

SMART CONTRACT AUDIT REPORT

for

Venus Token Converter

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1 Introduction

Given the opportunity to review the design document and related smart contract source code of the Venus Token Converter protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

1.1 About Venus Token Converter

The Venus protocol enables a complete algorithmic money market protocol on BNB Smart Chain. Venus enables users to utilize their cryptocurrencies by supplying collateral to the protocol that may be borrowed by pledging over-collateralized cryptocurrencies. The audited Venus Token Converter supports the conversion from the generated protocol incomes and distribution to different targets (RiskFund, Prime, VTreasury, or XVSVault). The basic information of the audited protocol is as follows:

Item	Description
Name	Venus
Website	https://venus.io/
Туре	EVM Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	September 27, 2023

Table 1.1: Basic Information of The Venus Token Converter

In the following, we show the Git repository of reviewed files and the commit hash value used in this audit. This audit covers on the changes from PR #9: https://github.com/VenusProtocol/protocol-reserve/pull/9.

https://github.com/VenusProtocol/protocol-reserve.git (cc3d674)

And here is the commit ID after fixes for the issues found in the audit have been checked in:

https://github.com/VenusProtocol/protocol-reserve.git (5d1c6b3)

1.2 About PeckShield

PeckShield Inc. [12] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

High Critical High Medium

High Medium

Low

High Low

High Medium

Low

High Medium

Low

Likelihood

Table 1.2: Vulnerability Severity Classification

1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [11]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

Table 1.3: The Full List of Check Items

Category	Check Item
	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
Basic Coding Bugs	Revert DoS
Dasic Couling Dugs	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
Advanced DeFi Scrutiny	Digital Asset Escrow
rataneed Der i Geraemi,	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
Additional Recommendations	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [10], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings.

1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during
	the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functional-
	ity that processes data.
Numeric Errors	Weaknesses in this category are related to improper calcula-
	tion or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like
	authentication, access control, confidentiality, cryptography,
	and privilege management. (Software security is not security
	software.)
Time and State	Weaknesses in this category are related to the improper man-
	agement of time and state in an environment that supports
	simultaneous or near-simultaneous computation by multiple
	systems, processes, or threads.
Error Conditions,	Weaknesses in this category include weaknesses that occur if
Return Values,	a function does not generate the correct return/status code,
Status Codes	or if the application does not handle all possible return/status
	codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper manage-
	ment of system resources.
Behavioral Issues	Weaknesses in this category are related to unexpected behav-
	iors from code that an application uses.
Business Logics	Weaknesses in this category identify some of the underlying
	problems that commonly allow attackers to manipulate the
	business logic of an application. Errors in business logic can
	be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used
	for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of
	arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written
	expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices
	that are deemed unsafe and increase the chances that an ex-
	ploitable vulnerability will be present in the application. They
	may not directly introduce a vulnerability, but indicate the
	product has not been carefully developed or maintained.

2 | Findings

2.1 Summary

Here is a summary of our findings after analyzing the design and implementation of the Venus Token Converter protocol. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings
Critical	0
High	3
Medium	1
Low	2
Informational	0
Total	6

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities that need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 3 high-severity vulnerabilities, 1 medium-severity vulnerability, and 2 low-severity vulnerabilities.

Title ID Severity Category **Status** PVE-001 High Public Exposure of State-Modifying Security Features Resolved Functions in RiskFundConverter Possible Underflow in RiskFundCon-**PVE-002** Medium **Business Logic** Resolved verter::postSweepToken() PVE-003 Numeric Errors Low Improved Precision in AbstractToken-Resolved Converter::getAmountIn() PVE-004 Low Trust Issue of Admin Keys Security Features Resolved **PVE-005** Resolved High Inconsistent Storage Layout For Risk-Coding Practices Fund Upgrade **PVE-006** RiskFundCon-Resolved High Incorrect Business Logic verter::postConversionHook() Logic

Table 2.1: Key Venus Token Converter Audit Findings

Besides recommending specific countermeasures to mitigate these issues, we also emphasize that it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms need to kick in at the very moment when the contracts are being deployed in mainnet. Please refer to Section 3 for details.

3 Detailed Results

3.1 Public Exposure of State-Modifying Functions in RiskFundConverter

• ID: PVE-001

• Severity: High

• Likelihood: High

• Impact: High

• Target: RiskFundConverter

• Category: Business Logic [8]

• CWE subcategory: CWE-841 [5]

Description

The audited Venus Token Converter contract converts and distributes the protocol incomes to different recipients. While examining the specific RiskFundConverter contract, we notice the exposure of a public function and this exposure may be abused to corrupt the intended token conversion functionality.

In the following, we show the code snippet from the exposed <code>postSweepToken()</code> routine. This routine is designed to be an internal helper for the public <code>sweepToken()</code> function. In other words, this routine should not be exposed to public. Its public exposure without any caller authentication will corrupt the internal asset states and thus cripple the token conversion functionality.

```
173
        function postSweepToken(address tokenAddress, uint256 amount) public override {
174
             uint256 balance = IERC20Upgradeable(tokenAddress).balanceOf(address(this));
175
             uint256 balanceDiff = balance - assetsReserves[tokenAddress];
176
177
             uint256 amountDiff = amount - balanceDiff;
178
             if (amountDiff > 0) {
179
                 address[] memory pools = getPools(tokenAddress);
180
                 uint256 assetReserve = assetsReserves[tokenAddress];
181
                 for (uint256 i; i < pools.length; ++i) {</pre>
182
                     uint256 poolShare = (poolsAssetsReserves[pools[i]][tokenAddress] *
                         EXP_SCALE) / assetReserve;
183
                     if (poolShare == 0) continue;
184
                     updatePoolAssetsReserve(pools[i], tokenAddress, amount, poolShare);
```

```
185 }
186 assetsReserves[tokenAddress] -= amountDiff;
187 }
188 }
```

Listing 3.1: RiskFundConverter::postSweepToken()

Recommendation Declare the above postSweepToken() routine as an internal function.

Status This issue has been fixed in the following commit: a1e6697.

3.2 Possible Underflow in RiskFundConverter::postSweepToken()

ID: PVE-002Severity: MediumLikelihood: LowImpact: High

• Target: RiskFundConverter

Category: Coding Practices [7]CWE subcategory: CWE-563 [4]

Description

In last Section 3.1, we have examined the public exposure of an internal helper routine. In this section, we further examine the logic of this helper and notice its logic needs to be improved.

Specifically, we show below the code snippet from the postSwepToken() routine. Within this routine, there is a assetsReserves state that keeps track of the amount for each asset that is transferred to ProtocolShareReserve combined (for all pools). And the information about the pool-specific asset amount in ProtocolShareReserve is saved in another state poolsAssetsReserves. It comes to our attention that when this helper postSweepToken() is called, the balanceDiff variable computes the amount of tokens that is free to sweep. If the requested amount is larger than balanceDiff, there is a need to reduce the pool-specific asset reserve to fulfill the sweep request. However, the calculation of amount - balanceDiff (line 177) may lead to a possible arithmetic underflow. Moreover, the pool-specific asset reserve needs to be updated as updatePoolAssetsReserve(pools[i], tokenAddress, amountDiff, poolShare), instead of the current approach of updatePoolAssetsReserve(pools[i], tokenAddress, amount, poolShare) (line 184).

```
function postSweepToken(address tokenAddress, uint256 amount) public override {
    uint256 balance = IERC20Upgradeable(tokenAddress).balanceOf(address(this));
    uint256 balanceDiff = balance - assetsReserves[tokenAddress];

uint256 amountDiff = amount - balanceDiff;
    if (amountDiff > 0) {
```

```
179
                 address[] memory pools = getPools(tokenAddress);
180
                 uint256 assetReserve = assetsReserves[tokenAddress];
181
                 for (uint256 i; i < pools.length; ++i) {</pre>
182
                     uint256 poolShare = (poolsAssetsReserves[pools[i]][tokenAddress] *
                         EXP_SCALE) / assetReserve;
183
                     if (poolShare == 0) continue;
184
                     updatePoolAssetsReserve(pools[i], tokenAddress, amount, poolShare);
185
                 }
186
                 assetsReserves[tokenAddress] -= amountDiff;
187
188
```

Listing 3.2: RiskFundConverter::postSweepToken()

Recommendation Revise the above postSweepToken() logic to properly keep track of the pool-specific reserve.

Status This issue has been fixed in the following commit: fd20124.

3.3 Improved Precision in AbstractTokenConverter::getAmountIn()

• ID: PVE-003

• Severity: Low

• Likelihood: Medium

• Impact: Low

• Target: AbstractTokenConverter

Category: Numeric Errors [9]

• CWE subcategory: CWE-190 [2]

Description

SafeMath is a widely-used Solidity math library that is designed to support safe math operations by preventing common overflow or underflow issues when working with uint256 operands. While it indeed blocks common overflow or underflow issues, the lack of float support in Solidity may introduce another subtle, but troublesome issue: precision loss. In this section, we examine one possible precision loss source that stems from the mixed multiplication (mul) and division (div).

In particular, we use the AbstractTokenConverter::getAmountIn() as an example. This routine is used to calculate the amount of input tokens sender would send on receiving the requested amountOutMantissa of output tokens.

```
function getAmountIn(

uint256 amountOutMantissa,

address tokenAddressIn,

address tokenAddressOut

) public returns (uint256 amountConvertedMantissa, uint256 amountInMantissa) {
```

```
427
             if (amountOutMantissa == 0) {
428
                 revert InsufficientInputAmount();
429
431
             ConversionConfig memory configuration = convertConfigurations[tokenAddressIn][
                 tokenAddressOut];
433
             if (!configuration.enabled) {
434
                 revert ConversionConfigNotEnabled();
435
437
             priceOracle.updateAssetPrice(tokenAddressIn);
438
             priceOracle.updateAssetPrice(tokenAddressOut);
440
             uint256 maxTokenOutReserve = balanceOf(tokenAddressOut);
441
             uint256 tokenInUnderlyingPrice = priceOracle.getPrice(tokenAddressIn);
442
             uint256 tokenOutUnderlyingPrice = priceOracle.getPrice(tokenAddressOut);
444
            /// amount of tokenAddressOut after including incentive \,
445
             uint256 conversionWithIncentive = MANTISSA_ONE + configuration.incentive;
446
             /// conversion rate after considering incentive(conversionWithIncentive)
447
             uint256 tokenInToOutConversion = (tokenInUnderlyingPrice *
                 conversionWithIncentive) / tokenOutUnderlyingPrice;
449
             amountInMantissa = ((amountOutMantissa * EXP_SCALE) / tokenInToOutConversion);
450
             amountConvertedMantissa = amountOutMantissa;
452
             /// If contract has less Liquity for tokenAddressOut than amountOutMantissa
453
             if (maxTokenOutReserve < amountOutMantissa) {</pre>
454
                 amountInMantissa = ((maxTokenOutReserve * EXP_SCALE) /
                     tokenInToOutConversion);
455
                 amountConvertedMantissa = maxTokenOutReserve;
456
            }
457
```

Listing 3.3: AbstractTokenConverter::getAmountIn()

We notice the calculation of the resulting amountInMantissa (line 449) involves mixed multiplication and devision. For improved precision, it is better to round-up the computation in favor of the protocol. Currently, it simply takes a round-down approach to compute the result. Note that the resulting precision loss may be just a small number, but it plays a critical role when certain boundary conditions are met. And it is always the preferred choice if we can avoid the precision loss as much as possible. Note the <code>getAmountOut()</code> routine can be similarly improved.

Recommendation Revise the above calculations to better mitigate possible precision loss.

Status This issue has been fixed in the following commits: 92342eb and 057dcee.

3.4 Trust Issue of Admin Keys

• ID: PVE-004

• Severity: Low

Likelihood: Low

Impact: Low

• Target: Multiple Contracts

• Category: Security Features [6]

• CWE subcategory: CWE-287 [3]

Description

In the new Venus Token Converter protocol, there is a privileged owner as well as the associated access control module, i.e., AccessControlledV8. They play a critical role in governing and regulating the system-wide operations (e.g., parameter setting and role authorization). They may also affect the flow of assets managed by this protocol. Our analysis shows that these privileged accounts need to be scrutinized. In the following, we examine the privileged accounts and their related privileged accesses in current contracts.

```
58
       function setConvertibleBaseAsset(address convertibleBaseAsset_) external onlyOwner {
59
            ensureNonzeroAddress(convertibleBaseAsset_);
60
           emit ConvertibleBaseAssetUpdated(convertibleBaseAsset, convertibleBaseAsset_);
61
           convertibleBaseAsset = convertibleBaseAsset_;
62
       }
63
64
       /// @dev Risk fund converter setter
65
       /// @param riskFundConverter_ Address of the risk fund converter
66
       /// @custom:event RiskFundConverterUpdated emit on success
67
       /// @custom:error ZeroAddressNotAllowed is thrown when risk fund converter address
           is zero
68
       function setRiskFundConverter(address riskFundConverter_) external onlyOwner {
69
           ensureNonzeroAddress(riskFundConverter_);
70
           emit RiskFundConverterUpdated(riskFundConverter, riskFundConverter_);
71
           riskFundConverter = riskFundConverter_;
72
       }
73
74
       /// @dev Shortfall contract address setter
75
       /// @param shortfallContractAddress_ Address of the auction contract
76
       /// @custom:error ZeroAddressNotAllowed is thrown when shortfall contract address is
77
       function setShortfallContractAddress(address shortfallContractAddress_) external
           ensureNonzeroAddress(shortfallContractAddress_);
78
79
           emit ShortfallContractUpdated(shortfall, shortfallContractAddress_);
80
           shortfall = shortfallContractAddress_;
```

Listing 3.4: Example Privileged Operations in the RiskFundV2 Contract

If the privileged admins are managed by a plain EOA account, this may be worrisome and pose counter-party risk to the exchange users. A multi-sig account could greatly alleviate this concern, though it is still far from perfect. Specifically, a better approach is to eliminate the administration key concern by transferring the role to a community-governed DAO. In the meantime, a timelock-based mechanism can also be considered as mitigation.

Moreover, it should be noted that current contracts have the support of being deployed behind a proxy. And there is a need to properly manage the proxy-admin privileges as they fall in this trust issue as well.

Recommendation Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

Status This issue has been resolved as the owner is managed by the Venus: Timelock contract deployed at 0x939bd8d64c0a9583a7dcea9933f7b21697ab6396.

3.5 Inconsistent Storage Layout For RiskFund Upgrade

• ID: PVE-005

• Severity: High

Likelihood: High

Impact: High

Target: RiskFund

• Category: Coding Practices [7]

• CWE subcategory: CWE-1041 [1]

Description

The current Venus Token Converter protocol will upgrade an earlier contract RiskFundV1 contract. Naturally, there is a requirement on the storage consistency between the old version and the new version. Our analysis shows that the new version has a different storage layout from the old version and this inconsistency needs to be resolved before the upgrade.

To elaborate, we show below the storage layout of of the new version RiskFundV2Storage. We notice it inherits five states from the earlier version, i.e., pancakeSwapRouter, minAmountToConvert, convertibleBaseAsset, shortfall, and poolReserves. However, our analysis on the old version indicates the following layout: convertibleBaseAsset, shortfall, pancakeSwapRouter, and minAmountToConvert. Beside the different layout, we also notice an extra poolReserves state that does not show up in the old version. Last but not least, for consistency, the RiskFundV2 contract inherits from Ownable2StepUpgradeable, AccessControlledV8, RiskFundV2Storage, and ReentrancyGuardUpgradeable. The

first Ownable2StepUpgradeable inheritance is redundant as it is already part of AccessControlledV8 inheritance.

```
32
   contract RiskFundV1Storage is ReserveHelpersStorage, MaxLoopsLimitHelpersStorage {
33
       /// @notice This state is deprecated, using it to prevent storage collision
34
       address private pancakeSwapRouter;
35
       /// @notice This state is deprecated, using it to prevent storage collision
36
       uint256 private minAmountToConvert;
37
38
       address public convertibleBaseAsset;
       address public shortfall;
39
40
41
       /// @notice Store base asset's reserve for specific pool
42
       mapping(address => uint256) public poolReserves;
43
  }
44
45 /// @title RiskFundV2Storage
46 /// @author Venus
47 /// @dev Risk fund V2 storage
48 contract RiskFundV2Storage is RiskFundV1Storage {
49
       /// @notice Risk fund converter address
50
       address public riskFundConverter;
```

Listing 3.5: Key Storage States Defined in RiskFundV2

Listing 3.6: Key Storage States Defined in RiskFundV1

Recommendation Ensure the storage consistency between the old version and the new version.

Status This issue has been fixed in the following commit: 74f9cb2.

3.6 Incorrect RiskFundConverter::postConversionHook() Logic

• ID: PVE-006

• Severity: High

Likelihood: High

• Impact: High

• Target: RiskFundConverter

• Category: Business Logic [8]

• CWE subcategory: CWE-841 [5]

Description

As mentioned in Section 3.1, the Venus Token Converter contract converts and distributes the protocol incomes to different recipients. While examining the token conversion logic, we notice there is a post-conversion hook and this hook logic needs to be revisited.

In the following, we show the code snippet from the related postConversionHook() routine inside the RiskFundConverter contract. This routine is designed to perform bookkeeping tasks after the conversion. However, we notice the current implementation mixes tokenInAddress with tokenOutAddress and updates the wrong asset reserve numbers as well. Specifically, the amountIn is expected to be the protocol incoming amount, which will be sent to the destinationAddress.

```
202
                                    function postConversionHook(
203
                                                     address tokenInAddress,
204
                                                     address tokenOutAddress,
205
                                                     uint256 amountIn,
206
                                                     uint256 amountOut
207
                                    ) internal override {
208
                                                     address[] memory pools = getPools(tokenInAddress);
209
                                                     uint256 assetReserve = assetsReserves[tokenInAddress];
210
                                                     for (uint256 i; i < pools.length; ++i) {</pre>
211
                                                                       uint256 poolShare = (poolsAssetsReserves[pools[i]][tokenInAddress] *
                                                                                        EXP_SCALE) / assetReserve;
212
                                                                      if (poolShare == 0) continue;
213
                                                                       updatePoolAssetsReserve(pools[i], tokenInAddress, amountIn, poolShare);
214
                                                                       uint256 poolAmountOutShare = (poolShare * amountOut) / EXP_SCALE;
215
                                                                       IRiskFund (destination \verb|Address|). update \verb|PoolState| (pools[i], token \verb|OutAddress|, token \verb|OutAddress|), token \verb|OutAddress|, to
                                                                                        poolAmountOutShare);
216
                                                     }
217
218
                                                     assetsReserves[tokenInAddress] -= amountIn;
219
```

Listing 3.7: RiskFundConverter::postSweepToken()

Recommendation Revise the above postConversionHook() routine to properly update the token amount. An updated version is shown in the following:

```
function postConversionHook(
```

```
203
                                                     address tokenInAddress,
204
                                                     address tokenOutAddress,
205
                                                     uint256 amountIn,
206
                                                     uint256 amountOut
207
                                    ) internal override {
208
                                                    address[] memory pools = getPools(tokenOutAddress);
209
                                                    uint256 assetReserve = assetsReserves[tokenOutAddress];
210
                                                    for (uint256 i; i < pools.length; ++i) {</pre>
211
                                                                     uint256 poolShare = (poolsAssetsReserves[pools[i]][tokenOutAddress] *
                                                                                     EXP_SCALE) / assetReserve;
212
                                                                     if (poolShare == 0) continue;
213
                                                                     updatePoolAssetsReserve(pools[i], tokenOutAddress, amountOut, poolShare);
214
                                                                     uint256 poolAmountInShare = (poolShare * amountIn) / EXP_SCALE;
215
                                                                     IRiskFund (destination \verb|Address|). update \verb|PoolState| (pools[i], token In \verb|Address|), update \verb|PoolState| (pools[i], token In Address|), update \verb|PoolState| (pools[i], t
                                                                                      poolAmountInShare);
216
217
218
                                                     assetsReserves[tokenOutAddress] -= amountOut;
219
```

Listing 3.8: RiskFundConverter::postSweepToken()

Status This issue has been fixed in the following commit: 15c252d.

4 Conclusion

In this audit, we have analyzed the design and implementation of the Venus Token Converter protocol, which allows for the conversion from the generated protocol incomes and distribution to different targets (RiskFund, Prime, VTreasury, or XVSVault). The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and addressed.

Meanwhile, we need to emphasize that Solidity-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



References

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