

1) Let S be a sequence of m integer pairs $\langle (x_1, y_1), (x_2, y_2), \dots, (x_m, y_m) \rangle$. Each of the values x_i and y_i , for all $1 \leq i \leq m$, is an integer between 1 and n . A domino sequence is a subsequence $\langle (x_{i_1}, y_{i_1}), (x_{i_2}, y_{i_2}), \dots, (x_{i_t}, y_{i_t}) \rangle$ such that $1 \leq i_1 < i_2 < \dots < i_t \leq m$ and, for all $1 \leq j < t$, $y_{i_j} = x_{i_{j+1}}$. Note that it isn't necessarily true that $i_{j+1} = i_j + 1$, that is, the elements of the domino sequence don't have to be consecutive in S , but they have to appear in the right order.

Example: For $S = \langle (1, 3), (4, 2), (3, 5), (2, 3), (3, 8) \rangle$, both $\langle (1, 3), (3, 5) \rangle$ and $\langle (4, 2), (2, 3), (3, 8) \rangle$ are domino sequences.

Use dynamic programming to find a longest domino sequence of S in $O(n + m)$ time. Argue briefly that the running time of your algorithm is indeed $O(n + m)$ and that its output is indeed a longest domino sequence of S .

2) You are given a set $X = \{x_1, x_2, \dots, x_n\}$ of points on the real line. Your task is to design a greedy algorithm that finds a smallest set of intervals, each of length 2, that contains all the given points.

Example: Suppose that $X = \{1.5, 2.0, 2.1, 5.7, 8.8, 9.1, 10.2\}$. Then the three intervals $[1.5, 3.5]$, $[4, 6]$, and $[8.7, 10.7]$ are length-2 intervals such that every $x \in X$ is contained in one of the intervals. Note that 3 is the minimum possible number of intervals because points 1.5, 5.7, and 8.8 are far enough from each other that they have to be covered by 3 distinct intervals. Also, note that my solution is not unique – for example, I can shift the middle interval $[4, 6]$ to the right, say to $[5.7, 7.7]$, without disturbing the other intervals, and we would still have an optimal solution.

(a) Suppose that elements of X are presented in increasing order. Write a greedy algorithm code, running in $O(n)$ time, for this problem.