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1:

$$y^2 = x^3 + x + 6 \mod 11, K_{prA} = 6, K_{pubB} = (5,9)$$

session $key = K_{prA}K_{pubB} = 6(5,9)$

we should use double and add algorithm to reach 6 in the fastest way

$$P \rightarrow 2P \rightarrow 3P \rightarrow 6P$$

$$P + P = (5,9) + (5,9)$$

$$P = Q \rightarrow s = \frac{3x_1^2 + 1}{2y_1} \mod 11 = 76(18)^{-1} = 10 \times 8 =$$

 $80 \ mod \ 11 = 3$

$$x_3 = s^2 - x_1 - x_2 = -1 \mod 11 = 10$$

$$y_3 = s(x_1 - x_3) - y_1 = 3 \times 6 - 9 = 9$$

$$2P = (10,9)$$

$$3P = 2P + P = (10.9) + (5.9)$$

$$P \neq Q \rightarrow s = \frac{y_2 - y_1}{x_2 - x_1} = 0$$

$$x_3 = 0 - 10 - 5 = -15 \mod 11 = -4 \mod 11 = 7$$

$$y_3 = -9 \mod 11 = 2$$

$$3P = (7,2)$$

$$6P = 3P + 3P = (7,2) + (7,2)$$

$$P = Q \rightarrow s = s = \frac{3x_1^2 + 1}{2y_1} \mod 11 = 148(4)^{-1} = 5 \times 3 = 15 \mod 11 = 4$$

$$x_3 = 16 - 7 - 7 = 2$$

$$y_3 = 4(7 - 2) - 2 = 7$$

$$6P = (2,7)$$

$$session \ key = (2,7)$$

2.1:

$$a = 2, b = 2$$

 $4a^3 + 27b^2 = 4(2)^3 + 27(2)^2 = 32 + 108 = 140 \mod 17 \equiv 4$

2.2:

$$(2,7) \neq (5,2) \rightarrow s = \frac{y_2 - y_1}{x_2 - x_1} \mod p \rightarrow$$

$$s = (2-7)(5-2)^{-1} \mod 17 = 12 \times 6 = 72 \mod 17 = 4$$

$$x_3 = s^2 - x_1 - x_2 \mod 17 \rightarrow x_3 = 16 - 2 - 5 = 9$$

$$y_3 = s(x_1 - x_3) - y_1 = 4(2-9) - 7 = 40 - 7 = 16 \rightarrow$$

$$(x_3, y_3) = (9,16)$$

2.3:

from the book example we know #E=19 now we should calculate upper and lower bounds:

upper bound :
$$P + 1 + 2\sqrt{P} = 17 + 1 + 2 \times 4.12 = 26.24$$

lower bound :
$$P + 1 - 2\sqrt{P} = 17 + 1 - 2 \times 4.12 = 9.76$$

 $9.76 \le 19 \le 26.24$

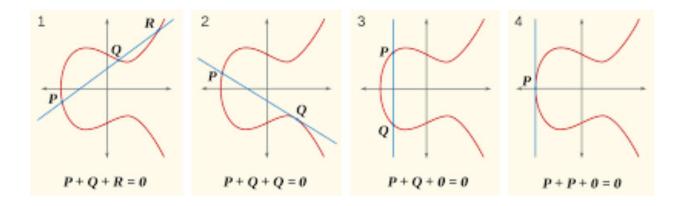
2.4:

since the group cardinality is 19, $\varphi(19) = 18$ which means all the elements except θ are primitive element.

3.1:

all the points on x-axis

3.2:



4:

$$p = 31, \alpha = 3, \beta = 6, x = 10, keys = (17,5)$$
 and $(13,5)$

4.1:

$$\alpha^{x} = \alpha^{10} = 3^{10} = 25$$

first verification step:

$$t = \beta^r . r^s \mod p$$

$$t_1 = 6^{17} \cdot 17^5 \mod 31 = 25$$

$$t_2 = 6^{13}.13^5 \mod 31 = 5$$

$$t_1 = \alpha^x$$
 , $t_2 \neq \alpha^x \rightarrow first\ signature\ is\ valid$

4.2:

$$k_E \in \{0,1,...,31-2\}$$
 such that $gcd(k_E,30) = 1$

so we can choose every number in the given range except 2,3,5,6,10,15, the overall amount is (29-0)+1-6=24

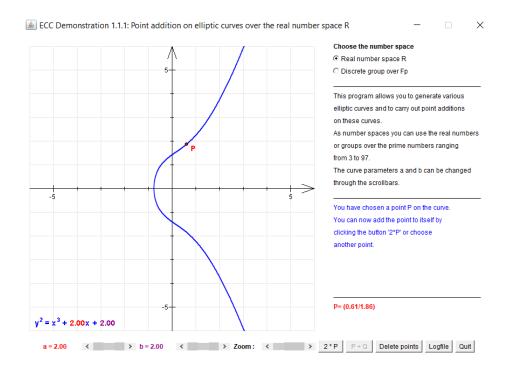
5:

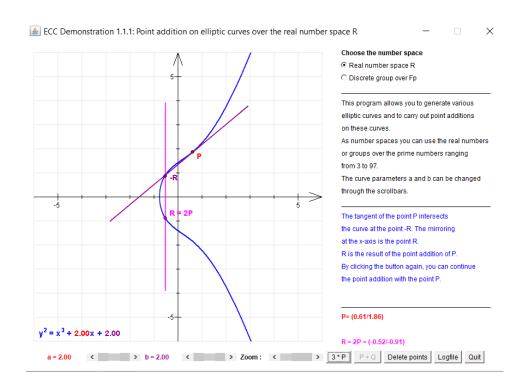
first Oscar receives public key (9797,131) then he chooses a random number smaller that 9797 ,s = 100 then Oscar computes $x \equiv s^e \rightarrow 100^{131} = 9190 \ mod \ 9797$ then Oscar sends (9190,100) to Alice so when Alice checks $x = s^e$ she verifies the signature

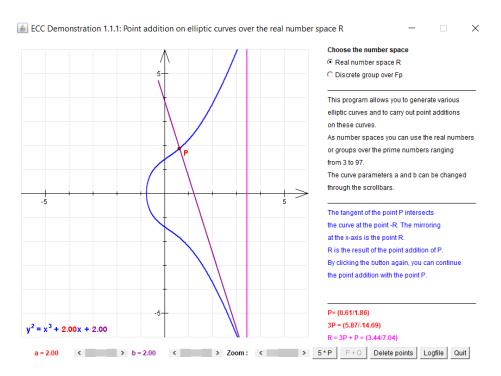
6:

1.

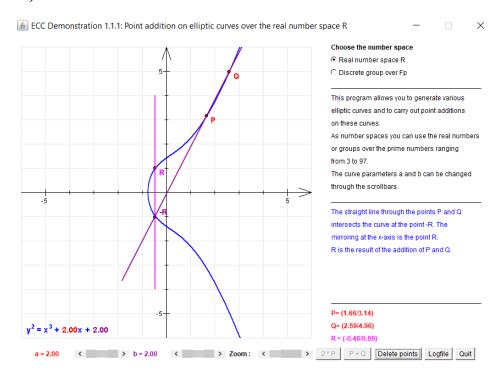
a)





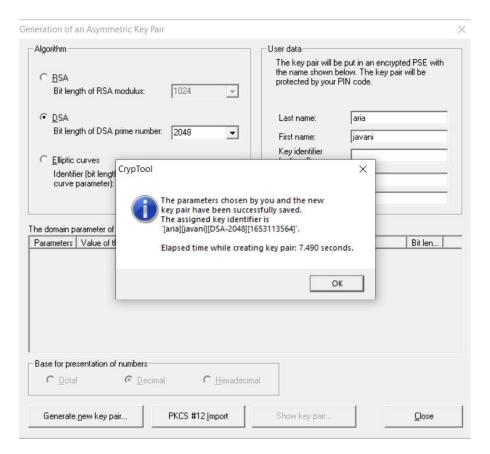


b)



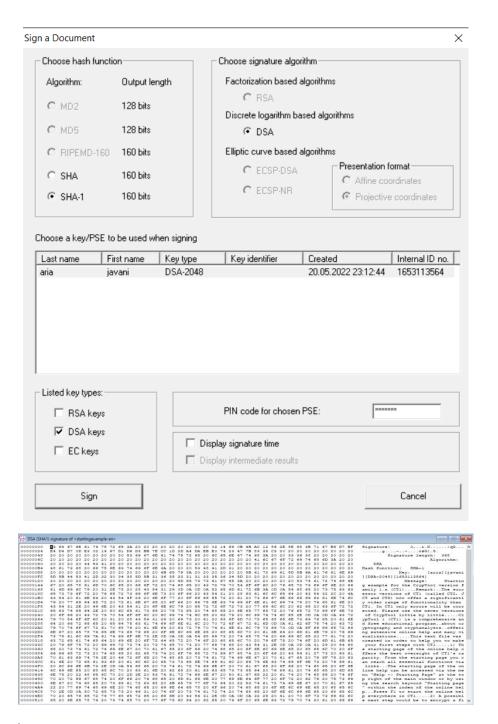
2.

a)

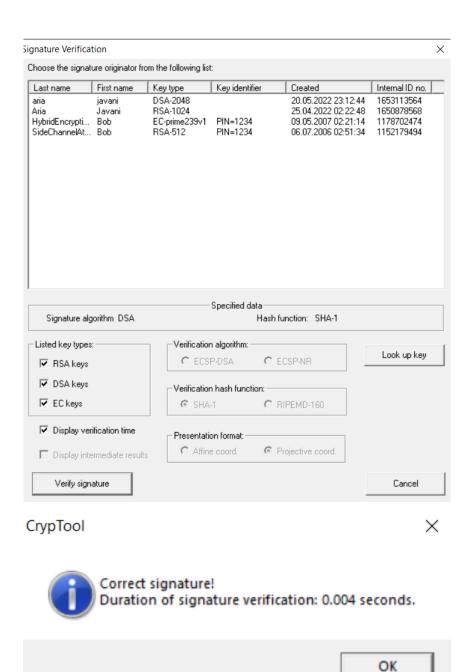


b)

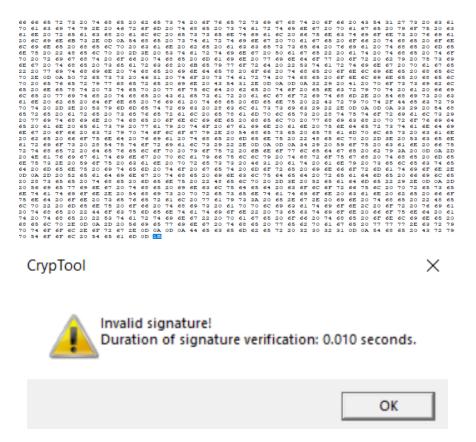
original text and signed text



c)



d)



since we changed the signature after verifying messages are not identical so this signature doesn't belong to this text.