

### Homework 3

**Question 1.** [20 marks] Solve the following exercises from **Week 3(b)** of the notes:

- (1) Exercise 1.1.6
- (2) Exercise 1.1.9
- (3) Exercise 1.1.11.

**Question 2.** [20 marks] Let  $\mathcal{L}$  be a first-order language with two unary function symbols  $\underline{f}$  and  $\underline{g}$ .

- (1) Find three  $\mathcal{L}$ -sentences  $\phi_1, \phi_2, \phi_3$ , such that:
  - In every  $\mathcal{L}$ -structure  $\mathcal{M}$ , we have that  $\mathcal{M} \models \phi_1$  if and only if  $\underline{f}$  and  $\underline{g}$  are constant.
  - In every  $\mathcal{L}$ -structure  $\mathcal{M}$ , we have that  $\mathcal{M} \models \phi_2$  if and only if  $\text{im}(\underline{f}) \cap \text{im}(\underline{g})$  contains exactly two elements.
  - In every  $\mathcal{L}$ -structure  $\mathcal{M}$ , we have that  $\mathcal{M} \models \phi_3$  if and only if  $\underline{f} = \underline{g}$ .

You don't have to give formal proofs that this is the case, but I suggest that you do.

- (2) Consider the following  $\mathcal{L}$ -formulas:

$$\begin{aligned} \psi_1 : \forall x \forall y \underline{f}(x) \doteq \underline{g}(y) \quad \psi_2 : \exists x \forall y \underline{f}(x) \doteq \underline{g}(y) \\ \psi_3 : \forall x \exists y \underline{f}(x) \doteq \underline{g}(y) \quad \psi_4 : \exists x \exists y \underline{f}(x) \doteq \underline{g}(y) \end{aligned}$$

Find  $\mathcal{L}$ -structures  $\mathcal{M}_1, \dots, \mathcal{M}_4$  and  $\mathcal{N}_1, \dots, \mathcal{N}_4$  such that  $\mathcal{M}_i \models \psi_i$  and  $\mathcal{N}_i \models \neg \psi_i$  for each  $i \leq 4$ .

**Question 3.** [20 marks] Let  $\mathcal{L}$  be a language with a single unary relation symbol  $\underline{P}$  and a single binary relation symbol  $\underline{R}$ . Consider the following  $\mathcal{L}$ -sentences:

$$\begin{aligned} \phi_1 : \exists x \forall y \exists ((\underline{P}(x) \rightarrow \underline{R}(x, y)) \wedge (\underline{P}(y) \wedge \neg \underline{R}(y, z))) \\ \phi_2 : \exists x \exists z ((\underline{R}(z, x) \rightarrow \underline{R}(x, z)) \rightarrow \forall y \underline{R}(x, y)) \\ \phi_3 : \forall y (\exists z \forall t \underline{R}(t, z) \wedge \forall x (\underline{R}(x, y) \rightarrow \neg \underline{R}(x, y))) \\ \phi_4 : \exists x \forall y ((\underline{P}(y) \rightarrow \underline{R}(y, x)) \wedge (\forall u)((\underline{P}(u) \rightarrow \underline{R}(u, y)) \rightarrow \underline{R}(x, y))) \end{aligned}$$

For each of these formulas determine whether or not it is satisfied in each of the following  $\mathcal{L}$ -structures:

- (1) A structure with universe  $\mathbb{N}$ , where  $\underline{R}$  is interpreted as the usual order  $\leq$  and  $\underline{P}$  is the set of all even numbers.

- (2) A structure with universe  $\mathcal{P}(\mathbb{N})$ , where  $\underline{R}$  is interpreted as  $\subseteq$  and  $\underline{P}$  is the set of all finite subsets of  $\mathbb{N}$ .
- (3) A structure with universe  $\mathbb{R}$  where  $\underline{R}$  is interpreted as the relation  $\{(x, y) \in \mathbb{R}^2 : x = y^2\}$  and  $\underline{P}$  is interpreted as the set of all rational numbers.

**Question 4.** [40 marks] In all the languages considered in this exercise  $\underline{R}$  is a binary relation symbol,  $\star, \oplus$  are binary function symbols, and  $\underline{c}, \underline{d}$  are constant symbols.

- (1) In each of the following five cases you are given a language  $\mathcal{L}_i$  and two structures  $\mathcal{M}_i$  and  $\mathcal{N}_i$ . For each  $i \leq 5$  find an  $\mathcal{L}_i$ -sentence which is true in  $\mathcal{M}_i$  but false in  $\mathcal{N}_i$ .

- |                                                                       |                                                 |                                                   |
|-----------------------------------------------------------------------|-------------------------------------------------|---------------------------------------------------|
| (1) $\mathcal{L}_1 = \{\underline{R}\}$                               | $\mathcal{M}_1 = (\mathbb{N}; \leq)$            | $\mathcal{N}_1 = (\mathbb{Z}; \leq)$              |
| (2) $\mathcal{L}_2 = \{\underline{R}\}$                               | $\mathcal{M}_2 = (\mathbb{Q}; \leq)$            | $\mathcal{N}_2 = (\mathbb{Z}; \leq)$              |
| (3) $\mathcal{L}_3 = \{\star\}$                                       | $\mathcal{M}_3 = (\mathbb{N}; \times)$          | $\mathcal{N}_3 = (\mathcal{P}(\mathbb{N}); \cap)$ |
| (4) $\mathcal{L}_4 = \{\underline{c}, \star\}$                        | $\mathcal{M}_4 = (\mathbb{N}; 1, \times)$       | $\mathcal{N}_4 = (\mathbb{Z}; 1, \times)$         |
| (5) $\mathcal{L}_5 = \{\underline{c}, \underline{d}, \oplus, \star\}$ | $\mathcal{M}_5 = (\mathbb{R}; 0, 1, +, \times)$ | $\mathcal{N}_5 = (\mathbb{Q}; 0, 1, +, \times)$   |

- (2) For each of the sentences below, in the language  $\{\underline{c}, \oplus, \star, \underline{R}\}$  find a structure which satisfies it and a structure which satisfies its negation:

- $\phi_1 : \quad \forall x \forall y \exists z (\neg y \doteq \underline{c} \rightarrow x \oplus (y \star z) \doteq \underline{c})$
- $\phi_2 : \quad \forall u \forall v \forall w \exists x (\neg x \doteq \underline{c} \rightarrow u \oplus (v \oplus x) \oplus (w \otimes x \star (x \star x)) \doteq \underline{c})$
- $\phi_3 : \quad \forall x \forall y \forall z (\underline{R}(x, x) \rightarrow \underline{R}(x \star z, y \star z))$
- $\phi_4 : \quad \forall x \forall y \forall z (\underline{R}(x, x) \wedge (\underline{R}(x, y) \wedge \underline{R}(y, z) \rightarrow \underline{R}(x, z)) \wedge (\underline{R}(x, y) \rightarrow \underline{R}(y, x)))$
- $\phi_5 : \quad \forall x \forall y (\underline{R}(x, y) \rightarrow \neg \underline{R}(y, x)).$