## Computability and Complexity Problem Set 10

Reductions and complexity classes

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- 1. Algorithmic versus functional reduction.
  - (a) Prove that for every set A,  $A \leq_a \bar{A}$ .
  - (b) Prove that for all sets A, B, if  $A \leq_f B$  and B is recursively enumerable, then A is recursively enumerable.
  - (c) Show an example for which the relation  $A \leq_f \bar{A}$  does not hold.
  - (d) Show an example for which the following relation does not hold: if  $A \leq_a B$  and B is recursively enumerable, then A is recursively enumerable.
- 2. Prove the following relations.
  - (a) Does  $A \leq_a B$  imply  $A \leq_f B$ ?
  - (b) Does  $A \leq_f B$  imply  $A \leq_a B$ ?
- 3. Let  $re = \{A | A \text{ is recursively enumerable}\}.$ 
  - (a) State the conditions to prove that HALT is re-complete wrt  $\leq_a$ .
  - (b) Prove the conditions.
- 4. Does the class  $DTIME(n^2)$  depend of your calculus model (e.g. Java language)? And what about the class P? Justify your answers.
- 5. Let A be a recursive set and f a total function from  $\mathbb N$  to  $\mathbb N$ . Is the following implication true? Explain.

$$A \in NTIME(f) \Rightarrow \exists c > 1 : A \in DTIME(c^f)$$

6. Let A be a recursive set and f a total function from  $\mathbb{N}$  to  $\mathbb{N}$ . Is the following implication true? Explain.

$$A \in DTIME(f) \Rightarrow A \in DSPACE(f)$$