

On Real Roots of Dense and Sparse Random Polynomials



$$f(x) = \underbrace{a_0 + a_1 x^3 + a_2 x^{11} + \dots + a_m x^{200}}_{\# \text{ of real roots.}} + \dots + a_{m-1} x^{198}$$

15 terms.

$$\left| f(x) = x^{100} + x^{10} - x \right.$$

Descartes sign rule:



If the # of sign changes is t in $f(x)$
the # of +ve real roots is $\#$
 t or $t-2$ or $t-4, \dots$

$$f(-x) \underset{-ve}{\sim}$$

Sparsity of f
is the no. of nonzero
monomials

For a K -sparse polynomials the total
no. of real roots is $\leq \boxed{2K-2}$.

$$(x - x^5 + \cancel{x^{10}}) \rightarrow +\infty$$

Ex:

$$f(x) = \frac{(x+1)}{2} \underbrace{(x^3 - 3x + 4)}_3 + \frac{(x^5 + 2)}{2} \underbrace{(x^5 - 4x^3 + 3)}_3$$

The Renz & Conjecture:- (Pascal Koiran)

$$F = \sum_{i=1}^m \left(\prod_{j=1}^{k_i} f_{ij} \right) \quad K := \max_i (k_i)$$

depth 4 $\sum \prod \sum \prod$

where f_{ij} are t sparse

then the # of real roots of F is $O((mKt)^c)$
equivalently $O(\text{poly}(mkt))$

$$O(m t^K)$$

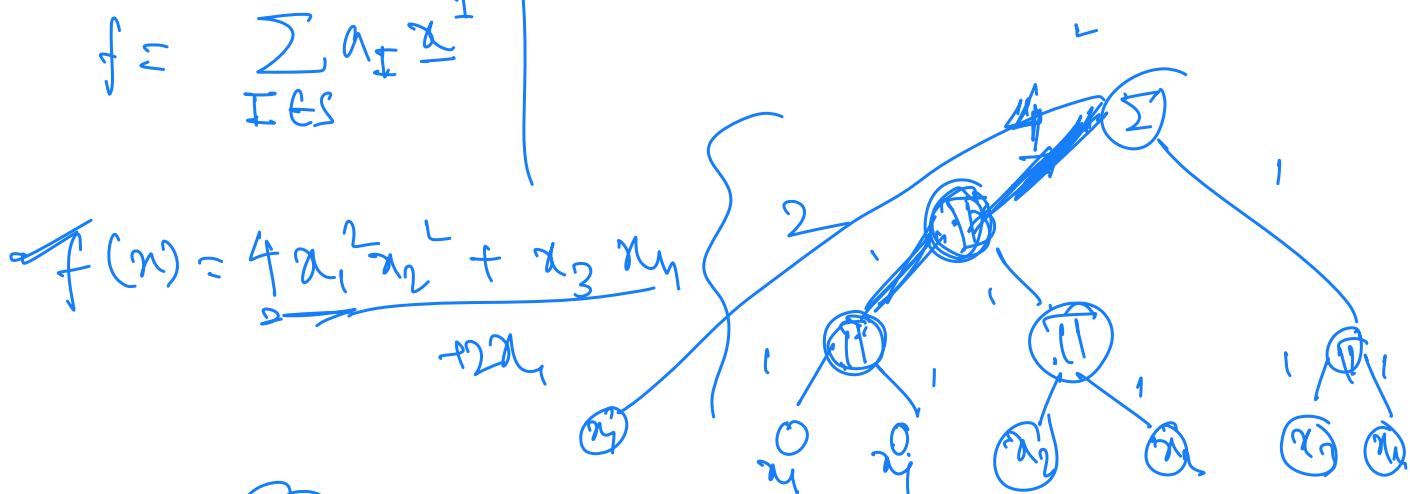
Arithmetic Circuits:

$$f \in F[a_0, \dots, a_n]$$

$$f = \sum_{I \in S} a_I x^I$$

No ds:

$$\Sigma ; \Pi$$



$\Sigma \Pi \Sigma \Pi$ arithmetic circuits.

$$\Sigma$$

$$\Pi$$

$$f_{ij} = \sum_{d_i=1}^t a_{ij} x^{d_i} \rightarrow \Sigma \Pi \text{ circuit}$$

depth -2

Bürgisser showed

→ The Real \mathcal{V} conjecture holds
in arrange case.

$$O(mK^2t)$$

$$\text{VP}^0 \neq \text{VNP}^0$$

$\Phi \leq \text{NP}$

$$\text{VP}_{\mathbb{C}} \neq \text{VNP}_{\mathbb{C}}$$

Coupled with some upper bounds

$$\boxed{\text{VP} \neq \text{VNP}}$$

Valiant's P

$$\sqrt{P} \stackrel{?}{=} VP$$

$$f(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

$$a_i \sim N(0,1)$$

E_n := Expected value of number of real zeros.

(1995)

Bürgisser showed the expected no. of real roots for K-th sparse poly. is $O(\sqrt{K} \log K)$

2020

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MPI

Jindal/Pandey/-

→ this is $O(\sqrt{K})$. AND it's tight.

Theorem 1: Let $S \subseteq \mathbb{N}$ be the support of f .

$$|S| = K. \quad S = \{e_1, \dots, e_K\}$$

$$f = \sum_{i=1}^K x_i^{e_i} \quad \text{if } f \neq 0$$

$$|E_S^K| \leq \frac{2}{\pi} \sqrt{K-1}$$

Thm 2: \exists a seqⁿ $S_n \subset \mathbb{N}$ $|S_n| = K+2$

s.t. $\forall K \geq 3$ $|E_{S_n}| \geq \frac{\pi - \sqrt{3}}{16\pi} \sqrt{K} + \frac{1}{4}$

Theorem (Kac '43) (Dense case)

$$f(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

$a_i \in N(0,1)$

$$E_n = \frac{1}{\pi} \int_{-\infty}^{\infty} \sqrt{\frac{1}{(t^2-1)^2} - \frac{(n+1)^2 t^{2n}}{(t^{2n+2}-1)^2}} dt$$

$$E_n \sim O(\log n)$$

Edelman-Kostlan's derivation of Kac's integral:-

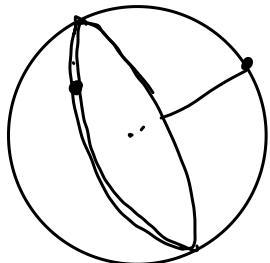
$$\begin{matrix} \vec{a} = \begin{pmatrix} a_0 \\ \vdots \\ a_n \end{pmatrix} : & \vartheta(t) = \begin{pmatrix} 1 \\ t \\ t^2 \\ \vdots \\ t^n \end{pmatrix} \in \mathbb{R}^{n+1} \\ f & \text{moment curve} \end{matrix} \quad t \in \mathbb{R}$$

t is a root iff $a \cdot \vartheta(t) = 0$

$$\alpha := a / \|a\| ; \quad \gamma(t) = \vartheta(t) / \| \vartheta(t) \|$$

Let S^n be the unit sphere on \mathbb{R}^{n+1} .

$\alpha, \gamma(t) \in S^n$.



Observation:

If $a \in N(0, I)$

the $\alpha = a/\|a\| \sim \text{Unif on } S^n$.

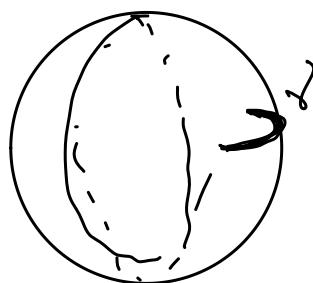
* density of a is independent of the orientation.

* Def^h: (Equator) $P_\perp := \{ \alpha \in S^n \mid \alpha \perp P \}$

Let γ be a curve on S^n ,

$\gamma_\perp := \{ P_\perp \mid P \in \gamma \}$

Def^h: The multiplicity of a point $Q \in \gamma_\perp$ is the number of equators in γ_\perp that contain Q [denoted by $m(Q)$]



Def^h: A curve on S^n is called "rectifiable" if the length is finite.

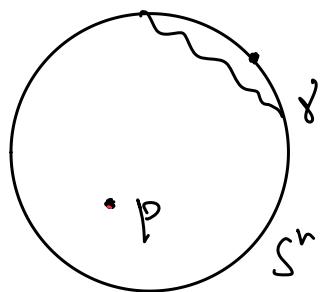
Lemma: for a rectifiable curve γ .

$$|\gamma_\perp| = \frac{|\gamma|}{\pi} \cdot A_n$$

area swept by equators includes overlapping)

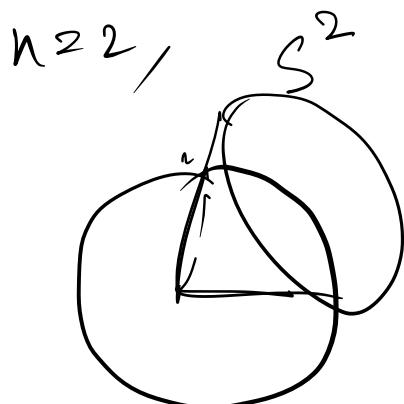
area of S^n .

$E_n :=$ expected # of real roots.



let $p \in S^n$

$m(p) =$ # real root of the poly. corresponding to p .



$$\gamma(t) = \begin{pmatrix} 1 \\ t \\ t^2 \end{pmatrix} \rightarrow \text{squashed cone}$$

$\gamma(t)$ is rectifiable

$$E_n = \int_{p \in S^n} m(p) \times \text{density of } p$$

$$= \frac{1}{A_n} \cdot \boxed{\int_{p \in S^n} m(p)}$$

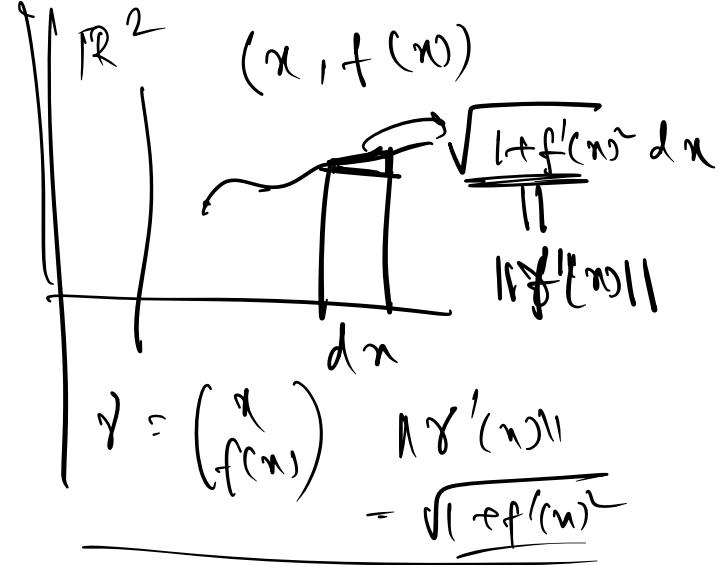
Area swept by $\gamma_1 = |\gamma_1|$

$$= \frac{|\gamma|}{\pi} \cdot A_n$$

$$\Rightarrow \boxed{E_n = \frac{|\gamma|}{\pi}}$$

$$E_n = \frac{1}{\pi} \int_{-\infty}^{\infty} \|\gamma'(t)\| dt$$

after doing messy algebra to revive Kac's integral.



Sparsity Conj

$f \in \mathbb{F}[x_0, \dots, x_n]$ with $\deg_{x_i}(f) < d$
 if $f = g \cdot h$ $g, h \in \mathbb{F}[x_0, \dots, x_n]$
 then $\|f\| \stackrel{o(\text{poly}(d))}{\geq} \|g\|$

$$s = \|f\| := \text{sparsity of } f.$$

$$\# \text{ real roots} \leq O(\sqrt{K})$$

$$\overline{(\# \text{ real roots})^2} \leq K \rightarrow \underline{\text{sparsity}}$$

$$\cancel{K \geq (\underset{\text{of } f}{\text{avg}} \# \text{ real roots})^2 \geq (\underset{\text{of } s}{\text{avg}} \# \text{ real roots})^2}$$

$$(n - \underset{\text{upper bound wrt } \|g\|}{\text{avg}} \# \text{ imaginary roots})^2$$

