindependent blockdevs, if this can be discerned, or randomly. This limits excessive I/O when blockdevs are added or removed from the pool, but maximizes the likelihood that up-to-date pool metadata is retrievable in case of failure.

To solve the second issue, Stratis monitors how large its most recent serialized metadata updates are, and increases the size of MDA areas on newly added devices when a fairly low threshold -%50 – is reached in comparison to the available MDAx size. This ensures that by the time sufficient blockdevs have been added to the pool to be in danger of serialized JSON data being too large, there are enough blockdevs with enlarged MDA space so that they can be used for MDA writes.

7 Delivery of Features

7.1 Stratis version 0.1

Simplest thing that does something useful

- 1. Create a pool
- 2. Destroy a pool
- 3. Create a filesystem
- 4. Create a filesystem from existing filesystem (a r/w snapshot)
- 5. Destroy a filesystem
- 6. List filesystems
- 7. Rename filesystems
- 8. List pools
- 9. Rename pools

- 10. List blockdevs in a pool
- 11. Redundancy level: none
- 12. D-Bus API
- 13. Command-line tool
- 14. Save/restore configuration across reboot
- 15. Initial disk labeling and on-disk metadata format
- 16. Error handling for missing, corrupted, or duplicate blockdevs in a pool

7.2 Stratis version 0.5

Add cache tier

- 1. List cache blockdevs in a pool
- 2. Add cache blockdevs
- 3. Remove cache blockdevs
- 4. Write-through caching only

7.3 Stratis version 1.0

Minimum Viable Product

- 1. Snapshot management: auto snaps, date-based culling, "promotion" from snap to "primary"
- 2. Monitor pool(s) for getting close to capacity, and do something (remount ro?) if dangerously full
- 3. Some sort of notification method to the user if pool is approaching capacity
- 4. Maintain filesystems: Grow a filesystem as it nears capacity
- 5. Maintain filesystems: Run fstrim periodically to release unused areas back to thinpool
- 6. Add and use an additional blockdev

7.4 Stratis version 2.0

Add Redundancy Support

- 1. Remove an existing blockdev
- 2. Redundancy level: raid1
- 3. Redundancy level: raid5
- 4. Redundancy level: raid6
- 5. Redundancy level: raid10
- 6. Cache redundancy level: raid1
- 7. Write-through caching enabled
- 8. thin/cache metadata validation/check (i.e. call thin check & cache check)
- 9. Quotas
- 10. Blockdev resize (larger)
- 11. Spares

7.5 Stratis version 3.0

Rough ZFS feature parity. New DM features needed.

- 1. Send/Receive Integrity checking (w/ self-healing only if on raid)
- 2. Raid scrub
- 3. Compression
- 4. Encryption
- 5. Dedupe[aq]
- 6. Raid write log (on ssd? To eliminate raid write hole)

7.6 Stratis version 4.0

Future features and evolution

- 1. Change a pool's redundancy level
- 2. Boot from a filesystem
- 3. Libstoragement integration
- 4. Multipath integration
- 5. Tag-based blockdev and filesystem classification/grouping
- 6. Mirroring across partitions within a pool, for multi-site or across hw failure domains (shelves/racks)

8 Open Questions

Initial filesystem sizing. Mkfs does different things depending on the size of the blockdev. If it's small then things will be suboptimal if we grow it substantially. Weigh this against too large, which would waste thinpool space (mkfs touches/allocates more thin blocks).

Alignment and tuning of sizes across layers. It would be great if the fs happens to write to a new location that allocates a thin block that it uses that entire block. Also, look at XFS allocation groups, they may work cross-purposes to thinpool by spreading files across the blkdev. Align as much as possible.

Behavior when thin pool is exhausted. Slow down? Switch all fs to readonly?

Should we separate cache into read devs and write devs? Write-back cache we may want to support redundancy (cache contains only copy of data until flush) whereas read cache redundancy serves little purpose, just reduces the total space available.

Snapshot semantics & management. Snapshots should be thought of as "secondary" and read-only until promoted to "primary" and read-write¹².

Quotas. If the default behavior is to grow a vol that nears its notional capacity, how do we limit it for vols (e.g. /var/log or time machine) where being limited in size is part of how they expect to operate?

Should we have a 'monitor' subcommand? What would it do?

Being able to identify and differentiate blockdevs within a pool is very important, given how many there can be. We should provide the user with system-generated info, and also allow the user to add their own descriptions of blockdevs.

Journald interactions

Nfs interactions

 $^{^{12}\}mathrm{See}\ \mathrm{GitHub}\ \mathrm{issue:https://github.com/stratis-storage/stratisd/issues/19}$