

Homework #3

CPE 512

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Part 1 Serial Test

CODE

• See Appendix A for the source code

OUTPUT

```
uahcls01@dmcvlogin1:Hw3> ./mm_mult_serial 4 6 3
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=1e-06 seconds
```

Part 2 Parallel Test

CODE

See Appendix B for the source code

OUTPUT NP = 1

```
uahcls01@dmcvlogin2:Hw3> mpirun -np 1 ./mm_mult_mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709

B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622

C matrix =
9889.42 18581 17272.7
```

```
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=1.7e-05 seconds
OUTPUT NP = 2
uahcls01@dmcvlogin2:Hw3> mpirun -np 2 ./mm_mult_mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=9.9e-05 seconds
OUTPUT NP = 3
uahcls01@dmcvlogin2:Hw3> mpirun -np 3 ./mm_mult_mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=0.000154 seconds
OUTPUT NP = 4
uahcls01@dmcvlogin2:Hw3> mpirun -np 4 ./mm mult mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
```

```
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=0.00015 seconds
OUTPUT NP = 5
uahcls01@dmcvlogin2:Hw3> mpirun -np 5 ./mm mult mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=0.004219 seconds
OUTPUT NP = 6
uahcls01@dmcvlogin2:Hw3> mpirun -np 6 ./mm_mult_mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
```

```
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=0.000299 seconds
OUTPUT NP = 7
uahcls01@dmcvlogin2:Hw3> mpirun -np 7 ./mm_mult_mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=0.004195 seconds
OUTPUT NP = 8
uahcls01@dmcvlogin2:Hw3> mpirun -np 8 ./mm mult mpi 4 6 3
A matrix =
48.3962 65.3245 15.0385 72.383 25.8898 46.0265
15.4881 50.6507 6.74602 71.0055 12.2209 77.5441
61.5452 31.5127 46.8515 89.4849 70.0342 57.3195
75.4144 83.5553 91.7832 7.74197 40.0845 11.1709
B matrix =
26.5416 83.9488 86.5328
51.0444 65.3442 85.2683
76.9977 49.0015 46.6826
12.2581 99.9706 40.1026
58.6347 47.2069 4.06732
37.0919 22.9082 82.6622
C matrix =
9889.42 18581 17272.7
7979.16 14392.3 15281.2
14179 23086.6 18811.4
16193.3 19210.5 19332
time=0.000227 seconds
```

I originally ran the program with np = 2, 4, and 8 first and then proceeded to go back and run the rest of the cases after reading the homework description more closely. My connection to the dmc seemed slow on the night I ran np = 1, 3, 5, 6, and 7 which I think was the cause of the substantial increase in completion time.

Part 3 Timing Analysis

SERIAL

```
Matrix Size = 256, time = 0.107185 seconds
Matrix Size = 512, time = 0.862933 seconds
Matrix Size = 768, time = 4.647209 seconds
Matrix Size = 1024, time = 26.243087 seconds
Matrix Size = 1280, time = 46.751129 seconds
Matrix Size = 1536, time = 106.47022 seconds
Matrix Size = 1792, time = 180.9898 seconds
Matrix Size = 2048, time = 371.17542 seconds
Matrix Size = 2304, time = 490.22311 seconds
Matrix Size = 2560, time = 665.90334 seconds
Matrix Size = 2816, time = 950.09028 seconds
Matrix Size = 3072, time = 1154.8234 seconds
Matrix Size = 3328, time = 1447.418 seconds
Matrix Size = 3584, time = 2131.6487 seconds
```

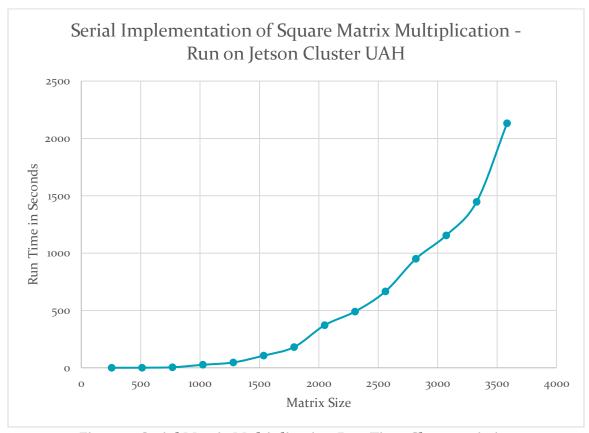


Figure 1. Serial Matrix Multiplication Run Time Characteristics

The base algorithm behaves as expected. We know that the algorithm used has an order notation of $O(n^3)$ and looking at **Figure 1** it can be seen that the run time is exponentially increasing as the matrix size increases.

PARALLEL

NP = 2

```
Matrix Size = 256, time = 0.069588 seconds
Matrix Size = 512, time = 0.510835 seconds
Matrix Size = 768, time = 2.1731 seconds
Matrix Size = 1024, time = 13.790599 seconds
Matrix Size = 1280, time = 23.50243 seconds
Matrix Size = 1536, time = 55.698489 seconds
Matrix Size = 1792, time = 97.612683 seconds
Matrix Size = 2048, time = 188.87857 seconds
Matrix Size = 2048, time = 244.05105 seconds
Matrix Size = 2560, time = 347.84285 seconds
Matrix Size = 2816, time = 474.08665 seconds
Matrix Size = 3072, time = 576.89958 seconds
Matrix Size = 3328, time = 736.07543 seconds
Matrix Size = 3584, time = 1175.1372 seconds
```

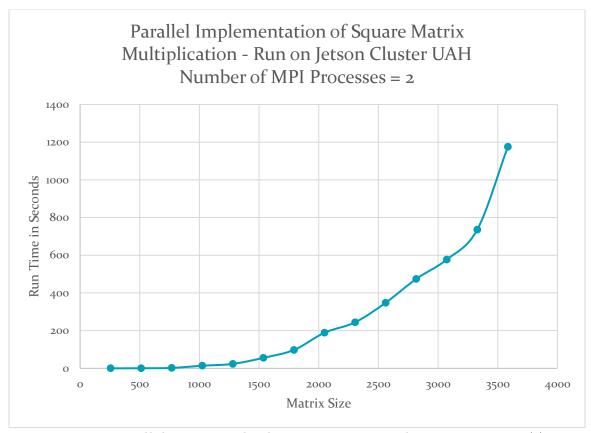


Figure 2. Parallel Matrix Multiplication Run Time Characteristics, NP(2)

Comparing **Figure 2** with **Figure 1** it seems that utilizing one more process and splitting the work evenly between the two has cut the completion time almost in half for every case, which is in line with what we would expect to see.

```
NP = 4
Matrix Size = 256, time = 0.054693 seconds
```

```
Matrix Size = 512, time = 0.289523 seconds Matrix Size = 768, time = 1.221861 seconds Matrix Size = 1024, time = 6.843977 seconds Matrix Size = 1280, time = 11.904123 seconds Matrix Size = 1536, time = 27.318558 seconds Matrix Size = 1792, time = 48.401298 seconds Matrix Size = 2048, time = 97.835424 seconds Matrix Size = 2304, time = 121.94584 seconds Matrix Size = 2560, time = 172.05514 seconds Matrix Size = 2816, time = 236.1049 seconds Matrix Size = 3072, time = 303.81367 seconds Matrix Size = 3328, time = 370.86322 seconds Matrix Size = 3584, time = 584.54626 seconds
```

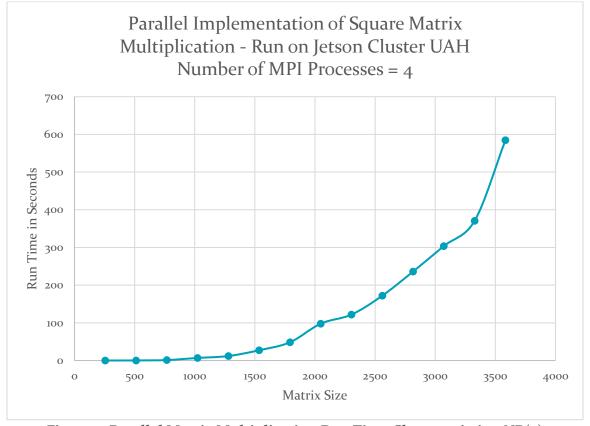


Figure 3. Parallel Matrix Multiplication Run Time Characteristics, NP(4)

NP = 8

```
Matrix Size = 256, time = 0.042516 seconds
Matrix Size = 512, time = 0.196104 seconds
Matrix Size = 768, time = 0.726653 seconds
Matrix Size = 1024, time = 3.569868 seconds
Matrix Size = 1280, time = 6.044348 seconds
Matrix Size = 1536, time = 13.962634 seconds
Matrix Size = 1792, time = 24.515898 seconds
Matrix Size = 2048, time = 49.485545 seconds
```

```
Matrix Size = 2304, time = 61.588715 seconds
Matrix Size = 2560, time = 88.130733 seconds
Matrix Size = 2816, time = 120.32 seconds
Matrix Size = 3072, time = 156.76207 seconds
Matrix Size = 3328, time = 187.06935 seconds
Matrix Size = 3584, time = 287.17158 seconds
```

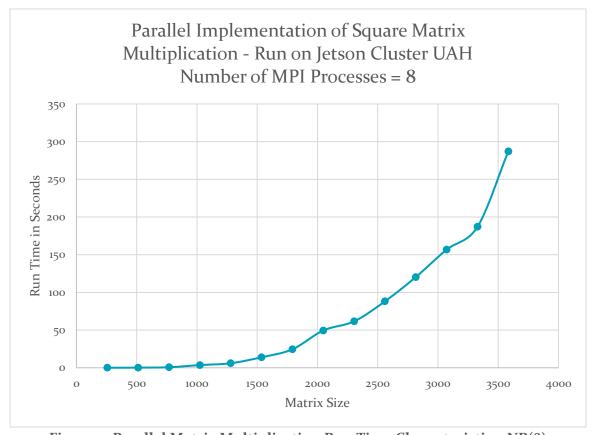


Figure 4. Parallel Matrix Multiplication Run Time Characteristics, NP(8)

Examining the above data, we can see a trend forming. Every time the number of processes is doubled we are seeing the completion time cut almost in half for each run. The above implementations are behaving just like the serial portion but because the work is being distributed and has parallelized the program, it is able to finish in a timelier manner. We can see from each graph, of each parallel run, that the run time is increasing at an exponential rate as matrix size increases which is consistent with the behavior of the serial program. The algorithm is executing $n*\frac{n*n}{p}$ times, where p is the number of processes. Therefore, the order notation for each implementation in an MPI process is $O(n^2)$ but overall the algorithm is still $O(n^3)$.

Below are the graphs for relative speed up and efficiency with regards to the serial version of the program.

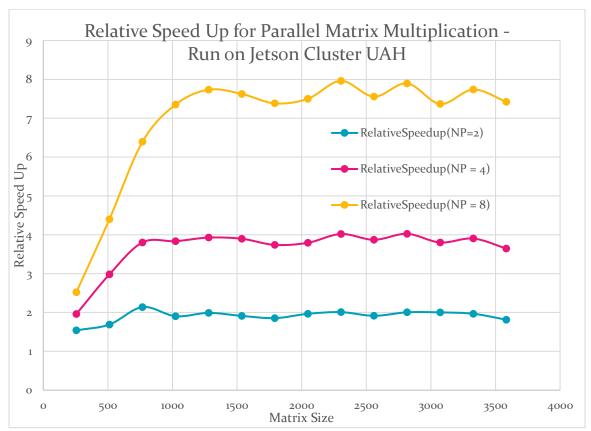


Figure 5. Parallel Matrix Multiplication Relative Speed Up, NP(2, 4, 8)

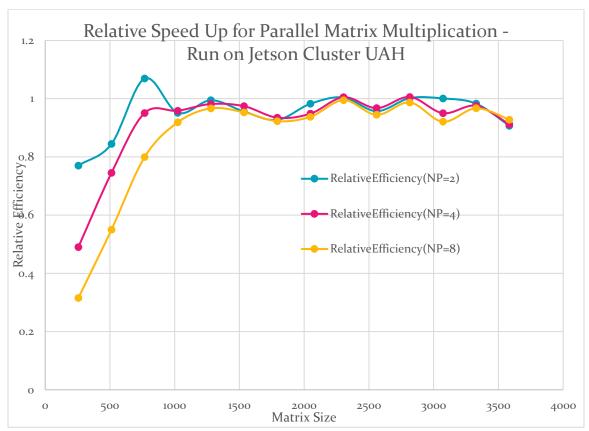


Figure 5. Parallel Matrix Multiplication Relative Efficiency, NP(2, 4, 8)

Appendix

A) SERIAL CODE using namespace std; #include <iostream> #include <iomanip> #include <sstream> #include <stdlib.h> #include <string.h> #include <sys/time.h> #define MX SZ 320 #define SEED 2397 /* random number seed */ #define MAX VALUE 100.0 $^{\prime \star}$ maximum size of array elements A, and B $^{\star \prime}$ /* copied from mpbench */ #define TIMER CLEAR (tv1.tv sec = tv1.tv usec = tv2.tv sec = tv2.tv usec = #define TIMER START gettimeofday(&tv1, (struct timezone*)0) #define TIMER ELAPSED ((tv2.tv usec-tv1.tv usec)+((tv2.tv sectv1.tv sec) *1000000)) #define TIMER STOP gettimeofday(&tv2, (struct timezone*)0) struct timeval tv1, tv2; This declaration facilitates the creation of a two dimensional dynamically allocated arrays (i.e. the lxm A array, the mxn B array, and the lxn C array). It allows pointer arithmetic to be applied to a single data stream that can be dynamically allocated. To address the element at row x, and column y you would use the following notation: A(x,y), B(x,y), or C(x,y), respectively. Note that this differs from the normal C notation if A were a two dimensional array of A[x][y] but is still very descriptive of the data structure. */ float *a, *b, *c; #define A(i,j) *(a+i*dim_m+j) #define B(i,j) *(b+i*dim n+j) #define C(i,j) *(c+i*dim n+j) Routine to retrieve the data size of the numbers array from the command line or by prompting the user for the information void get index size(int argc,char *argv[],int *dim l,int *dim m,int *dim n) { if(argc!=2 && argc!=4) { cout<<"usage: mm mult serial [l dimension] <m dimension n dimmension>" << endl; exit(1);else { if (argc == 2) { *dim l = *dim n = *dim m = atoi(argv[1]);else {

*dim l = atoi(argv[1]);

```
*dim m = atoi(argv[2]);
         *dim n = atoi(argv[3]);
      }
   if (*dim 1 \le 0 \mid | *dim n \le 0 \mid | *dim m \le 0) {
      cout<<"Error: number of rows and/or columns must be greater than 0"
          << endl;
      exit(1);
}
  Routine that fills the number matrix with Random Data with values
  between 0 and MAX VALUE
  This simulates in some way what might happen if there was a
  single sequential data acquisition source such as a single file
void fill_matrix(float *array,int dim_m,int dim_n)
   int i,j;
   for(i=0;i<dim m;i++) {</pre>
     for (j=0;j<dim_n;j++) {
        array[i*dim n+j]=drand48()*MAX VALUE;
   }
}
  Routine that outputs the matrices to the screen
void print matrix(float *array,int dim m,int dim n)
   int i,j;
   for(i=0;i<dim m;i++) {</pre>
      for (j=0; j < \dim n; j++) {
        cout << array[i*dim n+j] << " ";</pre>
      cout << endl;
}
  MAIN ROUTINE: summation of a number list
int main( int argc, char *argv[])
   float dot_prod;
   int dim 1, dim n, dim m;
   int i,j,k;
   get matrix sizes
   get_index_size(argc,argv,&dim_l,&dim_m,&dim_n);
   // dynamically allocate from heap the numbers in the memory space
   // for the a,b, and c matrices
   a = new (nothrow) float[dim l*dim m];
  b = new (nothrow) float[dim_m*dim_n];
  c = new (nothrow) float[dim_l*dim_n];
   if(a==0 || b==0 || c==0) {
```

```
cout <<"ERROR: Insufficient Memory" << endl;</pre>
  exit(1);
 }
   initialize numbers matrix with random data
 srand48(SEED);
 fill matrix(a,dim l,dim m);
 fill matrix(b, dim m, dim n);
  output numbers matrix
*/
// cout << "A matrix =" << endl;
// print_matrix(a,dim_l,dim_m);
// cout << endl;
//cout << "B matrix =" << endl;
 //print matrix(b,dim m,dim n);
 //cout << endl;</pre>
 Start recording the execution time
 TIMER CLEAR;
 TIMER START;
 // multiply local part of matrix
 for (i=0;i<dim 1;i++) {
    for (j=0; j<\overline{dim} n; j++) {
       dot_prod = \overline{0.0};
       for (k=0; k<dim_m; k++) {
          dot_prod += A(i,k)*B(k,j);
       C(i,j) = dot prod;
 }
   stop recording the execution time
 TIMER_STOP;
 //cout << "C matrix =" << endl;
 //print_matrix(c,dim_l,dim_n);
 //cout << endl;
 cout << "time=" << setprecision(8) << TIMER_ELAPSED/1000000.0</pre>
      << " seconds" << endl;
```

}

B) PARALLEL CODE

```
/* Matrix Matrix Multiplication Program Example -- serial version */
/* September 2016 -- B. Earl Wells -- University of Alabama */
                      in Huntsville
// mm mult serial.cpp
// compilation:
// gnu compiler
//
     g++ mm mult serial.cpp -o mm mult serial -O3 -lm
// Note: to compile a parallel MPI program version which is named
// mm mult mpi.cpp
//
   then execute the following command
//
       gnu compiler
//
         mpic++ mm mult mpi.cpp -o mm mult MPI gnu -lm -O3
/*
  This program is designed to perform matrix matrix multiplication
  A \times B = C, where A is an lxm matrix, B is a m \times n matrix and
  {\tt C} is a 1 x n matrix. The program is designed to be a template
  serial program that can be expanded into a parallel multiprocess
  and/or a multi-threaded program.
  The program randomly assigns the elements of the A and B matrix
  with values between 0 and a MAX VALUE. It then multiples the
  two matrices with the result being placed in the C matrix.
  The program prints out the A, B, and C matrices.
  The program is executed using one or three command line parameters.
  These parameters represent the dimension of the matrices. If only
  one parameter is used then then it is assumed that square matrices are
  to be created and multiplied together that have the specified
  dimension. In cases where three command line parameters are entered
  then the first parameter is the {\bf l} dimension, the second the {\bf m}, and
  the third is the n dimension.
  To execute:
  mm mult serial [l parameter] <m parameter n parameter>
  Editted by Kyle Ray
  October 12, 2017
  Homework #3
  Changed the serial version of the program to utilize the MPI library to
  perform the matrix multiplication in parallel on a number of MPI processes
  passed in on the command line.
using namespace std;
#include <iostream>
#include <iomanip>
#include <sstream>
#include <stdlib.h>
#include <string.h>
#include <sys/time.h>
//#include <time.h>
#include <mpi.h>
#define MX SZ 320
                         /* random number seed */
#define SEED 2397
#define MAX VALUE 100.0 ^{\prime \star} maximum size of array elements A, and B ^{\star \prime}
```

```
/* copied from mpbench */
#define TIMER CLEAR
                        (tv1.tv sec = tv1.tv usec = tv2.tv sec = tv2.tv usec =
0)
#define TIMER START
                        gettimeofday(&tv1, (struct timezone*)0)
#define TIMER ELAPSED
                         ((tv2.tv usec-tv1.tv usec)+((tv2.tv sec-
tv1.tv sec) *1000000))
#define TIMER STOP
                        gettimeofday(&tv2, (struct timezone*)0)
struct timeval tv1, tv2;
// Defines so that I can compile the code in visual studio
//#define srand48(s) srand(s)
//#define drand48() (((double)rand())/((double)RAND MAX))
This declaration facilitates the creation of a two dimensional
dynamically allocated arrays (i.e. the lxm A array, the mxn B
array, and the lxn C array). It allows pointer arithmetic to
be applied to a single data stream that can be dynamically allocated.
To address the element at row x, and column y you would use the
following notation: A(x,y), B(x,y), or C(x,y), respectively.
Note that this differs from the normal C notation if A were a
two dimensional array of A[x][y] but is still very descriptive
of the data structure.
* /
float *a, *b, *c;
float *group a, *group c;
#define A(i, \overline{j}) * (a+i*d\overline{m} m+j)
\#define B(i,j) * (b+i*dim n+j)
#define C(i,j) *(c+i*dim n+j)
   Routine to retrieve the matrix dimensions of the arrays from the
   command line.
void get index size(int argc, char *argv[], int *dim 1, int *dim m, int *dim n,
int rank) {
  if (argc != 2 && argc != 4) {
    if (rank == 0)
      cout << "usage: mm mult serial [l dimension] <m dimension n dimmension>"
        << endl;
      MPI Finalize();
      exit(1);
    }
  else {
    if (argc == 2) {
      *dim l = *dim n = *dim m = atoi(argv[1]);
    else {
      *dim_l = atoi(argv[1]);
      *dim m = atoi(argv[2]);
      *dim n = atoi(argv[3]);
  if (rank == 0)
    if (*dim 1 \le 0 | | *dim n \le 0 | | *dim m \le 0) {
      cout << "Error: number of rows and/or columns must be greater than 0"
        << endl;
      MPI Finalize();
```

```
exit(1);
   }
 }
}
  Routine that fills the number matrix with Random Data with values
  between 0 and MAX VALUE
  This simulates in some way what might happen if there was a
  single sequential data acquisition source such as a single file
void fill matrix(float *array, int dim m, int dim n)
 int i, j;
 for (i = 0; i < dim m; i++) {
   for (j = 0; j < \overline{dim} n; j++) {
     array[i*dim n + j] = drand48()*MAX VALUE;
 }
}
  Routine that outputs the matrices to the screen
void print matrix(float *array, int dim m, int dim n)
  int i, j;
 for (i = 0; i < dim m; i++) {
   for (j = 0; j < \dim n; j++) {
     cout << array[i*dim_n + j] << " ";
   cout << endl;</pre>
  }
}
// Routine to get the base number of multiplies for a ceratin process
int getBaseMult(int num mults, int numtasks, int rank)
 // Calculate the base number of multiplies for each task
 int base = num mults / numtasks;
  // Calculate the extra if it is not an even distribution
 int extra = num mults % numtasks;
  // If there are any extra assign them to the first tasks up to rank of extra
 if (extra != 0)
   // If rank is less than the number of extra items, then this process gets
an extra multiply to process
   if (rank < extra)</pre>
     base = (num mults / numtasks) + 1;
 return base;
}
/* ONE-TO-ALL SCATTER ROUTINE
Routine to divide the number of rows to each process based on the number
of multiplies that each process must perform. Each process will receive the
row up to number of multiplies in a sequential order. This method will also
keep up
```

```
with the current column slider and pass that to the process so it knows where
start it's set of multiplications.
The partial arrays for each process are stored in the group a array.
void scatter(float* a, float* b, float *group a, int root, int rank, int
numtasks,
  int dim 1, int dim m, int dim n, int* start column)
 MPI Status status;
 int type = 234;
  // How many multiplies will the entire operation require?
  int num mults = 0;
  if (\dim 1 == 1)
   num \overline{\text{mults}} = \text{dim } n;
  else \overline{i}f (dim_n == 1)
   num_mults = dim_l;
    num mults = dim l * dim n;
  // Variables to keep up with what row and column we are reading from
  int begin row = 0;
  int begin column = 0;
  // Root process does all of the work
  if (rank == root)
    // Loop over the MPI tasks
    for (int mpi task = 0; mpi task < numtasks; mpi task++)</pre>
      // Get the number of multiplies
      int base = getBaseMult(num_mults, numtasks, mpi_task);
      // Each task gets at least one row to work with
      int num rows = 1;
      // Variables to keep up with navigating the matrices
      int count = base;
      int curr col = begin column;
      // Calculate the number of rows that we need to send to the process.
      while (count > 0)
        // Get the current distance from the end of dim n
        int diff = dim_n - curr_col;
        // If we have more multiplies to process then left for this row, we
must add another row
        if (count > diff)
          num rows++;
          count -= diff;
          curr col = 0;
          continue;
        count -= dim n;
      }
      // If the current loop iteration doesn't correspond to the root then we
must send the column and number of rows to the process
```

```
if (mpi task != root)
        MPI_Send(&begin_column, 1, MPI_INT, mpi_task, type, MPI COMM WORLD); //
send the column slider location
       MPI_Send(&num_rows, 1, MPI_INT, mpi_task, type, MPI_COMM_WORLD); //
send the number of rows for correct sizing
      // Local Buffer variables
      float* local a = new float[num rows*dim m];
      // Reset our column check and multiply count
      curr col = begin column;
      count = base;
      // Row assignment
      for (int r = 0; r < num rows; r++)
        // Fill up the local matrix with values from the main A matrix
        for (int t = 0; t < dim m; t++)
          // If this is the root, go ahead and store it in the buffer
          if (mpi task == root)
           group a[r*dim m + t] = a[begin row*dim m + t];
          else // Store it in the local that will be sent to the other
processes
           local a[r*dim m + t] = a[begin row*dim m + t];
          }
        }
        // Calculate the distance from the end of dim n for this set of row
multiplications
        int diff = dim n - curr col;
        // If we still have more to process we must add the next row to this
processes variables
        if (diff <= count)
         begin row++;
          count -= diff;
          curr col = 0;
       }
      }
      // Column Assignment
      // This logic will keep up with what column each process should start
performing calculations
      // Base is the updated base number of multiplies each process must
perform
      int temp_column = base % dim_n; // Get the leftover after applying number
of multiplies
      begin column += temp column; // Update the current column index
      if (begin column >= dim n) // Account for wrapping around dim n
       begin column = begin column % dim n;
      // Send the data to the other processes
      if (mpi_task != root)
```

```
MPI Send(local a, num rows*dim m, MPI FLOAT, mpi task, type,
MPI COMM WORLD);
     }
      delete[] local a;
  else
    // Calculate the base number of multiplies for each task
    int base = getBaseMult(num mults, numtasks, rank);
    // I must send the number of rows as well
    int num rows = 0;
    // Receive smaller arrays as well as the starting dim n column from the
root process
    MPI_Recv(start_column, 1, MPI_INT, root, type, MPI_COMM_WORLD, &status);
    MPI_Recv(&num_rows, 1, MPI_INT, root, type, MPI_COMM_WORLD, &status);
    MPI_Recv(group_a, num_rows*dim_m, MPI_FLOAT, root, type, MPI COMM WORLD,
&status);
}
// All to one gather routine
// Each process will send their calculated sub matrix back to the root process
void gather(float* c, float* group c, int num mults, int root, int rank, int
numtasks, int dim 1, int dim n)
 MPI Status status;
 int type = 123;
  // Calculate the base number of multiplies for each task
  int base = getBaseMult(num mults, numtasks, rank);
  if (rank == root)
    int curr ind = 0;
    // Piece back together the matrix
    for (int mpi task = 0; mpi task < numtasks; mpi task++)</pre>
      if (mpi task == root)
        for (int i = 0; i < base; i++)
          // Copy what the root has
          c[curr_ind] = group_c[i];
          curr ind++;
        }
      else
        // Receive from the processes
        // Calculate the base number of multiplies for each task
        int base = getBaseMult(num mults, numtasks, mpi task);
        float* temp = new float[base];
        MPI Recv(temp, base, MPI FLOAT, mpi task, type, MPI COMM WORLD,
&status);
```

```
for (int i = 0; i < base; i++)
          //cout << "Group_c Item " << temp[i] << endl;
          c[curr_ind] = temp[i];
         curr ind++;
        delete[] temp;
    }
  }
  else
    // Send the matrix to the root
   MPI Send(group c, base, MPI FLOAT, root, type, MPI COMM WORLD);
}
  MAIN ROUTINE: summation of a number list
int main(int argc, char *argv[])
 float dot prod;
 int dim_l, dim_n, dim_m;
 int i, \bar{j}, k;
 int num mults, group size, num group;
 int numtasks, rank, num;
 int start column;
 MPI Status status;
 // Main Routine
 MPI Init(&argc, &argv); // initalize MPI environment
 MPI Comm size (MPI COMM WORLD, &numtasks); // get total number of MPI
processes
 MPI_Comm_rank(MPI_COMM_WORLD, &rank); // get unique task id number
 // get matrix sizes
  get index size(argc, argv, &dim 1, &dim m, &dim n, rank);
  // The root process fills the matrices and then passes them to the othe
processes
 if (rank == 0)
    // dynamically allocate from heap the numbers in the memory space
   // for the a,b, and c matrices
    a = new (nothrow) float[dim l*dim m];
   b = new (nothrow) float[dim m*dim n];
    c = new (nothrow) float[dim_l*dim_n];
    if (a == 0 || b == 0 || c == 0)
     cout << "ERROR: Insufficient Memory 1" << endl;</pre>
     MPI Abort (MPI COMM WORLD, 1);
       initialize numbers matrix with random data
```

```
* /
   srand48 (SEED);
   fill_matrix(a, dim_l, dim_m);
    fill_matrix(b, dim_m, dim_n);
     output numbers matrix
    //cout << "A matrix =" << endl;
   //print matrix(a, dim l, dim m);
   //cout << endl;
   //cout << "B matrix =" << endl;
   //print matrix(b, dim m, dim n);
   //cout << endl;
 else
   b = new (nothrow) float[dim_m*dim_n];
  // Start recording the execution time
 if (rank == 0)
   TIMER CLEAR;
   TIMER_START;
 MPI Bcast(b, dim m*dim n, MPI FLOAT, 0, MPI COMM WORLD);
 // broad cast the data size, which is really the number of multiplies
 if (\dim 1 == 1)
  num \overline{\text{mults}} = \text{dim } n;
 else if (\dim_n == 1)
   num mults = dim 1;
   num mults = dim l * dim n;
 int base = getBaseMult(num mults, numtasks, rank);
 // Each process has a local array set for local calculations
 group a = new (nothrow) float[dim l*dim m];
 group c = new (nothrow) float[dim l*dim n];
 start column = 0;
 if (group_a == 0 \mid | group_c == 0)
   cout << "ERROR: Insufficient Memory 2" << endl;</pre>
   MPI Abort (MPI COMM WORLD, 1);
 // Scatter the Data
 // The root process needs to scatter the correct amount of data to each
 scatter(a, b, group a, 0, rank, numtasks, dim 1, dim m, dim n,
&start column);
 // Each process will start working on the data here
 int startIndex = 0;
 int row = 0;
 int col = start_column;
```

```
for(int i = 0; i < base; i++)
   group c[i] = 0;
    for (\overline{int} j = 0; j < \dim m; j++)
     group_c[i] += group_a[dim_m*row + j] * b[j*dim_n + col];
    // Keep up with the column so we know when to bump the row
    if (start column != 0 \&\& (start column % (dim n - 1) == 0))
     start column = 0;
     row++;
   else if ((start column + 1) != dim n) // can't exceed the dim n for current
column
     start_column++;
   col = start column;
  // Gather
 gather(c, group c, num mults, 0, rank, numtasks, dim l, dim n);
 Start recording the execution time
 // TIMER CLEAR;
 // TIMER START;
     stop recording the execution time
  //TIMER STOP;
 if (rank == 0)
   TIMER STOP;
   //cout << "C matrix =" << endl;
   //print matrix(c, dim l, dim n);
   //cout << endl;</pre>
   cout << "time=" << setprecision(8) << TIMER ELAPSED/1000000.0</pre>
       << " seconds" << endl;
  }
 // if (rank == 0)
  // cout << "Made it to the cleanup" << endl;</pre>
  // Clear out memory
 if (rank == 0)
   delete[] a;
   delete[] b;
   delete[] c;
 delete[] group_a;
  delete[] group_c;
```

```
// Terminate MPI Program -- perform necessary MPI housekeeping
// clear out all buffers, remove handlers, etc.
MPI_Finalize();
```

C) RUN TIME CHARACTERISTICS ON ONE GRAPH

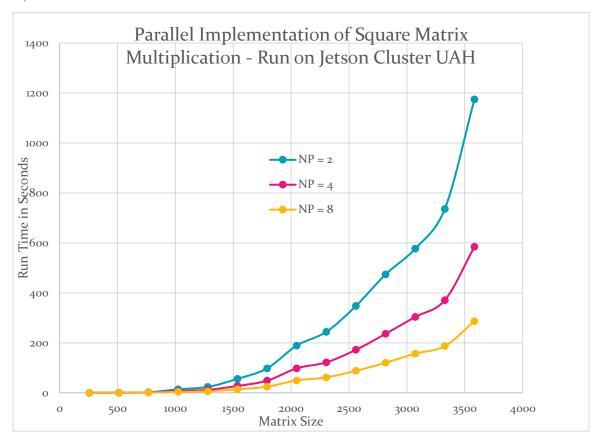


Figure 6. Parallel Matrix Multiplication Run Time Characteristics, NP(2,4,8)