

## Homework 1

**Question 1.** [30 marks] Solve the following exercises:

- Exercise [1.4.5](#)
- Exercise [1.4.11](#)
- Exercise [1.5.4](#)

A large part of the course going forward will involve *proofs by induction*. Many of you may have not seen these before, and the following two questions are here to make sure everyone is on the same page.

But first of all, *what* is a “proof by induction”? Well... This is a relatively long story, but here is the gist of it. Often-times we construct a set  $\mathcal{I}$  of mathematical objects in the following way: (a) We describe the “basic objects” of  $\mathcal{I}$ . (b) We describe a way of building larger objects of  $\mathcal{I}$  from the objects we have already built. This is called an **inductive** (or **recursive**) definition.

Then, to prove that some property is true of all elements of  $\mathcal{I}$ , it suffices to prove that: (a) It is true of the “basic objects” of  $\mathcal{I}$ . (b) If it is true of some objects, then it is true of all larger objects built from them.

That’s really abstract, so the next two exercises will try to illustrate it.

**Question 2.** [30 marks] We can build the set of natural numbers  $\mathbb{N}$  by induction as follows: (a)  $0 \in \mathbb{N}$  (that’s the basic object). (b) If  $n \in \mathbb{N}$  then  $n + 1 \in \mathbb{N}$  (that how to build larger objects from smaller ones). If we then wish to prove that some property  $P(n)$  is true of all  $n \in \mathbb{N}$  we have to show that: (a)  $P(0)$  is true. (b) If  $P(n)$  is true, then so is  $P(n + 1)$ . Now, on with the question:

- (1) Prove by induction that for all  $n \in \mathbb{N}$  the following statement is true:

$$\sum_{i=0}^n i^3 = \frac{n^2(n+1)^2}{4}.$$

- (2) Prove by induction that for all  $n \in \mathbb{N}$  the following statement is true:

$$9^n - 1 \text{ is divisible by } 8.$$

If  $\mathcal{I}$  is defined by induction, then we can also use induction to define functions on  $\mathcal{I}$  – to define a function  $f$  on all of  $\mathcal{I}$  we just need to define  $f$  on the “basic objects” and provided we know how to compute  $f$  on simpler objects, we need to define  $f$  on larger ones built from them.

We can build more things by induction, not just the natural numbers. This will occupy us a lot in the next chapter, and the next exercise is a sort-of warm up.

**Question 3.** [40 marks] We define a set  $\mathcal{S}$  of strings of symbols, inductively, as follows: (a)  $\mathbf{a} \in \mathcal{S}$ , and  $\mathbf{b} \in \mathcal{S}$  (these are the basic objects). (b) If  $s, t \in \mathcal{S}$  then:  $(s) \in \mathcal{S}$ ,  $(st) \in \mathcal{S}$  (where  $st$  denotes the concatenation of  $s$  and  $t$ ), and  $([s] \uparrow [t]) \in \mathcal{S}$  (so in this example we have three ways of building objects from smaller ones). For example:

$$\mathbf{a} \in \mathcal{S}, (\mathbf{a}) \in \mathcal{S}, (\mathbf{aa}) \in \mathcal{S}, ((\mathbf{a})\mathbf{a}) \in \mathcal{S}, ([\mathbf{a}] \uparrow [(\mathbf{ba})]) \in \mathcal{S}.$$

We can inductively define a function  $\mathbf{bra} : \mathcal{S} \rightarrow \mathbb{N}$  as follows  $\mathbf{bra}(\mathbf{a}) := 0$ ,  $\mathbf{bra}(\mathbf{b}) := 0$ . Once  $\mathbf{bra}(s), \mathbf{bra}(t)$  have been defined, we set:

$$\begin{aligned}\mathbf{bra}((s)) &:= \mathbf{bra}(s) + 2 \\ \mathbf{bra}((st)) &:= \mathbf{bra}(s) + \mathbf{bra}(t) + 2 \\ \mathbf{bra}([s] \uparrow [t]) &:= \mathbf{bra}(s) + \mathbf{bra}(t) + 6\end{aligned}$$

- (1) Prove that for all  $s \in \mathcal{S}$ , we have that  $\mathbf{bra}(s)$  is even.

Let  $s, t \in \mathcal{S}$ . We can define a function  $\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]} : \mathcal{S} \rightarrow \mathcal{S}$  as follows  $\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}(\mathbf{a}) := s$ ,  $\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}(\mathbf{b}) := t$ , and:

$$\begin{aligned}\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}((s)) &:= (\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}(s)) \\ \mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}((st)) &:= (\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}(s)\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}(t)) \\ \mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}([s] \uparrow [t]) &:= ([\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}(s)] \uparrow [\mathbf{Sub}_{[s/\mathbf{a}, t/\mathbf{b}]}(t)])\end{aligned}$$

- (2) Prove that for all  $s \in \mathcal{S}$ ,  $\mathbf{Sub}_{[\mathbf{a}/\mathbf{a}, \mathbf{b}/\mathbf{b}]}(s)$  does not contain the letter  $\mathbf{b}$ .  
(3) Prove that for all  $s \in \mathcal{S}$ ,  $\mathbf{Sub}_{[\mathbf{bb}/\mathbf{a}, \mathbf{aa}/\mathbf{b}]}(s)$  contains an even number of occurrences of each of  $\mathbf{a}$  and  $\mathbf{b}$ .

Finally, define a function  $\mathbf{let} : \mathcal{S} \rightarrow \mathbb{N}$  as follows:  $\mathbf{let}(\mathbf{a}) = 1$ ,  $\mathbf{let}(\mathbf{b}) = 1$ , and:

$$\begin{aligned}\mathbf{let}((s)) &:= \mathbf{let}(s) \\ \mathbf{let}((st)) &:= \mathbf{let}(s) + \mathbf{let}(t) \\ \mathbf{let}([s] \uparrow [t]) &:= \mathbf{let}(s) + \mathbf{let}(t)\end{aligned}$$

- (4) Show that for all  $s \in \mathcal{S}$ ,  $\mathbf{bra}(\mathbf{Sub}_{[(\mathbf{a})/\mathbf{a}, (\mathbf{b})/\mathbf{b}]}(s)) = \mathbf{bra}(s) + 2\mathbf{let}(s)$ .

[*Hint.* I am asking here for a “formal” proof by induction. Clearly write out what you need to prove. This should guide you.]