

# Belief and Opinion Dynamics and Aggregation in Multi-Agent Systems

## Part 2: Opinion Dynamics and Aggregation

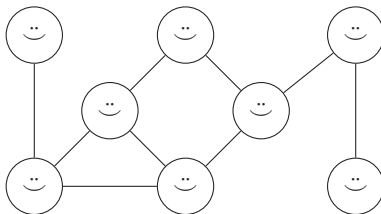
Umberto Grandi · Emiliano Lorini · Arianna Novaro



IJCAI-2020 Tutorial (website)

# Election day in Influenceville

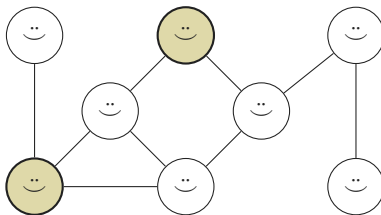
The inhabitants of *Influenceville* are notoriously open to new ideas.



Each day, if 50% (or more!) of an influenceviller's friends think  $a$ , the influenceviller will immediately change her opinion to  $a$  as well.

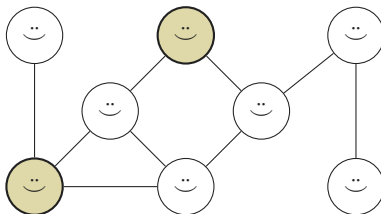
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One day, two politically active influencevillers decide they will vote for the **green candidate** in the upcoming election.



# Election day in Influenceville

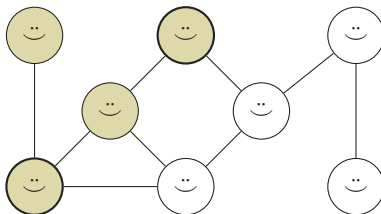
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*After how many days will **all** the influencevillers choose to vote for the green candidate (the one above is day zero)?*

# Election day in Influenceville

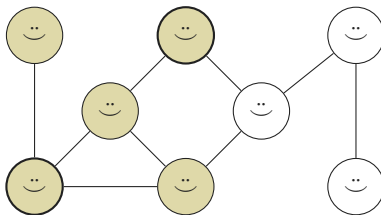
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Day 1

# Election day in Influenceville

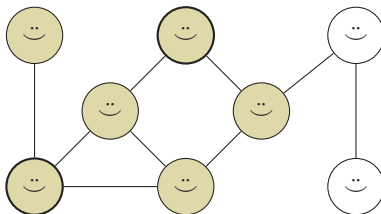
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Day 2

## Election day in Influenceville

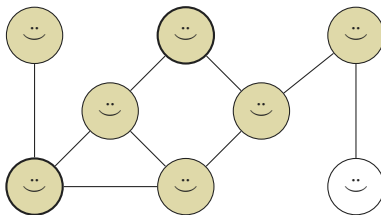
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Day 3

## Election day in Influenceville

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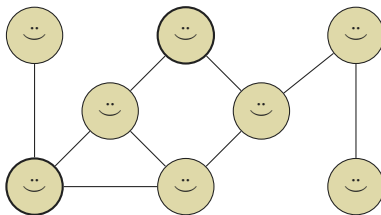


Day 4



## Election day in Influenceville

One day, two politically active influencevillers decide they will vote for the **green candidate** in the upcoming election.



Day 5

# Plan for this tutorial

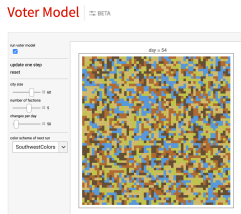
- ☐ Review of opinion diffusion and social influence models
- ☐ Judgment and goal aggregation
- ☐ Propositional opinion diffusion
- ☐ Diffusion of constrained opinions
- ☐ Strategic opinion diffusion

# **Review of opinion diffusion and social influence models**

# Voter model

Well-known models, widely studied and used in simulations:

- At any point in time **two voters meet** and one of the two, again at random, takes the **opinion of the other**



Play with a simple voter model on <https://demonstrations.wolfram.com/VoterModel/>!

- R. A. Holley, T. M. Liggett. Ergodic Theorems for Weakly Interacting Infinite Systems and the Voter Model. *The Annals of Probability*. 1975.
- P. Clifford, A. Sudbury. A model for spatial conflict. *Biometrika*. 1973.

# Weighted averages

Another well-known and studied model:

- ▶ Agents have **opinions** in  $[0, 1]$ , forming a vector  $o$
- ▶ A stochastic matrix  $G$  represents the **intensity** of **mutual influence** among agents
- ▶ Multiply  $G$  with  $o$  to obtain the influenced opinions, as the weighted sums of the opinions of one's influencers

The linearity of influence and its clean mathematical formulation allows for the use of **fixed-point theorems** from linear algebra.

- M. H. DeGroot. Reaching a consensus. *Journal of the American Statistical Association*. 1974.
- K. Lehrer, C. Wagner. Applications of the Consensus Model. Rational Consensus in Science and Society. *Philosophical Studies Series in Philosophy*. 1981.

# Independent cascades

A probabilistic model of influence **mixing weights and probability**:

- ▶ Agents are on a directed network and have **binary** 0/1 opinions
- ▶ Each edge  $e$  has an influence probability  $p_e$
- ▶ All agents with 0-opinion flip one coin for each edge that points to them from an agent with opinion 1, biased with the influence probability of the edge

The well-known paper below studies algorithms to **find a set of seeds** of fixed size which **maximise** the spread of the 1-opinion.

- D. Kempe, J. Kleinberg, E. Tardos. Maximizing the spread of influence through a social network. *ACM SIGKDD International Conference on Knowledge discovery and data mining*. 2003.

# Our approach

We started from the less well-studied **threshold models** (Granovetter, Schelling, 1978) and from work on **eliciting social influence from agents' choices** (Grabisch and Rusinowska), and embarked on the study of diffusion of **complex opinions**:

- ▶ opinions as binary choices over interconnected issues, or preferences over alternatives
- ▶ social influence modelled through **aggregation functions** for each agent (decentralised)
- ▶ include elements of strategic reasoning into models of diffusion (intentional-unintentional)

# Judgment and goal aggregation



# Judgment aggregation

An elegant framework to represent **binary collective decisions** over a set of interrelated **issues** having an underlying **integrity constraint**.



	valid	breached	liable
Judge 1	✓	✓	✓
Judge 2	×	✓	×
Judge 3	✓	×	×
Majority	✓	✓	×



*The doctrinal paradox: do you see why?*

- L.A. Kornhauser, L.G. Sager. The One and the Many: Adjudication in Collegial Courts. *Calif. Law Review*. 1993.
- C. List, P. Pettit. Aggregating Sets of Judgments: An Impossibility Result. *Economics and Philosophy*. 2002.
- U. Endriss. Judgment Aggregation. In *Handbook of Computational Social Choice*. 2016.

# Binary aggregation framework

- ▶ A set  $\mathcal{N}$  of  $n$  agents deciding over a set  $\mathcal{I}$  of  $m$  binary issues
- ▶ An integrity constraint  $\Gamma$  over the issues in  $\mathcal{I}$
- ▶ The individual ballot  $B_a \in \{0, 1\}^m$  of agent  $a \in \mathcal{N}$
- ▶ A profile  $\mathbf{B} = (B_1, \dots, B_n)$  with the ballots of the  $n$  agents
- ▶ An aggregation rule  $F$  associating ballot(s) to a profile

	$p_1$	$p_2$	$\dots$	$p_m$
$B_1$	1	1	$\dots$	0
$B_2$	0	1	$\dots$	0
$\dots$	$\dots$	$\dots$	$\dots$	$\dots$
$B_n$	1	0	$\dots$	0
$F(\mathbf{B})$	1	1	$\dots$	0

- U. Grandi, U. Endriss. Binary Aggregation with Integrity Constraints. IJCAI-2011.

# Binary aggregation rules

- **Majority**: accept an issue if ( $> / \geq$ ) 50% of agents accept it
- **Unanimity**: accept/reject an issue if all agents accept/reject it

	$p_1$	$p_2$	$p_3$
$B_1$	1	1	0
$B_2$	0	1	0
$B_3$	1	1	0
Majority( $\mathbf{B}$ )	1	1	0
Unanimity( $\mathbf{B}$ )	0	1	0

If the majority outcome is not consistent with the integrity constraint  $\Gamma$ , more sophisticated rules needed (e.g., *Kemeny* rule).

- J. Lang, M. Slavkovik. Judgment aggregation rules and voting rules. ADT-2013.

# Axiomatic analysis of rules

How can we analyse aggregation rules, to choose one to use?

⇒ By checking whether they satisfy some **desirable properties** (called *axioms* in the literature), expressed formally.

*How to capture formally our intuition for a “good” property?*

# Axiomatic analysis of rules

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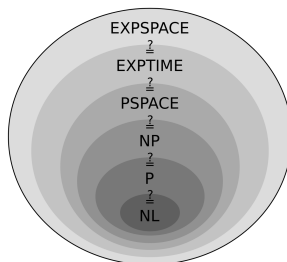
*How to capture formally our intuition for a “good” property?*

- ▶ **Anonymity**: it should not matter which agent submitted a certain ballot to compute the outcome of the rule
- ▶ **Unanimity**: if every agent agrees on an issue (or on a whole ballot), the outcome of the rule should agree with them
- ▶ **Monotonicity**: if an issue is currently accepted in the outcome, it should still be accepted if more support is given to it

# Computational analysis of rules

How can we analyse aggregation rules, to choose one to use?

⇒ By studying the **computational complexity** of calculating their outcome (i.e., the *winner determination problem*)



# Goal-based voting



*M: Movie night?*



*T: Take-away?*



*C: Playing cards?*

- ▶ Aggregation of **goals** instead of judgments ( $\Rightarrow$  collective **plan**)
- ▶ Balance **expressivity**, **compactness** and **complexity**:
  - Propositional formula  $\gamma_i$  instead of ballot  $B_i$  for  $i \in \mathcal{N}$
  - Restrictions on the language of goals to improve results
  - Need to generalize known voting rules to the new input

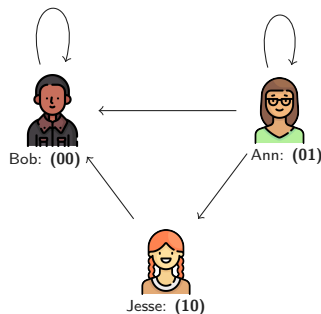
• A. Novaro, U. Grandi, D. Longin, E. Lorini. Goal-Based Collective Decisions: Axiom. & Complexity. IJCAI-2018.

# Propositional opinion diffusion



# Opinion diffusion on a network

Agents are connected to each other over a network and they can **influence** one another on their opinions over some binary issues.

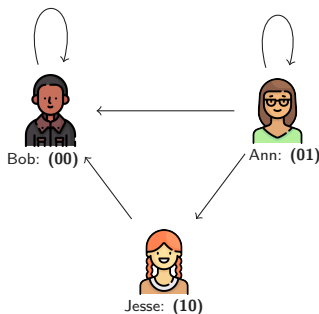


- U. Grandi, E. Lorini, L. Perrussel. Propositional Opinion Diffusion. AAMAS-2015.

# Propositional Opinion Diffusion (POD)

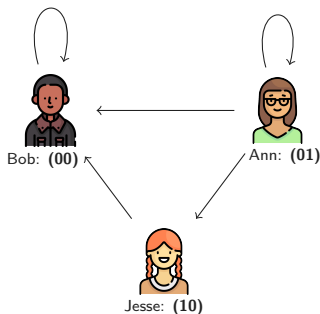
Extension of the binary aggregation framework with addition of:

- ▶ an **influence network**:  $E \subseteq \mathcal{N} \times \mathcal{N}$
- ▶ the **influencers** for each agent  $i$ :  $Inf(i) = \{j \mid (j, i) \in E\}$



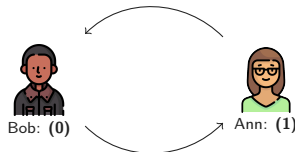
# Synchronous iterative diffusion

$$B_i^t = \begin{cases} B_i^{t-1} & \text{if } Inf(i) = \emptyset \\ F_i(B_{Inf(i)}^{t-1}) & \text{otherwise} \end{cases}$$



# POD termination

POD **terminates** on a class of graphs  $\mathcal{E}$  if there does not exist an **infinite** sequence of effective updates starting from any initial opinion profile on any graph  $G \in \mathcal{E}$ .



*Simplest non-terminating graph*

- If  $F_i$  satisfies ballot-**monotonicity** for all  $i \in \mathcal{N}$ , then POD terminates on the class of **directed acyclic graphs with loops** after at most  $\text{diam}(E) + 1$  number of steps.

# Sufficient and necessary conditions

Related work found **necessary and sufficient** conditions for POD-termination:

- ▶  $F_i$  are independent, monotonic, and responsive, the graph  $G$  is serial, **and**
- ▶ *[intricate conditions in terms of winning and losing coalitions of  $F_i$  interlocking on  $G$ ]*

- Z. Christoff, D. Grossi. Stability in Binary Opinion Diffusion. LORI-2017.

# Diffusion of constrained opinions

## Example

Four individuals need to decide on whether to build a **skyscraper** (S), a new road (R), or a **hospital** (H). Law says that if S and H are built then R should also be built.



(Hosp and SkyS) implies Road

Voter 1:  
Y N N

Voter 2:  
N N Y

Voter 3:  
Y Y Y

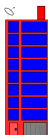


Voter 4:  
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*What happens if voter 4 is influenced on **all three issues**?*

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Voter 4:  
N N N

*What happens if voter 4 is influenced on **all three issues**?*  
*What happens if voter 4 is influenced on **one issue at the time**?*



# POD vs. proposition-wise updates

Individuals update using aggregation functions  $F$  on **all issues**:

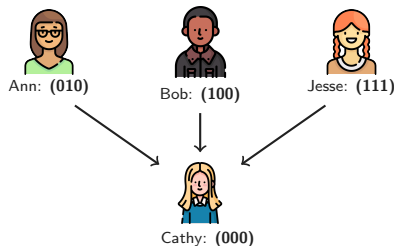
- ▶  $\text{POD}_F$  is a transformation function that updates the opinion of a subset of individuals **on all issues at the same time** (provided the update is consistent with the constraint  $\Gamma$ ).

Or on **subsets of issues** of limited size:

- ▶  $\text{PWOD}_F^k$  is a transformation function that updates the opinion of a subset of individuals **on a subset of issues of size at most  $k$**  (provided the update is consistent with  $\Gamma$ ).

## Example (continued)

An influence network between four agents, with  $\Gamma = (S \wedge H \rightarrow R)$ :

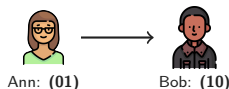


If  $F_{\text{Cathy}}$  is the **strict majority rule** and  $B$  the above profile, then:

- ▶  $\text{POD}_F(B) = \{B\}$ ,  $B$  is a **termination profile**
- ▶  $\text{PWOD}_F^1(B) = \{(010, 100, 111, \mathbf{010}), (010, 100, 111, \mathbf{100}), B\}$ .
- ▶  $\text{PWOD}_F^2(B) = \text{PWOD}_F^1(B)$

## Problematic example

Let there be two issues and  $\Gamma = p \text{ XOR } q = \{01, 10\}$



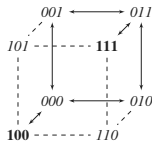
Whatever the **unanimous**  $F$ :

- ▶  $\text{POD}_F(\mathbf{B}) = \{\mathbf{B}, \mathbf{B}'\}$  where  $B'_1 = B'_2 = (0, 1)$
- ▶  $\text{PWOD}_F^1(\mathbf{B}) = \{\mathbf{B}\}$

**Question:** Can we characterise the set of integrity constraints on which  $\text{PWOD}_F^k$ -reachability corresponds to  $\text{POD}_F$ -reachability?

# Reachability result

An integrity constraint  $\Gamma$  is *k-geodetic* iff for all models  $B$  and  $B'$  of  $\Gamma$ , at least one of the shortest paths connecting them is composed of nodes that are all models of  $\Gamma$ . A non-example:



- Let  $\Gamma$  be an integrity constraint. Any profile  $B'$  that is  $\text{POD}_F$ -reachable from a  $\Gamma$ -consistent initial profile  $B$  is also  $\text{PWOD}_F^k$ -reachable from  $B$  if and only if  $\Gamma$  is *k-geodetic*.

# Examples of 1-geodetic constraints

**Preferences.** Let  $a > b$  be a set of binary questions for candidates  $a, b, c, \dots$ . The constraints are that of **transitivity**, **completeness** and **anti-symmetry**.

- Set of constraints is 1-geodetic: two distinct linear orders always differ on at least one adjacent pair.

- Brill, Elkind, Endriss, Grandi. Pairwise Diffusion of Preference Rankings in Social Networks. IJCAI-2016.

**Budget constraints.** Enumerate all combinations of items that exceed a given budget. They are *negative formulas*, i.e., one DNF representation only has negative literals: a sufficient condition for 1-geodeticity.

- O. Ekin, P. L. Hammer, A. Kogan. On Connected Boolean Functions. *Discrete Mathematics*, 1999.

# Cost of constraints and termination

- **Cost of constraints:** Can we quantify the gain in terms of influence that is given by allowing updates on  $k$  issues?

**Answer:** the **influence gap** is the sum of the distances between each individual's opinion and the aggregated one of its influencers. We show that this figure for  $\text{POD}_F$  is larger than for  $\text{PWOD}_F^k$  and give precise bounds.

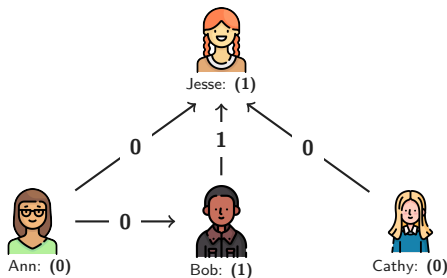
- **Termination:** Can we find conditions on the graph and aggregator that guarantee termination?

**Answer:** preliminary findings on complete graphs and DAGs. For arbitrary graphs we have to assume consistent aggregation of influencers' opinions. Ongoing: general termination results.

# Strategic opinion diffusion

# Exerting influence over opinions

- Extension of POD where agents have **individual goals** and they choose whether to **exert their influence** on the issues.

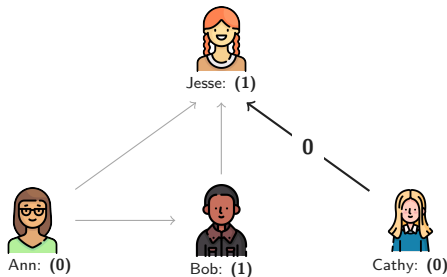


- U. Grandi, E. Lorini, A. Novaro, L. Perrussel. Games of Influence. *Journal of Logic and Computation*. 2021.



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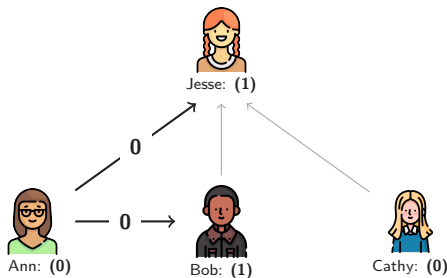
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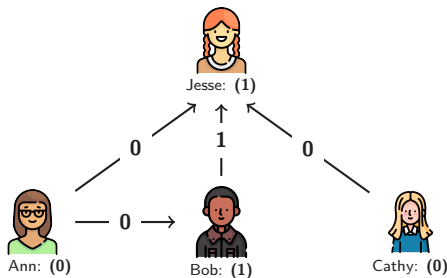
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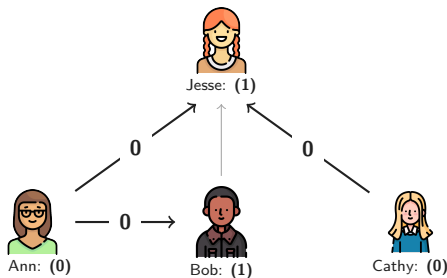
# Unanimous aggregation of opinions

- Agents need to **aggregate** the opinions of their **influencers**
- ⇒ **Unanimity** of those influencers who exert their opinion



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# Influence games

Agents first decide whether to **exert their influence** on some issues, then they **update their opinions** as per those of their influencers.

- ▶ An **influence game** is a tuple  $(\mathcal{N}, \mathcal{I}, E, S_0, F_i, \gamma_i)$
- ▶ Each  $\gamma_i$  is a formula in a language of Linear Temporal Logic
  - *Example of issue-wise consensus goal:*

$$\Diamond \Box \bigwedge_{p \in J} ((\bigwedge_{i \in C} \text{op}(i, \{p\})) \vee (\bigwedge_{i \in C} \neg \text{op}(i, \{p\})))$$

for  $C \subseteq N$  and  $J \subseteq \mathcal{I}$ .

# Game-theoretic and computational complexity results for influence games

- ▶ If  $E$  is a graph such that  $\text{Inf}(i) \neq \emptyset$  for all  $i \in \mathcal{N}$  and  $\gamma_i$  is the **issue-wise consesus** goal for all  $i \in \mathcal{N}$  and  $J \subseteq \mathcal{I}$ , then there is a **Nash equilibrium** for any initial state  $S_0$ .
- ▶ Checking if a profile is a Nash equilibrium (**Nash-membership**) for influence games with unanimous aggregation rule is **in PSPACE** for memory-less strategies.

# Conclusions

# Summary

- ▶ Election day in Influenceville: *hard-threshold cascade model*
- ▶ A partial review of opinion diffusion and social influence models
- ▶ Judgment aggregation: *doctrinal paradox*, *binary aggregation with constraints*, *aggregation rules* (majority, unanimity), *axioms* (anonymity, unanimity, monotonicity); Goal-based voting
- ▶ Propositional opinion diffusion (POD): *iterative diffusion process*, *convergence results*
- ▶ Diffusion of constrained opinions: *reachability result* and *k-geodetic constraints*
- ▶ Influence games: *exerting one's influence*, *unanimous aggregation*, *game-theory and complexity results*



# Future work

- ▶ Investigate the effects of opinion diffusion on voting:
  - Preliminary results in simulation show that the frequency of Condorcet winners increases after preference diffusion is run
  - Recent work explored control by an external player of the result of an election via opinion diffusion
- ▶ Go further in the analysis of constrained opinion diffusion obtaining general termination results
- ▶ Develop a full-fledged model of strategic opinion diffusion, in which agents can lie about their opinions towards the satisfaction of their goals

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