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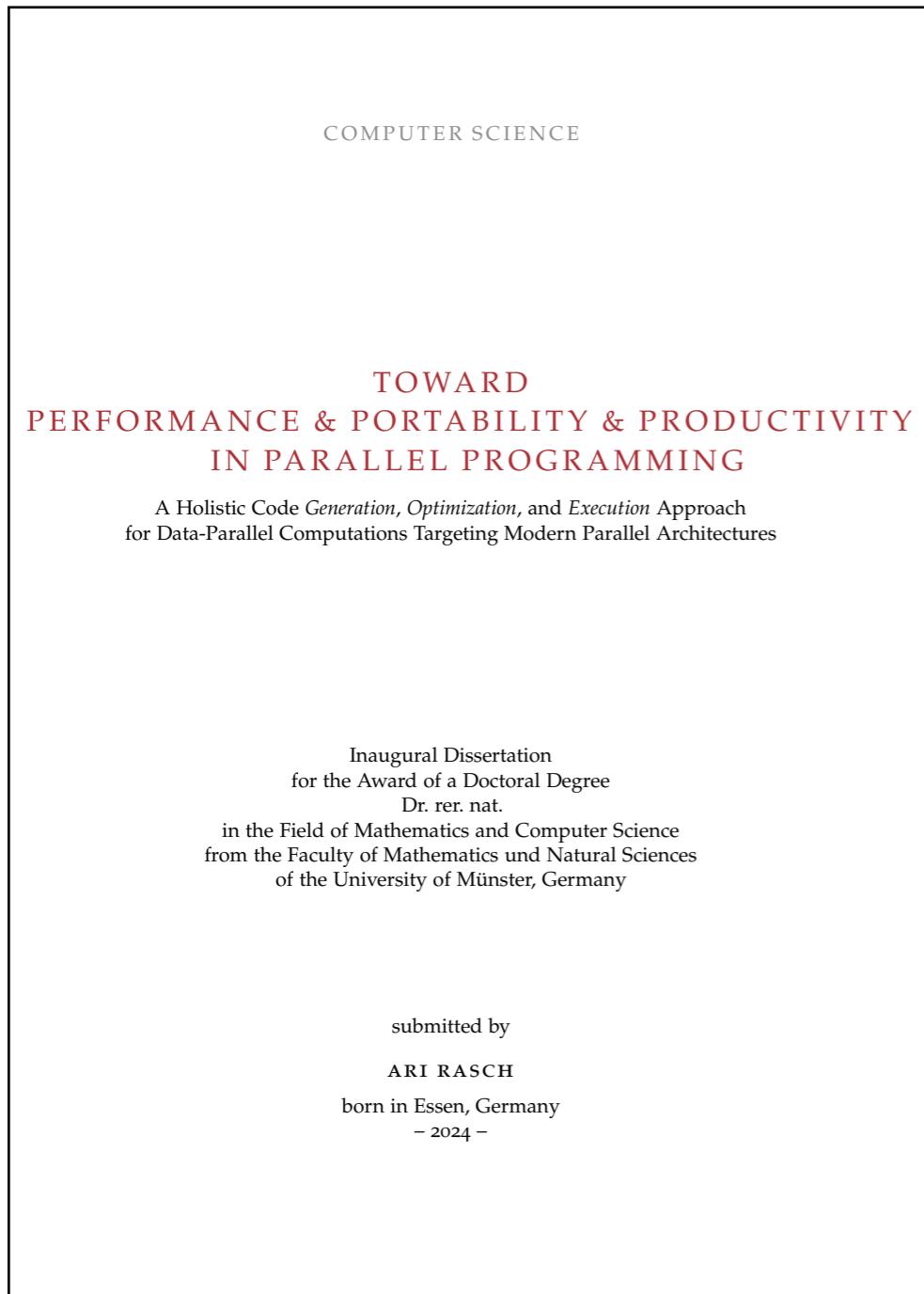
SC24
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Toward **Performance & Portability & Productivity** in Parallel Programming

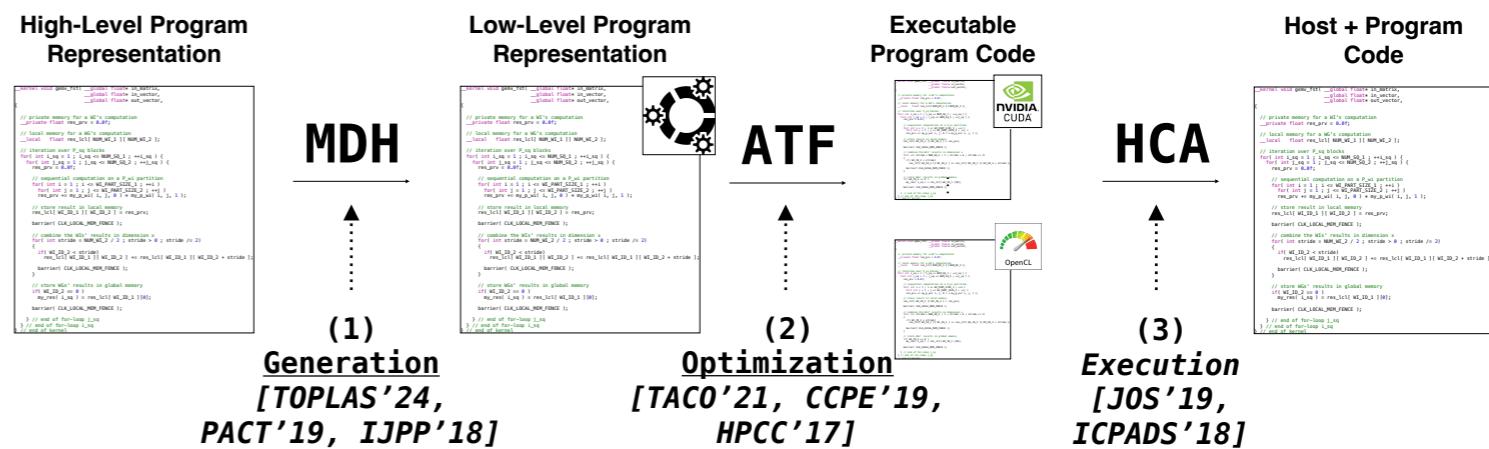
A Holistic *Code Generation, Optimization, and Execution* Approach
for Data-Parallel Computations Targeting Modern Parallel Architectures

Ari Rasch
University of Münster, Germany

Introductory Remark



This thesis describes three major (sub-)projects:



The sub-project complement each other to a holistic approach to code

Generation & Optimization & Execution

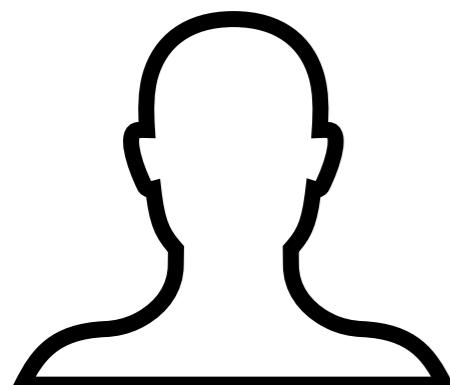
Major Contribution

Medium Contribution

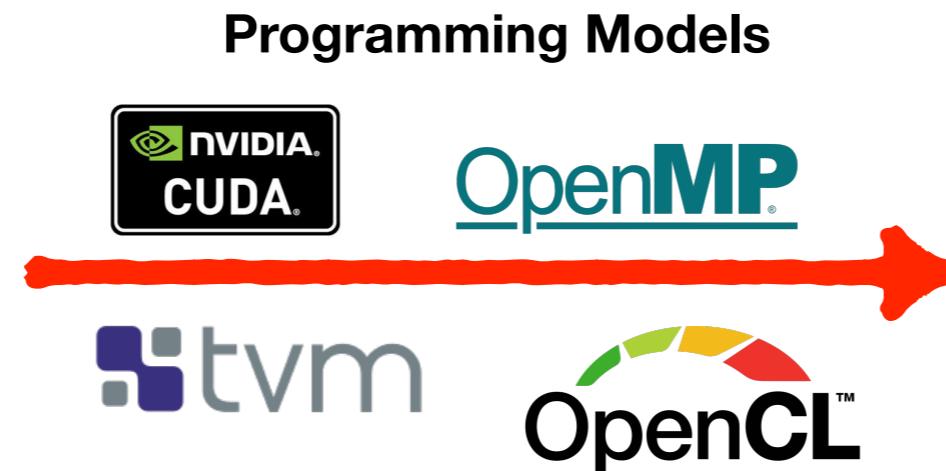
Minor Contribution

Parallel Programming in Today's World

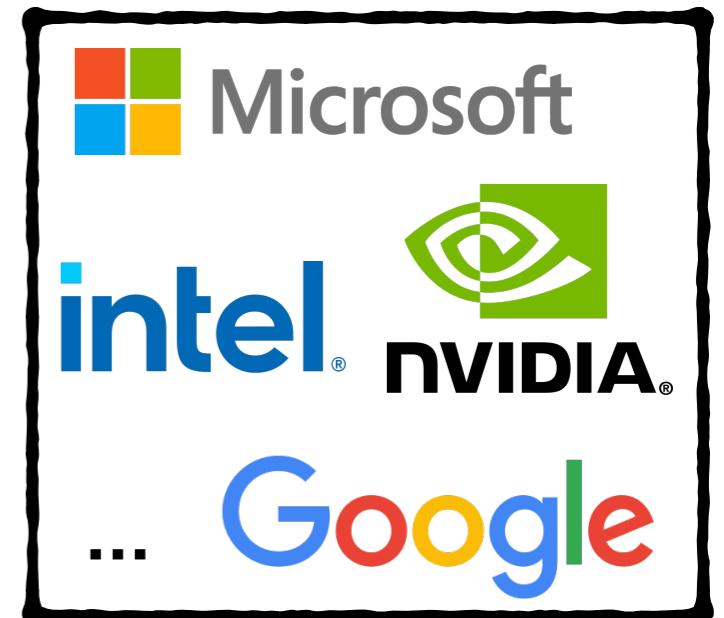
Parallel programming is hard:



Domain Scientist



Parallel Architectures



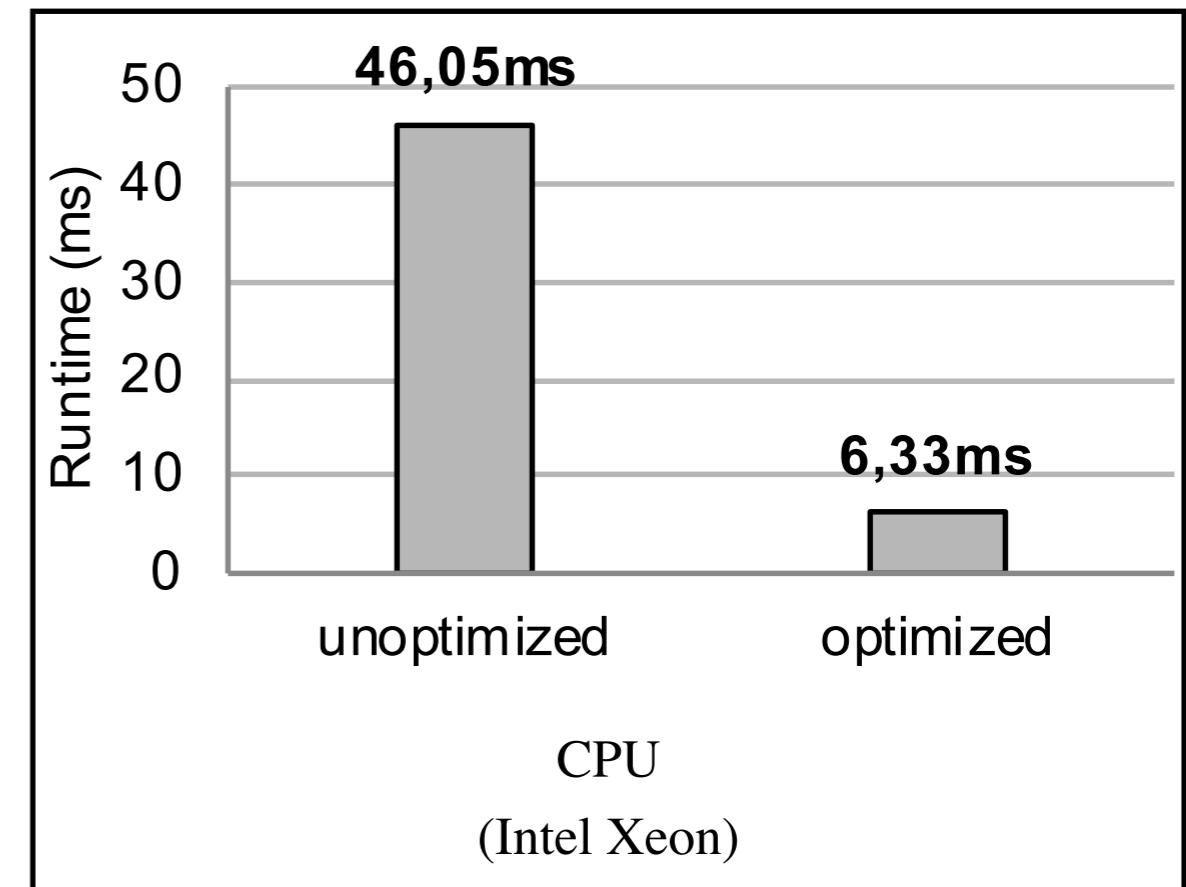
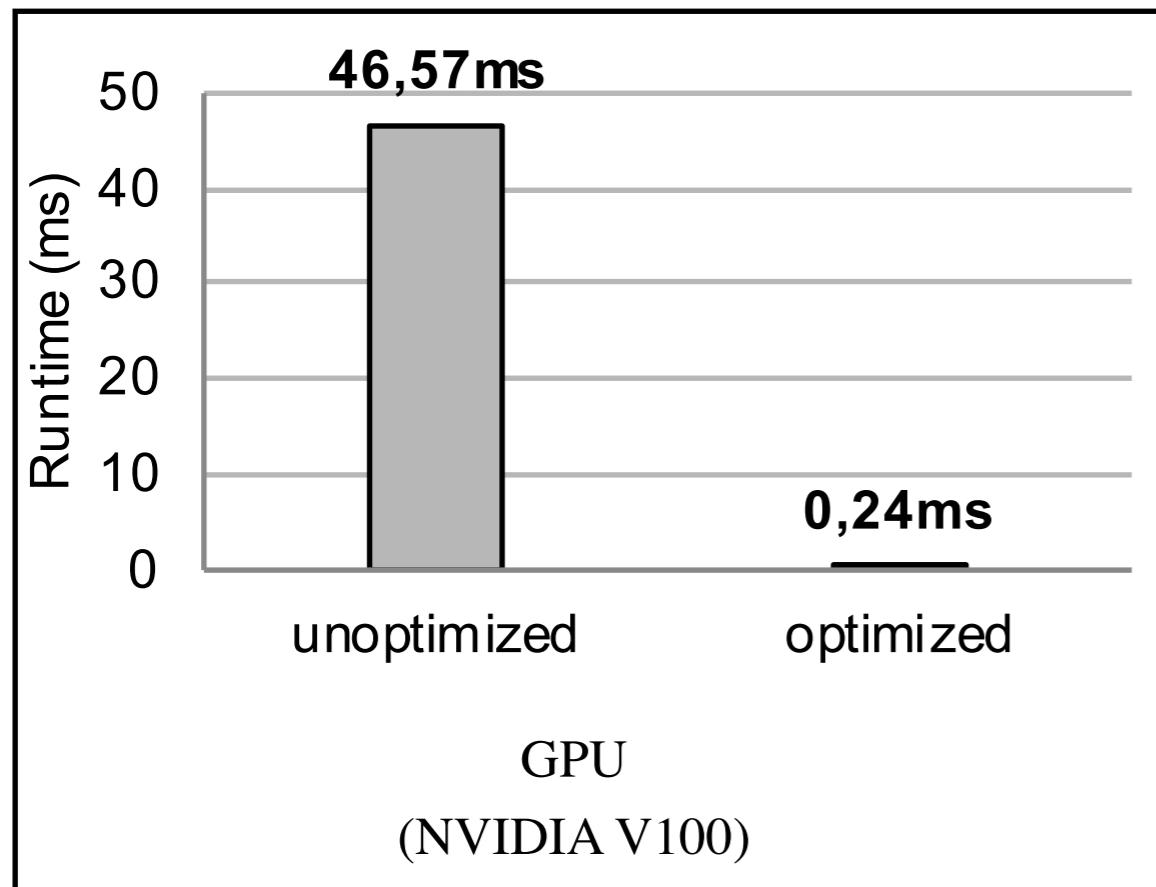
- + domain knowledge
- but lacks hardware & optimization details

Struggle with
simultaneously achieving
Performance & Portability & Productivity

- + tremendous performance
- require advanced & specific optimizations

Challenges: Performance & Portability & Productivity

The *Performance* challenge:

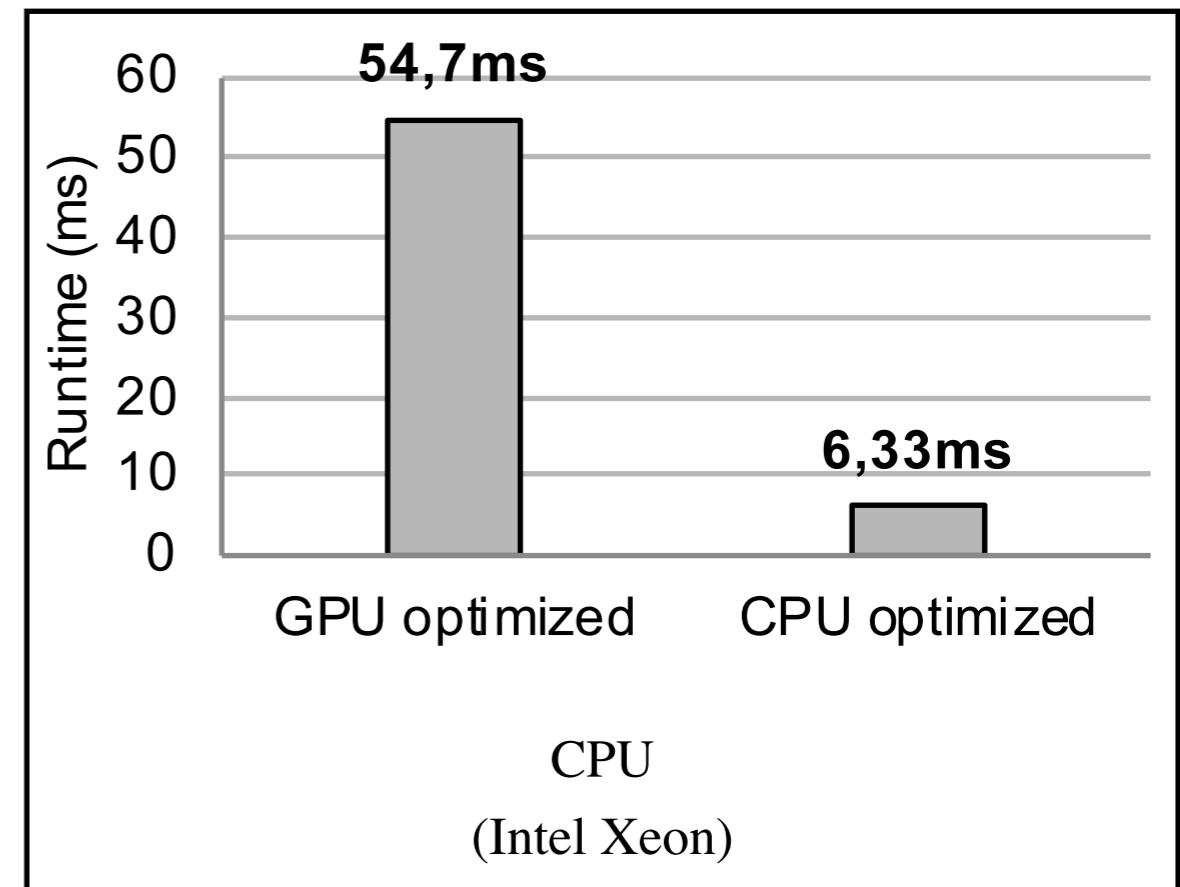
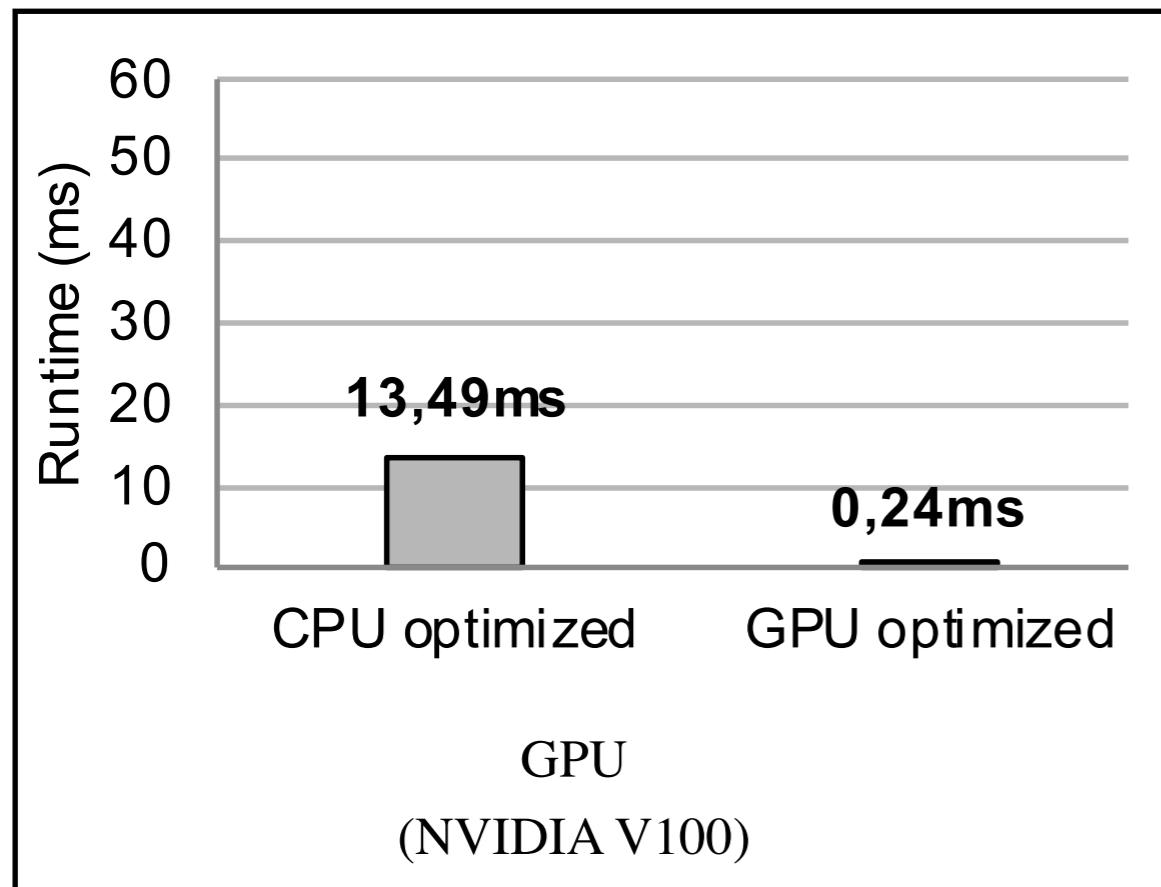


Runtime (lower is better) of unoptimized vs optimized matrix multiplication on GPU (left) and CPU (right).

High *Performance* requires *complex optimizations*

Challenges: Performance & Portability & Productivity

The Portability challenge:



Runtime (lower is better) of GPU/CPU-optimized matrix multiplication
on GPU (left) and CPU (right).

High Portability requires architecture(/data)-specific optimizations

Challenges: Performance & Portability & Productivity

The Productivity challenge:

```
1 __kernel void MatMul( __global const float A[M][K] ,
2                      __global const float B[K][N] ,
3                      __global      float C[M][N] )
4 {
5     int i = get_global_id(0);
6     int j = get_global_id(1);
7
8     for( int k=0 ; k<K ; ++k )
9         C[i][j] += A[i][k] * B[k][j];
10 }
```

Naive OpenCL implementation of matrix multiplication

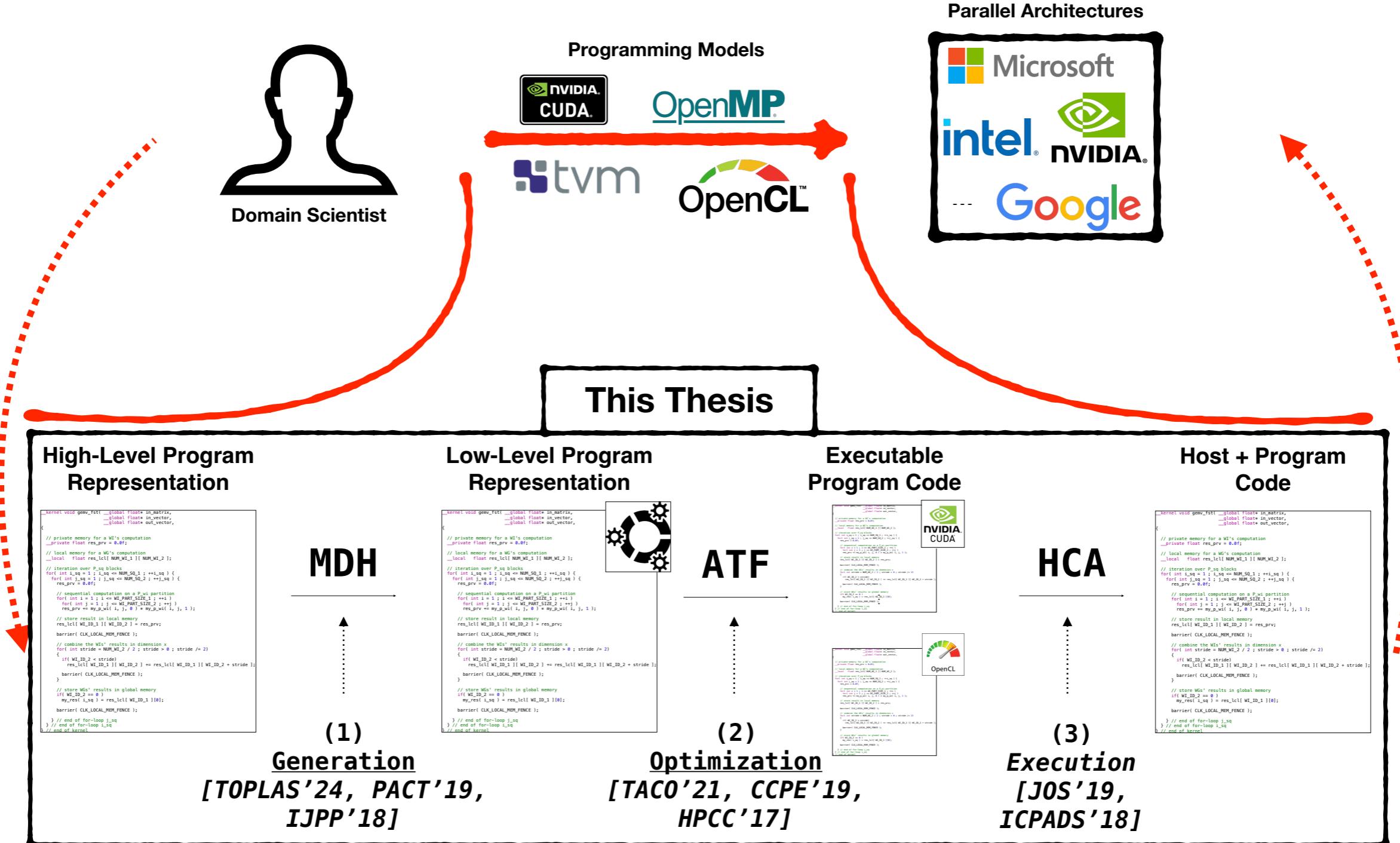
```
1   __kernel void MatMul( /* ... */ )
2   {
3       const size_t i_wg_l_1 = get_group_id(2);
4       // ... 5 lines skipped
5
6       __private TYPE_TS res_p[/*...*/][/*...*/];
7       {
8           // ... 7 lines skipped
9               for (size_t p_iteration_l_1 = 0; p_iteration_l_1 < (2);
10                   ++p_iteration_l_1) {
11                   for (size_t p_iteration_l_2 = 0; p_iteration_l_2 < (1)
12                       ; ++p_iteration_l_2) {
13                       size_t p_iteration_r_1 = 0;
14                       res_p[p_step_l_1][((p_iteration_l_1) * 1 + 0)][(0)][
15                           p_step_l_2][((p_iteration_l_2) * 1 + 0)] = f(
a[((l_step_l_1 * (32 / 1) + ((p_step_l_1 *
1 + (((p_iteration_l_1) * 1 + 0)) / 1) * 1
+ i_wi_l_1 * 1 + (((p_iteration_l_1) * 1 +
0)) % 1))) / 1) * (64 * 1) + i_wg_l_1 * 1 +
(((p_step_l_1 * (2 + (((p_iteration_l_1) *
1 + 0)) / 1) * 1 + i_wi_l_1 * 1 + ((((
p_iteration_l_1) * 1 + 0)) % 1))) % 1))) *
1024 + (((l_step_r_1 * (2 / 1) + ((((
p_step_r_1 * (1 + (((p_iteration_r_1) * 1 +
0)) / 1) * 1 + i_wi_r_1 * 1 + ((((
p_iteration_r_1) * 1 + 0)) % 1))) / 1) * (2
* 1) + i_wg_r_1 * 1 + (((p_step_r_1 * (1) +
((p_iteration_r_1) * 1 + 0)) / 1) * 1 +
i_wi_r_1 * 1 + (((p_iteration_r_1) * 1 + 0)
) % 1))) % 1))),,
14           // ... 107 lines skipped
15       }
```

Optimized OpenCL implementation of matrix multiplication

High ***Productivity*** requires ***automatic optimization***

Contributions of this Thesis

This thesis introduces a novel, holistic approach to ***Generating*** & ***Optimizing*** & ***Executing*** code:



The ultimate goal of **MDH+ATF+HCA** is to simultaneously achieve
Performance & **Portability** & **Productivity**

Outline

This talk(/thesis) is structured into three main parts:

High-Level Program Representation

```
_kernel void gemv_fst( __global float* in_matrix,
                      __global float* in_vector,
                      __global float* out_vector,
```

// private memory for a WI's computation
`__private float res_prv = 0.0f;`

// local memory for a WG's computation
`_local float res_lcl[NUM_WI_1][NUM_WI_2];`

// iteration over P_sq blocks
`for(int i_sq = 1 ; i_sq <= NUM_SO_1 ; ++i_sq) {`
 `for(int j_sq = 1 ; j_sq <= NUM_SO_2 ; ++j_sq) {`
 `res_prv += res_lcl[i_sq][j_sq] * my_p_wi[i, j, 1];`

// sequential computation on a P-wi partition
 `for(int i = 1 ; i <= WI_PART_SIZE_1 ; ++i)`
 `for(int j = 1 ; j <= WI_PART_SIZE_2 ; ++j)`
 `res_prv += my_p_wi[i, j, 0] * my_p_wi[i, j, 1];`

// store result in local memory
`res_lcl[WI_ID_1][WI_ID_2] = res_prv;`

barrier(CLK_LOCAL_MEM_FENCE);

// combine the WIs' results in dimension x
`for(int stride = NUM_WI_2 / 2 ; stride > 0 ; stride /= 2)`
`{`
 `if(WI_ID_2 < stride)`
 `res_lcl[WI_ID_1][WI_ID_2] += res_lcl[WI_ID_1][WI_ID_2 + stride];`

barrier(CLK_LOCAL_MEM_FENCE);

// store WGs' results in global memory
`if(WI_ID_2 == 0)`
 `my_res[i_sq] = res_lcl[WI_ID_1][0];`

barrier(CLK_LOCAL_MEM_FENCE);

} // end of for-loop i_sq
} // end of for-loop i_sq
} // end of kernel

MDH

(1)
Generation
*[TOPLAS'24,
 PACT'19, IJPP'18]*

1. Part: How to *generate* automatically optimizable (auto-tunable) code?

Low-Level Program Representation

```
_kernel void gemv_fst( __global float* in_matrix,
                      __global float* in_vector,
```

// private memory for a WI's computation
`__private float res_prv = 0.0f;`

// local memory for a WG's computation
`_local float res_lcl[NUM_WI_1][NUM_WI_2];`

// iteration over P_sq blocks
`for(int i_sq = 1 ; i_sq <= NUM_SO_1 ; ++i_sq) {`
 `for(int j_sq = 1 ; j_sq <= NUM_SO_2 ; ++j_sq) {`
 `res_prv += res_lcl[i_sq][j_sq] * my_p_wi[i, j, 1];`

// sequential computation on a P-wi partition
 `for(int i = 1 ; i <= WI_PART_SIZE_1 ; ++i)`
 `for(int j = 1 ; j <= WI_PART_SIZE_2 ; ++j)`
 `res_prv += my_p_wi[i, j, 0] * my_p_wi[i, j, 1];`

// store result in local memory
`res_lcl[WI_ID_1][WI_ID_2] = res_prv;`

barrier(CLK_LOCAL_MEM_FENCE);

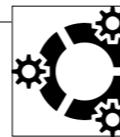
// combine the WIs' results in dimension x
`for(int stride = NUM_WI_2 / 2 ; stride > 0 ; stride /= 2)`
`{`
 `if(WI_ID_2 < stride)`
 `res_lcl[WI_ID_1][WI_ID_2] += res_lcl[WI_ID_1][WI_ID_2 + stride];`

barrier(CLK_LOCAL_MEM_FENCE);

// store WGs' results in global memory
`if(WI_ID_2 == 0)`
 `my_res[i_sq] = res_lcl[WI_ID_1][0];`

barrier(CLK_LOCAL_MEM_FENCE);

} // end of for-loop j_sq
} // end of for-loop i_sq



ATF

(2)
Optimization
*[TACO'21, CCPE'19,
 HPCC'17]*

2. Part: How to automatically *optimize* (auto-tune) code?

Executable Program Code

```
_kernel void gemv_fst( __global float* in_matrix,
```

// private memory for a WI's computation
`__private float res_prv = 0.0f;`

// local memory for a WG's computation
`_local float res_lcl[NUM_WI_1][NUM_WI_2];`

// iteration over P_sq blocks
`for(int i_sq = 1 ; i_sq <= NUM_SO_1 ; ++i_sq) {`
 `for(int j_sq = 1 ; j_sq <= NUM_SO_2 ; ++j_sq) {`
 `res_prv += res_lcl[i_sq][j_sq] * my_p_wi[i, j, 1];`

// sequential computation on a P-wi partition
 `for(int i = 1 ; i <= WI_PART_SIZE_1 ; ++i)`
 `for(int j = 1 ; j <= WI_PART_SIZE_2 ; ++j)`
 `res_prv += my_p_wi[i, j, 0] * my_p_wi[i, j, 1];`

// store result in local memory
`res_lcl[WI_ID_1][WI_ID_2] = res_prv;`

barrier(CLK_LOCAL_MEM_FENCE);

// combine the WIs' results in dimension x
`for(int stride = NUM_WI_2 / 2 ; stride > 0 ; stride /= 2)`
`{`
 `if(WI_ID_2 < stride)`
 `res_lcl[WI_ID_1][WI_ID_2] += res_lcl[WI_ID_1][WI_ID_2 + stride];`

barrier(CLK_LOCAL_MEM_FENCE);

// store WGs' results in global memory
`if(WI_ID_2 == 0)`
 `my_res[i_sq] = res_lcl[WI_ID_1][0];`

barrier(CLK_LOCAL_MEM_FENCE);

} // end of for-loop j_sq
} // end of for-loop i_sq



HCA

(3)
Execution
*[JOS'19,
 ICPADS'18]*

```
_kernel void gemv_fst( __global float* in_matrix,
```

// private memory for a WI's computation
`__private float res_prv = 0.0f;`

// local memory for a WG's computation
`_local float res_lcl[NUM_WI_1][NUM_WI_2];`

// iteration over P_sq blocks
`for(int i_sq = 1 ; i_sq <= NUM_SO_1 ; ++i_sq) {`
 `for(int j_sq = 1 ; j_sq <= NUM_SO_2 ; ++j_sq) {`
 `res_prv += res_lcl[i_sq][j_sq] * my_p_wi[i, j, 1];`

// sequential computation on a P-wi partition
 `for(int i = 1 ; i <= WI_PART_SIZE_1 ; ++i)`
 `for(int j = 1 ; j <= WI_PART_SIZE_2 ; ++j)`
 `res_prv += my_p_wi[i, j, 0] * my_p_wi[i, j, 1];`

// store result in local memory
`res_lcl[WI_ID_1][WI_ID_2] = res_prv;`

barrier(CLK_LOCAL_MEM_FENCE);

// combine the WIs' results in dimension x
`for(int stride = NUM_WI_2 / 2 ; stride > 0 ; stride /= 2)`
`{`
 `if(WI_ID_2 < stride)`
 `res_lcl[WI_ID_1][WI_ID_2] += res_lcl[WI_ID_1][WI_ID_2 + stride];`

barrier(CLK_LOCAL_MEM_FENCE);

// store WGs' results in global memory
`if(WI_ID_2 == 0)`
 `my_res[i_sq] = res_lcl[WI_ID_1][0];`

barrier(CLK_LOCAL_MEM_FENCE);

} // end of for-loop j_sq
} // end of for-loop i_sq

OpenCL

Code Generation via MDH



Overview Getting Started Code Examples Publications Citations Contact



Multi-Dimensional Homomorphisms (MDH)

An Algebraic Approach Toward Performance & Portability & Productivity for Data-Parallel Computations

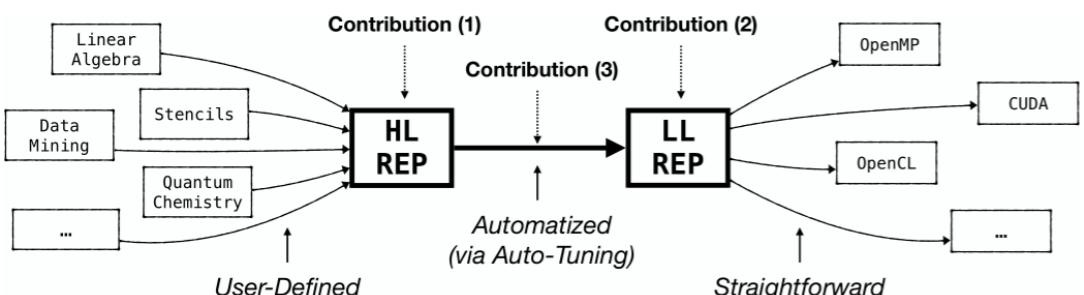
Overview

The approach of **Multi-Dimensional Homomorphisms (MDH)** is an algebraic formalism for systematically reasoning about *de-composition* and *re-composition* strategies of data-parallel computations (such as linear algebra routines and stencil computations) for the memory and core hierarchies of state-of-the-art parallel architectures (GPUs, multi-core CPU, multi-device and multi-node systems, etc).

The MDH approach (formally) introduces:

1. *High-Level Program Representation* (*Contribution 1*) that enables the user conveniently implementing data-parallel computations, agnostic from hardware and optimization details;
2. *Low-Level Program Representation* (*Contribution 2*) that expresses device- and data-optimized de- and re-composition strategies of computations;
3. *Lowering Process* (*Contribution 3*) that fully automatically lowers a data-parallel computation expressed in its high-level program representation to an optimized instance in its low-level representation, based on concepts from automatic performance optimization (a.k.a. *auto-tuning*), using the *Auto-Tuning Framework (ATF)*.

The MDH's low-level representation is designed such that *Code Generation* from it (e.g., in *OpenMP* for CPUs, *CUDA* for NVIDIA GPUs, or *OpenCL* for multiple kinds of architectures) becomes straightforward.



Our *Experiments* report encouraging results on GPUs and CPUs for MDH as compared to state-of-practice approaches, including NVIDIA *cuBLAS/cuDNN* and Intel *oneMKL/oneDNN* which are hand-optimized libraries provided by vendors.

<https://mdh-lang.org>

ACM TOPLAS 2024

(De/Re)-Composition of Data-Parallel Computations via Multi-Dimensional Homomorphisms

ARI RASCH, University of Muenster, Germany

Data-parallel computations, such as linear algebra routines and stencil computations, constitute one of the most relevant classes in parallel computing, e.g., due to their importance for deep learning. Efficiently de-composing such computations for the memory and core hierarchies of modern architectures and re-composing the computed intermediate results back to the final result—we say *(de/re)-composition* for short—is key to achieve high performance for these computations on, e.g., GPU and CPU. Current high-level approaches to generating data-parallel code are often restricted to a particular subclass of data-parallel computations and architectures (e.g., only linear algebra routines on only GPU or only stencil computations), and/or the approaches rely on a user-guided optimization process for a well-performing (de/re)-composition of computations, which is complex and error prone for the user.

We formally introduce a systematic (de/re)-composition approach, based on the algebraic formalism of Multi-Dimensional Homomorphisms (MDHs). Our approach is designed as general enough to be applicable to a wide range of data-parallel computations and for various kinds of target parallel architectures. To efficiently target the deep and complex memory and core hierarchies of contemporary architectures, we exploit our introduced (de/re)-composition approach for a correct-by-construction, parametrized cache blocking, and parallelization strategy. We show that our approach is powerful enough to express, in the same formalism, the (de/re)-composition strategies of different classes of state-of-the-art approaches (scheduling-based, polyhedral, etc.), and we demonstrate that the parameters of our strategies enable systematically generating code that can be fully automatically optimized (auto-tuned) for the particular target architecture and characteristics of the input and output data (e.g., their sizes and memory layouts). Particularly, our experiments confirm that via auto-tuning, we achieve higher performance than state-of-the-art approaches, including hand-optimized solutions provided by vendors (such as NVIDIA cuBLAS/cuDNN and Intel oneMKL/oneDNN), on real-world datasets and for a variety of data-parallel computations, including linear algebra routines, stencil and quantum chemistry computations, data mining algorithms, and computations that recently gained high attention due to their relevance for deep learning.

CCS Concepts: • Computing methodologies → Parallel computing methodologies; Machine learning;
• Theory of computation → Program semantics; • Software and its engineering → Compilers;

Additional Key Words and Phrases: Code generation, data parallelism, auto-tuning, GPU, CPU, OpenMP, CUDA, OpenCL, linear algebra, stencils computation, quantum chemistry, data mining, deep learning

A full version of this article is provided by Rasch [2024], which presents our novel concepts with all of their formal details. In contrast to the full version, this article relies on a simplified formal foundation for better illustration and easier understanding. We often refer the interested reader to Rasch [2024] for formal details that should not be required for understanding the basic ideas and concepts of our approach.

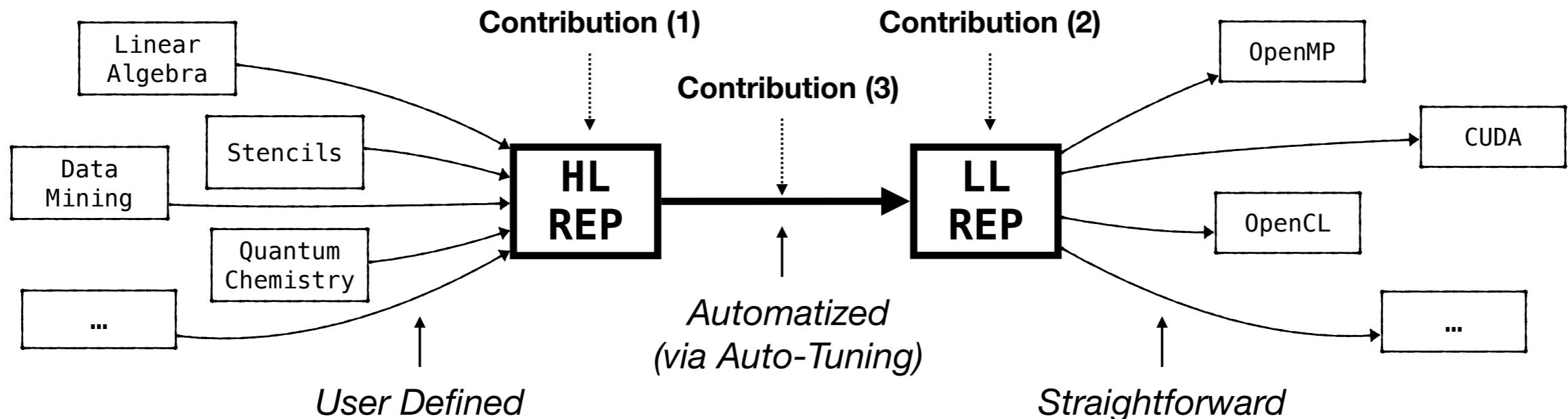
This work was funded by the Deutsche Forschungsgemeinschaft (DFG, German Research Foundation)—project PPP-DL (470527619).

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Code Generation via MDH



MDH is a (formal) framework for expressing & optimizing data-parallel computations:



1. **Contribution 1 (HL-REP):** defines *data parallelism* & introduces *higher-order functions* for expressing data-parallel computations, agnostic from hardware and optimization details while still capturing high-level information relevant for generating high-performing code
2. **Contribution 2 (LL-REP):** allows expressing and reasoning about optimizations for the memory and core hierarchies of contemporary parallel architectures & generalizes these optimizations to apply to arbitrary combinations of data-parallel computations and architectures
3. **Contribution 3 (\rightarrow):** introduces a *structured optimization process* – for arbitrary combinations of data-parallel computations and parallel architectures – that allows *fully automatic* optimization (auto-tuning)

Code Generation via MDH



Example: MatVec expressed in MDH

```
MatVec<T∈TYPE| I,K∈ℕ> := out_view<T>( w:(i,k)↔(i) ) ∘  
                                md_hom<I,K>( *, (#+,+) ) ∘  
                                inp_view<T,T>( M:(i,k)↔(i,k) , v:(i,k)↔(k) )
```

High-Level Representation of MatVec

What is happening here:

- `inp_view` captures the accesses to input data
- `md_hom` expresses the data-parallel computation
- `out_view` captures the accesses to output data

¹We can generate such MDH expressions also automatically from straightforward (annotated) code in Python or C

Code Generation via MDH



md_hom	f	\otimes_1	\otimes_2	\otimes_3	\otimes_4
Dot	*	+	/	/	/
MatVec	*	++	+	/	/
MatMul	*	++	++	+	/
MatMul ^T	*	++	++	+	/
bMatMul	*	++	++	++	+

Views	inp_view			out_view		
	A	B	C			
Dot	(k) \mapsto (k)	(k) \mapsto (k)	(k) \mapsto ()			
MatVec	(i, k) \mapsto (i, k)	(i, k) \mapsto (k)	(i, k) \mapsto (i)			
MatMul	(i, j, k) \mapsto (i, k)	(i, j, k) \mapsto (k, j)	(i, j, k) \mapsto (i, j)			
MatMul ^T	(i, j, k) \mapsto (k, i)	(i, j, k) \mapsto (j, k)	(i, j, k) \mapsto (j, i)			
bMatMul	(b, i, j, k) \mapsto (b, i, k)	(b, i, j, k) \mapsto (b, k, j)	(b, i, j, k) \mapsto (b, i, j)			

1) Linear Algebra Routines

md_hom	f	\otimes_1	\otimes_2
MBBS	id	$\text{++}_{\text{prefix-sum}}(+)$	+

Views	inp_view		out_view	
	A	Out		
MBBS	(i, j) \mapsto (i, j)	(i) \mapsto (i)		

8) Maximum Bottom Box Sum

md_hom	f	\otimes_1	\otimes_2
Jacobi1D	J _{1D}	++	/
Jacobi2D	J _{2D}	++	++

Views	inp_view		out_view	
	I	0		
Jacobi1D	(i) \mapsto (i+0), (i) \mapsto (i+1), (i) \mapsto (i+2)	(i) \mapsto (i)		
Jacobi2D	(i, j) \mapsto (i, j+1), (i, j) \mapsto (i+1, j), ...	(i, j) \mapsto (i, j)		

3) Jacobi Stencils

md_hom	f	\otimes_1
map(f)	f	++
reduce(\oplus)	id	\oplus
reduce(\oplus, \otimes)	(x) \mapsto (x, x)	(\oplus, \otimes)

Views	inp_view			out_view		
	I	0 ₁	0 ₂			
map(f)	(i) \mapsto (i)	(i) \mapsto (i)	/			
reduce(\oplus)	(i) \mapsto (i)	(i) \mapsto ()	/			
reduce(\oplus, \otimes)	(i) \mapsto (i)	(i) \mapsto ()	(i) \mapsto ()			

6) Map/Reduce Patterns

md_hom	f	\otimes_1	\otimes_2	\otimes_3	\otimes_4	\otimes_5	\otimes_6	\otimes_7	\otimes_8	\otimes_9	\otimes_{10}
Conv2D	*	++	++	+	+	/	/	/	/	/	/
MCC	*	++	++	++	++	+	+	+	/	/	/
MCC_Capsule	*	++	++	++	++	+	+	+	++	++	+

2) Convolution Stencils

Views	inp_view		out_view	
	I	F	0	
Conv2D	(p, q, r, s) \mapsto (p+r, q+s)	(p, q, r, s) \mapsto (r, s)	(p, q, r, s) \mapsto (p, q)	
MCC	(n, p, ...) \mapsto (n, p+r, q+s, c)	(n, p, ...) \mapsto (k, r, s, c)	(n, p, ...) \mapsto (n, p, q, k)	
MCC_Capsule	(n, p, ...) \mapsto (n, p+r, q+s, c, cm, ck)	(n, p, ...) \mapsto (k, r, s, c, ck, cn)	(n, p, ...) \mapsto (n, p, q, k, cm, cn)	

4) Probabilistic Record Linkage

Views	inp_view		out_view	
	N	E	M	
PRL	wght	++	maxPRL	
PRL	(i, j) \mapsto (i)	(i, j) \mapsto (j)	(i, j) \mapsto (i)	

7) Scan Pattern

md_hom	f	\otimes_1	Views	inp_view	out_view
scan(\oplus)	id	$\text{++}_{\text{prefix-sum}}(\oplus)$	scan(\oplus)	(i) \mapsto (i)	(i) \mapsto (i)

5) Histogram

Views	inp_view			out_view		
	Bins	Elems	Out			
Histo	f _{Histo}	++	+	Histo	(b, e) \mapsto (b)	(b, e) \mapsto (e)

MDH is capable of expressing a wide range of data-parallel computations from popular domains

Code Generation via MDH



Performance Evaluation: (via runtime comparison)

Deep Learning	NVIDIA Ampere GPU									
	ResNet-50				VGG-16				MobileNet	
	Training		Inference		Training		Inference		Training	Inference
	MCC	MatMul	MCC	MatMul	MCC	MatMul	MCC	MatMul	MCC	MCC
TVM+Ansor	1.00	1.26	1.05	2.22	0.93	1.42	0.88	1.14	0.94	1.00
PPCG	3456.16	8.26	-	7.89	1661.14	7.06	5.77	5.08	2254.67	7.55
PPCG+ATF	3.28	2.58	13.76	5.44	4.26	3.92	9.46	3.73	3.31	10.71
cuDNN	0.92	-	1.85	-	1.22	-	1.94	-	1.81	2.14
cuBLAS	-	1.58	-	2.67	-	0.93	-	1.04	-	-
cuBLASEx	-	1.47	-	2.56	-	0.92	-	1.02	-	-
cuBLASLt	-	1.26	-	1.22	-	0.91	-	1.01	-	-



Highlights only

MDH speedup over

- TVM: **0.88x – 2.22x**
- PPCG: **2.58x – 13.76x**
- (cuBLAS/cuDNN: **0.91x – 2.67x**)

Deep Learning	Intel Skylake CPU									
	ResNet-50				VGG-16				MobileNet	
	Training		Inference		Training		Inference		Training	Inference
	MCC	MatMul	MCC	MatMul	MCC	MatMul	MCC	MatMul	MCC	MCC
TVM+Ansor	1.53	1.05	1.14	1.20	1.97	1.14	2.38	1.27	3.01	1.40
Pluto	355.81	49.57	364.43	13.93	130.80	93.21	186.25	36.30	152.14	75.37
Pluto+ATF	13.08	19.70	170.69	6.57	3.11	6.29	53.61	8.29	3.50	25.41
oneDNN	0.39	-	5.07	-	1.22	-	9.01	-	1.05	4.20
oneMKL	-	0.44	-	1.09	-	0.88	-	0.53	-	-
oneMKL(JIT)	-	6.43	-	8.33	-	27.09	-	9.78	-	-



MDH speedup over

- TVM: **1.05 – 3.01x**
- Pluto: **6.29x – 364.43x**
- (oneMKL/oneDNN: **0.39x – 9.01x**)

Case Study “Deep Learning” for which most competitors are highly optimized (most challenging for us!)

Significantly higher speedups for other case studies,
e.g., >170x over TVM on GPU already for straightforward dot products

Code Optimization via ATF



Overview Getting Started Code Examples Publications Citations Contact



Auto-Tuning Framework (ATF)

Efficient Auto-Tuning of Parallel Programs with Constrained Tuning Parameters

Overview

The **Auto-Tuning Framework (ATF)** is a general-purpose auto-tuning approach: given a program that is implemented as generic in performance-critical program parameters (a.k.a. *tuning parameters*), such as sizes of tiles and numbers of threads, ATF fully automatically determines a hardware- and data-optimized configuration of such parameters.

Key Feature of ATF

A key feature of ATF is its support for **Tuning Parameter Constraints**. Parameter constraints allow auto-tuning programs whose tuning parameters have so-called *interdependencies* among them, e.g., the value of one tuning parameter has to evenly divide the value of another tuning parameter.

ATF's support for parameter constraints is important: modern parallel programs target novel parallel architectures, and such architectures typically have deep memory and core hierarchies thus requiring constraints on tuning parameters, e.g., the value of a tile size tuning parameter on an upper memory layer has to be a multiple of a tile size value on a lower memory layer.

For such parameters, ATF introduces novel concepts for **Generating & Storing & Exploring** the search spaces of constrained tuning parameters, thereby contributing to a substantially more efficient overall auto-tuning process for such parameters, as confirmed in our **Experiments**.

Generality of ATF

For wide applicability, ATF is designed as generic in:

1. The target program's **Programming Language**, e.g., C/C++, CUDA, OpenMP, or OpenCL. ATF offers *pre-implemented cost functions* for conveniently auto-tuning C/C++ programs, as well as CUDA and OpenCL kernels which require host code for their execution which is automatically generated and executed by ATF's pre-implemented CUDA and OpenCL cost functions. ATF also offers a pre-implemented *generic cost function* that can be used for conveniently auto-tuning programs in any other programming language different from C/C++, CUDA, and OpenCL.

<https://atf-tuner.org>

ACM TACO 2021

Efficient Auto-Tuning of Parallel Programs with Interdependent Tuning Parameters via Auto-Tuning Framework (ATF)

ARI RASCH and RICHARD SCHULZE, University of Muenster, Germany
MICHEL STEUWER, University of Edinburgh, United Kingdom
SERGEI GORLATCH, University of Muenster, Germany

Auto-tuning is a popular approach to program optimization: it automatically finds good configurations of a program's so-called tuning parameters whose values are crucial for achieving high performance for a particular parallel architecture and characteristics of input/output data. We present three new contributions of the Auto-Tuning Framework (ATF), which enable a key advantage in *general-purpose auto-tuning*: efficiently optimizing programs whose tuning parameters have *interdependencies* among them. We make the following contributions to the three main phases of general-purpose auto-tuning: (1) ATF *generates* the search space of interdependent tuning parameters with high performance by efficiently exploiting parameter constraints; (2) ATF *stores* such search spaces efficiently in memory, based on a novel chain-of-trees search space structure; (3) ATF *explores* these search spaces faster, by employing a multi-dimensional search strategy on its chain-of-trees search space representation. Our experiments demonstrate that, compared to the state-of-the-art, general-purpose auto-tuning frameworks, ATF substantially improves generating, storing, and exploring the search space of interdependent tuning parameters, thereby enabling an efficient overall auto-tuning process for important applications from popular domains, including stencil computations, linear algebra routines, quantum chemistry computations, and data mining algorithms.

CCS Concepts: • General and reference → Performance; • Computer systems organization → Parallel architectures; • Software and its engineering → Parallel programming languages;

Additional Key Words and Phrases: Auto-tuning, parallel programs, interdependent tuning parameters

ACM Reference format:

Ari Rasch, Richard Schulze, Michel Steuwer, and Sergei Gorlatch. 2021. Efficient Auto-Tuning of Parallel Programs with Interdependent Tuning Parameters via Auto-Tuning Framework (ATF). *ACM Trans. Archit. Code Optim.* 18, 1, Article 1 (January 2021), 26 pages.

<https://doi.org/10.1145/3427093>

This is a new paper, not an extension of a conference paper.

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Code Optimization via ATF



Advantage of ATF over state-of-the-art general-purpose AT approaches:

ATF finds values of performance-critical parameters with

interdependencies among them

via optimized processes to

generating & storing & exploring

the spaces of interdependent parameters

→ We illustrate ATF by comparing it to MIT's ***OpenTuner*** [PACT'14] & ***CLTune*** [MCSoC'15] which is the foundation of many related approaches (e.g., **KernelTuner** & **KTT**).

Note: BaCO [ASPLOS'23] & KTT recently adopted the ATF techniques to also efficiently handle interdependent tuning parameters.

Code Optimization via ATF



How does ATF achieve its efficiency for *interdependent tuning parameters*:

ATF introduces
parameter constraints

```
tuner.addParameter( "tp_1", T1 );
tuner.addParameter( "tp_2", T2 );
// ...
tuner.addConstraint(
  [](T1 tp_1, T2 tp_2, ... ) -> bool
  { /* ... */ }
```

CLTune constraints

```
tuner.addParameter( "tp_1", R1, [](T1 tp_1) -> bool { /* ... */ } );
tuner.addParameter( "tp_2", R2, [](T2 tp_2) -> bool { /* ... */ } );
```

ATF constraints

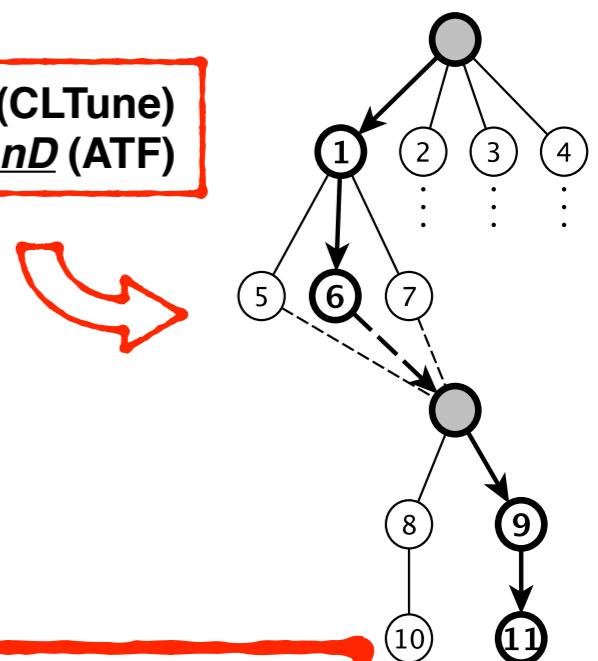
ATF introduces the
CoT (Chain-of-Trees) space structure

```
SP := [ (1,1) | (2,1) | (2,2) | ... ]
```

CLTune search space



ATF
CoT search space



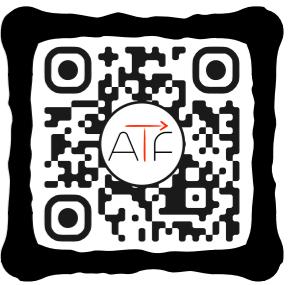
Defined on:
search space (CLTune)
vs. *parameters* (ATF)

verbose & 1D (CLTune)
vs. *compact & nD* (ATF)

ATF introduces a novel
constraint design and search **space structure**
to efficiently **generate** & **store** & **explore** constrained search spaces

Highlights only

Code Optimization via ATF



ATF is able to auto-tune **modern** parallel computations, e.g., for **GPUs & CPUs**:

Stencil

ATF is able to auto-tune CONV to:

>40x higher performance than CONV+**CLTune** on CPU **>10⁴x** higher performance than CONV+**CLTune** on GPU

>3x higher performance than **Intel MKL-DNN** on CPU **>15x** higher performance than **NVIDIA cuDNN** on GPU



Linear Algebra

ATF is able to auto-tune GEMM to:

>2x higher performance than GEMM+**CLTune** on CPU **>120x** higher performance than GEMM+**CLTune** on GPU

>2x higher performance than **Intel MKL** on CPU **>2x** higher performance than **NVIDIA cuBLAS** on GPU



Quantum Chemistry

ATF is able to auto-tune CCSD(T) to:

>2x higher performance than **Facebook TC** on GPU CLTune **fails!** (too high search space generation time)

Data Mining

ATF is able to auto-tune PRL to:

>1.6x higher performance than PRL+**CLTune** on CPU **>1.07x** higher perform. than PRL+**CLTune** on GPU

OpenTuner fails for all studies

Code Execution via HCA



Overview Contact

HCA

Host Code Abstraction (HCA)
*A High-Level Abstraction for Host Code Programming
Designed for Distributed, Heterogeneous Systems*

Overview
The **Host Code Abstraction (HCA)** is a high-level programming abstraction that simplifies implementing and optimizing so-called host code which is required in modern parallel programming approaches (e.g., **CUDA** and **OpenCL**) to execute code on the devices of distributed, heterogeneous systems.
More details will follow soon!

Contact

<https://hca-project.org>

Journal of Supercomputing 2019

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<https://doi.org/10.1007/s11227-019-02829-2>



dOCAL: high-level distributed programming with OpenCL and CUDA

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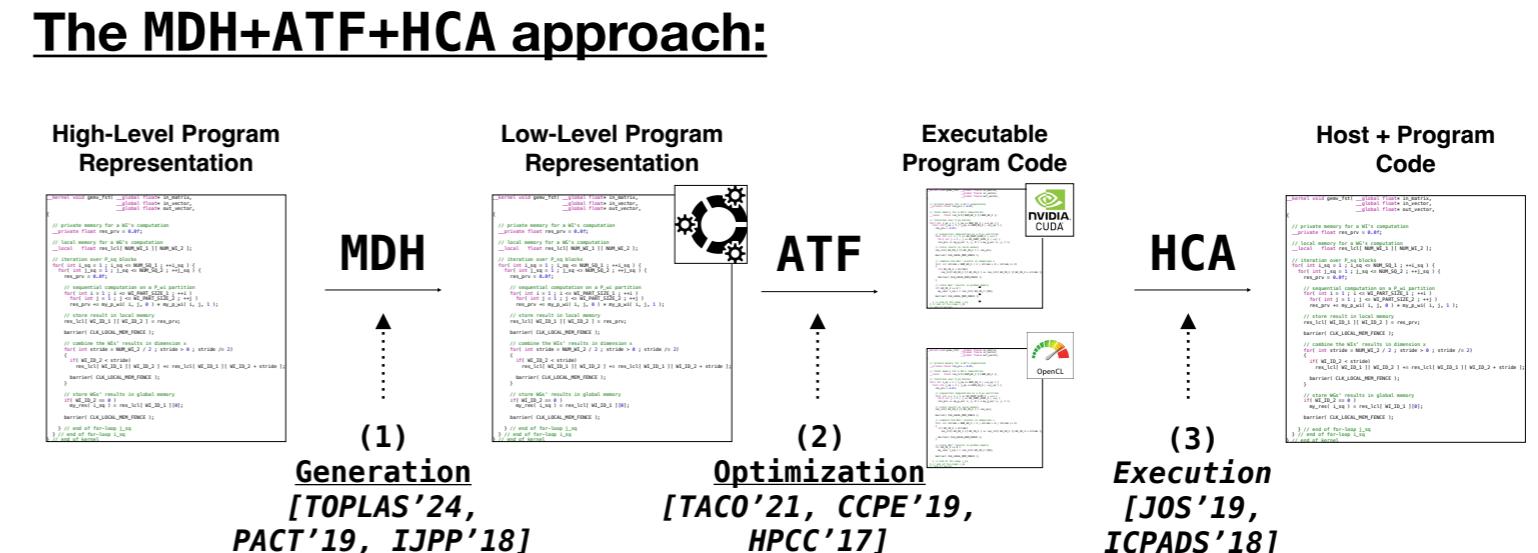
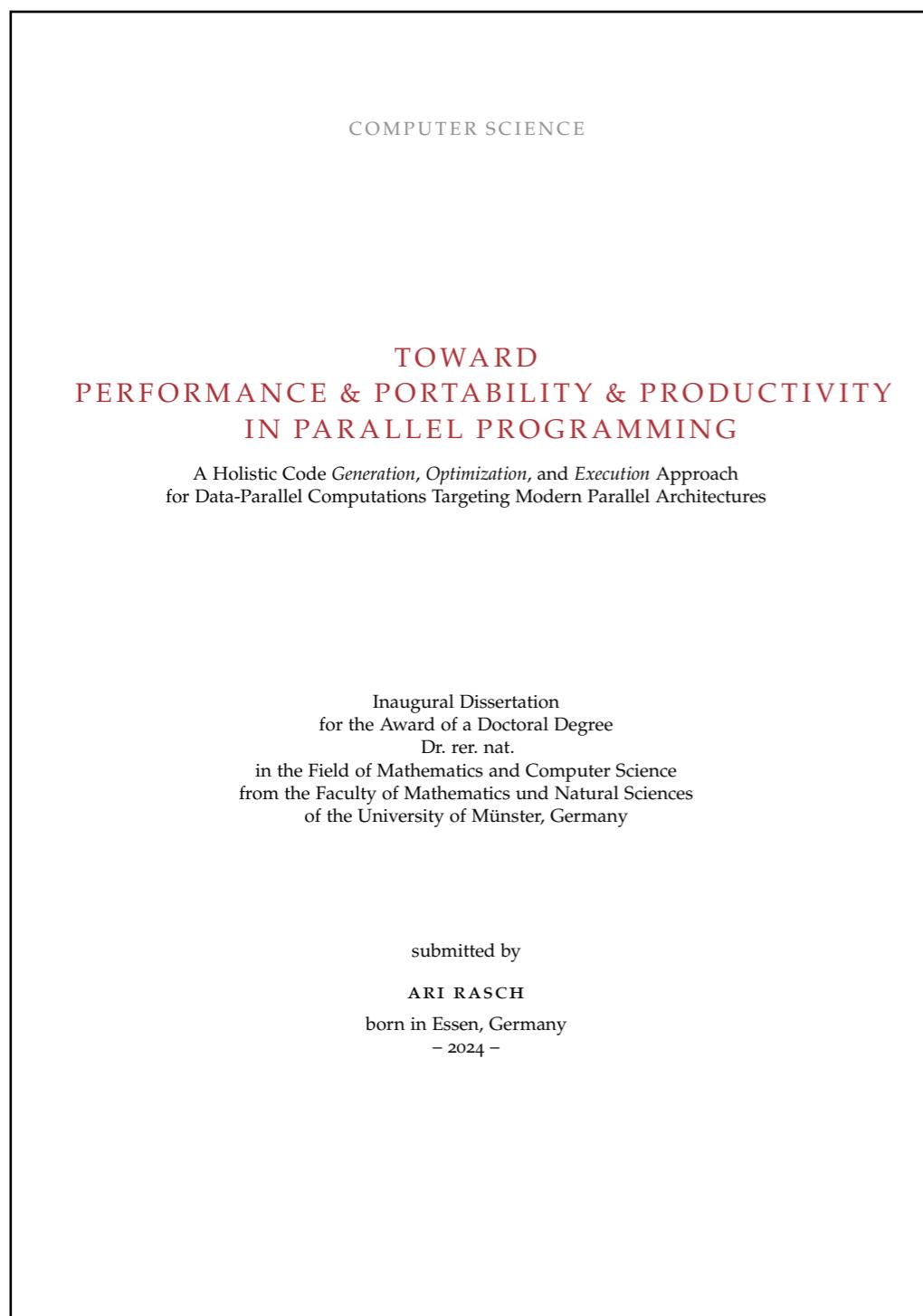
Abstract

In the state-of-the-art parallel programming approaches OpenCL and CUDA, so-called host code is required for program's execution. Efficiently implementing host code is often a cumbersome task, especially when executing OpenCL and CUDA programs on systems with multiple nodes, each comprising different devices, e.g., multi-core CPU and graphics processing units; the programmer is responsible for explicitly managing node's and device's memory, synchronizing computations with data transfers between devices of potentially different nodes and for optimizing data transfers between devices' memories and nodes' main memories, e.g., by using pinned main memory for accelerating data transfers and overlapping the transfers with computations. We develop distributed OpenCL/CUDA abstraction layer (dOCAL)—a novel high-level C++ library that simplifies the development of host code. dOCAL combines major advantages over the state-of-the-art high-level approaches: (1) it simplifies implementing both OpenCL and CUDA host code by providing a simple-to-use, high-level abstraction API; (2) it supports executing arbitrary OpenCL and CUDA programs; (3) it allows conveniently targeting the devices of different nodes by automatically managing node-to-node communications; (4) it simplifies implementing data transfer optimizations by providing different, specially allocated memory regions, e.g., pinned main memory for overlapping data transfers with computations; (5) it optimizes memory management by automatically avoiding unnecessary data transfers; (6) it enables interoperability between OpenCL and CUDA host code for systems with devices from different vendors. Our experiments show that dOCAL significantly simplifies the development of host code for heterogeneous and distributed systems, with a low runtime overhead.

Skipped for brevity

Summary

The **MDH+ATF+HCA** approach achieves **Performance & Portability & Productivity** for data-parallel computations targeting modern parallel architectures:



- The three sub-projects — **MDH** & **ATF** & **HCA** — complement each other to a holistic code *Generation & Optimization & Execution* approach
- There are many (promising) future directions for **MDH** & **ATF** & **HCA** (one part of thesis dedicated to FW)



MDH



ATF



HCA