

# Numerical Solution of PDEs using the Finite Element Method

A Poisson solver - deal.II step-3

**Luca Heltai** <[luca.heltai@sissa.it](mailto:luca.heltai@sissa.it)>

International School for Advanced Studies ([www.sissa.it](http://www.sissa.it))

Mathematical Analysis, Modeling, and Applications ([math.sissa.it](http://math.sissa.it))

Master in High Performance Computing ([www.mhpc.it](http://www.mhpc.it))

SISSA mathLab ([mathlab.sissa.it](http://mathlab.sissa.it))



# Aims for this Lecture

- First introduction into assembly of sparse linear systems
- Translation of weak form to assembly loops
- Applying boundary conditions
- Using linear solvers
- Post-processing and visualisation



# Reference material

- Tutorials
  - Step-3  
[https://dealii.org/current/doxygen/deal.II/step\\_3.html](https://dealii.org/current/doxygen/deal.II/step_3.html)
- Documentation
  - [https://www.dealii.org/current/doxygen/deal.II/group\\_FE\\_vs\\_Mapping\\_vs\\_FEValues.html](https://www.dealii.org/current/doxygen/deal.II/group_FE_vs_Mapping_vs_FEValues.html)
  - [https://www.dealii.org/current/doxygen/deal.II/group\\_UpdateFlags.html](https://www.dealii.org/current/doxygen/deal.II/group_UpdateFlags.html)





# Recap of Poisson Problem

Variational, continuous problem, infinite dimensional space:

$$\int_{\Omega} \nabla u \cdot \nabla v = \int_{\Omega} f v \quad \forall v \in H_0^1(\Omega)$$

Variational, discrete problem, finite dimensional space:

$$\int_{\Omega} \nabla u_h \cdot \nabla v_h = \int_{\Omega} f v_h \quad \forall v_h \in V_h \subset H_0^1(\Omega)$$



# Recap of Poisson Problem

$$\int_{\Omega} \nabla u_h \cdot \nabla v_h = \int_{\Omega} f v_h \quad \forall v_h \in V_h = \text{span}\{v_i\}_{i=1}^N$$



$$A_{ij} u^j = F_i \quad u_h := u^i v_i$$

$$A_{ij} := \int_{\Omega} \nabla v_j \nabla v_i \quad F_i := \int_{\Omega} f v_i$$



# Split Assembly on cells

$$A_{ij} := \int_{\Omega} \nabla v_j \cdot \nabla v_i d\Omega \quad F_i := \int_{\Omega} f v_i d\Omega$$

$$A_{ij} = \sum_m \int_{T_m} \nabla v_j \cdot \nabla v_i dT_m$$

To make this efficiently, we need a smart way to map  
local dofs to global dofs



# Split Assembly on cells

$$A_{ij} = \sum_m \int_{T_m} \nabla v_j \cdot \nabla v_i dT_m = \sum_m \int_{\hat{T}} [(\nabla v_j) \circ F_m] \cdot [(\nabla v_i) \circ F_m] J_m d\hat{T}$$

$$v_i \circ F_m|_{T_m} = \sum_{\alpha} P_{mi\alpha} \hat{v}_{\alpha}$$

$$P_{mi\alpha} = \begin{cases} 1 & \text{if local dof } \alpha \text{ on element } T_m \text{ maps to global dof } i \\ 0 & \text{otherwise} \end{cases}$$





# Split Assembly on cells

$$A_{ij} = \sum_m \int_{T_m} \nabla v_j \cdot \nabla v_i dT_m = \sum_m \int_{\hat{T}} [(\nabla v_j) \circ F_m] \cdot [(\nabla v_i) \circ F_m] J_m d\hat{T}$$

$$A_{ij} = \sum_m \sum_{\alpha} \sum_{\beta} P_{mi\alpha} \int_{\hat{T}} [(DF_m^{-T} \hat{\nabla} \hat{v}_{\alpha})] \cdot [DF_m^{-T}(\hat{\nabla} \hat{v}_{\beta})] J_m d\hat{T} P_{mj\beta}$$

$$A_{ij} = \sum_m \sum_{\alpha} \sum_{\beta} \sum_q P_{mi\alpha} [(DF_m^{-T} \hat{\nabla} \hat{v}_{\alpha})](\hat{x}_q) \cdot [DF_m^{-T}(\hat{\nabla} \hat{v}_{\beta})](\hat{x}_q) J_m(\hat{x}_q) w_q P_{mj\beta}$$

$$m \in [0, N_{\text{cell}}) \quad \alpha, \beta \in [0, N_{\text{localdofs}}) \quad i, j \in [0, N_{\text{dofs}}) \quad q \in [0, N_{\text{qpoints}})$$





# Local VS global matrix

$$A_{ij} = \sum_m \int_{T_m} \nabla v_j \cdot \nabla v_i dT_m = \sum_m \int_{\hat{T}} [(\nabla v_j) \circ F_m] \cdot [(\nabla v_i) \circ F_m] J_m d\hat{T}$$

$$A_{ij} = \sum_m \sum_{\alpha} \sum_{\beta} \sum_q P_{mi\alpha} [(DF_m^{-T} \hat{\nabla} \hat{v}_{\alpha})](\hat{x}_q) \cdot [DF_m^{-T}(\hat{\nabla} \hat{v}_{\beta})](\hat{x}_q) J_m(\hat{x}_q) w_q P_{mj\beta}$$

$$a_{m \alpha\beta} := \sum_q [(DF_m^{-T} \hat{\nabla} \hat{v}_{\alpha})](\hat{x}_q) \cdot [DF_m^{-T}(\hat{\nabla} \hat{v}_{\beta})](\hat{x}_q) J_m(\hat{x}_q) w_q$$

$$A = \sum_m P_m^T a_m P_m$$



# Local VS global right-hand-side

$$F_i = \sum_m \int_{T_m} f v_i dT_m = \sum_m \int_{\hat{T}} [f \circ F_m] [v_i \circ F_m] J_m d\hat{T}$$

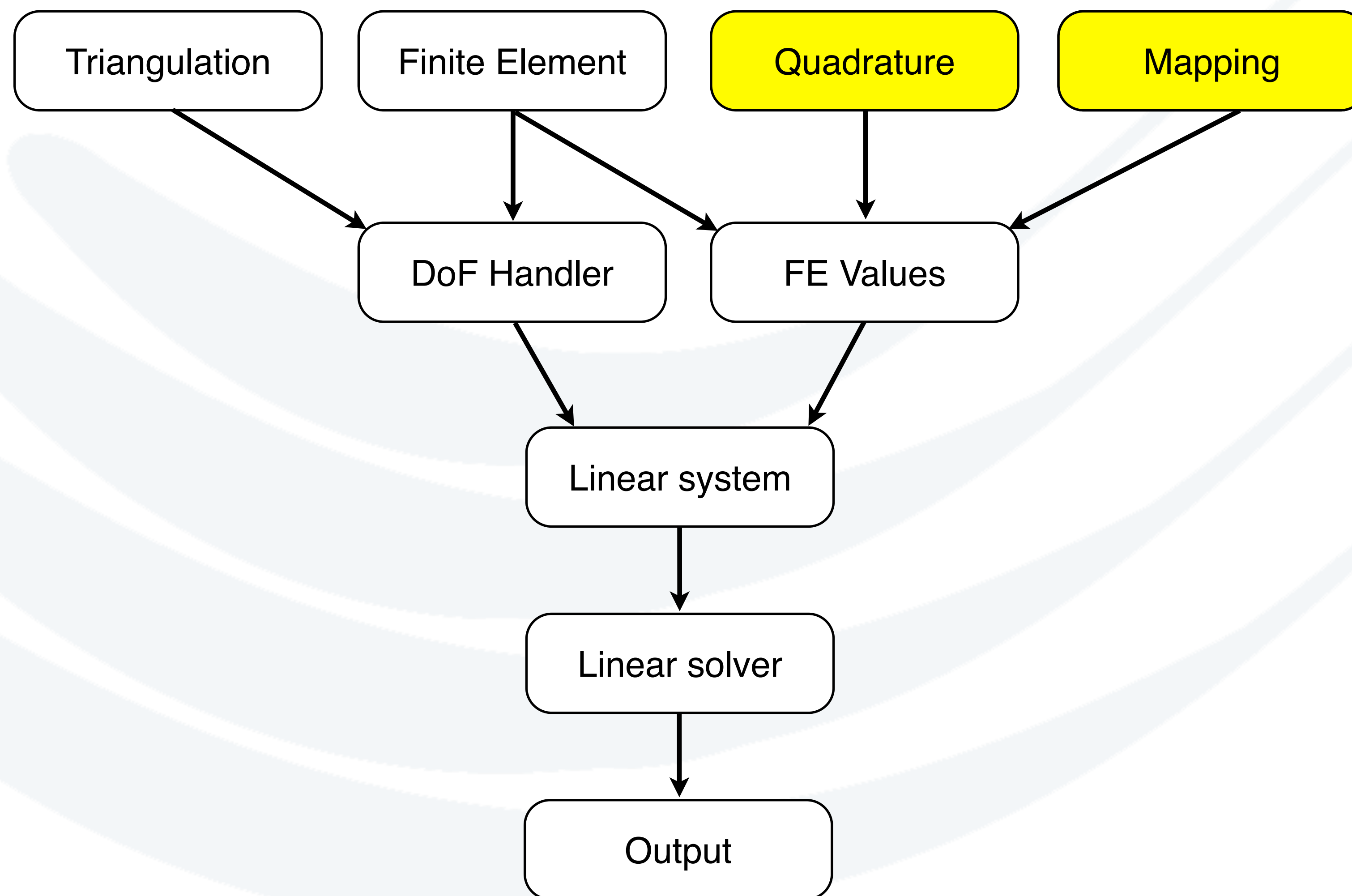
$$F_i = \sum_m \sum_{\alpha} \sum_q P_{mi\alpha} [f \circ F_m](\hat{x}_q) \hat{v}_{\alpha}(\hat{x}_q) J_m(\hat{x}_q) w_q$$

$$f_m_{\alpha} := \sum_{\alpha} \sum_q [f \circ F_m](\hat{x}_q) \hat{v}_{\alpha}(\hat{x}_q) J_m(\hat{x}_q) w_q$$

$$F = \sum_m P_m^T f_m$$



# Structure of a prototypical FE problem



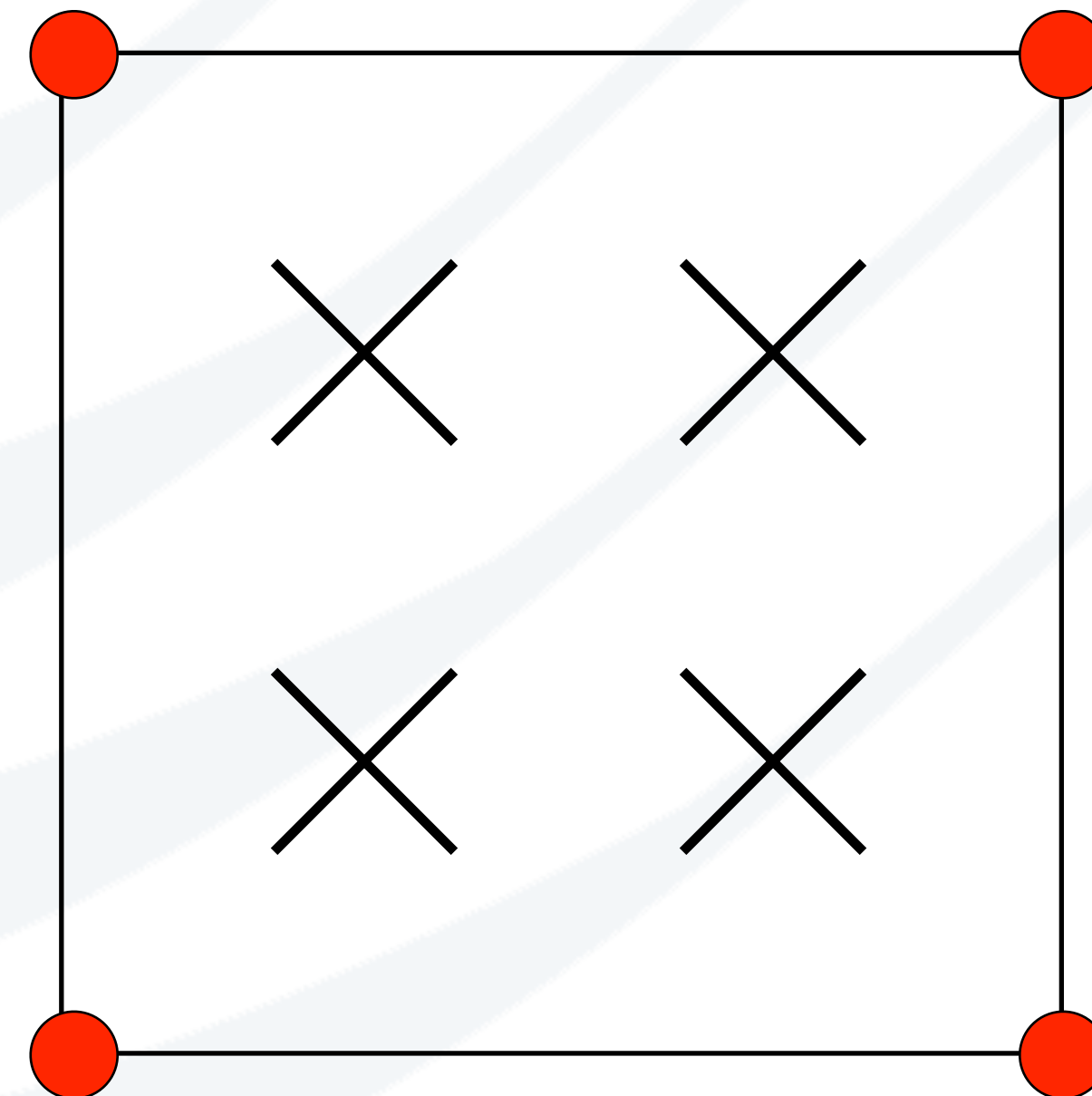




# Integration on a cell: the Quadrature classes

- n-Order Gauss quadrature
- Other rules
  - Gauss Lobatto
  - Simpson
  - Trapezoidal
  - Midpoint
  - A few others
- Anisotropic
- Tensor product

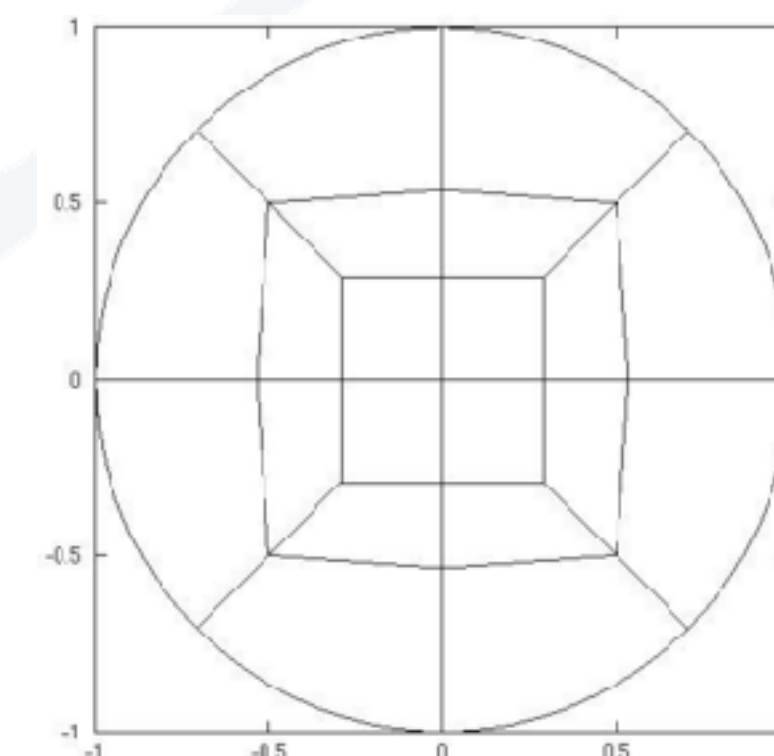
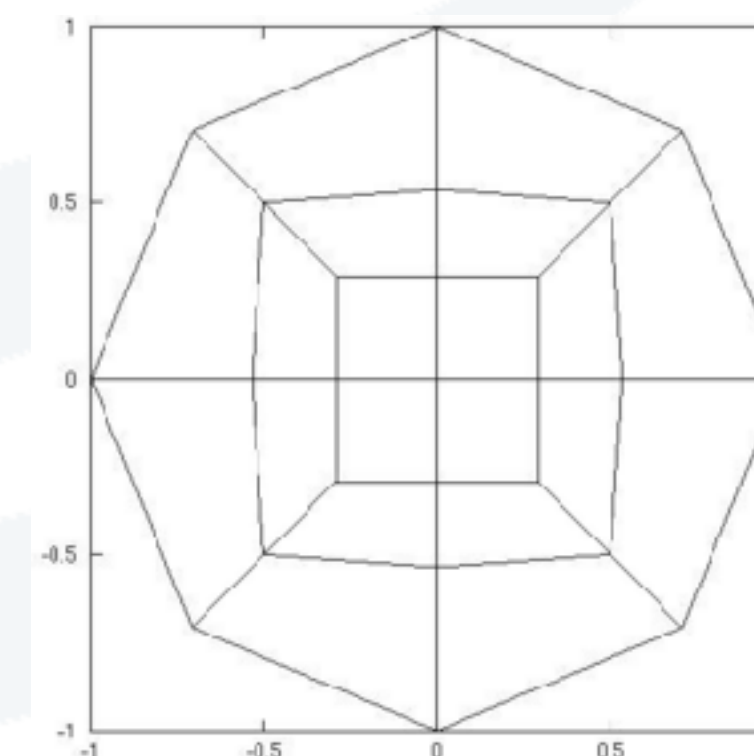
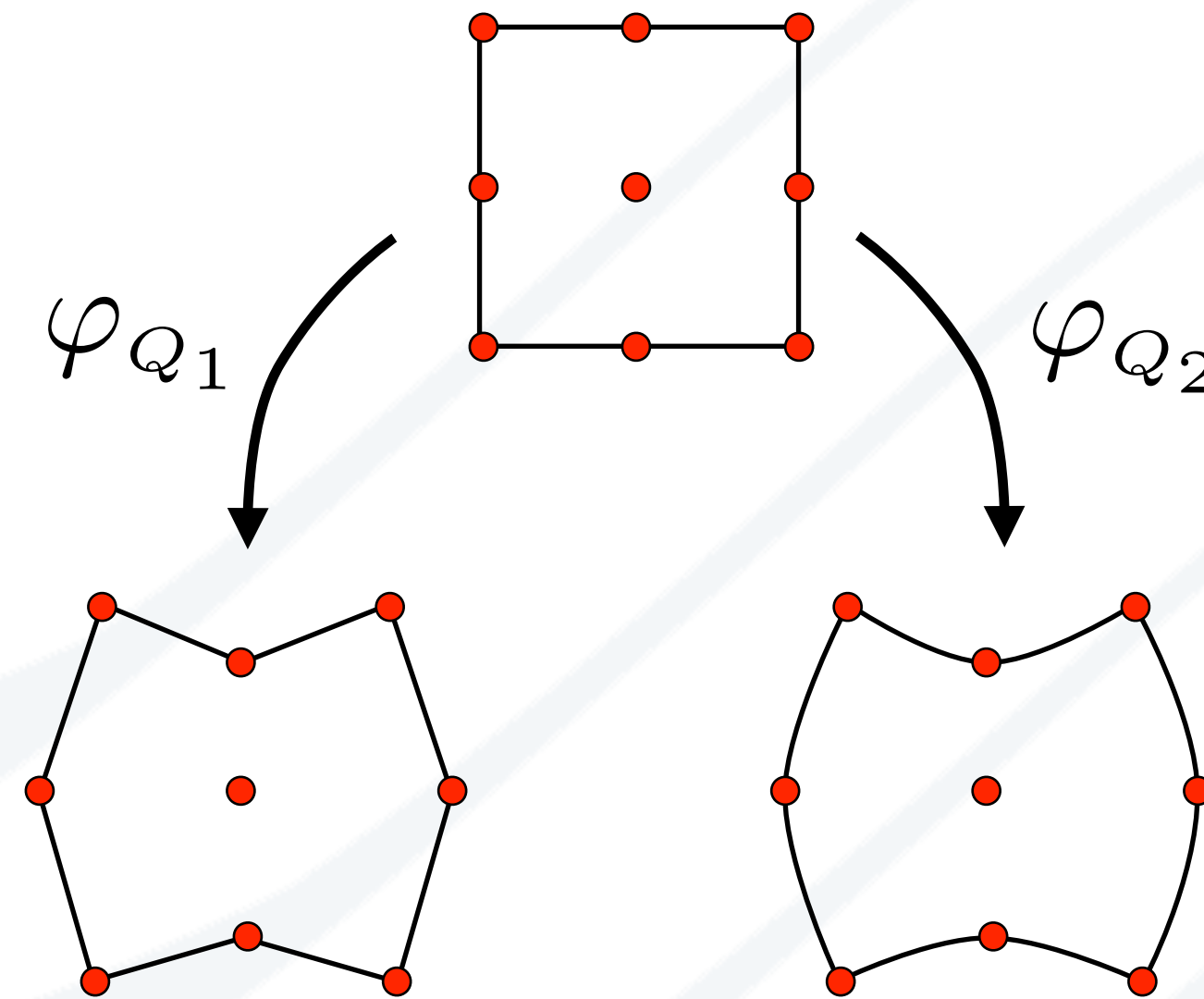
FE\_Q<2>(1)





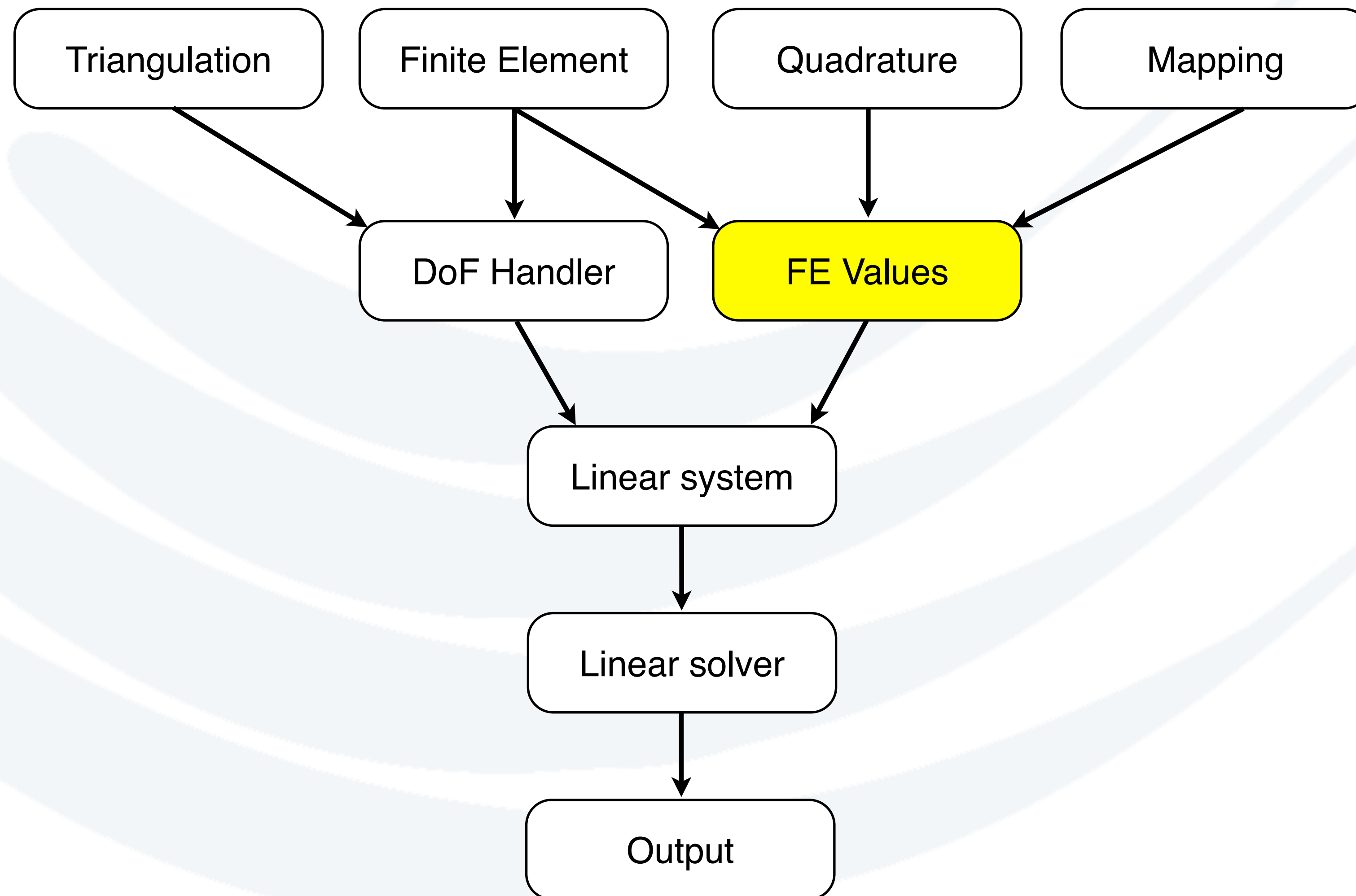
# Integration on a cell: the Mapping classes

- n-order mappings
- Increase accuracy of:
  - Integration schemes
  - Surface basis vectors
- Lagrangian / Eulerian
  - Latter useful for fluid and contact problems, data visualisation
- Boundary and interior manifolds





# Structure of a prototypical FE problem







# Integration on a cell: the FEValues class

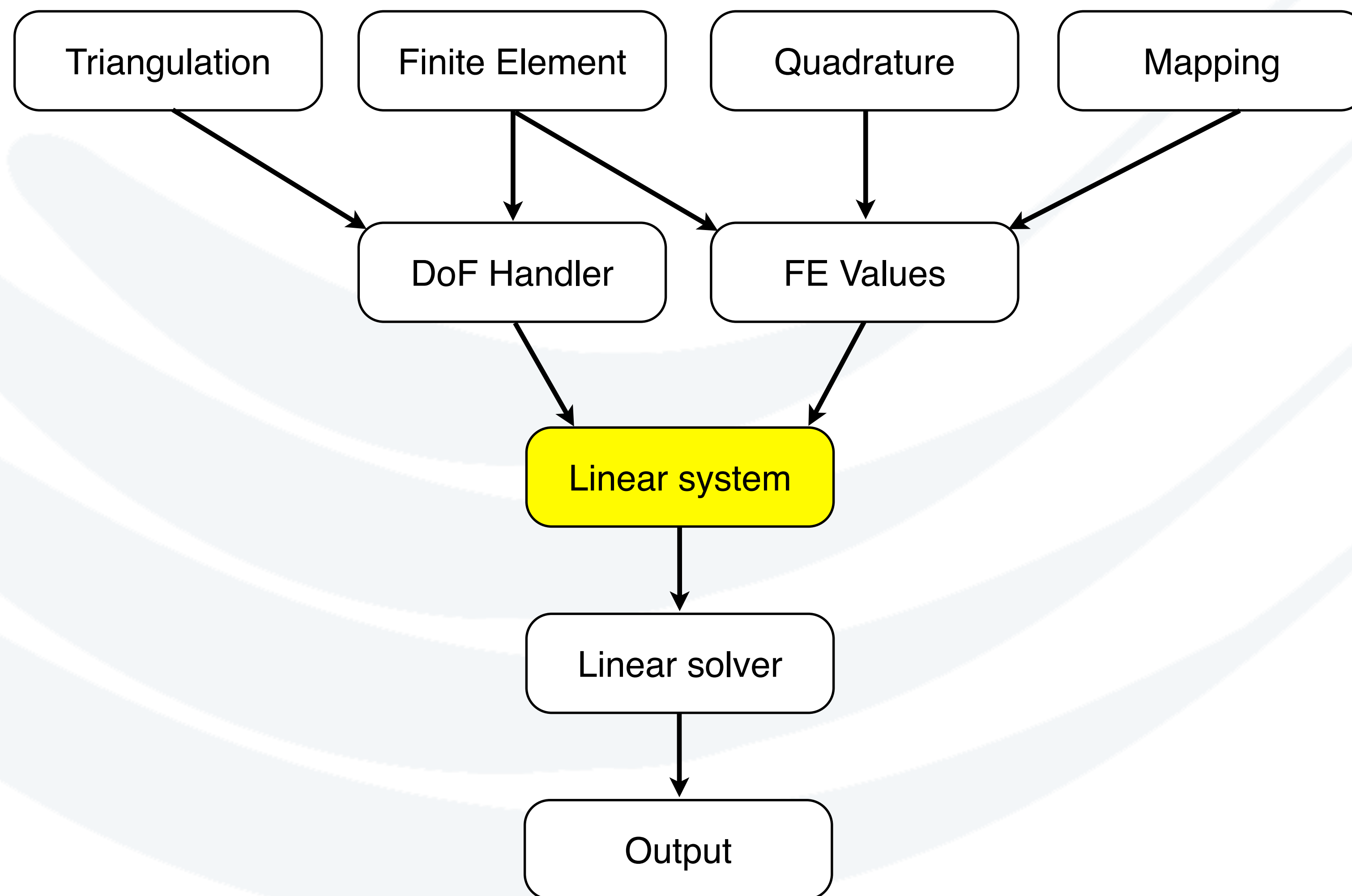
- Object that helps perform integration
- Combines information of:
  - Cell geometry
  - Finite-element system
  - Quadrature rule
  - Mappings
- Can provide:
  - Shape function data
  - Quadrature weights and mapping jacobian at a point
  - Normal on face surface
  - Covariant/contravariant basis vectors
- More ways it can help:
  - Object to extract shape function data for individual fields
  - Natural expressions when coding
- Low level optimisations

$$a_{IJ} := \sum_q [(DF_m^{-T} \hat{\mathbf{V}} \hat{\mathbf{v}}_I)](\hat{\mathbf{x}}_q) \cdot [DF_m^{-T}(\hat{\mathbf{V}} \hat{\mathbf{v}}_J)](\hat{\mathbf{x}}_q) J_m(\hat{\mathbf{x}}_q) w_q$$

```
cell_matrix(I,J) +=  
    * fe_values.shape_grad (I, q_point)  
    * fe_values.shape_grad (J, q_point)  
    * fe_values.JxW (q_point);
```



# Structure of a prototypical FE problem





# Sparse linear systems

- Minimise data storage
  - Evaluate grid connectivity
- Functions to help set up
  - Connectivity
  - Constraints
- Minimal access times
  - Direct manipulation of (non-zero) entries
  - Matrix-vector operations
    - Skip over zero-entries
- Types
  - Unity (monolithic, contiguous)
  - Block sparse structures
- Sub-organisation (e.g. component-wise)

$$[K] \{d\} = \{F\}$$

$$\begin{bmatrix} K_{11} & K_{12} \\ K_{21} & K_{22} \end{bmatrix} \begin{Bmatrix} d_1 \\ d_2 \end{Bmatrix} = \begin{Bmatrix} F_1 \\ F_2 \end{Bmatrix}$$

$$\begin{aligned} & \cdot \left( K_{11} - K_{12} K_{22}^{-1} K_{21} \right) d_1 \\ & \quad = F_1 - K_{12} K_{22}^{-1} F_2 \\ & \cdot d_2 = K_{22}^{-1} (F_2 - K_{21} d_1) \end{aligned}$$





# Solving Poisson's equation

- Demonstration: Step-3  
[https://www.dealii.org/current/doxygen/deal.II/step\\_3.html](https://www.dealii.org/current/doxygen/deal.II/step_3.html)  
<http://www.math.colostate.edu/~bangerth/videos.676.10.html>
- Key points
  - Local assembly + quadrature rules
  - Distribution of local contributions to the global linear system
  - Application of boundary conditions
  - Solving a linear system
  - Output for visualisation

