## 【没选修实验课的同学专用的编程题】--Miller-Rabin

用C/C++实现Miller-Rabin素数测试算法。请提交固定链接的MarkDown页面。选修实验课的同学也可以考虑实现,不过就不强制要求。

```
// https://github.com/gou4shi1/oj/blob/master/poj1811.cpp
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <iostream>
#include <time.h>
using namespace std;
typedef long long 11;
const int INF = 0x7f7f7f7f;
const int maxn = 100 + 10;
// ret = a * b ( mod n )
11 mul mod ( 11 a, 11 b, 11 n ) {
   a %= n;
   b %= n;
   ll ret = 0;
   11 \text{ tmp} = a;
   while (b) {
       if (b & 1 ) {
           ret += tmp;
           if ( ret > n )
               ret -= n;
       }
        tmp <<= 1;
       if (tmp > n)
           tmp -= n;
       b >>= 1;
   return ret;
}
// ret = a^n ( mod MOD )
ll pow mod ( ll a, ll n, ll mod ) {
   ll ret = 1;
   11 tmp = a % mod;
   while (n) {
       if ( n & 1 )
           ret = mul mod ( ret, tmp, mod );
       tmp = mul_mod ( tmp, tmp, mod );
       n >>= 1;
   return ret;
}
```

```
/*
* witness whether a^(n-1) = 1 \pmod{n}
* n-1 = x*2<sup>t</sup>, 中间判断
* /
bool witness ( ll a, ll n, ll x, ll t ) {
  ll ret = pow mod (a, x, n);
   11 last = ret;
   for ( int i = 1; i <= t; ++i ) {
       ret = mul mod ( ret, ret, n );
       if ( ret == 1 && last != 1 && last != n - 1 )
          return true;
       last = ret;
   }
    if ( ret != 1 )
       return true;
   else
      return false;
}
/*
* 合数可能被判成素数, 概率(1/4)^NUM
* 素数一定返回true
* n-1 = x*2<sup>t</sup>, 中间判断
* /
bool Miller Rabin ( ll n ) {
   if ( n < 2 )
       return false;
   if (n == 2)
       return true;
   if (!(n & 1))
       return false;
   11 x = n - 1, t = 0;
   while (!(x & 1))
      x >>= 1, ++t;
   srand ( time ( NULL ) );
   const static int NUM = 8;
   for ( int i = 0; i < NUM; ++i ) {
       ll a = rand () % ( n - 1 ) + 1;
       if ( witness ( a, n, x, t ) )
          return false;
   return true;
}
ll factor[ maxn ]; //素因数分解结果
int tot; //素因子个数
// gcd要保证a>b否则死循环
11 gcd ( ll a, ll b ) {
   11 t;
```

```
while (b) {
     t = a;
      a = b;
      b = t % b;
   return a >= 0 ? a : -a;
}
/*
* 分解x, c是常数参数?算导上是-1
* 生日悖论来提高概率
* 可能返回n本身
* /
ll polard rho ( ll n, ll c ) {
  11 i = 1, k = 2;
   srand ( time ( NULL ) );
   11 x = rand () % (n - 1) + 1;
   11 y = x;
   while ( 1 ) {
      ++i;
       x = ( mul mod ( x, x, n ) + c ) % n;
      11 d = gcd (y - x, n);
      if ( d != 1 && d != n )
         return d;
      if (y == x)
         return n;
      if ( i == k )
         y = x, k += k;
  }
void findfac ( ll n ) {
   if (n == 1)
      return;
   if ( Miller Rabin ( n ) ) {
      factor[tot++] = n;
      return;
   }
   const static int k = 107;
   11 p = n;
   int c = k;
   while (p \ge n)
     p = polard rho (p, c--);
   findfac ( p );
   findfac ( n / p );
}
/*
* poj1811
* 输入n(2<n<2<sup>54</sup>),如果是素数,输出prime
* 否则输出最小的素因子
* /
```

```
int main () {
 #ifdef LOCAL
    freopen ( "in", "r", stdin );
    // freopen("out", "w", stdout);
 #endif
    int T;
    scanf ( "%d", &T );
    while ( T-- ) {
       11 n;
        scanf ( "%lld", &n );
        if ( Miller_Rabin ( n ) )
          printf ( "Prime\n" );
        else {
           tot = 0;
           findfac ( n ); // factor其实就是所有的因子了
           11 ans = factor[ 0 ];
           for ( int i = 1; i < tot; ++i )
              ans = min ( ans, factor[ i ] );
           printf ( "%lld\n", ans );
    }
   return 0;
```