

# INFO-F-405 : Security

## Encryption and statistical analysis

The goal of this session is to learn about encryption and decryption as well as about first cryptographic tools and basic techniques of cryptanalysis. We are going to use shift ciphers during this exercises.

### 1 Mono-alphabetic encryption

One of the first ciphers was the Caesar's cipher, it is a mono-alphabetic shift cipher. Long time ago this cipher was used with shift of 3 letters in the alphabet and it was enough to secure messages, since most of people did not know how to read at all. This method became completely obsolete with the discovery of an attack based on frequency analysis.

There exist other mono-alphabetic ciphers like substitution cipher, but we are going to use a special case which is the shift cipher.

#### 1.1 Encryption and decryption

##### 1.1.1 Encryption

The encryption is relatively strait forward. Let us represent letters of the latin alphabet by numbers between 0 and 25, in other words:

$$A \rightarrow 0, B \rightarrow 1, \dots, Z \rightarrow 25$$

Now we can write an encryption function  $E_k(x)$  for the shift cipher in  $\mathbb{Z}_{26}$  with a secret key  $k$ , where  $k \in [0; 25]$ , this value  $k$  represents the amplitude of the shift. Here is the definition of the encryption function:

$$E_k(x) = (x + k) \bmod 26$$

##### 1.1.2 Decryption

The decryption function is the inverse of the encryption function  $E_k(x)$ , we will note it as  $D_k(x)$ . We can write the decryption function:

$$D_k(x) = (x - k) \bmod 26$$

## Exercise 1

A cipher with  $\mathcal{M} = \mathcal{C}$  for certain keys  $k \in \mathcal{K}$  is called *involutory* if its encryption and decryption procedures become identical, i.e. for all  $m \in \mathcal{M} : E_k(m) = D_k(m)$ . Under which keys  $k \in [0, 25]$  the above Shift Cipher becomes involutory?

### 1.2 Attacks

#### 1.2.1 Brute-force attack

The exhaustive key search a.k.a. brute-force is very easy to implement and in this case it is also very easy (and fast) to execute since we only have to test 26 different values for the key. Afterwards a human can easily detect which decryption was successful. Sometimes one can immediately spot the key while looking at the ciphertext by finding small frequently used words like “the”, “of” or “to”.

## Exercise 2

Write a program that uses brute-force attack in order to find the secret key that decrypts the following message:

FbZR crbcyr jrne Fhcrezna cnwnznf. Fhcrezna jrnef Puhpx Abeevf cnwnznf.

Here is another ciphertext:

'Yvccf, nfigu!' - zj fev fw kyv wzijk kyzexj gvfgcv kip  
kf gizek nyve kyvp cvrie r evn gifxirddzex crexlrzv.

#### 1.2.2 Known-Plaintext Attacks

In one of the previous exercises we performed cryptanalysis by knowing only the ciphertexts. Here we consider known-plaintext attacks where the analyst may know plaintext/ciphertext pairs.

## Exercise 3

- (a) How many pairs of plaintext/ciphertext characters must be known in order to determine the key of the Shift Cipher?
- (b) Give an answer to question (a) for the Vigenère Cipher.

#### 1.2.3 Statistical analysis

One of the first statistical attacks on this cipher is an attack based on frequency analysis. The idea is to calculate the frequency of each letter of the ciphertext and to compare them to the reference table that contains frequencies of letter use in a given language. Using this technique we can find the correct shift i.e. the secret key that was used. Table 1 gives frequencies of letters for a text written in English.

## Exercise 4

## Exercice 1

$k=0 \rightarrow$  pas de shift

$k=13$  car double shift:  $13 + 13 \pmod{26} = 0$

## Exercice 3

a) seulement 1 : exemple :  $\overset{\text{cipher}}{\uparrow} C = d$   $\overset{\text{plaintext message}}{\uparrow} m = a$

$$\Rightarrow k = C - m = 4 - 1 = 3$$

b) la taille de la clé

A	B	C	D	E	F	G	H	I	J	K	L	M
8.17	1.49	2.78	4.25	12.70	2.23	2.02	6.09	6.97	0.15	0.77	4.03	2.41
N	O	P	Q	R	S	T	U	V	W	X	Y	Z
6.75	7.51	1.93	0.09	5.99	6.33	9.06	2.76	0.98	2.36	0.15	1.97	0.07

Table 1: Relative frequencies of letter in a text in English (in %).

Write a program that uses frequency analysis in order to find the key. Test it on the previous exercise and compare the result with the previous one (that used the brute-force attack).

### 1.3 Kasiski examination/Classical ciphers

**Kasiski Method.** This test allows to get a good approximation of the size of the key. The idea is to search patterns of 2, 3 or 4 letters (N-grams) that appear regularly. Once we find some patterns, we can calculate the distance between them. Once it is done, we can calculate the *gcd* between these distances and approximate the size of the key.

**Vigenère cipher.** The Vigenère cipher is a polyalphabetic cipher, its development was a very important step in the history of cryptography. The Vigenère cipher was not broken for the period of 200–300 years after its development.

#### Exercise 5

The following ciphertext was generated using the Vigenère cipher from the plaintext in English language alphabet containing only capital letters (ABC...XYZ). That is,  $\#c_i = \#m_i + \#k_j \pmod{26}$  where  $k = k_1 \dots k_d$  is the key,  $m = m_1 \dots m_n$  is the plaintext, and  $c = c_1 \dots c_n$  is the ciphertext. Vigenère cipher assumes the following numbering of alphabet letters:  $\#A = 0, \#B = 1, \dots, \#Y = 24, \#Z = 25$ .

LNRKE RRLWZ HCUEG ZAQEO PQLGY EBDY QWFRS YLLCG TMVEY DNBKE  
 RRLBX SMIHG VLGIE GQTFH LYQDM ISSEM YUILH SQRWC VAGOE BXPRR  
 TFHSS QTGOL UHFMX NBWHC VEYRF EUECQ ALGWC OIETH ZHNCD TFWHC  
 ZATHS GQOSU YCOLM ZSSEM YUILH

- (1) The most frequent digramms in the ciphertext are HC, HS and RR, each occurs three times. Find the length  $d$  of the key using the Kasiski Method.
- (2) Assume that digramm HC decrypts to HE and HS decrypts to ES. Determine the letters of the key  $k$ .

#### Playfair cipher.

- was used in WW I by British army
- substitutes digrams (more efficient than substituting single characters)
- key  $k$  is a 5x5 matrix filled with alphabet letters (without J)

Let  $m = m_0, \dots, m_n$ .  $Enc(k, m) = c_0, \dots, c_n$  according to the rules:

1.  $m_i, m_{i+1}$  in one row:  $c_i$  right to  $m_i$ ,  $c_{i+1}$  right to  $m_{i+1}$
2.  $m_i, m_{i+1}$  in one column:  $c_i$  below  $m_i$ ,  $c_{i+1}$  below  $m_{i+1}$
3.  $m_i, m_{i+1}$  span a rectangle:  $c_i$  opposite to  $m_i$ ,  $c_{i+1}$  opposite to  $m_{i+1}$  (within the same row)
4.  $m_i = m_{i+1}$  or  $|m|$  is odd: inject or append some pre-specified dummy letter.

### Exercise 6

Compute the plaintext from the following Playfair ciphertext using the  $5 \times 5$ -key matrix below.

FK SM VS WV ST QC

A	B	L	Z	E
S	Y	F	C	O
P	M	W	D	G
T	I	R	K	V
N	Q	U	H	X

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A
C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C
E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D
F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E
G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F
H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G
I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H
J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I
K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J
L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K
M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L
N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M
O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N
P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P
R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q
S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R
T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T
V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U
W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V
X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W
Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X
Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y

Figure 1: Vigenère square.