CS205 C/ C++ Programming - Project2

樊顺

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1. Part 1 Requirements

- redirect the output to the file.
- Calculate the matrix multiple as double and float separately.
- Improving the speed of calculating.

2. Part 2 The Implement of the Requirements

2.1. arguments process

For the input data, since I don't know the matrix size of the input file, I used the function: int getNumber(string name); to get the number in the input parameter to determine the size of the matrix in the input file. The code of this function show as follow:

```
int getNumber(string name)
{
    int len_s = name.size();
    int i=0, j=0;
    while (i < len_s)
    {
        if (name[i] >= '0'&& name[i] <= '9')
        {
            j = i;
            int len = 0;
            while (name[i] >= '0'&& name[i] <= '9')
            {
                i++;
                len++;
            }
            string s0 = name.substr(j, len);//获取子串
            int num=0;//数字字符串转换为整型数字
            stringstream s1(s0);
            s1 >> num;
```

```
return num;
}
else
{
   i++;
}
return 0;
}
```

2.2. redirect the output to the file

In this part, I use the <fstream> to implement the redirect the output to the file. And when finish the output, close the stream. The code of output show as follow:

2.3. Calculate the matrix multiple as double and float separately

In this task, I use the double type to store all input numbers in two vectors. After the calculate, the result also store as double in vector.

In the basic implement, I use the Two-dimensional array vector<vector<double>>> to store the input matrix.

Due to storage structure, the calculate was always using the double to calculate. And when the result output to the file, I make sure that the precision of all output data is 5 decimal places.

2.4. Improving the speed of calculating

2.4.1. Brute force solution

In this implement, I use three for loop to calculate two matrix multiple. The code is:

```
{
    vector<vector<double>> result;
    for (int i = 0; i < a.size(); i++)
    {
        vector<double> temp;
        for (int j = 0; j < b.size(); j++)
        {
            double res = 0;
            for (int k = 0; k < a.size(); k++)
            {
                res += a[i][k] * b[k][j];
            }
            temp.push_back(res);
        }
        result.push_back(temp);
}
return result;
}</pre>
```

• This method is the basic and slowest implementation. The time cost in calculate two matrix(the size both are 2048 * 2048) is 147s.

```
./matmul mat-A-2048.txt mat-B-2048.txt out2048.txt
Read in First file used time :0.709465s
Read in second file used time :0.710928s
Calculate two matrix multiple used time :147.305s
Output the result to file used time :2.674s
```

2.4.2. Reverse cycle order

• In this Implementation, I change the cycle order, make the first loop the k, the loop j (which means this method value taken from the a array is first calculated with the data in the column of the b array.).

The code is:

```
{
    res += s * b[k][j];
}
    temp.push_back(res);
}
result.push_back(temp);
}
return result;
}
```

• For this method, the performance of the program has been significantly improved, greatly reducing the time consumed by the calculation.

```
./matmul mat-A-2048.txt mat-B-2048.txt out2048.txt

Read in First file used time :0.704428s

Read in second file used time :0.699405s

Calculate two matrix multiple used time :142.224s

Calculate two matrix multiple with Reverse cycle order used time :32.7107s

Output the result to file used time :2.664s

Output the result to out2 file used time :2.586s
```

Analysis

In the analysis, it was found that the reasons for the inefficiency of matrix multiplication were on the one hand due to the complexity of the $O(n^3)$ algorithm, and on the other hand due to the discontinuity of memory access, which led to the low hit rate of the cache in the storage space.

Assume the definition of a jump uncle to evaluate the discontinuity of memory access.

Assuming that the matrix is stored in rows.

For a loop like ijk, after each number in result matrix is counted, it jumps once, and jumps back to the beginning after the end of a line. For a matrix of size n, the number of hops is $n^3 + n^2 + n$ times.

For a loop like ikj, it calculates the matrix row by row, so the result matrix will only jump once after each row is calculated. During this period, the b matrix jumped n times in total, that is, during the calculation of one row, it jumped n+1 times. The result matrix has a total of n rows, so the total number of jumps is $n^2 + n$, which greatly reduces the number of jumps and improves the operating efficiency.