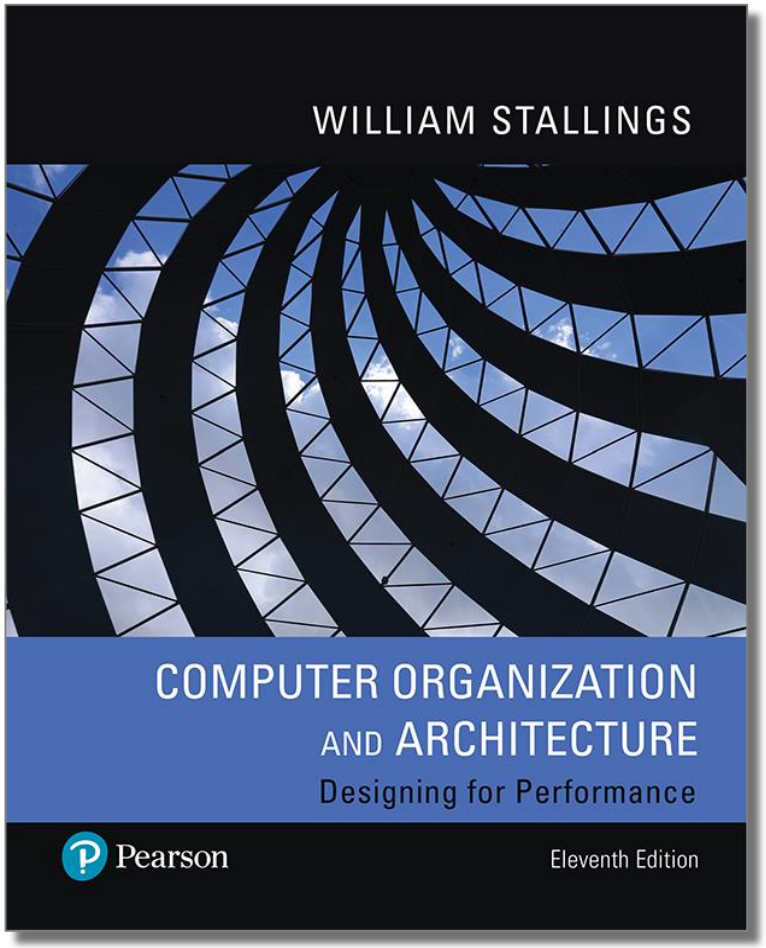
Computer Organization and Architecture

Designing for Performance

11th Edition



Chapter 14

Instruction Sets: Addressing Modes and Formats

Addressing Modes

- So far, we focused on what an instruction does
- In this chapter, we will focus on to specify the operands and operations of instructions
- How is the address of an operand specified?
- How are the bits of an instruction organized to define the operand addresses and operation of that instruction.
- The address field of an instruction is typically relatively small.
- We want to reference a large range of locations in the main memory

Addressing Modes

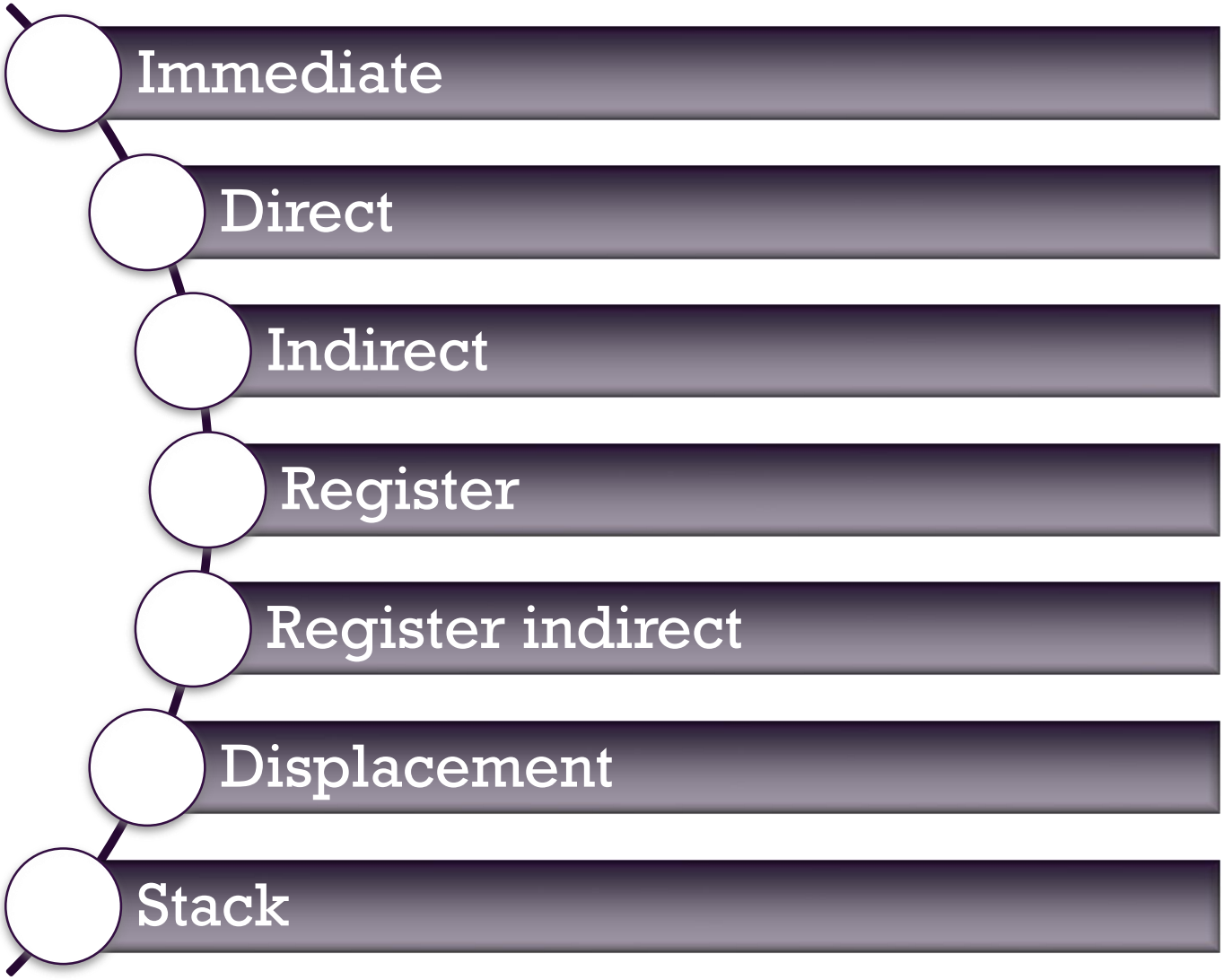


Figure 14.1
Addressing Modes

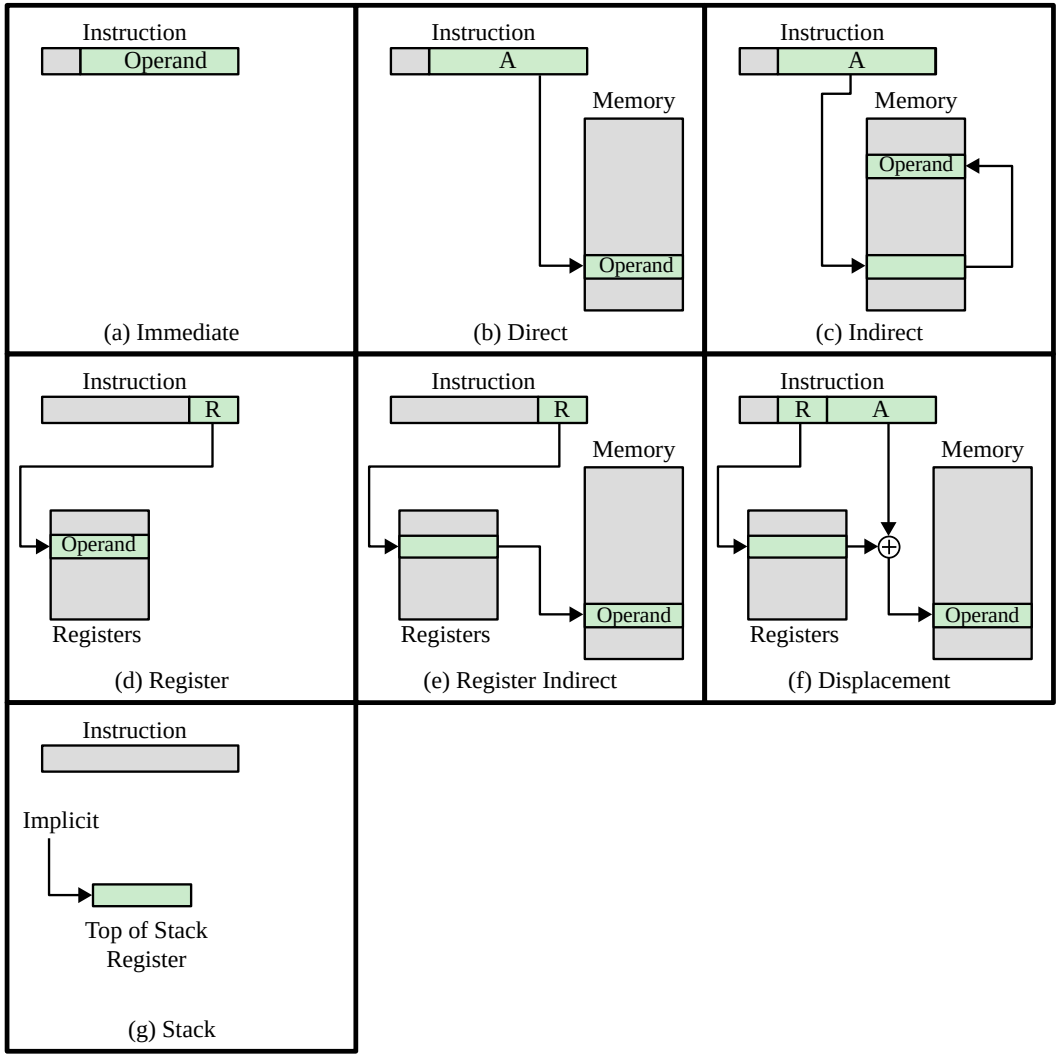


Table 14.1
Basic Addressing Modes

Mode	Algorithm	Principal Advantage	Principal Disadvantage
Immediate	Operand = A	No memory reference	Limited operand magnitude
Direct	EA = A	Simple	Limited address space
Indirect	EA = (A)	Large address space	Multiple memory references
Register	EA = R	No memory reference	Limited address space
Register indirect	EA = (R)	Large address space	Extra memory reference
Displacement	EA = A + (R)	Flexibility	Complexity
Stack	EA = top of stack	No memory reference	Limited applicability

Immediate Addressing

- Simplest form of addressing
- Operand = A
- This mode can be used to define and use constants or set initial values of variables
  - Typically the number will be stored in twos complement form
  - The leftmost bit of the operand field is used as a sign bit
- Advantage:
  - No memory reference other than the instruction fetch is required to obtain the operand, thus saving one memory or cache cycle in the instruction cycle
- Disadvantage:
  - The size of the number is restricted to the size of the address field, which, in most instruction sets, is small compared with the word length

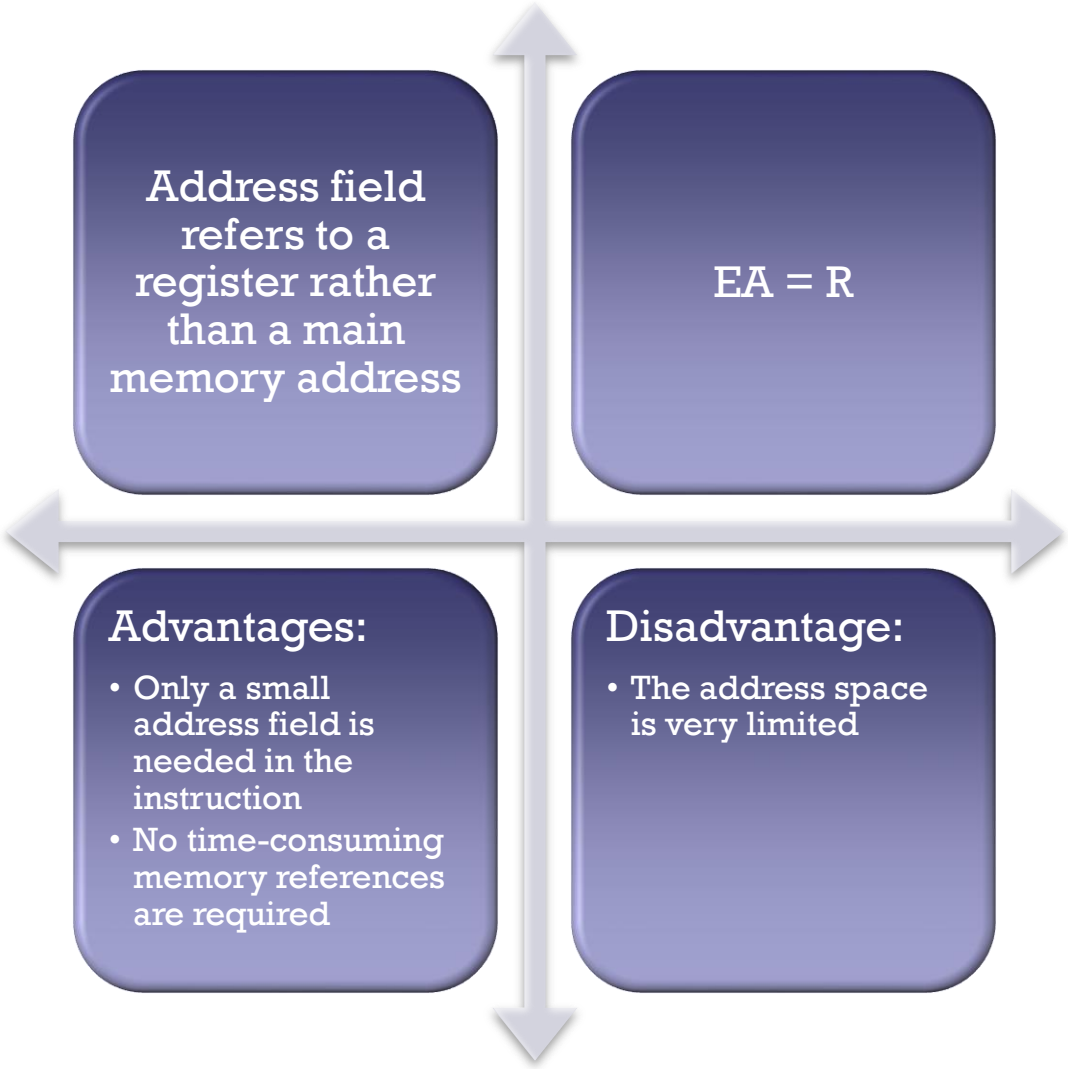
Direct Addressing

- Address field contains the effective address of the operand
- Effective address (EA) = address field (A)
- Was common in earlier generations of computers
- Requires only one memory reference and no special calculation
- Limitation is that it provides only a limited address space

Indirect Addressing

- Reference to the address of a word in memory which contains a full-length address of the operand
- EA = (A)
  - Parentheses are to be interpreted as meaning *contents of*
- Advantage:
  - For a word length of *N* an address space of 2^N is now available
- Disadvantage:
  - Instruction execution requires two memory references to fetch the operand
    - One to get its address and a second to get its value
- A rarely used variant of indirect addressing is multilevel or cascaded indirect addressing
  - EA = (. . . (A) . . .)
  - Disadvantage is that three or more memory references could be required to fetch an operand

Register Addressing



Register Indirect Addressing

- Analogous to indirect addressing
  - The only difference is whether the address field refers to a memory location or a register
- EA = (R)
- Address space limitation of the address field is overcome by having that field refer to a word-length location containing an address
- Uses one less memory reference than indirect addressing

Displacement Addressing

- Combines the capabilities of direct addressing and register indirect addressing
- EA = A + (R)
- Requires that the instruction have two address fields, at least one of which is explicit
  - The value contained in one address field (value = A) is used directly
  - The other address field refers to a register whose contents are added to A to produce the effective address
- Most common uses:
  - Relative addressing
  - Base-register addressing
  - Indexing

Relative Addressing

The implicitly referenced register is the program counter (PC)

- The next instruction address is added to the address field to produce the EA
- Typically the address field is treated as a twos complement number for this operation
- Thus the effective address is a displacement relative to the address of the instruction

Exploits the concept of locality

Saves address bits in the instruction if most memory references are relatively near to the instruction being executed

Base-Register Addressing

- The referenced register contains a main memory address and the address field contains a displacement from that address
- The register reference may be explicit or implicit
- Exploits the locality of memory references
- Convenient means of implementing segmentation
- In some implementations a single segment base register is employed and is used implicitly
- In others the programmer may choose a register to hold the base address of a segment and the instruction must reference it explicitly

Indexing

- The address field references a main memory address and the referenced register contains a positive displacement from that address
- The method of calculating the EA is the same as for base-register addressing
- An important use is to provide an efficient mechanism for performing iterative operations
- Autoindexing
  - Automatically increment or decrement the index register after each reference to it
  - EA = A + (R)
  - (R) □(R) + 1
- Postindexing
  - Indexing is performed after the indirection
  - EA = (A) + (R)
- Preindexing
  - Indexing is performed before the indirection
  - EA = (A + (R))

Stack Addressing

- A stack is a linear array of locations
  - Sometimes referred to as a *pushdown list* or *last-in-first-out queue*
- A stack is a reserved block of locations
  - Items are appended to the top of the stack so that the block is partially filled
- Associated with the stack is a pointer whose value is the address of the top of the stack
  - The stack pointer is maintained in a register
  - Thus references to stack locations in memory are in fact register indirect addresses
- Is a form of implied addressing
- The machine instructions need not include a memory reference but implicitly operate on the top of the stack

Figure 14.2
x86 Addressing Mode Calculation

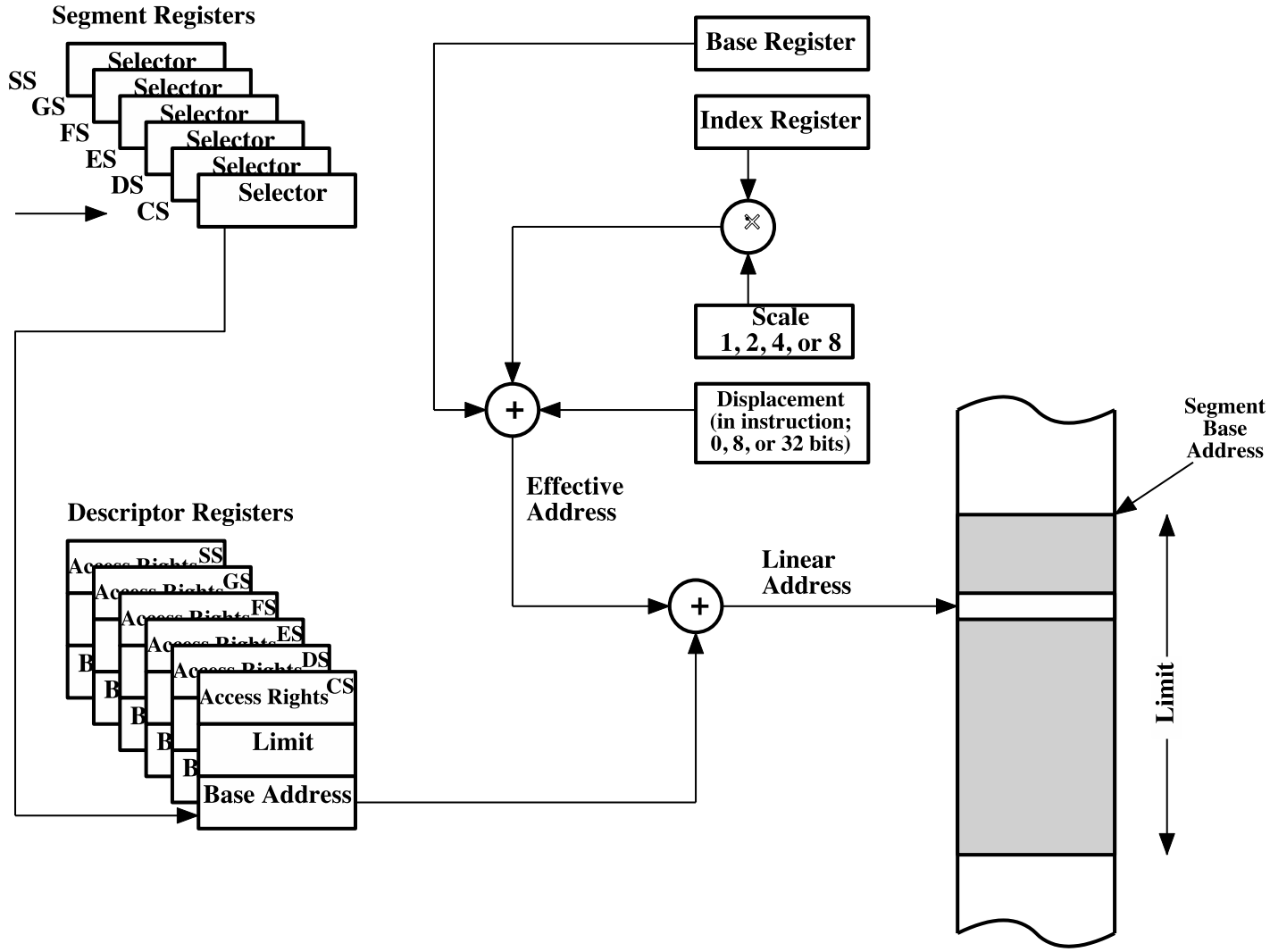


Table 14.2
x86 Addressing Modes

Mode	Algorithm
Immediate	Operand = A
Register Operand	LA = R
Displacement	LA = (SR) + A
Base	LA = (SR) + (B)
Base with Displacement	LA = (SR) + (B) + A
Scaled Index with Displacement	LA = (SR) + (I) × S + A
Base with Index and Displacement	LA = (SR) + (B) + (I) + A
Base with Scaled Index and Displacement	LA = (SR) + (I) × S + (B) + A
Relative	LA = (PC) + A

LA = linear address
(X) = contents of X
SR = segment register
PC = program counter
A = contents of an address field in the instruction

R = register
B = base register
I = index register
S = scaling factor

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