Ad Hoc Networks

The IMPORTANT framework for analyzing the Impact of Mobility on Performance Of RouTing protocols for Adhoc NeTworks

Fan Bai [a,*], Narayanan Sadagopan [b], Ahmed Helmy [a]

[a] *Department of Electrical Engineering, University of Southern California, Los Angeles, CA 90089, USA*
[b] *Department of Computer Science, University of Southern California, Los Angeles, CA 90089, USA*

Abstract

A Mobile Ad hoc Network (MANET) is a collection of wireless mobile nodes forming a temporary network without using any existing infrastructure. Since not many MANETs are currently deployed, research in this area is mostly simulation based. Random Waypoint is the commonly used mobility model in these simulations. Random Waypoint is a simple model that may be applicable to some scenarios. However, we believe that it is not sufficient to capture some important mobility characteristics of scenarios in which MANETs may be deployed. Our framework aims to evaluate the impact of different mobility models on the performance of MANET routing protocols. We propose various protocol independent metrics to capture interesting mobility characteristics, including spatial and temporal dependence and geographic restrictions. In addition, a rich set of parameterized mobility models is introduced including Random Waypoint, Group Mobility, Freeway and Manhattan models. Based on these models several 'test-suite' scenarios are chosen carefully to span the metric space. We demonstrate the utility of our test-suite by evaluating various MANET routing protocols, including DSR, AODV and DSDV. Our results show that the protocol performance may vary drastically across mobility models and performance rankings of protocols may vary with the mobility models used. This effect can be explained by the interaction of the mobility characteristics with the connectivity graph properties. Finally, we attempt to decompose the reactive routing protocols into mechanistic "building blocks" to gain a deeper insight into the performance variations across protocols in the face of mobility.
© 2003 Elsevier B.V. All rights reserved.

Keywords: Mobile ad hoc network; Performance; Mobility; Routing protocols

1. Introduction

A Mobile Ad hoc NETwork (MANET) is a collection of wireless nodes communicating with each other in the absence of any infrastructure. Classrooms, battlefields and disaster relief activities are a few scenarios where MANETs can be used. MANET research is gaining ground due to the ubiquity of small, inexpensive wireless communicating devices. Since, not many MANETs have been deployed, most of this research is simulation based. These simulations have several parameters including the mobility model and the communicating traffic pattern. In this paper, we

---
* Corresponding author. Tel.: +1-213-740-2685; fax: +1-213-740-9803.
E-mail addresses: fbai@usc.edu (F. Bai), narayans@usc.edu (N. Sadagopan), helmy@usc.edu (A. Helmy).

focus on the impact of mobility models on the performance of MANET routing protocols. We acknowledge that the communicating traffic pattern also has a significant impact on the routing protocol performance and merits a study on its own. However, as in most studies in this area, in order to isolate the effect of mobility, we fix the communicating traffic pattern to consist of randomly chosen source–destination pairs with long enough session times.

Mobility pattern, in many previous studies was assumed to be Random Waypoint. In the current network simulator (ns-2) distribution, the implementation of this mobility model is as follows: at every instant, a node randomly chooses a destination and moves towards it with a velocity chosen uniformly randomly from $[0, V_{max}]$, where V_{max} is the maximum allowable velocity for every mobile node [1]. Most of the simulations using the Random Waypoint model are based on this standard implementation. For the rest of the paper, we refer to this basic implementation as the Random Waypoint model.

In the future, MANETs are expected to be deployed in myriads of scenarios having complex node mobility and connectivity dynamics. For example, in a MANET on a battlefield, the movement of the soldiers will be influenced by the commander. In a city-wide MANET, the node movement is restricted by obstacles or maps. The node mobility characteristics are very application specific. Widely varying mobility characteristics are expected to have a significant impact on the performance of the routing protocols like DSR [5], DSDV [6] and AODV [7]. Random Waypoint is a well-designed and commonly used mobility model, but we find it is insufficient to capture those characteristics, such as

1. Spatial dependence of movement among nodes.
2. Temporal dependence of movement of a node over time.
3. Existence of barriers or obstacles constraining mobility.

In this study, we focus on the impact of the above-mentioned mobility characteristics on protocol performance. While doing so, we propose a generic framework to systematically analyze the impact of mobility on the performance of routing protocols for MANETs. This analysis attempts to answer the following questions:

1. Whether and to what degree mobility affects routing protocol performance?
2. If the answer to 1 is yes, why?
3. If the answer to 1 is yes, how?

To answer *Whether*, the framework evaluates the performance of these routing protocols over different mobility patterns that capture some of the characteristics listed above. The mobility models used in our study include the Random Waypoint, Group Mobility [8], Freeway and Manhattan. To answer *Why*, we propose some protocol independent metrics such as mobility metrics and connectivity graph metrics. Mobility metrics aim to capture some of the aforementioned mobility characteristics. Connectivity graph metrics aim to study the effect of different mobility patterns on the connectivity graph of the mobile nodes. It has also been observed in previous studies that under a given mobility pattern, routing protocols like DSR, DSDV and AODV perform differently [9–11]. This is possibly because each protocol differs in the basic mechanisms or ''building blocks'' it uses. For example, DSR uses route discovery, while DSDV uses periodic updates. To answer *How*, we want to investigate the effect of mobility on some of these ''building blocks'' and how they impact the protocol performance as a ''whole''.

In order to conduct our research and answer the above questions systematically, we propose a framework for analyzing the Impact of Mobility on the Performance Of RouTing protocols in Adhoc NeTworks (IMPORTANT). Through this framework we illustrate how modeling mobility is important in affecting routing performance and understanding the mechanism of ad hoc routing protocols. As shown in Fig. 1, our framework focuses on the following aspects: mobility models, the metrics for mobility and connectivity graph characteristics, the potential relationship between mobility and routing performance and the analysis of impact of mobility on building blocks of ad hoc routing protocols.
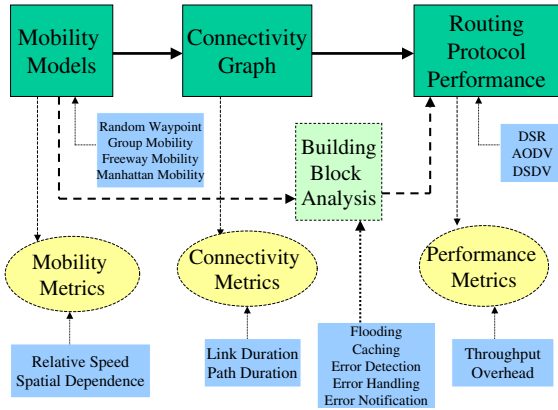
Fig. 1. IMPORTANT framework.

The rest of this paper is organized as follows. Section 2 gives a brief description of the related work and elaborates our contribution. Section 3 discusses some limitations of the Random Waypoint model and motivates part of our framework. Section 4 presents our proposed metrics to capture characteristics of mobility and the connectivity graph between the mobile nodes. Section 5 describes the mobility models used and introduces two new models, the Freeway mobility model and the Manhattan mobility model. Results of our simulation experiments are presented and discussed in Section 6. The analysis of the impact of mobility on protocol building blocks is discussed in Section 7. Finally, our conclusions from this study and planned future work are listed in Section 8.

2. Related work

Extensive research has been done in modeling mobility for MANETs. In this section, we mainly focus on experimental research in this area. This research can be broadly classified as follows based on the methodology used.

2.1. Random Waypoint based performance comparisons

Much of the initial research was based on using Random Waypoint as the underlying mobility model and Constant Bit Rate (CBR) traffic consisting of randomly chosen source–destination pairs as the traffic pattern. Routing protocols like DSR [5], DSDV [6], AODV [7] and TORA [12] were mainly evaluated based on the following metrics: packet delivery ratio (ratio of the number of packets received to the number of packets sent) and routing overhead (number of routing control packets sent). Ref. [9] concluded that on-demand protocols such as DSR and AODV performed better than table driven ones such as DSDV at high mobility rates, while DSDV performed quite well at low mobility rates. Ref. [10] performed a comparison study of the two on-demand routing protocols: DSR and AODV, using the performance metrics of packet delivery ratio and end to end delay. It observed that DSR outperforms AODV in less demanding situations, while AODV outperforms DSR at heavy traffic load and high mobility. However, the routing overhead of DSR was found to be lesser than that of AODV. In the above studies, focus was given on performance evaluation, while parameters investigated in the mobility model were change of maximum velocity and pause time. In our work, however, we design our test suites very carefully to pick scenarios that span a much larger set of mobility characteristics. Not only do we use Random Waypoint but also other mobility models such as Reference Point Group Mobility (RPGM) [8], Freeway and Manhattan in our evaluation of the performance of routing protocols.

2.2. Scenario based performance comparisons

Random Waypoint is a simple model that is easy to analyze and implement. This has probably been the main reason for the widespread use of this model for simulations. Realizing that Random Waypoint is too general a model, recent research has started focusing on alternative mobility models and protocol independent metrics to characterize them. Ref. [13] conducted a scenario based performance analysis of the MANET protocols. It proposed models for a few "realistic" scenarios such as a conference, event coverage and disaster relief. To differentiate between scenarios used, the study introduced the relative motion of the mobile nodes as a mobility metric. Their conclusions

about the performance of proactive and reactive protocols were similar to [9]. Ref. [11] used a mobility model in which each node computes its next position based on a probability distribution. This model does not allow significant changes in direction between successive instants. It concluded that proactive protocols perform better than reactive ones in terms of packet delivery ratio and end-to-end delay. However, reactive protocols were seen to incur a lower routing overhead. Ref. [8] introduced the RPGM model, which is one of the mobility models used in this study. Rate of link changes was used to characterize a few group mobility patterns as well as Random Waypoint. It observed that the rate of link change for Random Waypoint was higher than that for RPGM. From experiments, it observed that protocols like AODV, DSDV and HSR [14] perform worse with Random Waypoint than with RPGM. Thus, it concluded that mobility models do matter and it is not sufficient to simulate protocols with only the "Random Walk" like models. Ref. [15] proposed a mobility framework that consisted of a Mobility Vector model which can be used to generate "realistic" movement patterns used in several varied applications. It proposed the Displacement Measure that is a normalization of the actual distance traveled by the geographic displacement as a metric to evaluate the different movement patterns including those generated by Random Waypoint, Random Walk, RPGM and Mobility Vector models. By experiments, it observed that Random Waypoint and Random Walk produced higher Displacement Measure as compared to the Mobility Vector model. It studied the effect of transmission range on throughput across different mobility models and concluded that as the transmission range is increased, the rate of link changes decreased and the throughput for all protocols increased. However, the link change rate does not seem to vary greatly across the different mobility models. As far as routing overhead was concerned, Mobility Vector was seen to produce a worse overhead than Random Waypoint. Our study is also framework based. However, we do not aim to provide a generic mobility model from which all "realistic" mobility patterns can be derived. Rather, our framework aims at systematically

studying the effect of mobility per se on performance of MANET routing protocols. The contributions of our proposed framework are threefold:

1. Focus on mobility characteristics such as spatial dependence, geographic restrictions and temporal dependence. Define mobility metrics that capture these characteristics. Choose mobility models that span the metric space and use them to evaluate the performance of routing protocols.
2. Define connectivity graph metrics. Study the interaction of mobility metrics and connectivity graph metrics and its effect on protocol performance.
3. Analyze the reasons for the differences in protocol performance as a "whole" by investigating the effect of mobility on "parts" that build the protocol.

3. Limitations of Random Waypoint

Random Waypoint model was introduced in [9] and is among the most commonly used mobility models in the MANET research community. In this model, at every instant, each mobile node chooses a random destination and moves towards it with a speed uniformly distributed in $[0, V_{max}]$, where V_{max} is the maximum allowable speed for a node. After reaching the destination, the node stops for a duration defined by the "pause time" parameter. After this duration, it again chooses a random destination and repeats the whole process again until the simulation ends.

The Random Waypoint model is widely accepted mainly due to its simplicity of implementation and analysis. However, we observe that the basic Random Waypoint model as used in most of the simulations is insufficient to capture the following mobility characteristics:

1. *Temporal dependency:* Due to physical constraints of the mobile entity itself, the velocity of mobile node will change continuously and gently instead of abruptly, i.e. the current velocity is dependent on the previous velocity. However, intuitively, the velocities at two different

time slots are independent in the Random Waypoint model.

2. *Spatial dependency:* The movement pattern of a mobile node may be influenced by and correlated with nodes in its neighborhood. In Random Waypoint, each mobile node moves independently of others.

3. *Geographic restrictions:* In many cases, the movement of a mobile node may be restricted along the street or a freeway. A geographic map may define these boundaries.

In our study, we focus on the above-mentioned characteristics. In the next section, we formally define metrics to capture some of these characteristics.

4. Metrics

To quantitatively and qualitatively analyze the impact of mobility on routing protocol performance, we make use of several protocol independent metrics and protocol performance metrics. The protocol independent metrics attempt to extract the characteristics of mobility and the connectivity graph between the mobile nodes. These metrics are then used to explain the impact of mobility on the protocol performance metrics. Those metrics can be broadly classified as

1. Mobility metrics.
2. Connectivity graph metrics.
3. Protocol performance metrics.

4.1. Terminology

Before formally defining the metrics, we introduce some basic terminology that will be used later in the paper:

1. $\overrightarrow{V_i}(t)$: velocity vector of node i at time t.
2. $|\overrightarrow{V_i}(t)|$: speed of node i at time t.
3. $\theta_i(t)$: angle made by $\overrightarrow{V_i}(t)$ at time t with the X-axis.
4. $\overrightarrow{a_i}(t)$: acceleration vector of node i at time t.
5. $x_i(t)$: X co-ordinate of node i at time t.
6. $y_i(t)$: Y co-ordinate of node i at time t.

7. $D_{i,j}(t)$: Euclidean distance between nodes i and j at time t.
8. $\mathrm{RD}(\vec{a}(t), \vec{b}(t'))$: relative direction (RD) (or cosine of the angle) between the two vectors $\vec{a}(t)$, $\vec{b}(t')$ is given by $\frac{\vec{a}(t) \cdot \vec{b}(t')}{|\vec{a}(t)| * |\vec{b}(t')|}$.
9. $\mathrm{SR}(\vec{a}(t), \vec{b}(t'))$: speed ratio (SR) between the two vectors $\vec{a}(t)$, $\vec{b}(t')$ is given by $\frac{\min |\vec{a}(t)|, |\vec{b}(t')|}{\max |\vec{a}(t)|, |\vec{b}(t')|}$.
10. R: transmission range of a mobile node.
11. N: number of mobile nodes.
12. T: simulation time.
13. *random( )*: returns a value uniformly distributed in the interval $[-1, 1]$.

4.2. Mobility metrics

We propose these metrics to differentiate the various mobility patterns used in our study. The basis of differentiation is the extent to which a given mobility pattern captures the characteristics of spatial dependence, temporal dependence and geographic restrictions. In addition to these metrics, we also use the relative speed metric that differentiates mobility patterns based on relative motion. This metric was proposed in [13].

1. *Degree of spatial dependence:* It is a measure of the extent of similarity of the velocities of two nodes that are not too far apart. Formally,

$$D_{\mathrm{spatial}}(i, j, t) = \mathrm{RD}(\overrightarrow{v_i}(t), \overrightarrow{v_j}(t)) \\ * \mathrm{SR}(\overrightarrow{v_i}(t), \overrightarrow{v_j}(t)).$$

The value of $D_{\mathrm{spatial}}(i, j, t)$ is high when the nodes i and j travel in more or less the same direction and at almost similar speeds. However, $D_{\mathrm{spatial}}(i, j, t)$ decreases if the relative direction or the speed ratio decreases.

As it is rare for a node's motion to be spatially dependent on a far off node, we add the condition that

$$D_{i,j}(t) > c_1 * R \quad \Rightarrow \quad D_{\mathrm{spatial}}(i, j, t) = 0,$$

where $c_1 > 0$ is a constant which will be determined during our experiments in 6.

Average degree of spatial dependence: It is the value of $D_{\mathrm{spatial}}(i, j, t)$ averaged over node pairs and time instants satisfying certain condition. Formally,

$$\overline{D}_{\text{spatial}} = \frac{\sum_{t=1}^{T} \sum_{i=1}^{N} \sum_{j=i+1}^{N} D_{\text{spatial}}(i,j,t)}{P},$$

where P is the number of tuples (i,j,t) such that $D_{\text{spatial}}(i,j,t) \neq 0$. Thus, if mobile nodes move independently of one another, then the mobility pattern is expected to have a smaller value for $\overline{D}_{\text{spatial}}$. On the other hand, if the node movement is co-ordinated by a central entity, or influenced by nodes in its neighborhood, such that they move in similar directions and at similar speeds, then the mobility pattern is expected to have a higher value for $\overline{D}_{\text{spatial}}$.

2. *Degree of temporal dependence:* It is a measure of the extent of similarity of the velocities of a node at two time slots that are not too far apart. It is a function of the acceleration of the mobile node and the geographic restrictions. Formally,

$$D_{\text{temporal}}(i,t,t') = \text{RD}(\vec{v_i}(t), \vec{v_i}(t')) \\ * \text{SR}(\vec{v_i}(t), \vec{v_i}(t')).$$

The value of $D_{\text{temporal}}(i,t,t')$ is high when the node travels in more or less the same direction and almost at the same speed over a certain time interval that can be defined. However, $D_{\text{temporal}}(i,t,t')$ decreases if the relative direction or the speed ratio decreases.

Arguing in a way similar to that for $D_{\text{spatial}}(i,j,t)$, we have the following condition

$$|t - t'| > c_2 \quad \Rightarrow \quad D_{\text{temporal}}(i,t,t') = 0,$$

where $c_2 > 0$ is a constant which will be determined during our experiments in Section 6.

Average degree of temporal dependence: It is the value of $D_{\text{temporal}}(i,t,t')$ averaged over nodes and time instants satisfying certain condition. Formally,

$$\overline{D}_{\text{temporal}} = \frac{\sum_{i=1}^{N} \sum_{t=1}^{T} \sum_{t'=1}^{T} D_{\text{temporal}}(i,t,t')}{P},$$

where P is the number of tuples (i,t,t') such that $D_{\text{temporal}}(i,t,t') \neq 0$ Thus, if the current velocity of a node is completely independent of its velocity at some previous time step, then the mobility pattern is expected to have a smaller value for $\overline{D}_{\text{temporal}}$. However, if the current velocity is strongly dependent on the velocity at some previous time step, then the mobility

pattern is expected to have a higher value for $\overline{D}_{\text{temporal}}$.

3. *Relative speed (RS):* We use the standard definition from physics, i.e.

$$\text{RS}(i,j,t) = |\vec{V_i}(t) - \overrightarrow{V_j(t)}|.$$

As in the case of $D_{\text{spatial}}(i,j,t)$, we add the following condition:

$$D_{i,j}(t) > c_3 * R \quad \Rightarrow \quad \text{RS}(i,j,t) = 0,$$

where $c_3 > 0$ is a constant which will be determined during our experiments in 6.

Average relative speed: It is the value of $\text{RS}(i,j,t)$ averaged over node pairs and time instants satisfying certain condition. Formally,

$$\overline{\text{RS}} = \frac{\sum_{i=1}^{N} \sum_{j=1}^{N} \sum_{t=1}^{T} \text{RS}(i,j,t)}{P},$$

where P is the number of tuples (i,j,t) such that $\text{RS}(i,j,t) \neq 0$.

4. *Geographic restrictions:* For this metric, we developed the notion of degree of freedom of points on a map. Degree of freedom of a point is the number of directions a node can go after reaching that point, but currently we do not have a good way of quantitatively aggregating this definition for the whole map. Thus, we do not quantitatively define the geographic restrictions, but we qualitatively include it in our study as will be seen in Section 5.

4.3. Connectivity graph metrics

Since routing protocol performance is in general affected by the network topology dynamics, we feel that it is useful to have metrics to analyze the effect of mobility on the connectivity graph between the mobile nodes. The connectivity graph metrics aim to study this effect. These metrics might also help in relating mobility metrics with protocol performance, which will be shown in Section 6.

The connectivity graph is the graph $G = (V, E)$, such that $|V| = N$ and at time t, a link $(i,j) \in E$ iff $D_{i,j}(t) \leqslant R$. Let $X(i,j,t)$ be an indicator random variable which has a value 1 iff there is a link between nodes i and j at time t. $X(i,j) = \max_{t=1}^{T}$

$X(i, j, t)$ be an indicator random variable which is 1 if a link existed between nodes i and j at any time during the simulation, 0 otherwise.

1. *Number of link changes:* Number of link changes for a pair of nodes i and j is the number of times the link between them transitions from "down" to "up". Formally,

$$\text{LC}(i, j) = \sum_{t=1}^{T} C(i, j, t),$$

where $C(i, j, t)$ is an indicator random variable such that $C(i, j, t) = 1$ iff $X(i, j, t-1) = 0$ and $X(i, j, t) = 1$ i.e. if the link between nodes i and j is down at time $t - 1$, but comes up at time t.
Average number of link changes: It is the value of $\text{LC}(i, j)$ averaged over node pairs satisfying certain condition. Formally,

$$\overline{\text{LC}} = \frac{\sum_{i=1}^{N} \sum_{j=i+1}^{N} \text{LC}(i, j)}{P},$$

where P is the number of pairs i, j such that $X(i, j) \neq 0$.

2. *Link duration:* For two nodes i and j, at time t_1, duration of the link (i, j) is the length of the longest time interval $[t_1, t_2]$ during which the two nodes are within the transmission range of each other. Moreover these two nodes are not within the transmission range at time $t_1 - \epsilon$ and time $t_2 + \epsilon$ for $\epsilon > 0$. Formally,

$$\text{LD}(i, j, t_1) = t_2 - t_1$$

iff $\forall t \quad t_1 \leqslant t \leqslant t_2, \quad \epsilon > 0 : X(i, j, t) = 1$ and $X(i, j, t_1 - \epsilon) = 0$ and $X(i, j, t_2 + \epsilon) = 0$. Otherwise, $\text{LD}(i, j, t_1) = 0$.
Average link duration: It is the value of $\text{LD}(i, j)$ averaged over all existing links for node pairs satisfying certain condition. Formally,

$$\overline{\text{LD}} = \frac{\sum_{t_1=0}^{T} \sum_{i=1}^{N} \sum_{j=i+1}^{N} \text{LD}(i, j, t_1)}{P},$$

where P is the number of tuples (i, j, t_1) such that $\text{LD}(i, j, t_1) \neq 0$.

3. *Path duration:* For a path $P = \{n_1, n_2, \ldots n_k\}$, consisting of k nodes, at time t_1, path duration is the length of the longest time interval $[t_1, t_2]$, during which each of the $k - 1$ links between the nodes exist. Moreover, at time $t_1 - \epsilon$ and

time $t_2 + \epsilon$, $\epsilon > 0$, at least one of the k links does not exist. Thus, path duration is limited by the duration of the links along its path. Specifically, at time t_1, path duration is the minimum of the durations of the $k - 1$ links $(n_1, n_2), (n_2, n_3), \ldots,$ (n_{k-1}, n_k) at time t_1. Formally,

$$\text{PD}(n_1, n_k, t_1) = \min_{1 \leqslant z \leqslant k-1} \text{LD}(n_z, n_{z+1}, t_1),$$

that is, $\text{PD}(n_1, n_k, t_1)$ is the shortest path duration between node n_1 and node n_k at time t_1.
Average path duration: It is the value of $\text{PD}(n_1, n_k, t_1)$ averaged over all existing paths for node pairs satisfying certain condition. Formally,

$$\overline{\text{PD}} = \frac{\sum_{t_1=0}^{T} \sum_{n_1=1}^{N} \sum_{n_k=n_1+1}^{N} \text{PD}(n_1, n_k, t_1)}{P},$$

where P is the number of tuples (n_1, n_k, t_1) such that $\text{PD}(n_1, n_k, t_1) \neq 0$.

4. *Path availability:* It is the fraction of time during which a path is available between two nodes i and j. The node pairs of interest are the ones that have communication traffic between them. Formally,

$$\text{PA}(i, j) = \begin{cases} \frac{\sum_{t=\text{start}(i,j)}^{T} A(i, j, t)}{T - \text{start}(i,j)} & \text{if } T - \text{start}(i, j) > 0, \\ 0 & \text{otherwise}, \end{cases} \tag{1}$$

where $A(i, j, t)$ is an indicator random variable which has a value 1 if a path is available from node i to node j at time t, and has a value 0 otherwise. $\text{start}(i, j)$ is the time at which the communication traffic between nodes i and j starts.
Average path availability: It is the value of $\text{PA}(i, j)$ averaged over node pairs satisfying certain condition. Formally,

$$\overline{\text{PA}} = \frac{\sum_{i=1}^{N} \sum_{j=i+1}^{N} \text{PA}(i, j)}{P},$$

where P is the number of pairs i, j such that $T - \text{start}(i, j) > 0$.

4.4. Protocol performance metrics

We evaluate the performance of the MANET routing protocols using the metrics of throughput

(ratio of the number of packets delivered to the number of packets sent) and routing overhead (number of routing control packets sent) as done in several previous studies in this area of research.

5. Mobility models

As mentioned in Section 1, Random Waypoint does not seem to capture the mobility characteristics of spatial dependence, temporal dependence and geographic restrictions. In the previous section, we defined mobility metrics that either qualitatively or quantitatively define these characteristics. To thoroughly study the effect of mobility on MANET protocol performance, we seek to evaluate the protocols over a rich set of mobility models that span the design space of the mobility metrics. Thus, apart from Random Waypoint, we use the following mobility models:

1. RPGM model
2. Freeway mobility model
3. Manhattan mobility model

Each of the above models has certain characteristics that are different from Random Waypoint, which will be shown by our metrics and simulations.

1. *RPGM model:* Ref. [8] introduced this model. Here, each group has a logical center (group leader) that determines the group's motion behavior. Initially, each member of the group is uniformly distributed in the neighborhood of the group leader. Subsequently, at each instant, every node has a speed and direction that is derived by randomly deviating from that of the group leader.
   Applications: Group mobility can be used in military battlefield communications where the commander and soldiers form a logical group. More applications are mentioned in [8].
   Important characteristics: Each node deviates its velocity (both speed and direction) randomly from that of the leader. The movement in group mobility can be characterized as follows:

(a) $|\vec{V}_{\text{member}}(t)| = |\vec{V}_{\text{leader}}(t)| + \text{random}(\ ) * \text{SDR} * \text{max\_speed}$

(b) $\theta_{\text{member}}(t) = \theta_{\text{leader}}(t) + \text{random}(\ ) * \text{ADR} * \text{max\_angle}$

where $0 \leqslant \text{SDR}, \text{ADR} \leqslant 1$. SDR is the speed deviation ratio and ADR is the angle deviation ratio. SDR and ADR are used to control the deviation of the velocity (magnitude and direction) of group members from that of the leader. max_speed and max_angle are used to specify the maximum deviation a group member can take. In our simulation, we set maximum speed for the group leader as the max_speed and set 180° as the max angle. Since the group leader mainly decides the mobility of group members, group mobility pattern is expected to have high spatial dependence for small values of SDR and ADR.

2. *Freeway mobility model:* We propose this new model to emulate the motion behavior of mobile nodes on a freeway. The freeway map used in our study is shown in Fig. 2(a).
   Applications: It can be used in exchanging traffic status or tracking a vehicle on a freeway.
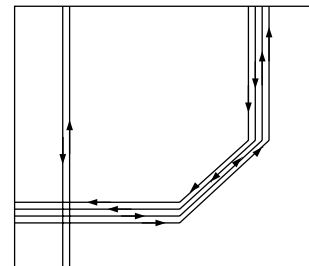
(a)
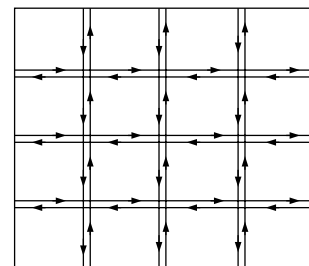
(b)

Fig. 2. Maps used in Freeway and Manhattan model. (a) Freeway map and (b) Manhattan map.

Important characteristics: In this model we use maps. There are several freeways on the map and each freeway has lanes in both directions. The differences between Random Waypoint and Freeway are the following:

(a) Each mobile node is restricted to its lane on the freeway.
(b) The velocity of mobile node is temporally dependent on its previous velocity.
(c) If two mobile nodes on the same freeway lane are within the safety distance (SD), the velocity of the following node cannot exceed the velocity of preceding node.

The inter-node and intra-node relationships involved are:

(a) $|\overrightarrow{V_i}(t+1)| = |\overrightarrow{V_i}(t)| + \text{random}(\ ) * |\overrightarrow{a_i}(t)|$
(b) $\forall i, \forall j, \forall t\ D_{i,j}(t) \leqslant \text{SD} \Rightarrow |\overrightarrow{V_i}(t)| \leqslant |\overrightarrow{V_j}(t)|$, if j is ahead of i in its lane.

Due to the above relationships, the Freeway mobility pattern is expected to have spatial dependence and high temporal dependence. It also imposes strict geographic restrictions on the node movement by not allowing a node to change its lane.

3. *Manhattan mobility model:* We introduce the Manhattan model to emulate the movement pattern of mobile nodes on streets defined by maps. The Manhattan map used in our study is shown in Fig. 2(b).
   Applications: It can be useful in modeling movement in an urban area where a pervasive computing service between portable devices is provided.
   Important characteristics: Maps are used in this model too. The map is composed of a number of horizontal and vertical streets. Each street has two lanes for each direction (north and south direction for vertical streets, east and west for horizontal streets). The mobile node is allowed to move along the grid of horizontal and vertical streets on the map. At an intersection of a horizontal and a vertical street, the mobile node can turn left, right or go straight. This choice is probabilistic: the probability of

moving on the same street is 0.5, the probability of turning left is 0.25 and the probability of turning right is 0.25.
The velocity of a mobile node at a time slot is dependent on its velocity at the previous time slot. Also, a node's velocity is restricted by the velocity of the node preceding it on the same lane of the street. The inter-node and intra-node relationships involved are the same as in the Freeway model.
Thus, the Manhattan mobility model is also expected to have high spatial dependence and high temporal dependence. It too imposes geographic restrictions on node mobility. However, it differs from the Freeway model in giving a node some freedom to change its direction.

Most of the mobility models mentioned above are parameterized. E.g. SDR and ADR are some of the parameters used in RPGM, while maps are important parameters in the Freeway and Manhattan models. Although we did not quantitatively define geographic restrictions in Section 4, we qualitatively include them in our study by using the Freeway and Manhattan models. Using a parameterized approach, we aim to get a good coverage of the design space of the proposed mobility metrics by producing a rich set of mobility patterns that can be used as a "test-suite" for further research.

6. Experiments

As a first step, we wanted to validate if our proposed metrics differentiate the mobility models. Once this was done, we focused on answering the following questions: *Whether* mobility affects protocol performance?, if yes, we attempt to answer the questions *Why?* and *How?* mentioned in Section 1.

6.1. Validating the mobility metrics

Our mobility scenario generator produced the different mobility patterns following the RPGM, Freeway and Manhattan models according to the

format required by *ns-2*. In all these patterns, 40 mobile nodes moved in an area of 1000 m × 1000 m for a period of 900 s. Random Waypoint mobility pattern was generated using the *setdest* tool which is a part of the *ns-2* distribution. For RPGM, we used two different mobility scenarios: single group of 40 nodes and four groups of 10 nodes each moving independently of each other and in an overlapping fashion. Both speed deviation ratio (SDR) and angle deviation ratio (ADR) were set to 0.1. For the Freeway and Manhattan models, the nodes were placed on the freeway lanes or local streets randomly in both directions initially. Their movement was controlled as per the specifications of the models. If a node moves beyond the boundary of the area it is re-inserted at the beginning position in a randomly chosen lane in the area. The maximum speed V_{max} was set to 1, 5, 10, 20, 30, 40, 50 and 60 m/s to generate different movement patterns for the same mobility model. On evaluating these patterns with our mobility metrics, we observed that some of the metrics were able to differentiate between the mobility patterns based on the characteristics we focused on, while the others failed.

Average relative speed ($\overline{RS}$): We experimented with different values of the constant c_3 mentioned in Section 4. For the value of $c_3 = 2$, $\overline{RS}$ could differentiate between the different mobility patterns very clearly. As seen in Fig. 3, $\overline{RS}$ has the lowest value for RPGM (single group and multiple
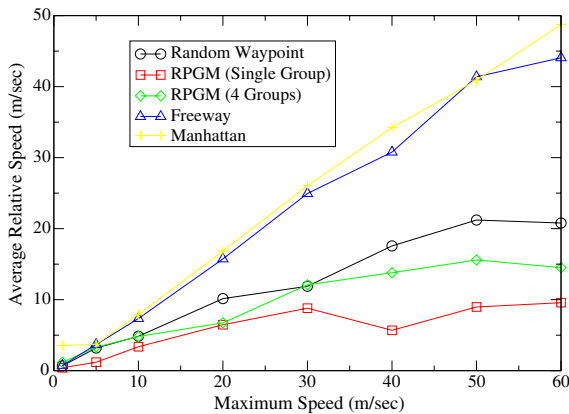
group mobility) as the nodes move together in a co-ordinated fashion with little deviation, while it has a medium value for Random Waypoint. Its value for the Freeway and Manhattan mobility patterns is the highest and almost twice that for Random Waypoint. This high value is because of the movement in opposite direction for both Freeway and Manhattan mobility patterns.

Average degree of spatial dependence ($\overline{D}_{spatial}$): We experimented with different values of the constant c_1 mentioned in Section 4. For the value of $c_1 = 2$, $\overline{D}_{spatial}$ could differentiate between the different mobility patterns very clearly. As seen in Fig. 4, $\overline{D}_{spatial}$ has a higher value for single group mobility (around 0.5) than that of multiple group mobility (about 0.35). However, for the Random Waypoint, Manhattan and Freeway, its value is almost 0. Intuitively, in RPGM, the group leader controls the movement of the mobile node and thus the mobility pattern has a high spatial dependence. Initially, we expected the Freeway and Manhattan mobility patterns to have a high spatial dependence as a node's movement is influenced by nodes before it in the lane. Due to the use of lanes in opposite directions in the map, the positive degree of spatial dependence of a node with nodes in the same direction cancels the negative degree of spatial dependence of the node with nodes traveling in the opposite direction.

Average degree of temporal dependence ($\overline{D}_{temporal}$): This metric could not differentiate be-
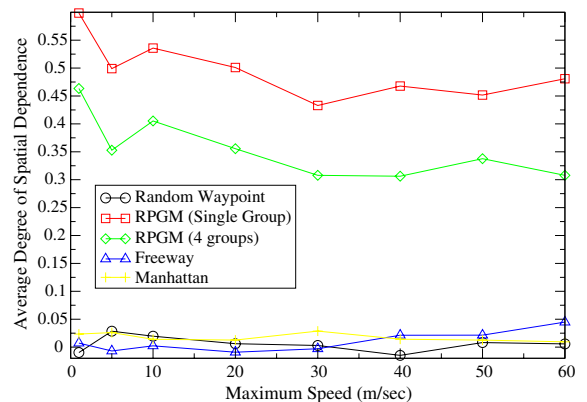

Fig. 3. Average relative speed.

Fig. 4. Average degree of spatial dependence.

tween the various mobility patterns used in our study. The usefulness of this metric is still under investigation.

In summary, $\overline{RS}$ and $\overline{D}_{\text{spatial}}$ are found to be useful mobility metrics in our study. Figs. 3 and 4 show that for each of these metrics, we had scenarios with relatively low values, medium values and relatively high values. Similarly, for geographic restrictions, the Freeway does not allow a node to change directions as freely as the Manhattan model. So, we believe that our "test-suite" has given a reasonably good coverage of the mobility metric space.

6.2. Validating the connectivity graph metrics

To study the effect of mobility on the connectivity graph, we evaluated the connectivity graphs resulting from the mobility patterns used in Section 6.1. We had the following observations about the connectivity graph metrics:

Average link duration ($\overline{LD}$): As seen in Fig. 5, $\overline{LD}$ has a higher value for single group and multiple groups than Random Waypoint. For the Freeway and Manhattan its value is similar to Random Waypoint or even worse. Since nodes in a group move at velocities that are deviated by a small fraction from the group leader, an already existing link between two nodes is expected to have a higher duration. The low value for the Freeway and Manhattan may be because of the opposite direction of motion and high relative speeds.

Average path duration ($\overline{PD}$): The path duration is counted as the interval between the time when the path is set up and the time when the path is broken. However, there can be potentially exponential paths between any specific source–destination pair. Analyzing the duration of all these paths might not be feasible. As a reasonable approximation, we define the path duration as the duration of the shortest path. The shortest path between the source and the destination is computed by the *Breadth First Search* (BFS) algorithm [19]. Similar to $\overline{LD}$, as shown in Fig. 6, RPGM (single group and multiple groups) has a higher $\overline{PD}$ value than Random Waypoint. The $\overline{PD}$ values for Manhattan and Freeway are similar to or slightly lower than Random Waypoint. Since each path is composed of several links, it is likely that the behavior of path duration will be determined by the behavior of link duration.

Average number of link changes ($\overline{LC}$): This metric was not able to differentiate between the several mobility patterns used in our study.

Average path availability ($\overline{PA}$): We use the BFS algorithm [19] on the snapshots of the network to calculate whether a path between a specific source–destination pair exists. For RPGM (single group), RPGM (mutiple group), Random Waypoint, Freeway and Manhattan models, $\overline{PA}$ is found to be around 100%, 92%, 97%, 99% and 95% respectively. In most cases, a path is available at least 95% of the time. Thus, the difference across the models was too small to be of any help.
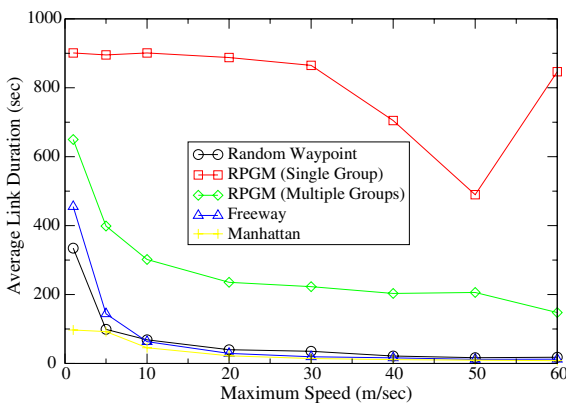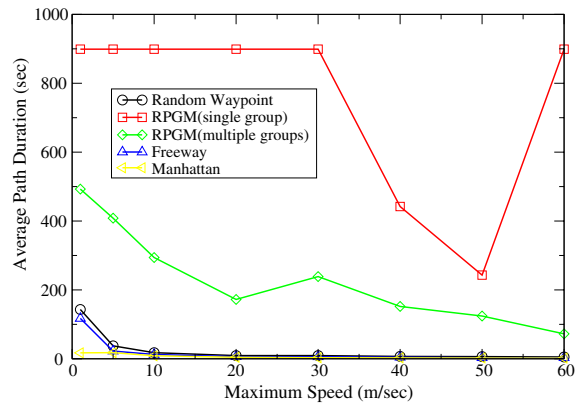

Fig. 5. Average link duration.

Fig. 6. Average path duration.

In summary, $\overline{LD}$ and $\overline{PD}$ are found to be useful metrics to differentiate the connectivity graph arising from the different mobility patterns used in our study.

6.3. Whether mobility affects protocol performance?

To evaluate the effect of mobility on the performance of protocols, we carried out simulations in the network simulator (*ns-2*) environment with the CMU Wireless Ad Hoc networking extension. The transmission range of the nodes was 250 m. The mobility patterns used were the same as those used to Section 6.1. The traffic pattern was generated by the *cbrgen* tool that is part of *ns-2* distribution. The traffic consisted of 20 CBR sources and 30 connections. The source–destination pairs were chosen at random. The data rate used was 4 packets/s and the packet size was 64 bytes.

To remove any effects due to randomness of the traffic pattern, we used different random seeds to generate three different traffic patterns having the same number of sources and connections. The results for each model (for a given V_{max}) are averaged over simulation runs using these three different traffic patterns.

We evaluated the performance of DSR, AODV and DSDV across this rich set of mobility models and observed that the mobility models may drastically affect protocol performance. We use DSR as an illustrative example.

DSR shows a difference of almost 40% in throughput from Manhattan to the RPGM (single group) model as seen from Fig. 7(a).

Also, there is an order of magnitude difference in the routing overhead of DSR across the various models as shown by Fig. 7(b). Similar performance differences were observed for other protocols used in our study. We observed that DSR, DSDV and AODV achieve the highest throughput and the least overhead with RPGM and incur high overhead and low throughput with both Freeway and Manhattan models. This is consistent with the observations made in [8] which evaluated the protocols using Random Waypoint and several other group mobility applications. However, we take a step further and attempt to analyze the reason for this performance difference in Section 6.4.
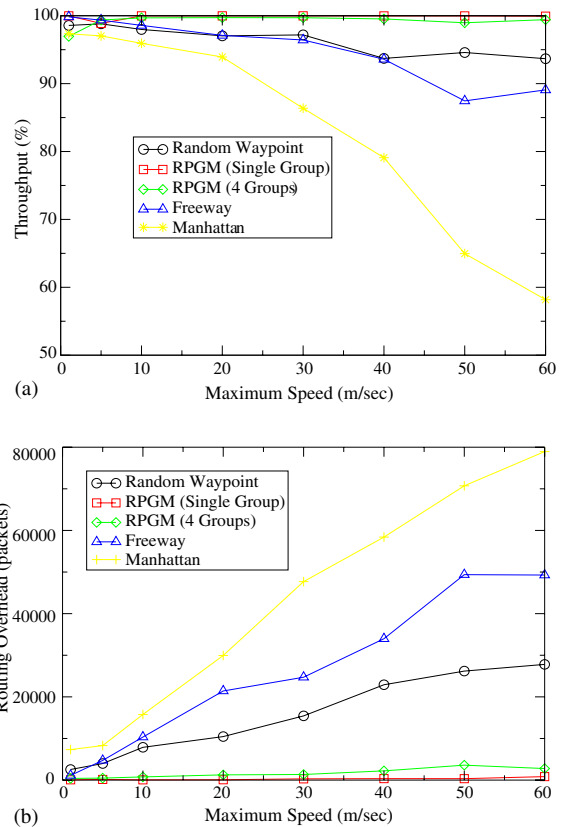

Fig. 7. DSR protocol performance across mobility models. (a) Throughput and (b) routing overhead.

Relative performance of protocols across mobility models: In this part, we investigated the effect of mobility on relative rankings of protocol performance. As shown in Figs. 8(a), 9(a), 10(a), 11(a), and 12(a), DSR seems to produce the highest throughput in most cases, while AODV seems to outperform DSR (by almost 31%) in the Manhattan model. As seen from Figs. 9(a) and 12(a), the relative ranking of AODV and DSDV in terms of throughput seems to depend on the underlying mobility model.

Also, DSR incurs the least routing overhead in most cases, while DSDV has a lower overhead than DSR in the Freeway and Manhattan models as shown in Figs. 11(b) and 12(b). The relative ranking of DSR and DSDV in terms of routing overhead seems to depend on the underlying mobility model as shown in Figs. 8(b), 9(b), 10(b), 11(b), and 12(b).
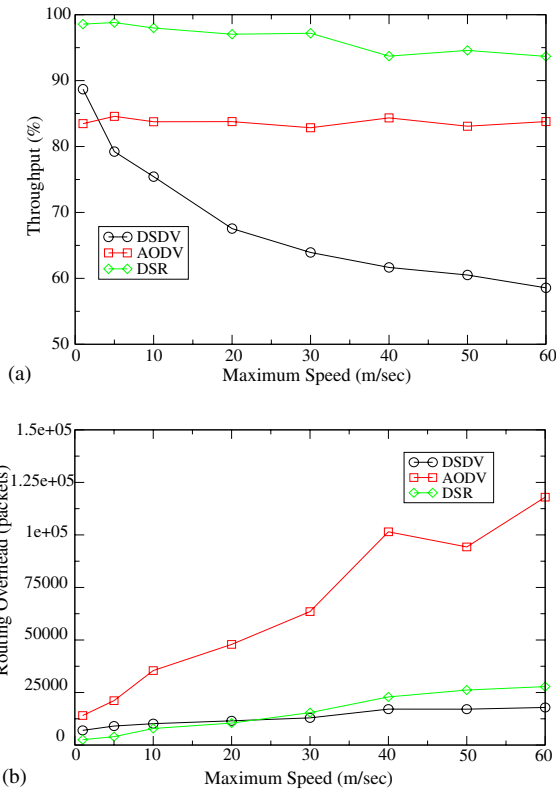
Fig. 8. Protocol performance for Random Waypoint model. Random Waypoint: throughput (a) and routing overhead (b).

Fig. 9. Protocol performance for RPGM (single group) model. RPGM (single group): throughput (a) and routing overhead (b).

Thus, we conclude that relative rankings of protocols may vary with the mobility model used. We also observe that DSDV achieves a higher throughput than AODV (by around 10%) in RPGM. Thus, in general it is not always true that on-demand protocols perform better than table driven ones in terms of throughput. Also, a protocol with the least overhead does not always produce the highest throughput. E.g. in the Freeway model, DSDV seems to have the least throughput and the least overhead.

Although, these results were somewhat expected, the quantitative analysis helped us gain a lot of insight to answer the next question.

6.4. Why mobility affects protocol performance?

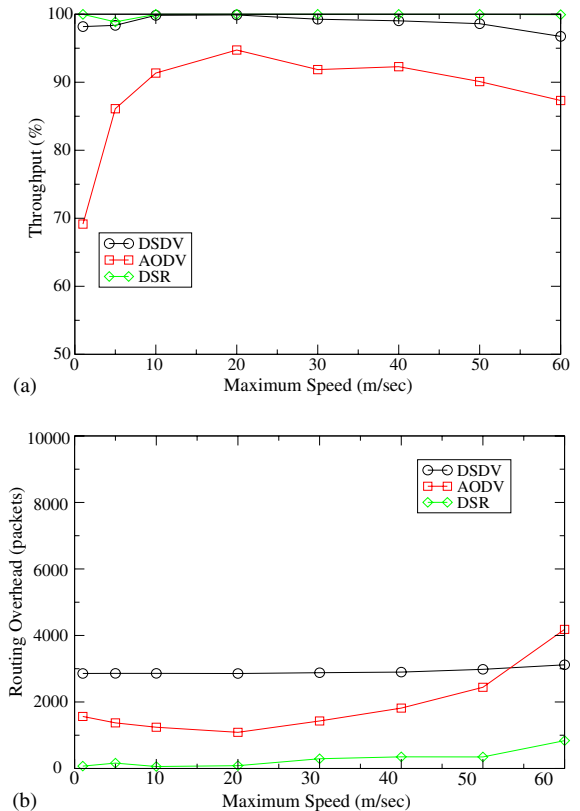First, the relationship between the mobility metrics and the performance metrics was unclear.

But after introducing the connectivity graph metrics, we were able to observe a very clear correlation between average degree of spatial dependence, average relative speed, average link duration, average path duration and protocol performance metrics. *The mobility pattern influences the connectivity graph which in turn influences the protocol performance.* The relationship is identified in Fig. 1.

In general, it was observed that DSR, DSDV and AODV had a higher throughput and lower overhead for the group mobility models than for the Random Waypoint model. At the same time, all the protocols had a higher throughput and lower overhead for Random Waypoint than the Freeway and Manhattan models. One plausible reason for this observation can be as follows:
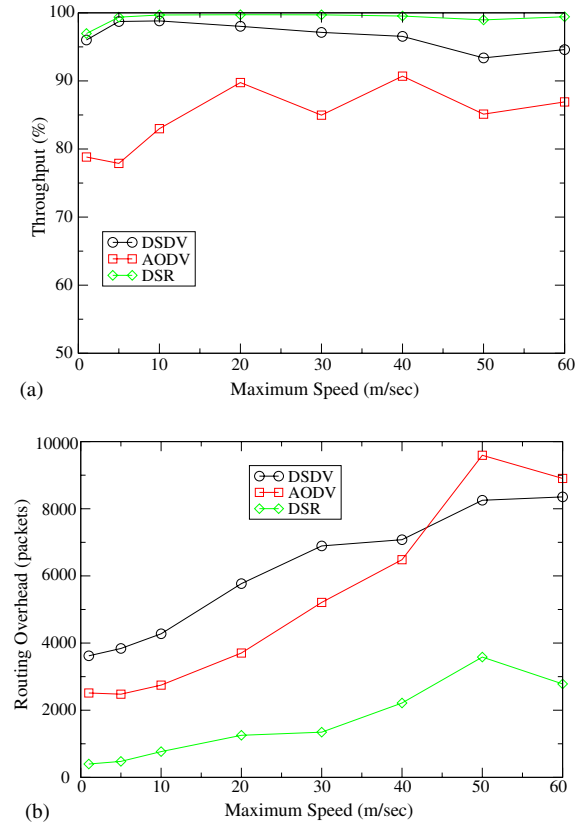
Fig. 10. Protocol performance for RPGM (four groups) model. RPGM (four groups): throughput (a) and routing overhead (b).

Fig. 11. Protocol performance for Freeway model. Freeway: throughput (a) and routing overhead (b).

1. With similar relative speed, between Random Waypoint and RPGM, high degree of spatial dependence (for RPGM) means higher link duration and correspondingly higher path duration, which in turn will result in higher throughput and lower routing overhead.
2. With the same degree of spatial dependency, between Freeway/Manhattan and Random Waypoint, high relative speed (for Freeway/Manhattan) means lower link duration and correspondingly lower path duration, which will result in lower throughput and higher overhead.

The above reasoning can be explained as follows: For a given relative speed, if a mobility pattern has a high degree of spatial dependence, an already existing link between two nodes is expected to remain stable for a longer period of time
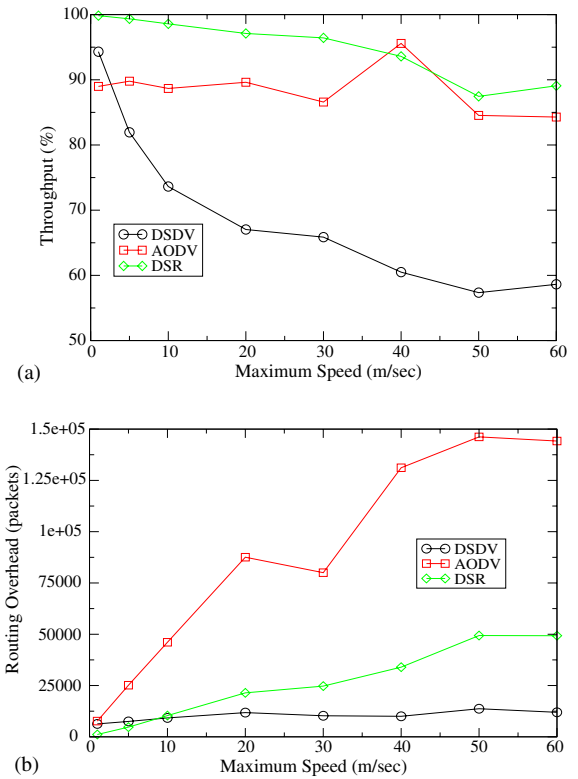
as the nodes are likely to move together, then the path between source and destination is relatively stable. Thus fewer packets will be dropped due to path breakage leading to higher throughput. At the same time, the control overhead is lower as little effort is needed to repair the seldom broken path. For a given spatial dependence, if a mobility pattern has a high relative speed, the nodes might move out of range more quickly. Thus an already existing link and the paths using this specific link may remain stable for a relatively shorter duration. This may lead to more packets being dropped due to path breakage, resulting in lower throughput. Higher control overhead is needed to repair the more frequently broken path. We also note that the Freeway and Manhattan mobility patterns have high relative speed and low degree of spatial dependence leading to the worst performance of all the protocols while using these models.
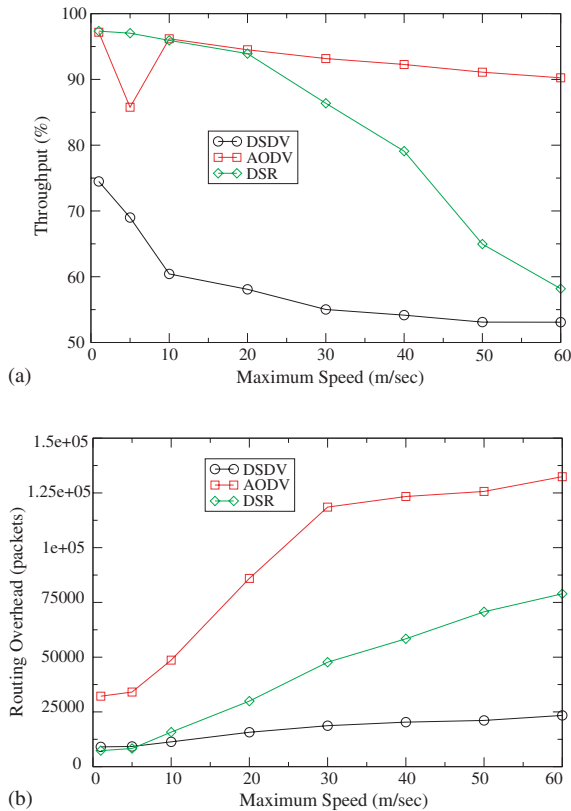
(a)

(b)

Fig. 12. Protocol performance for Manhattan model. Manhattan: throughput (a) and routing overhead (b).

7. How mobility affects protocol performance? Analysis of building blocks

Unlike the conventional evaluation studies [16–18], we pursue our analysis beyond the ''whole protocol'' level and attempt to answer *How* mobility affects protocol performance by looking into the ''parts'' that constitute the MANET routing protocols. We propose an approach to systematically decompose a protocol into its functional mechanistic ''building blocks''. Each building block can be thought of as a parameterized ''black box''. The parameter settings define the behavior of each block, while the behavior of building blocks and the nature of interaction between the building blocks defines the behavior of the protocol as a ''whole''. We use the analysis of reactive protocols as an example to illustrate this ap-

proach. In this section, we carry out a preliminary analysis of the impact of mobility on two reactive routing protocols after identifying the basic building blocks of MANET reactive routing protocols and their parameter setting. Thus we can extract the relative merits of different parameter settings and achieve a better understanding of various building blocks of MANET routing protocols, which will serve as a solid cornerstone for development of more efficient MANET routing protocols.

Fig. 13(a) and (b) shows the building block architecture for DSR and AODV respectively, Fig. 13(c) shows a generalized building block architecture for reactive MANET protocols.

7.1. Design choices of building blocks for DSR and AODV

First we discuss the design choices (parameter settings) of the identified building blocks of reactive MANET routing protocols and specific parameter settings for DSR and AODV. We pose some questions about the utility of the various design choices made by these protocols. In Section 7.2, we attempt to answer these questions.

The mechanism of reactive MANET routing protocols such as DSR and AODV [1] is composed of two major phases.

7.1.1. Route setup phase

Route discovery is the major mechanism in this phase. It is initiated if there is no cached route available to the destination. This mechanism consists of the following building blocks:

1. *Flooding building block:* The flooding building block takes responsibility to distribute the route request messages within the network. Here, the key parameter is *the range of flooding*, generally

---

[1] Current AODV implementation in *ns-2* (version ns-2.1b8a) and IETF specification adopt the expanding ring search for query flooding, localized rediscovery for error handling and source-specific error notification, which may be different with the original AODV paper. In our study, if not particularly specified, we refer the implementation in *ns-2* (version ns-2.1b8a) as the standard AODV mechanism.
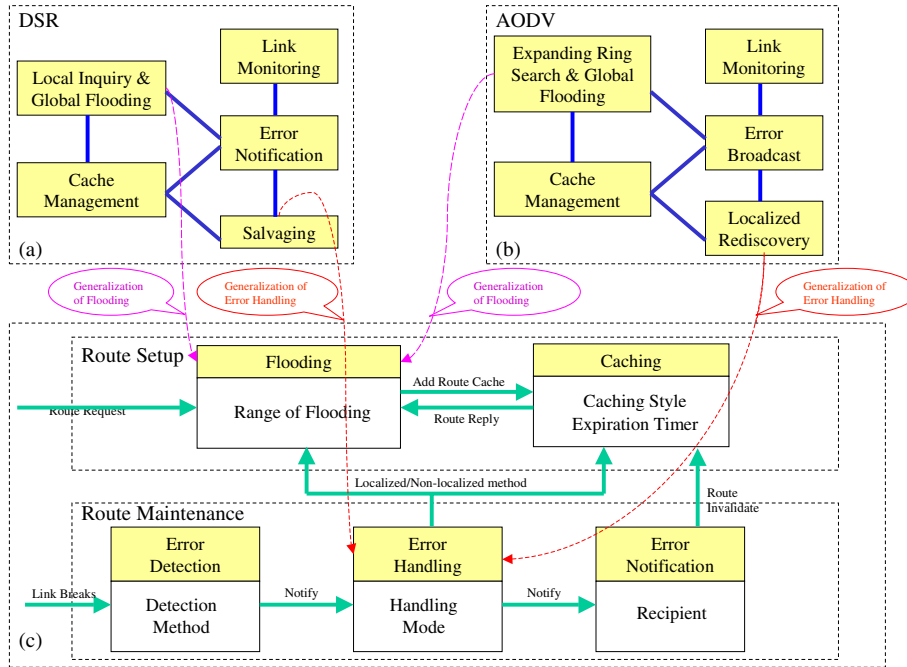
Fig. 13. Diagram of building block framework for reactive protocols.

described by TTL field in the IP header. If the TTL is set to the network diameter, a global flooding is done. One optimization is to implement the localized controlled flooding before global flooding. This may be useful if the probability of finding an appropriate cache in the neighborhood is high.

For the *range of flooding*, both DSR and AODV use the two-step controlled flooding. DSR conducts a non-propagating direct-neighborhood inquiry ($TTL = 1$) before the global flooding ($TTL = D$, D is network diameter). Similarly, AODV uses the expanding ring search ($TTL = 1, 3, 5, 7$) before the global flooding is initiated. Here, we want to answer the following question: *How useful are non-propogating route requests?*

2. *Caching building block:* The caching building block helps to efficiently and promptly provide the route to the destination without referring to the destination every time. Several parameters affect the behavior of the caching building block. One of those parameters is *whether ag-*

gressive caching is allowed, i.e. whether multiple cache entries are allowed for the same destination and whether a node can cache the route information it overhears? Generally speaking, aggressive caching scheme increases the possibility of finding an appropriate route without re-initiating a route discovery.

DSR uses *aggressive caching*, while AODV does not. For caching, we are interested in the following questions: *How useful is caching? and Is aggressive caching better than non-aggressive caching?*

7.1.2. Route maintenance phase

Route maintenance phase takes the responsibility of detecting broken links and repairing the corresponding routes. This phase is made up of the following building blocks:

1. *Error detection building block:* It is used to monitor the status of the link of a node with its immediate neighbors. Several methods to monitor the link status between neighbors can be used:

MAC level acknowledgement, network-layer explicit Hello message or network-layer passive overhearing scheme. Here, the parameter is *the mode of error detection used*.

For *mode of error detection*, both DSR and AODV can use either of the three choices. Since all these schemes are very similar, we do not investigate this building block in our analysis.

2. *Error handling building block:* It finds alternative routes to replace an invalid route after a broken link is detected. One of the parameters to this block is *what recovery scheme should be used*. In a localized recovery, the node detecting the broken link will attempt to find an alternative route in its own cache or do a localized flooding before asking the source to re-initiate the route discovery.

For *localized recovery*, in DSR, on detecting a broken link, the upstream node will first search its cache to replace the invalid route (this scheme is called salvaging), although the found alternative route may also be invalid in some scenarios. While in AODV, the upstream node detecting the broken link will initiate a localized flooding to find the route to the destination. For this building block, we are interested in the following question: *Which is a better scheme for localized error handling: cache lookup or localized flooding?*

3. *Error notification building block:* It is used to notify the nodes in the network about invalid routes. The key parameter to this building block is *the recipient of the error message*. Either only the source is notified or the entire network is notified.

For *recepient of error notification*, both DSR and AODV notify the error to the source. Since both DSR and AODV use the same parameter setting, we do not investigate this block during our simulations.

Besides these three questions about the design choices, we are also interested at the plausible explanation for the observation we made in Section 6: DSR outperforms AODV in most mobility scenarios except the Freeway and Manhattan model with high mobility.

7.2. Performance evaluation

We identified parts of the network simulator (*ns-2*) code [1] which implement these building blocks and profiled them during our simulations. The simulation setting is exactly the same as Section 6.

7.2.1. Flooding

Fig. 14(a) and (b) shows the *ratio of non-propagating route requests to the total number of route requests* issued by the DSR and AODV respectively. This metric measures the likelihood of finding a route to the destination from the source's neighborhood. The result shows that non-propagating route request is frequently used (more than 30% for DSR and more than 10% for AODV in most scenarios).
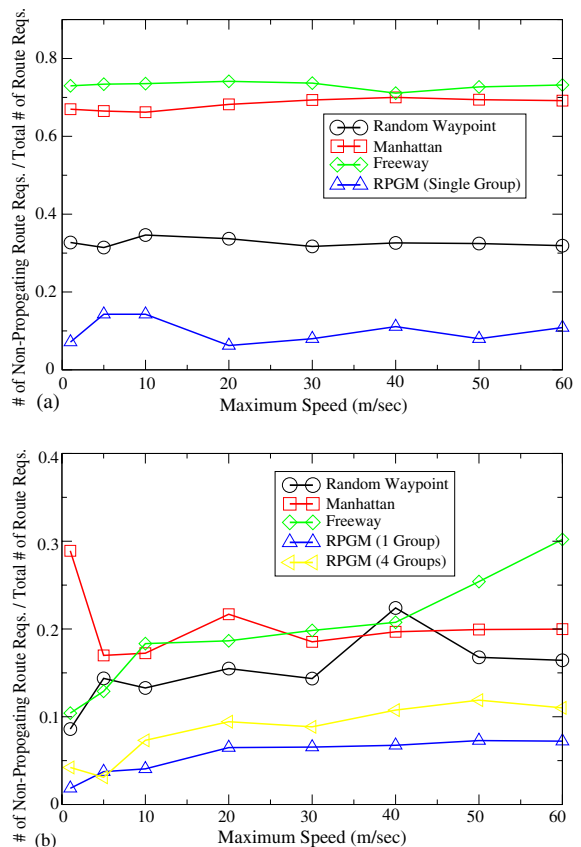

Fig. 14. Ratio of non-propogating route requests to total number of route requests. (a) DSR and (b) AODV.

It is observed that this ratio increases from Random Waypoint to Manhattan to Freeway mobility models. This is because the geographic constraints on movement are greater in Freeway than Manhattan, which are in turn greater than in Random Waypoint. Thus, the likelihood of finding a route to the destination from the source's neighbors increases from Random Waypoint to Manhattan to Freeway.

On the other hand, the ratio for DSR is almost twice as large as that for AODV across all mobility models. A possible reason for this comes from the fact that DSR uses aggressive caching as compared to AODV. When such a caching scheme is coupled with the mechanism of non-propagating route requests, it translates to low routing overhead and high throughput as was shown in our study and several other comparative studies.

Thus, it seems that caching has a significant impact on the performance of DSR and AODV. Hence we study it next.

7.2.2. Caching

To measure the effectiveness of caching, we evaluate the *ratio of the number of route replies coming from the cache to the total number of route replies*.

Fig. 15(a) and (b) shows that this ratio is high for Random Waypoint, Manhattan and Freeway models, which implies that most of the route replies for these mobility models come from the cache (more than 80% in most mobility scenarios).

The difference in the ratio for DSR and AODV is greater than 20% for all mobility models. DSR uses aggressive caching as compared to AODV. Thus, the likelihood of a route reply coming from a cache is higher in DSR than in AODV. Thus, fewer route requests will be needed and thus the routing overhead of DSR is lower than AODV as we observed in Section 6. Thus, aggressive caching seems to be a good design choice.

To completely evaluate the caching strategy, we also need to examine the validity of the cache entries. We evaluate the *ratio of invalid cache entries to the total number of cache entries* for DSR.

As shown by Fig. 16, the ratio increases from RPGM (around 10%) to Random Waypoint to Freeway (around 60%) to Manhattan (around
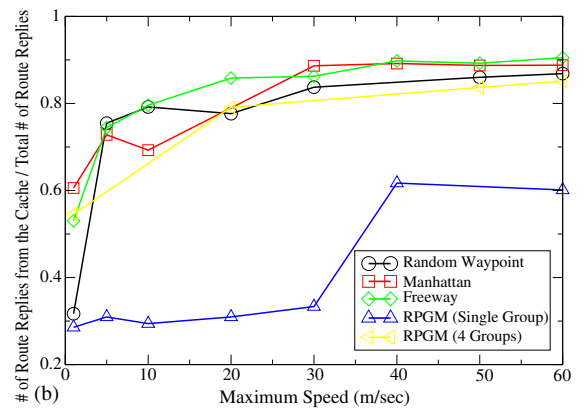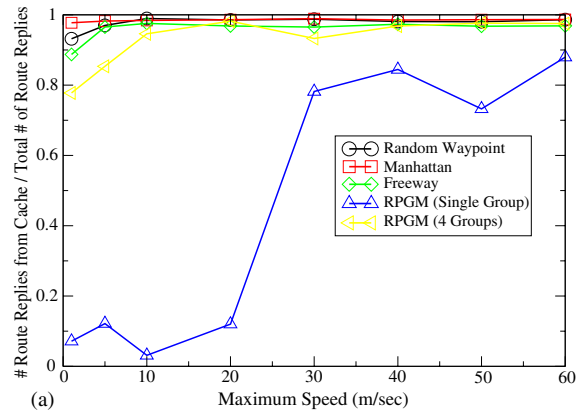

Fig. 15. Ratio of the number of route replies from the cache to the total number of route replies. (a) DSR and (b) AODV.
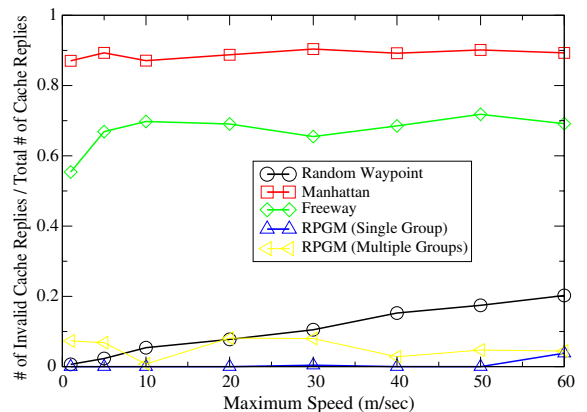

Fig. 16. Ratio of the number of invalid route replies from the cache to the total number of route replies from the cache.

80%) mobility models. Thus caching may have adverse effects in mobility models with a high relative speed which may lead to cache invalidation. Packets may be sent on invalid routes which might lead to packets being dropped and route request retries. This leads to a lower throughput and higher overhead for DSR for the Freeway and Manhattan models as was shown in our study.

On the other hand, in mobility models with very high relative speed like Manhattan and Freeway, AODV seems to achieve as good a throughput as DSR (and sometimes better). AODV does not use aggressive caching, thus the ratio of the number of route replies coming from the cache to the total number of route replies is lesser for AODV than DSR. Thus, the likelihood of getting invalid routes from the cache is lesser for AODV than for DSR. This may explain why AODV outperforms DSR in Freeway and Manhattan models with high mobility.

Moreover, at high relative speeds, the number of routes broken is greater. Thus, a protocol which has a better error handling mechanism at higher relative speeds might perform better in such situations. This line of reasoning leads us to evaluate the next building block of interest—error handling.

7.2.3. Error handling

To study the effectiveness of error handling, we focus on localized error handling. We evaluate the *ratio of the number of localized error handling to the total number of route errors* for both DSR and AODV. For DSR, we notice that salvaging accounts for less than 2% of the total number of route errors. Moreover, if we take invalid cache entries into account, the effect of salvaging on the protocol performance is further lowered. On the other hand, in AODV, a route request is initiated by the upstream node which detects the broken link if it is closer to the destination. As Fig. 17 shows, the ratio is between 40% and 50% for Freeway and Manhattan models. Moreover the routes obtained by this mechanism are more up to date than those from the cache salvaging in DSR. This is another factor which explains the better performance of AODV as compared to DSR in the Freeway and Manhattan models.
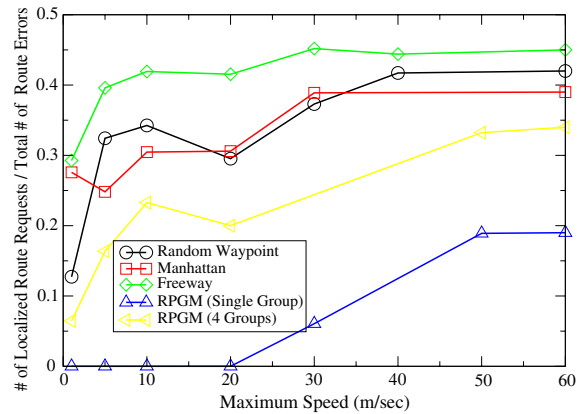

Fig. 17. Ratio of the number of localized route recovery requests to the total number of route errors for AODV.

7.3. Discussion

The above study of the building blocks has given us greater insight into the design of the reactive routing protocols for MANETs. Decomposing a protocol into building blocks and evaluating these building blocks have shown us the scenarios in which the chosen parameters can give a better performance. From the above study, we learnt the following principles of protocol design:

1. Caching helps reduce the protocol overhead. However, whether aggressive caching should be used depends on the scenarios in which the protocol will be deployed. For low mobility scenarios, aggressive caching might be useful, while for higher mobility scenarios, the more stale cache entries incurred by aggressive caching might affect the protocol throughput.
2. Non-propagating route requests, when combined with caching also reduce the protocol overhead. If caching is widely done in the network, it may be more advantageous to do non-propagating route requests (or expanding ring search) than globally flooding the route request. In DSR, due to aggressive caching, it may be more useful to do expanding ring search (from the source) on a route error than doing a global flooding (from the source). Again this might work well only for low mobility scenarios.

3. The nature of localized error handling also has a significant impact on protocol performance. Re-initiating a route request from an intermediate node can be more advantageous than doing a local cache lookup in high mobility scenarios, while a cache lookup might be more advantageous for low mobility scenarios.

Thus, no particular parameter setting of these building blocks is the most optimal for all scenarios. This further strengthens our conclusion that there is no clear winner among the protocols across all mobility scenarios.

8. Conclusions and future work

In this paper, we proposed a framework to analyze the impact of mobility pattern on routing performance of mobile ad hoc network in a systematic manner. In our study, we observe that the mobility pattern does influence the performance of MANET routing protocols. This conclusion is consistent with the observation of previous studies. But unlike previous studies that compared different ad hoc routing protocols, there is no clear winner among the protocols in our case, since different mobility patterns seem to give different performance rankings of the protocols. We hope that our "test-suite" of mobility models can be incorporated into the current scenarios used to test the MANET routing protocols.

Moreover, we observe that the mobility pattern influences the connectivity graph that in turn influences the protocol performance. In addition, we did a preliminary investigation of the common building blocks of MANET routing protocols, the effect of mobility on these building blocks and how they influence the protocol as a "whole".

In the future, we plan to study the impact of our "test-suite" on the performance of other ad hoc network protocols like multicast ad hoc, geographic routing protocols. This study would help us understand the impact of mobility more deeply and clearly. We believe that several parameters such as traffic patterns, node density and initial placement pattern of nodes may affect the routing performance and need to investigate them further.

9. Uncited references

[2–4].

References

[1] L. Breslau, D. Estrin, K. Fall, S. Floyd, J. Heidemann, A. Helmy, P. Huang, S. McCanne, K. Varadhan, Y. Xu, H. Yu, Advances in network simulation, IEEE Computer 33 (5) (2000) 59–67.

[2] F. Bai, N. Sadagopan, A. Helmy, Important: a framework to systematically analyze the impact of mobility on performance of routing protocols for adhoc networks, in: INFOCOM 2003, April 2003.

[3] N. Sadagopan, F. Bai, B. Krishnamachari, A. Helmy, Paths: analysis of path duration statistics and their impact on reactive manet routing protocols, in: MobiHoc 2003, June 2003.

[4] F. Bai, N. Sadagopan, A. Helmy, Brics: a building-block approach for analyzing routing protocols in ad hoc networks—a case study of reactive routing protocols, in: USC-CS-TR-02-775, November 2002.

[5] D.B. Johnson, D.A. Maltz, J. Broch, DSR: the dynamic source routing protocol for multi-hop wireless ad hoc networks, in: C. Perkins (Ed.), Ad Hoc Networking, Addison-Wesley, Reading, MA, 2001, pp. 139–172.

[6] C.E. Perkins, P. Bhagwat, Highly dynamic destination sequenced distance vector routing (DSDV) for mobile computers, in: ACM SIGCOMM, 1994, pp. 234–244.

[7] C. Perkins, Ad hoc on demand distance vector (AODV) routing, internet draft, draft-ietf-manet-aodv-00.txt.

[8] X. Hong, M. Gerla, G. Pei, C.-C. Chiang, A group mobility model for ad hoc wireless networks, in: ACM/IEEE MSWiM, August 1999.

[9] J. Broch, D.A. Maltz, D.B. Johnson, Y.-C. Hu, J. Jetcheva, A performance comparison of multi-hop wireless ad hoc network routing protocols, in: Proceedings of the Fourth Annual ACM/IEEE International Conference on Mobile Computing and Networking, ACM, October 1998.

[10] S.R. Das, C.E. Perkins, E.M. Royer, Performance comparison of two on-demand routing protocols for ad hoc networks, in: INFOCOM, March 2000.

[11] S.R. Das, R. Castaneda, J. Yan, R. Sengupta, Comparative performance evaluation of routing protocols for mobile, ad hoc networks, October 1998, pp. 153–161.

[12] V.D. Park, M.S. Corson, Temporally-ordered routing algorithm (TORA) version 1: functional specification, internet-draft, draft-ietf-manet-tora-spec-01.txt, August 1998.

[13] P. Johansson, T. Larsson, N. Hedman, B. Mielczarek, M. Degermark, Scenario-based performance analysis of routing protocols for mobile ad-hoc networks, in: International Conference on Mobile Computing and Networking (MobiCom'99), 1999, pp. 195–206.

[14] G. Pei, M. Gerla, X. Hong, C.-C. Chiang, A wireless hierarchical protocol with group mobility, in: IEEE WCNC, September 1999.

[15] X. Hong, T. Kwon, M. Gerla, D. Gu, G. Pei, A mobility framework for ad hoc wireless networks, in: ACM Second International Conference on Mobile Data Management (MDM), January 2001.

[16] D.A. Maltz, J. Broch, J. Jetcheva, D.B. Johnson, The effects of on-demand behavior in routing protocols for multi-hop wireless ad hoc networks, IEEE Journal on Selected Areas in Communications, Special Issue on Mobile and Wireless Networks, 17 (8) (1999) 1439–1453.

[17] W. Su, S.-J. Lee, M. Gerla, Mobility prediction in wireless networks, in: IEEE MILCOM, October 2000.

[18] S.-J. Lee, M. Gerla, C.-K. Toh, A simulation study of table-driven and on-demand routing protocols for mobile ad hoc networks, IEEE Network 13 (4) (1999) 48–54.

[19] T.H. Cormen, C.E. Leiserson, R.L. Rivest, C. Stein, Introduction to Algorithms, McGraw-Hill, New York, 1998.

Narayanan Sadagopan is a Ph.D. student in the Computer Science Department at the University of Southern California (USC) since Summer 2001. He completed his M.S. in Computer Science at USC (2001) and B.E. in Computer Science and Engineering from National Institute of Technology (NIT), Trichy, India (1998).

He is interested in mathematical modeling and algorithms. Currently, he is applying these techniques to routing in adhoc and sensor networks. URL: http://www-scf.usc.edu/~nsadagop.

Ahmed Helmy received his Ph.D. in Computer Science (1999), M.S. in Electrical Engineering (1995) from the University of Southern California, M.S. in Engineering Mathematics (1994) and B.S. in Electronics and Communications Engineering (1992) from Cairo University, Egypt.

Since 1999, he has been an Assistant Professor of Electrical Engineering at the University of Southern California. In 2002, Dr. Helmy received the National Science Foundation (NSF) CAREER Award, in 2000 he received the USC Zumberge Award, and in 2002 he received the best paper award from the IEEE/IFIP International Conference on Management of Multimedia Networks and Services (MMNS). In 2000, he founded—and is currently directing—the wireless networking laboratory at USC. His current research interests lie in the areas of protocol design and analysis for mobile ad hoc and sensor networks, mobility modeling, design and testing of multicast protocols, IP micro-mobility, and network simulation. URL: http://ceng.usc.edu/~helmy.

Fan Bai received B.S. in Automation Engineering from Tsinghua University, Beijing, China in 1999. Since 1999, he has been pursuing the Ph.D. degree at Department of Electrical Engineering at University of Southern California, California.

His current research interests include protocol design for mobile ad hoc network and wireless sensor network. URL: http://www-scf.usc.edu/~fbai.