Lecture 3: Process I

- Process is a program in execution.
 - It contains <u>every accounting information</u> of that running program, e.g.,
 - Current program counter
 - Accumulated running time
 - The list of files that are currently opened by that program
 - The page table

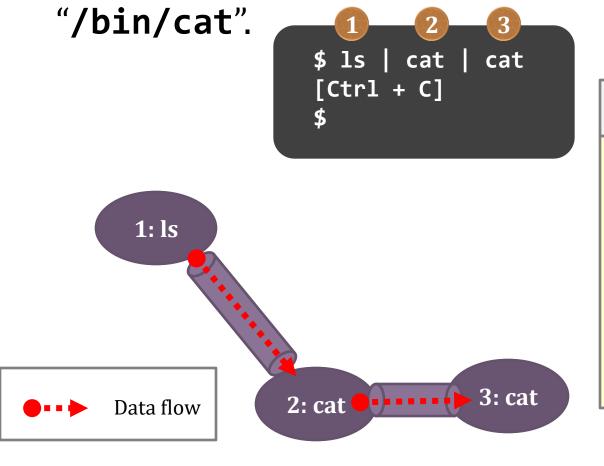
• ...

https://en.wikipedia.org/wiki/Process_control_block

```
$ 1s | cat | cat
[Ctrl + C]
$
```

- The command involves three processes.
- It will stop early if I send a signal to interrupt it.
- Its progress is determined by the scheduler.
- The three processes cooperate to give useful output.

- What are those two "cats"?
 - 2 different processes using the same code



```
If you don't know what a cat is.

#include <stdio.h>

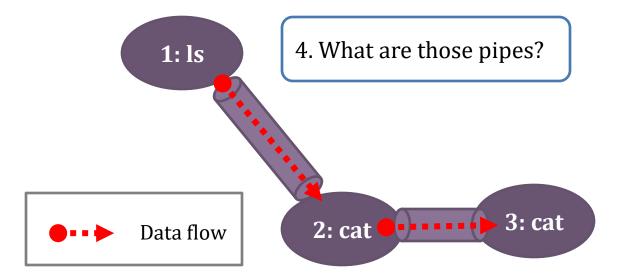
int main(void) {
   int c;
   while ( 1 ) {
      c = getchar();
      if( c == EOF )
           break;
      putchar(c);
   }
}
```

Our Roadmap

1. How to distinguish the two cats?

2. Who (and how to) create the processes?

3. Which should run first?



5. What if "**1s**" is feeding data too fast? Will the "**cat**" feels *full and dies*?!

Process identification

- How can we identify processes from one to another?
 - Each process is given an unique ID number, and is called the process ID, or the PID.
 - The system call, getpid(), prints the PID of the calling process.

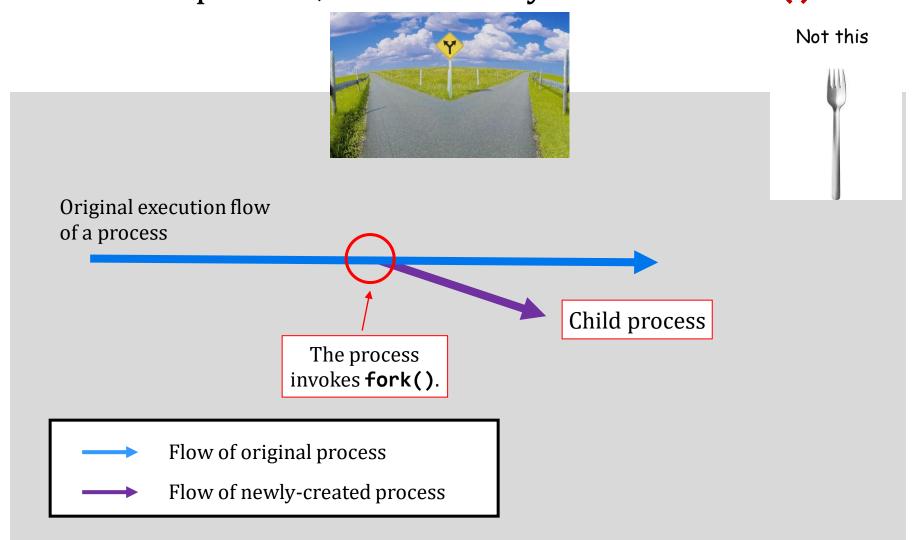
```
#include <stdio.h> // printf()
#include <unistd.h> // getpid()

int main(void) {
    printf("My PID is %d\n", getpid());
}
```

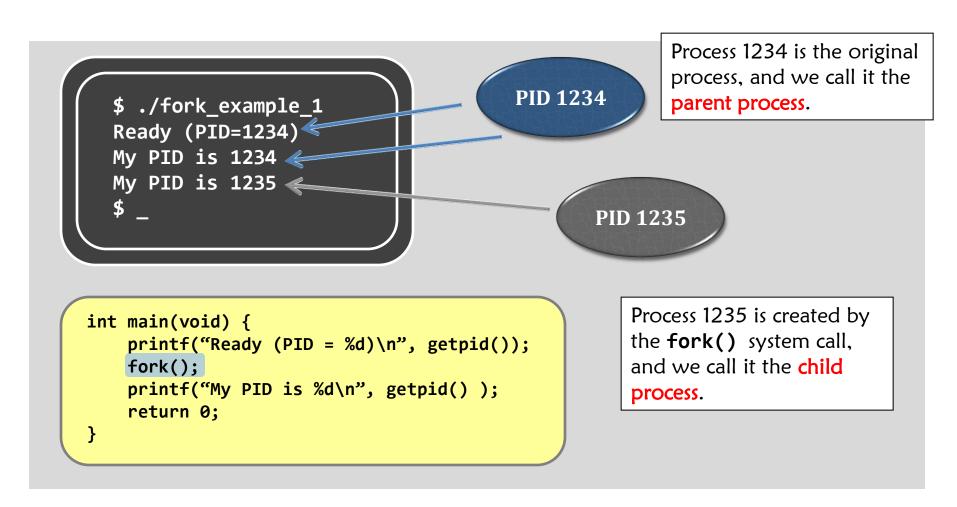
```
$ ./getpid
My PID is 1234
$ ./getpid
My PID is 1235
$ ./getpid
My PID is 1237
```

Process creation

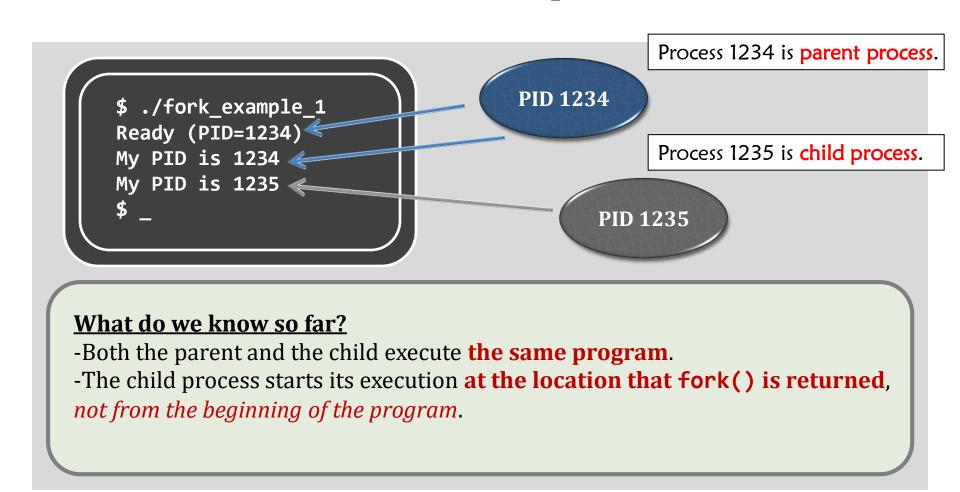
To create a process, we use the system call fork().



So, how do fork() and the processes behave?



So, how do fork() and the processes behave?



```
int main(void) {
      int result;
      printf("before fork ...\n");
      result = fork();
      printf("result = %d.\n", result);
      if(result == 0) {
8
        printf("I'm the child.\n");
9
        printf("My PID is %d\n", getpid());
10
11
      else {
        printf("I'm the parent.\n");
12
        printf("My PID is %d\n", getpid());
13
14
15
16
      printf("program terminated.\n");
17
```

```
$ ./fork_example_2
before fork ...
```

PID 1234

```
int main(void) {
      int result;
      printf("before fork ...\n");
      result = fork();
      printf("result = %d.\n", result);
      if(result == 0) {
        printf("I'm the child.\n");
9
        printf("My PID is %d\n", getpid());
10
11
      else {
        printf("I'm the parent.\n");
12
        printf("My PID is %d\n", getpid());
13
14
15
16
      printf("program terminated.\n");
17
```

```
./fork_example_2
before fork ...
```

PID 1234 fork() PID 1235

Let there be only **ONE CPU**. Then...

- Only one process is allowed to be executed at one time.
- However, we can't predict which process will be chosen by the OS.
- That is controlled by the OS's **scheduler**.

NOTE THIS

In this example, we assume that the parent, PID 1234, runs first, after the **fork()** call.

```
int main(void) {
      int result;
      printf("before fork ...\n");
      result = fork();
      printf("result = %d.\n", result);
      if(result == 0) {
        printf("I'm the child.\n");
        printf("My PID is %d\n", getpid());
10
11
      else {
12
        printf("I'm the parent.\n");
13
        printf("My PID is %d\n", getpid());
14
15
16
      printf("program terminated.\n");
17
```

```
$ ./fork_example_2
before fork ...
result = 1235
```

Important

For parent, the return value of **fork()** is the PID of the created child.

PID 1234 (running)

PID 1235 (waiting)

```
int main(void) {
      int result;
      printf("before fork ...\n");
      result = fork();
      printf("result = %d.\n", result);
      if(result == 0) {
        printf("I'm the child.\n");
        printf("My PID is %d\n", getpid());
10
11
      else {
12
        printf("I'm the parent.\n");
        printf("My PID is %d\n", getpid());
13
14
15
      printf("program terminated.\n");
17
```

```
$ ./fork_example_2
before fork ...
result = 1235
I'm the parent.
My PID is 1234
program terminated.
```

PID 1234 (dead)

PID 1235 (waiting)

```
int main(void) {
      int result;
      printf("before fork ...\n");
      result = fork();
      printf("result = %d.\n", result);
      if(result == 0) {
        printf("I'm the child.\n");
        printf("My PID is %d\n", getpid());
10
11
      else {
12
        printf("I'm the parent.\n");
13
        printf("My PID is %d\n", getpid());
14
15
      printf("program terminated.\n");
17
```

```
$ ./fork_example_2
before fork ...
result = 1235
I'm the parent.
My PID is 1234
program terminated.
result = 0
Important
```

For child, the return

value of **fork()** is **0**.

PID 1234 (dead)

PID 1235 (running)

```
int main(void) {
      int result;
      printf("before fork ...\n");
      result = fork();
      printf("result = %d.\n", result);
      if(result == 0) {
        printf("I'm the child.\n");
        printf("My PID is %d\n", getpid());
10
11
      else {
12
        printf("I'm the parent.\n");
13
        printf("My PID is %d\n", getpid());
14
15
      printf("program terminated.\n");
17
```

```
$ ./fork_example_2
before fork ...
result = 1235
I'm the parent.
My PID is 1234
program terminated.
result = 0
I'm the child.
My PID is 1235
program terminated.
$ _
```

PID 1234 (dead)

PID 1235 (dead)

- fork() behaves like "cell division".
 - It creates the child process by cloning from the parent process, including all user-space data, e.g.,

| Cloned items | Descriptions |
|-------------------------------------|---|
| Program counter [CPU register] | That's why they both execute from the same line of code after fork() returns. |
| Program code [File & Memory] | They are sharing the same piece of code. |
| Memory | Including local variables, global variables, and dynamically allocated memory. |
| Opened files [Kernel's internal] | If the parent has opened a file "A", then the child will also have file "A" opened automatically. |

- However...
 - fork() does not clone the following...
 - Note: PCB is in the kernel space.

| Distinct items | Parent | Child |
|------------------------|---------------------------|--|
| Return value of fork() | PID of the child process. | 0 |
| PID | Unchanged. | Different, not necessarily be "Parent PID + 1" |
| Parent process | Unchanged. | Parent. |
| Running time | Cumulated. | Just created, so should be 0. |
| [Advanced] File locks | Unchanged. | None. |

- process creation.
- program execution.



fork() can only duplicate...

- If a process can only <u>duplicate itself</u> and <u>always</u> runs the same program, it's not quite meaningful
 - how can we execute other programs?

- We want CHANGE!
 - Meet the exec*() system call family.

execl() – a member of the exec system call family (and the family has 6 members).

```
int main(void) {
                                                            $ ./exec_example
                                                            before execl ...
  printf("before execl ...\n");
  execl("/bin/ls", "/bin/ls", NULL);
  printf("after execl ...\n");
                                       Arguments of the exec1() call
  return 0;
                                       1<sup>st</sup> argument: the program name, "/bin/ls" in the
                                       example.
                                       2<sup>nd</sup> argument: argument[0] to the program.
                                       3<sup>rd</sup> argument: argument[1] to the program.
```

execl() – a member of the exec system call family (and the family has 6 members).

```
int main(void) {
   printf("before execl ...\n");
   execl("/bin/ls", "/bin/ls", NULL);
   printf("after execl ...\n");
   return 0;
}
```

```
$ ./exec_example
before execl ...
exec_example
exec_example.c
```

What is the output?

The same as **the output of running "1s"** in the shell.

Example #1: run the command "/bin/ls"

execl("/bin/ls", "/bin/ls", NULL);

| Argument Order | Value in above example | Description |
|-------------------|------------------------|---|
| 1 | "/bin/ls" | The file that the programmer wants to execute. |
| 2 | "/bin/ls" | When the process switches to "/bin/ls", this string is the program argument[0]. |
| 3 | NULL | This states the end of the program argument list. |

Example #2: run the command "/bin/ls -1"

```
execl("/bin/ls", "/bin/ls", "-1", NULL);
```

| Argument Order | Value in above example | Description |
|-------------------|------------------------|---|
| 1 | "/bin/ls" | The file that the programmer wants to execute. |
| 2 | "/bin/ls" | When the process switches to "/bin/ls", this string is the program argument[0]. |
| 3 | "-1" | When the process switches to "/bin/ls", this string is the program argument[1]. |
| 4 | NULL | This states the end of the program argument list. |

execl() – a member of the exec system call family (and the family has 6 members).

```
int main(void) {
  printf("before execl ...\n");
  execl("/bin/ls", "/bin/ls", NULL);
  printf("after execl ...\n");
  return 0;
}
```

WHAT?!
The shell prompt appears!

```
$ ./exec_example
before execl ...
exec_example
exec_example.c
$ _
```

The output says:

- (1) The gray code block **is not** reached!
- (2) The process is **terminated**!

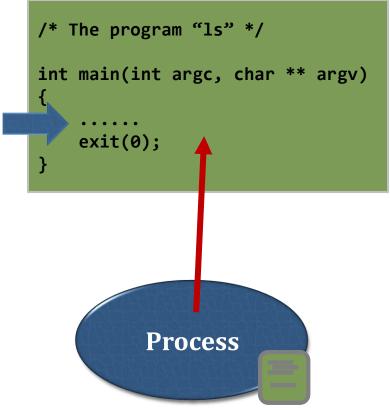
WHY IS THAT?!

The exec system call family is not simply a function that "invokes" a command.

```
int main(void) {
   printf("before execl ...\n");
   execl("/bin/ls", "/bin/ls", NULL);
   printf("after execl ...\n");
   return 0;
                                                        Process
Originally, the process is executing the
program "exec_example".
```

The exec system call family is not simply a function that "invokes" a command.

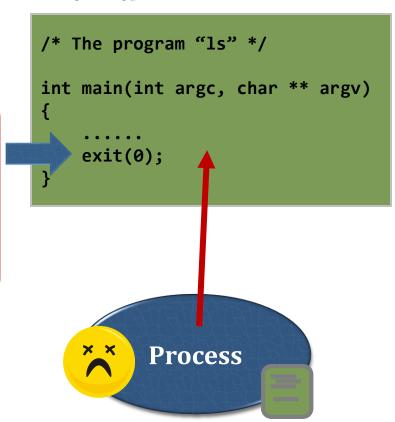
```
int main(void) {
   execl("/bin/ls", "/bin/ls", NULL);
   printf("after execl ...\n");
   return 0;
execl() changes the execution from
"exec_example" to "/bin/ls"
```



The exec system call family is not simply a function that "invokes" a command.

The "return" or the "exit()" statement in "/bin/ls" will terminate the process...

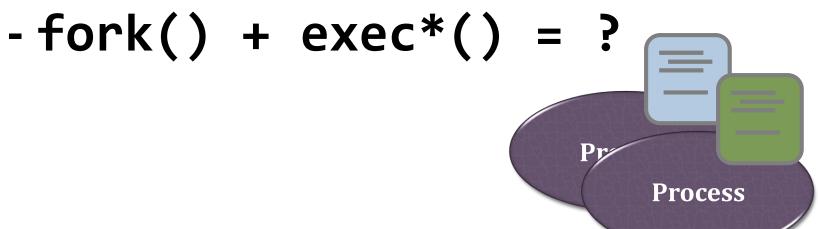
Therefore, it is certain that the process cannot go back to the old program!



- The process is changing the code that is executing and <u>never</u> returns to the original code.
 - The last two lines of codes are therefore not executed.
- The process that calls an exec* system call will replace userspace info, e.g.,
 - Program Code
 - Memory: local variables, global variables, and dynamically allocated memory;
 - Register value: e.g., the program counter;
- But, the kernel-space info of that process is preserved, including:
 - PID;
 - Process relationship;
 - etc.

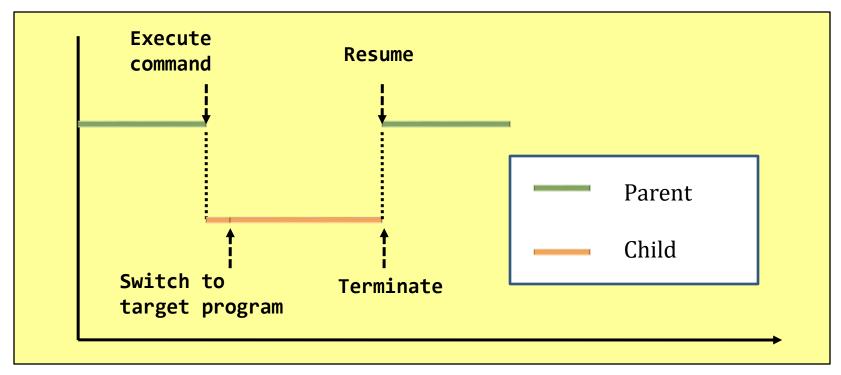
~ reverse takeover in stock market

- process creation.
- program execution.



When fork() meets exec*()...

- To implement the core part of a shell,
- To implement the C library call system()
- **•••**



fork() + exec*() = system()?

```
int system ver CS302(const char *cmd str) {
        if(cmd str == -1)
 2
 3
            return -1;
 4
        if(fork() == 0) {
 5
            execl(cmd_str, cmd_str, NULL);
            fprintf(stderr,
                "%s: command not found\n", cmd_str);
            exit(-1);
 8
 9
        return 0;
10
    }
11
                                                         $ ./system_implement_1
12
    int main(void) {
                                                         before...
13
        printf("before...\n\n");
                                                         Makefile
14
        system_ver_CS302("/bin/ls");
                                                         system implement 1
15
        printf("\nafter...\n");
                                                         system implement 1.c
16
        return 0;
17 }
```

fork() + exec*() = system()?!

```
int system ver CS302(const char *cmd str) {
        if(cmd str == -1)
 2
 3
            return -1;
        if(fork() == 0) {
 4
 5
            execl(cmd_str, cmd_str, NULL);
            fprintf(stderr,
               "%s: command not found\n", cmd_str);
            exit(-1);
8
9
        return 0;
10
   }
11
12
    int main(void) {
13
        printf("before...\n\n");
14
        system_ver_CS302("/bin/ls");
15
        printf("\nafter...\n");
16
        return 0;
17 }
```

Some strange cases may happen some times

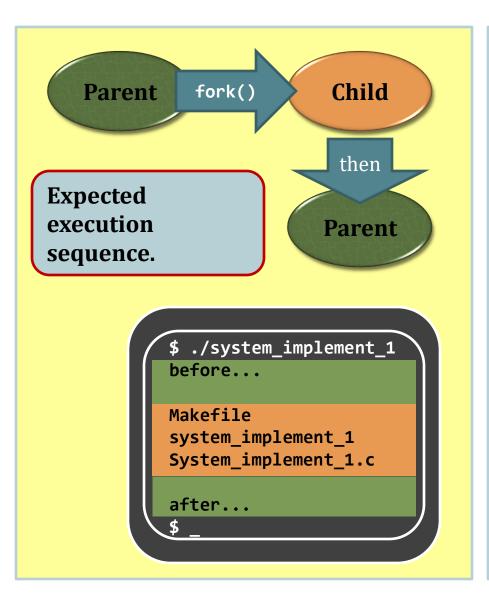
```
$ ./system_implement_1
before...

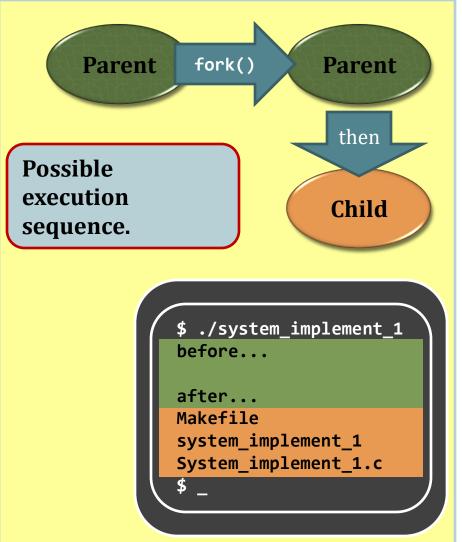
after...
Makefile
system_implement_1
system_implement_1.c
$ _
```

fork() + exec*() = system()...

```
int system ver CS302(const char *cmd str) {
        if(cmd str == -1)
                                                            Let's re-color the program!
             return -1:
 4
        if(fork() == 0) {
                                                                        Parent process
 5
             execl(cmd_str, cmd_str, NULL);
             fprintf(stderr,
                                                                        Child process
                "%s: command not found\n", cmd_str);
                                                                        Both processes
             exit(-1);
 8
 9
        return 0;
10
11
                                                            $ ./system_implement_1
12
    int main(void) {
                                                           before...
13
        printf("before...\n\n");
14
        system_ver_CS302("/bin/ls");
                                                           after...
15
        printf("\nafter...\n");
                                                           Makefile
16
        return 0;
                                                            system_implement_1
17 }
                                                            system implement 1.c
```

fork() + exec*() = system()...





 It is very weird to allow different execution orders.

- How to let the child to execute first?
 - But...we can't control the OS scheduler

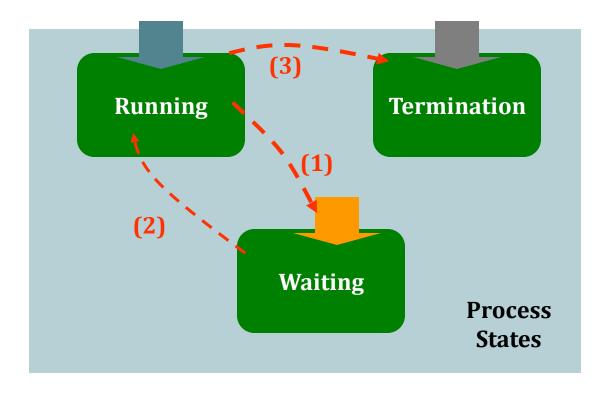
- Then, our problem becomes...
 - When to suspend the execution of the parent process?
 - Now to wake the parent up after the child is terminated?

fork()+ exec*() + wait() = system()

```
int system ver CS302(const char *cmd str) {
        if(cmd str == -1)
 2
 3
            return -1;
 4
        if(fork() == 0) {
            execl("/bin/sh", "/bin/sh",
                   "-c", cmd str, NULL);
 6
            fprintf(stderr,
                "%s: command not found\n", cmd_str);
            exit(-1);
 8
        wait(NULL);
9
10
        return 0;
                                                          $ ./system_implement_2
11
    }
                                                          before...
12
13
    int main(void) {
                                                         Makefile
14
        printf("before...\n\n");
                                                          system_implement_2
15
        system_ver_CS302("/bin/ls");
                                                          System_implement_2.c
16
        printf("\nafter...\n");
                                                          after...
17
        return 0;
18 }
```

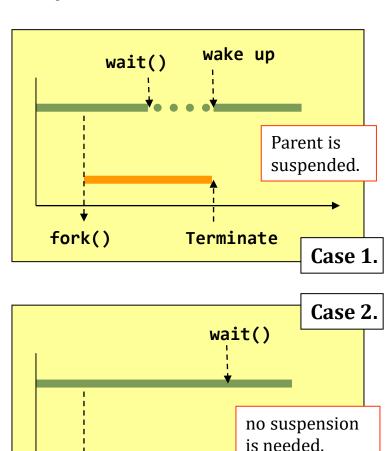
Process Life Cycle (user-space)

```
int main(void) {
   int x = 1;
   getchar();
   return x;
}
```



wait() - user-space

- wait() system call
 - suspend the calling process to waiting state and return (wakes up) when
 - one of its child processes changes from
 - **running to terminated.**
 - Or a <u>signal</u> is received (will cover)
 - return immediately (i.e., does nothing) if
 - It has no children
 - Or a child terminates before the parent calls wait for



Terminate

fork()

wait() VS waitpid()

| wait() | | waitpid() |
|---------------------------------------|----|--|
| Wait for any one of the children. | VS | Depending on the parameters, waitpid() will wait for a particular child only. |
| Detect child termination only. | | Depending on the parameters, waitpid() can detect multiple child's status change |

wait() also has a very
important hidden task
 (will cover next)

Summary

- A new process is created by fork()
 - Who is the first process?
- A process is a program being brought by exec to the memory
 - has state (initial state= ready)
 - waiting for the OS to schedule the CPU to run it
- Can a process execute more than one program?
 - Yes, keeps on calling the exec system call family
- You now know how system() C <u>library call</u> is implemented by <u>syscalls</u> fork(), exec(), and wait()

Thank You!

exec*() – arguments explained

- Environment variables
 - A set of strings maintained by the shell.

```
int main(int argc, char **argv, char **envp) {
   int i;
   for(i = 0; envp[i]; i++)
       printf("%s\n", envp[i]);
   return 0;
}
```

The "**envp" variable is an array of string A string is an array of characters

```
$ ./envp
SHELL=/bin/bash
PATH=.....

* _
```

exec*() – arguments explained

- Environment variables
 - A set of strings maintained by the shell.
 - Quite a number of programs will read and make use of the environment variable.

| Variable name | Description |
|---------------|--|
| SHELL | The path to the shell that you're using. |
| PWD | The full path to the directory that you're currently on. |
| HOME | The full path to your home directory. |
| USER | Your login name. |
| EDITOR | Your default text editor. |
| PRINTER | Your default printer. |