

# Computability and Complexity

## Problem Set 5

### Fixed point theorem and applications

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1. The **S** property is defined as follows :

$$\forall k \exists S \text{ total \& computable} : \varphi_k(x, y) = \varphi_{S(x)}(y)$$

Prove that the **S** property is a particular case of S-m-n (i.e. prove that S-m-n implies **S** for  $m = n = 1$ ).

2. Using the fixed point theorem, show that there exists a program  $P_n$  such that  $P_n$  terminates only for input  $n$ . (Hint : use the function  $g(n, x) = 1$  if  $x = n$ ,  $\perp$  otherwise together with the **S** property)
3. Using the fixed point theorem, show that there exists a program  $P_n$  that always outputs  $n$  (i.e. that prints its source code).
4. Prove Rice's theorem using the fixed point theorem.  
(Hint : define the function  $f(x) = i$  if  $x \in A$ ,  $j$  if  $x \in \overline{A}$ , with  $i \in \overline{A}$  and  $j \in A$ )
5. Prove that  $K = \{n \in \mathbb{N} \mid \varphi_n(n) \neq \perp\}$  is not recursive using the fixed point theorem.

### Challenge

Show that, for any computable total function  $f$ , there exist an infinity of  $k$ 's such that  $\varphi_k = \varphi_{f(k)}$ .

(Hint : Show that if it was not the case, we could find a computable total function that would not satisfy the fixed point theorem)

1. S-m-n  $\rightarrow \exists f$  totale calculable

$$\forall k: f_k(x, y) = f_{f(k, x)}(y)$$

S  $\rightarrow \forall k \exists f$  totale calculable :  $f_k(x, y) = f_{f(k, x)}(y)$

$$f(k, x) = f_S(k, x) = f_{k'}(x)$$

↳. Soit  $g(n, x) = 1 \begin{cases} \text{si } n = x \\ \perp \text{ sinon} \end{cases}$

- Avec S,  $\exists f$  totale calculable :  $f_{g(n, x)} = f_{S(n)}(x)$

- Point fixe :  $\exists p \text{ p.m. } n : f_n = f_{S(n)}$

$$\rightarrow f_n(x) = f_{S(n)}(x) = f_{g(n, x)}$$

$\rightarrow g$  calculable  $\rightarrow p$  existe

3.  $f(n, x) = x^n$  calculable

- avec S,  $\exists f$  totale calculable :  $f_{f(n, x)} = f_{S(n)}(x)$

- pt fixe :  $\exists p \text{ p.m. } n : f_n = f_{S(n)}$

$$\rightarrow \exists n : f_{f(n, x)} = f_{S(n)}(x) = f_n(x)$$

$\rightarrow f$  calculable  $\rightarrow p$  existe

4. Soit  $i \in A$ ,  $j \in \bar{A}$

$$\rightarrow f(x) = \begin{cases} j & \text{si } x \in A \\ i & \text{si } x \in \bar{A} \end{cases}$$

Si  $A$  récursif,  $f$  calculable

Point fixe:  $\exists k; \varphi_k = \varphi_{g(k)}$

Si  $k \in A$ ;  $\varphi_k = \varphi_j$

Si  $k \in \bar{A}$ ;  $\varphi_k = \varphi_i$