

Hands on - 6

3) The average run time complexity of the non random pivot version of Quicksort

Let the $T(n)$ be time complexity of sorting array size n .

→ The pivot divides array into two approximately equal halves.

$$T(n) = T(n/2) + T(n/2) + O(n)$$

$$T(n) = 2T(n/2) + O(n)$$

by masters theorem.

$$f(n) = O(n)$$

$$f(n) = O(n^{\log_b a - \epsilon}) \quad \forall \text{ for } \epsilon > 0$$

$$T(n) = O(n^{\log_b a}) = O(n^{\log_2 2}) = O(n)$$

recursive depth of tree is $\log_2 n$. Total becomes

$$T(n) = O(n \log n)$$

∴ Avg case: $O(n \log n)$